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(54) **METHODS AND CIRCUITRY FOR BUILT-IN SELF-TESTING OF CONTENT ADDRESSABLE MEMORIES**

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(21) Appl. No.: **11/208,134**

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(22) Filed: **Aug. 19, 2005**

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Reissue of:

(64) Patent No.: **6,609,222**  
Issued: **Aug. 19, 2003**  
Appl. No.: **09/654,197**  
Filed: **Sep. 1, 2000**

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U.S. Applications:

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Primary Examiner—Shelly A Chase

(51) **Int. Cl.**  
**G01R 31/28** (2006.01)

(57) **ABSTRACT**

(52) **U.S. Cl.** ..... **714/733**

(58) **Field of Classification Search** ..... **714/733,**  
**714/30, 365, 49**

See application file for complete search history.

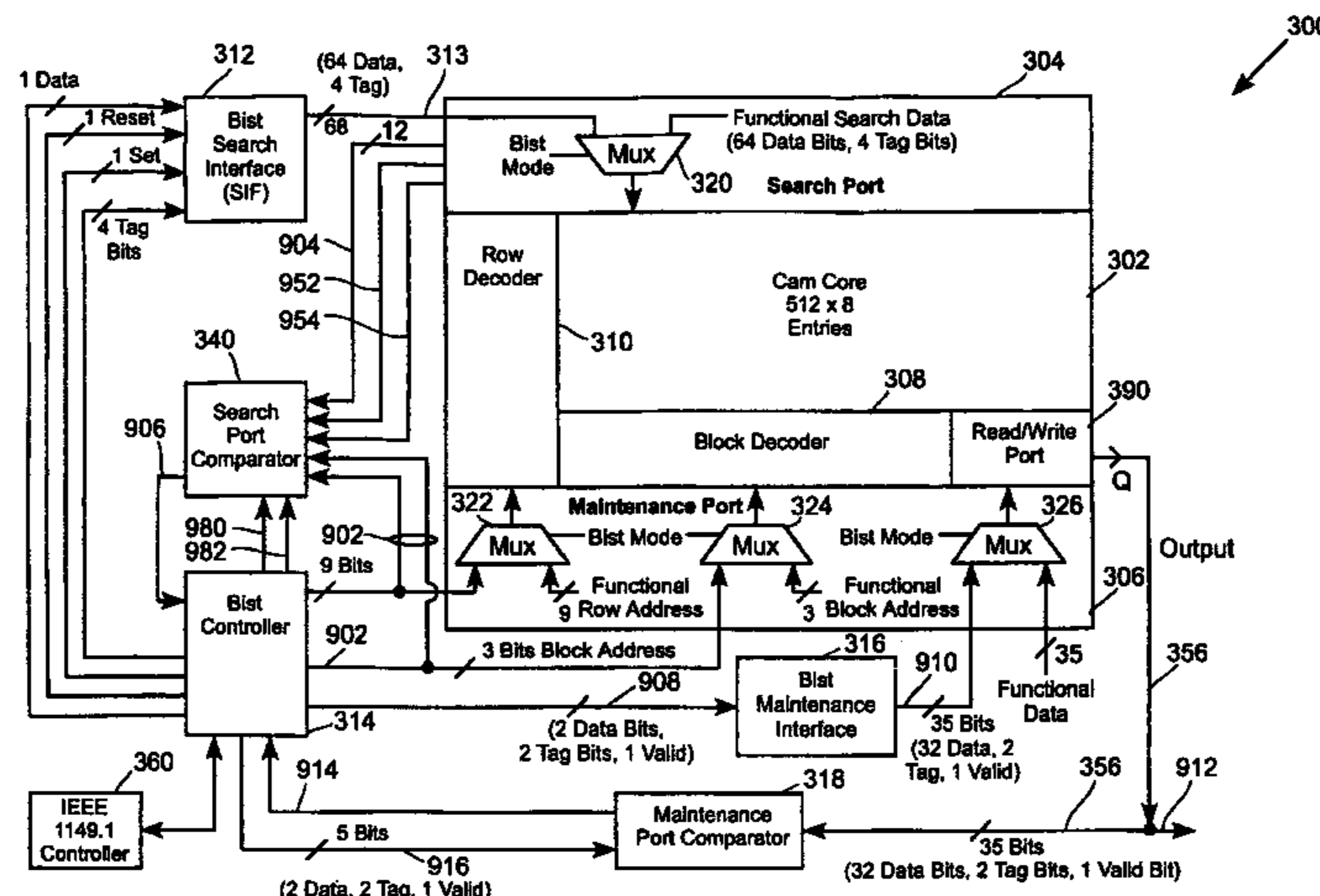
Methods for built-in self-test (BIST) testing and circuitry for testing a content addressable memory (CAM) core are provided. In one example, the BIST circuit includes a search port for enabling searches of the CAM core and a maintenance port for enabling addressing of locations of the CAM core. The maintenance port includes writing logic for writing to locations of the CAM core. The BIST circuit also includes a BIST controller for coordinating BIST testing of the CAM core. The BIST controller is capable of performing a BIST search on the CAM core on every cycle through the search port and performing a BIST write at selected times to the CAM core. Thus, the BIST write is capable of being performed in a same cycle as the BIST search permitting at-speed BIST. The BIST controller, performs BIST testing in a manner that limits the number of rows in the CAM that match at any given cycle, thus allowing a low-power BIST operation. The BIST controller can also be configured to coordinate simultaneous BIST testing of two or more CAM cores.

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**38 Claims, 27 Drawing Sheets**



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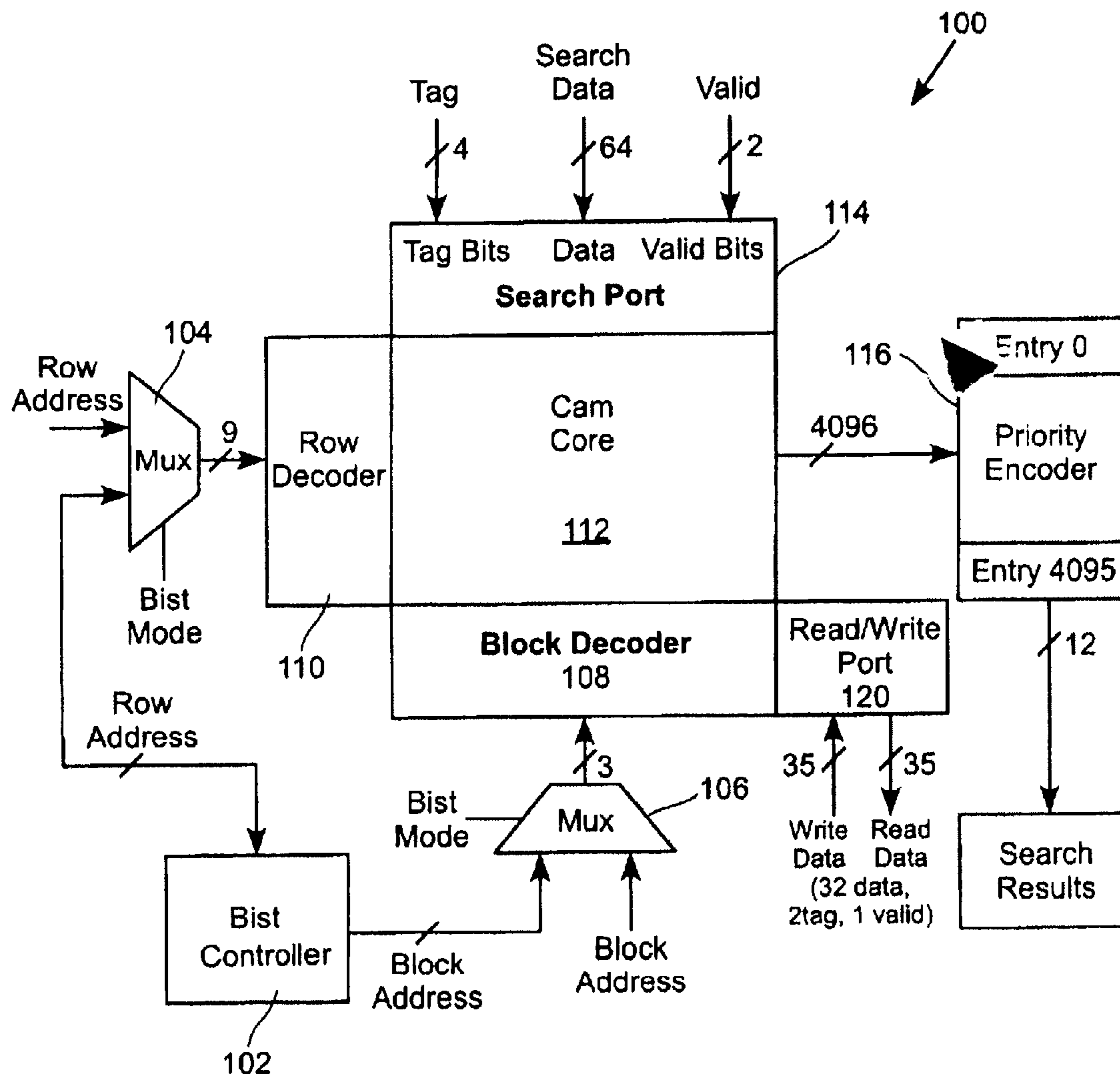


FIG. 1

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Entry 0	Higher Priority	NO
Entry 1		NO
Entry 2		NO
Entry 3		YES
Entry 4		NO
Entry 5		NO
Entry 6		YES
Entry 7	Lower Priority	NO
<b>Priority Encoder</b>		<b>Match?</b>

→ Search Result

FIG. 2A

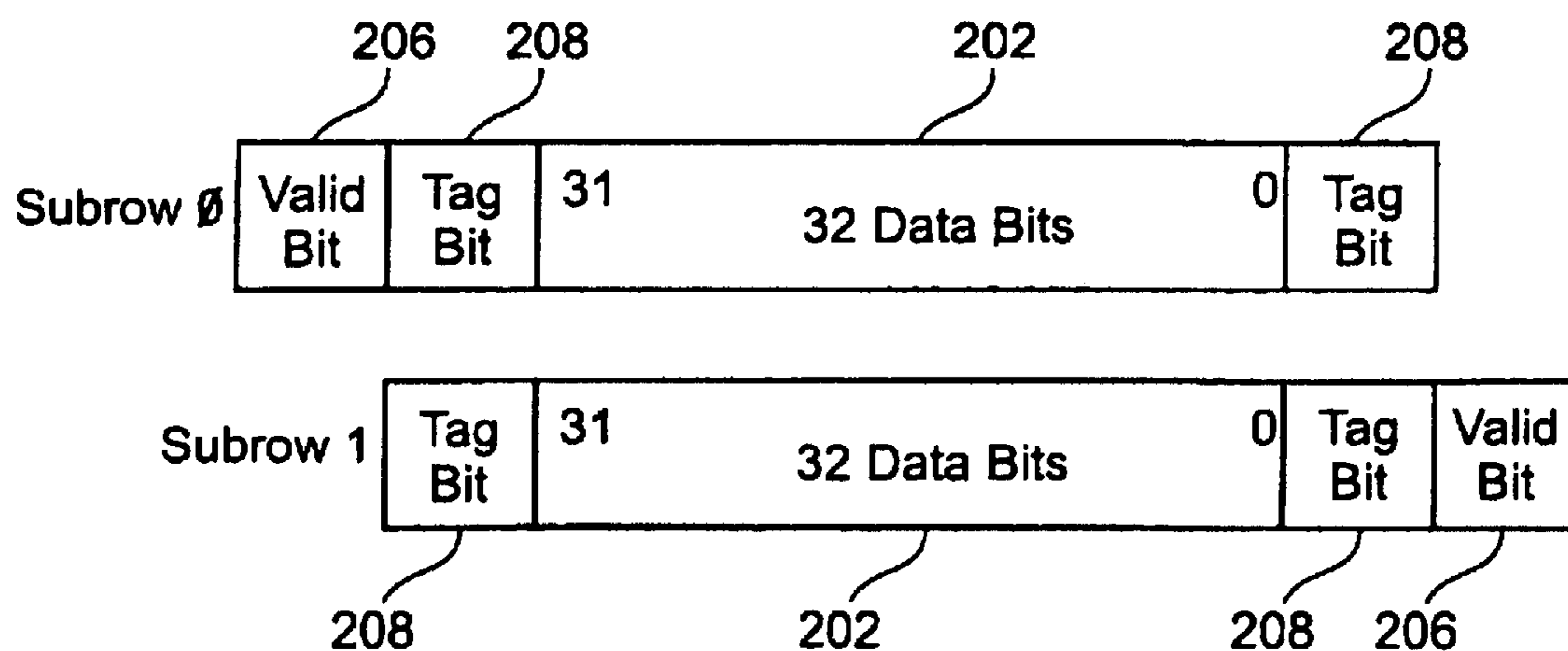


FIG. 2B

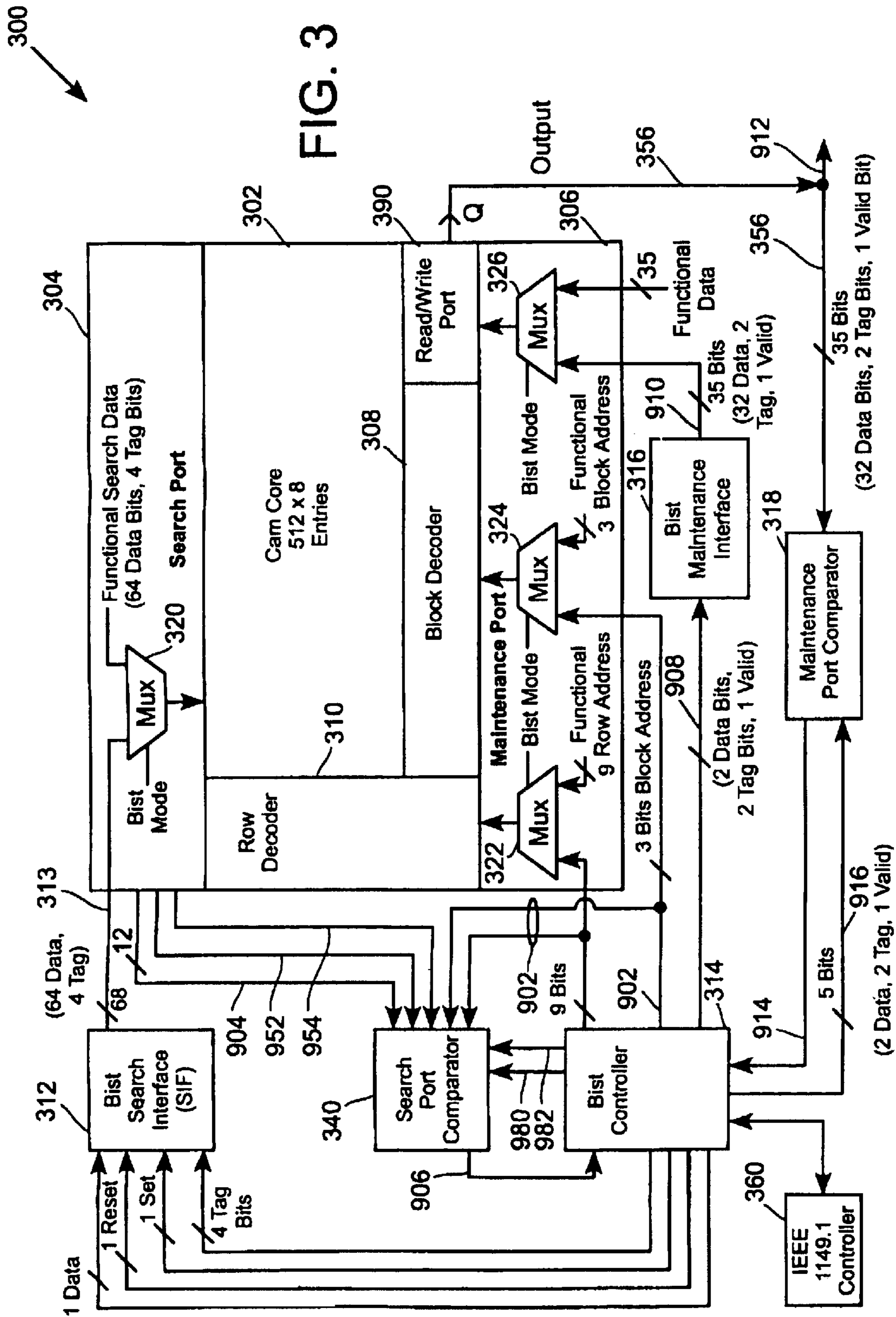


FIG. 3

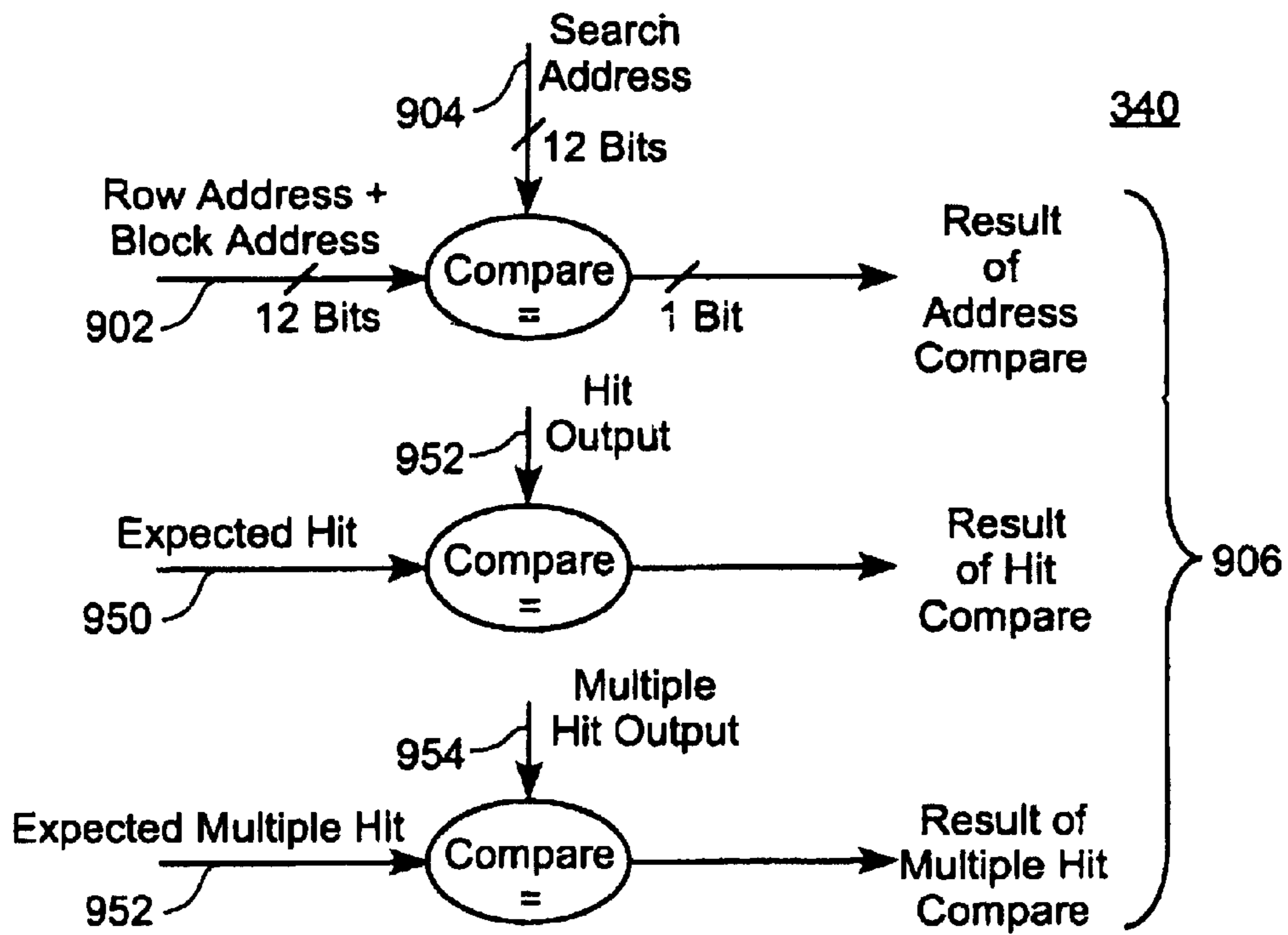


FIG. 4A

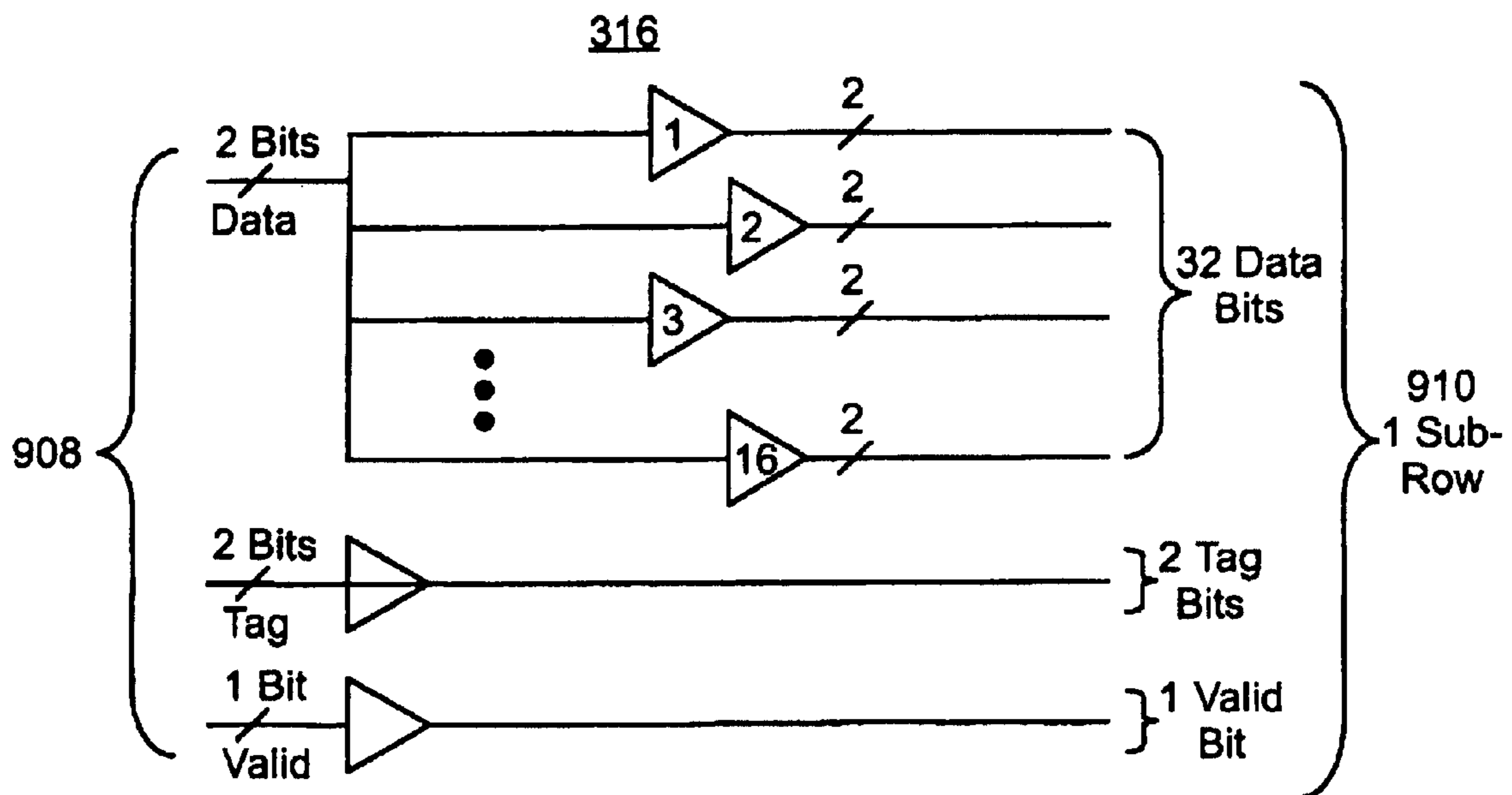


FIG. 4B

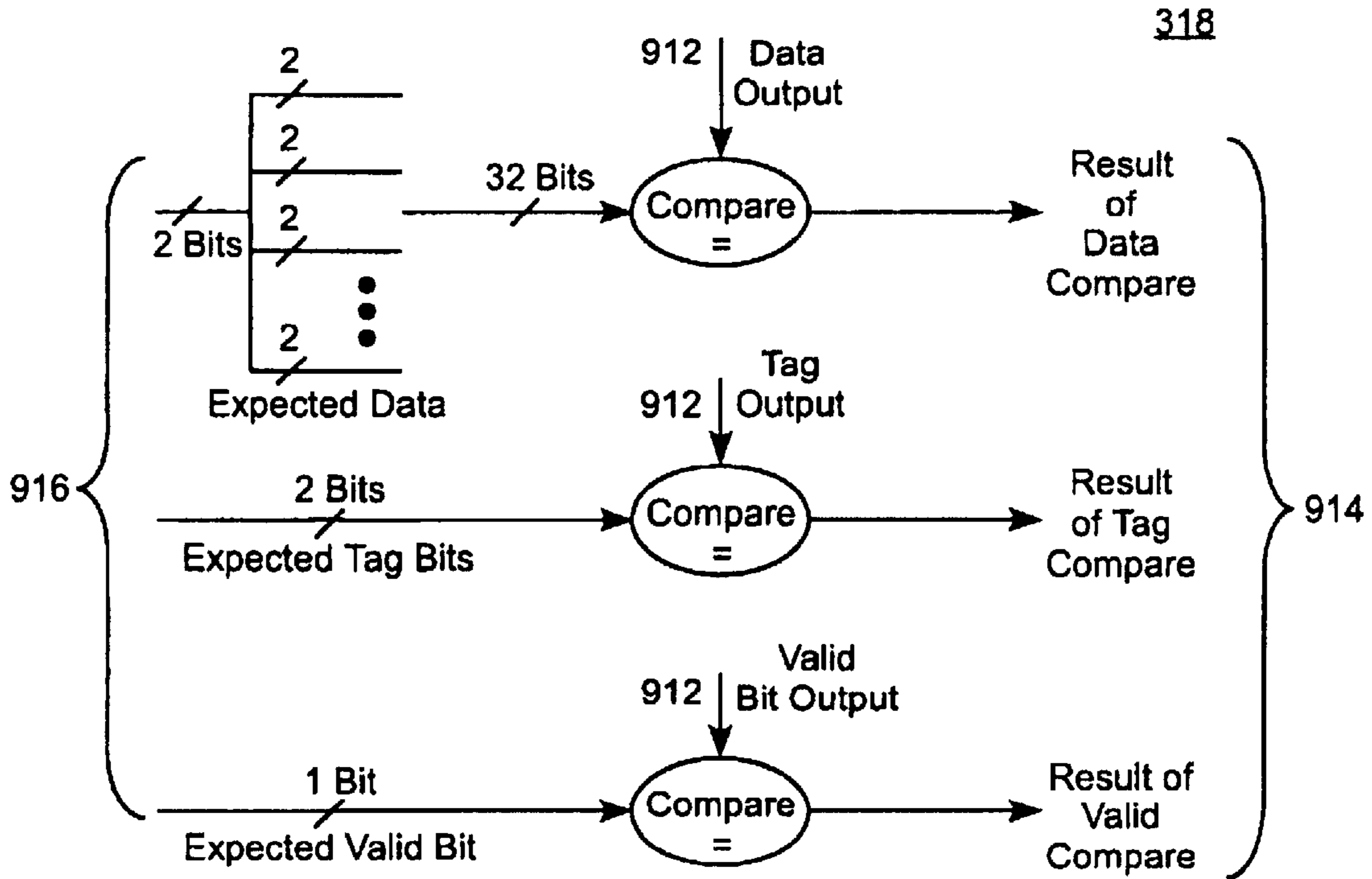


FIG. 4C

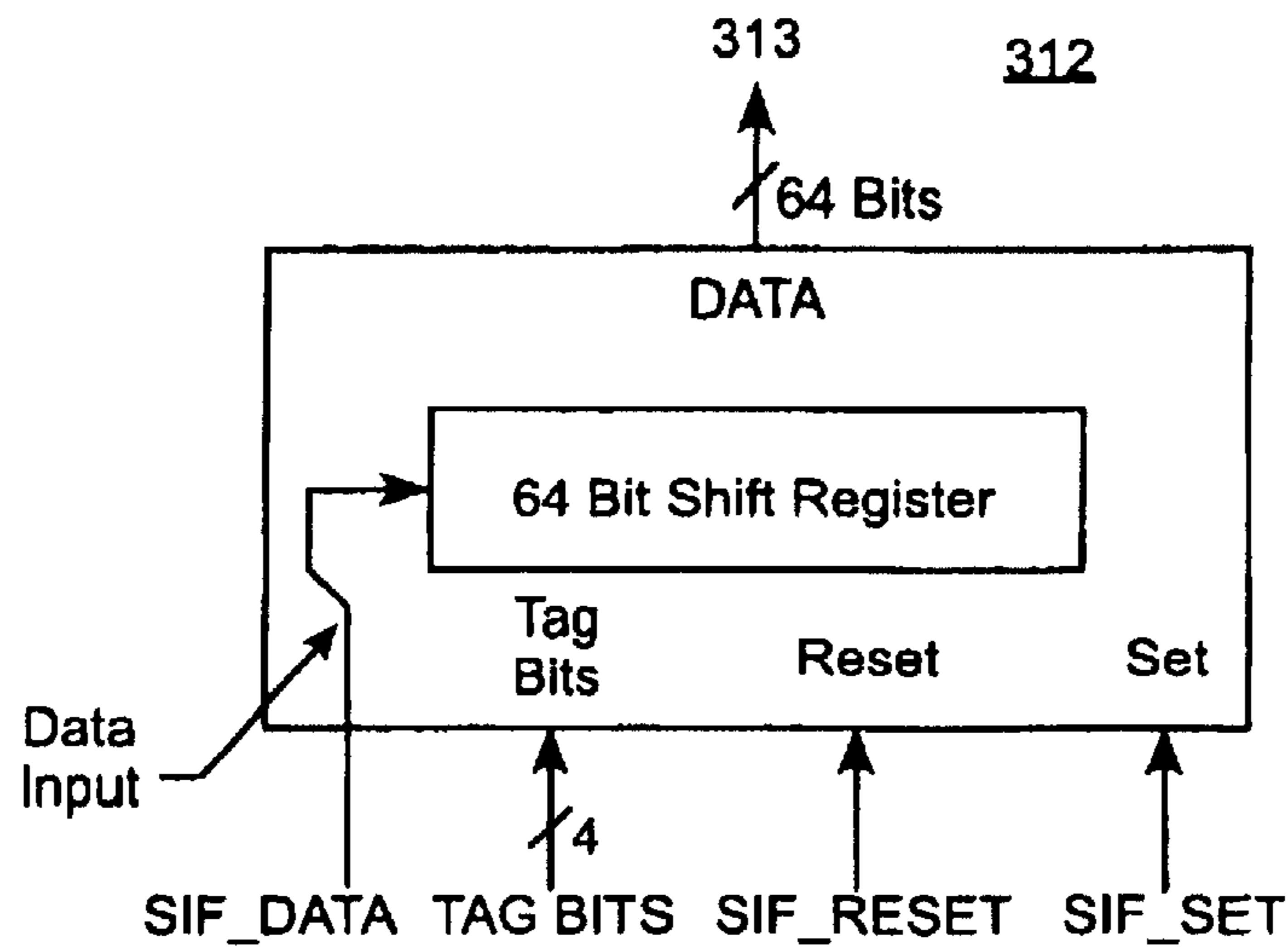


FIG. 4D

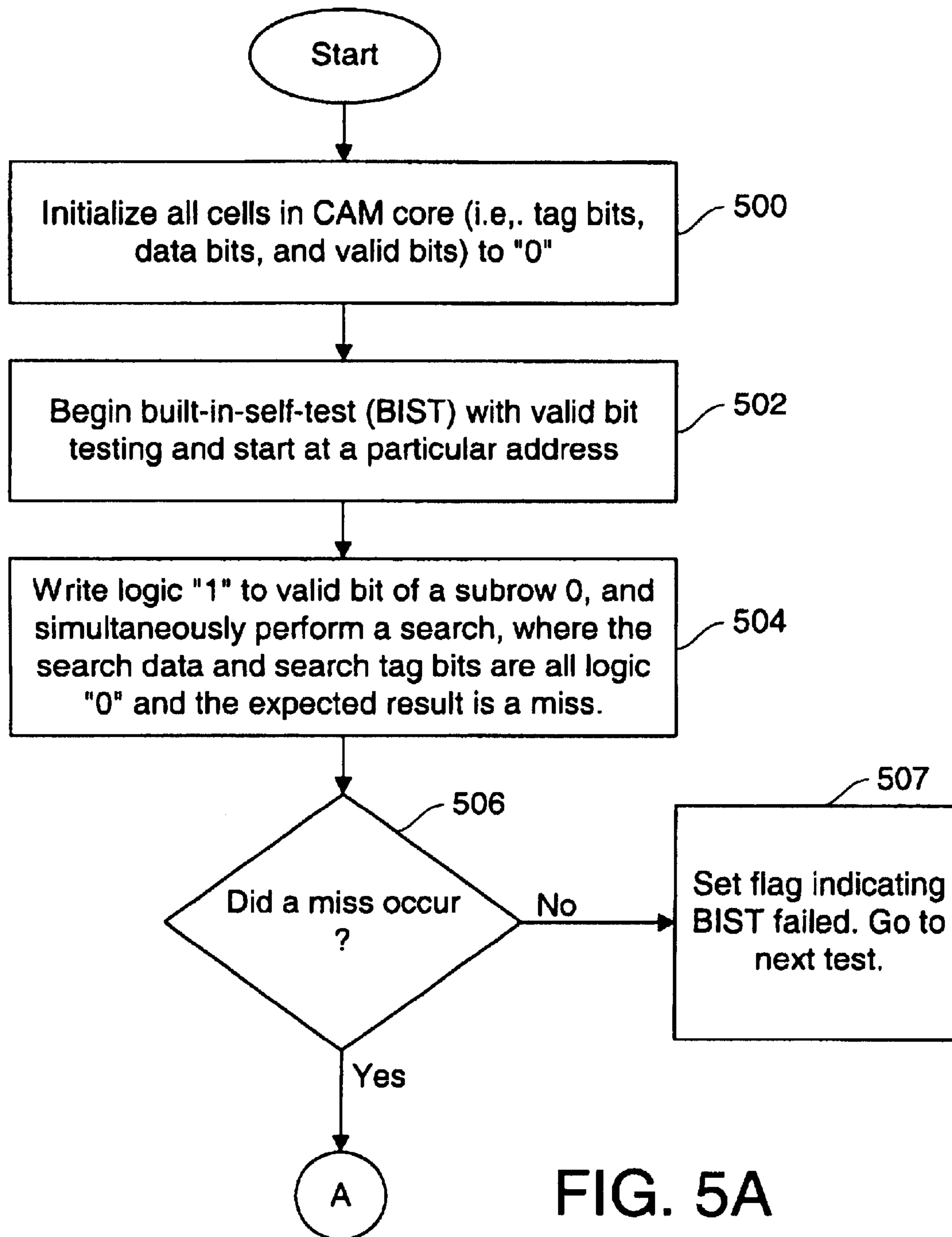


FIG. 5A



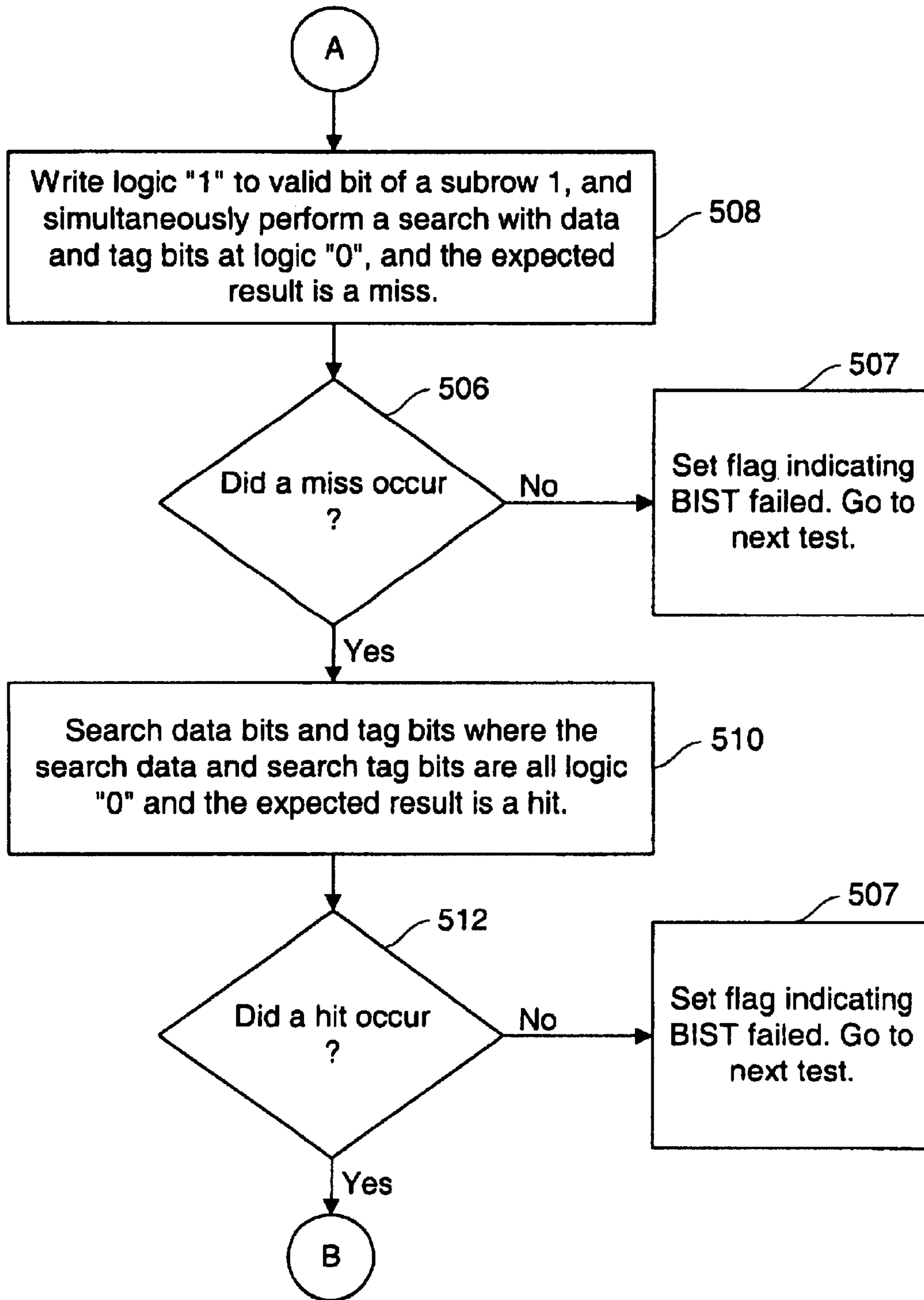


FIG. 5B

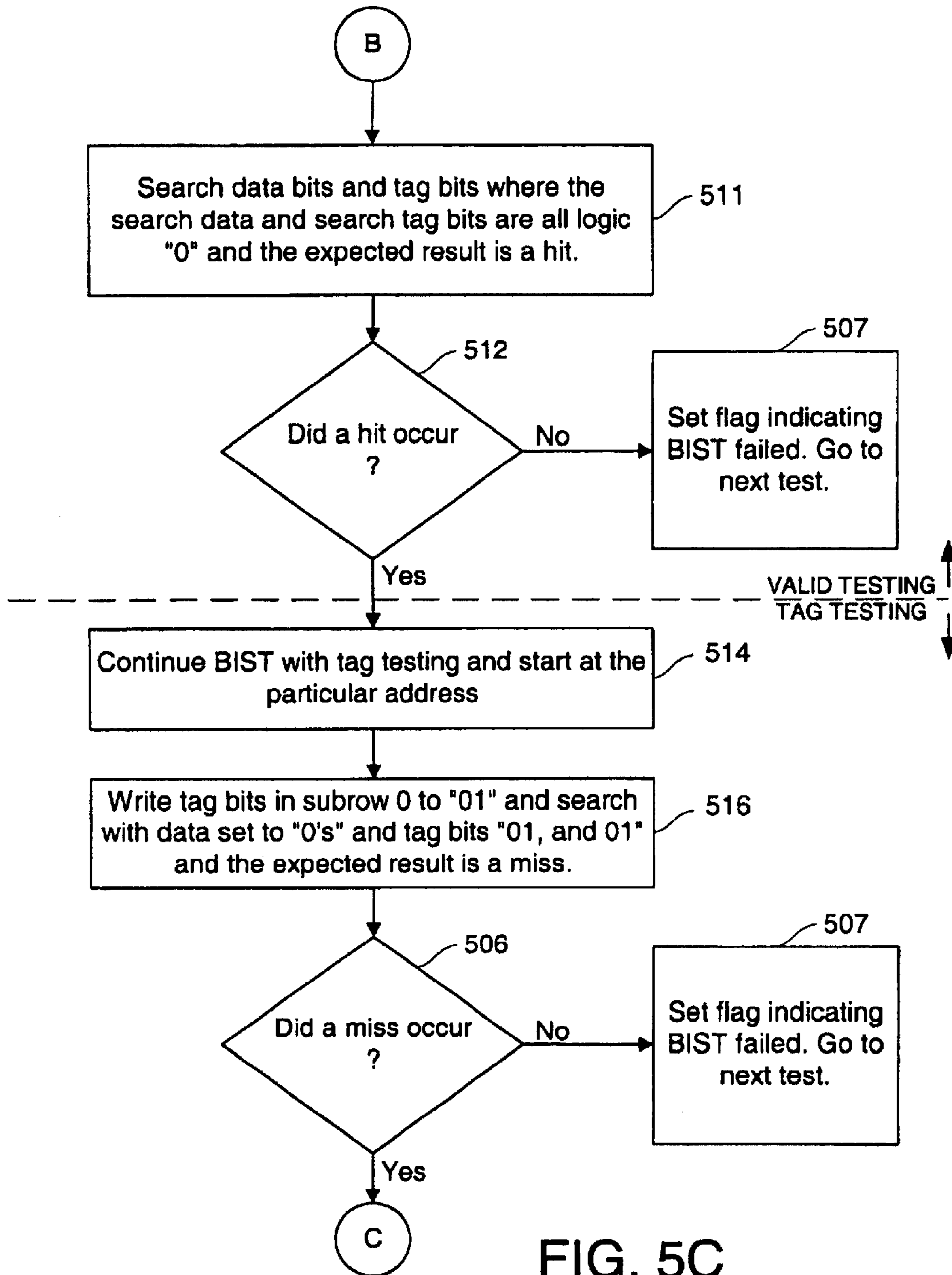


FIG. 5C

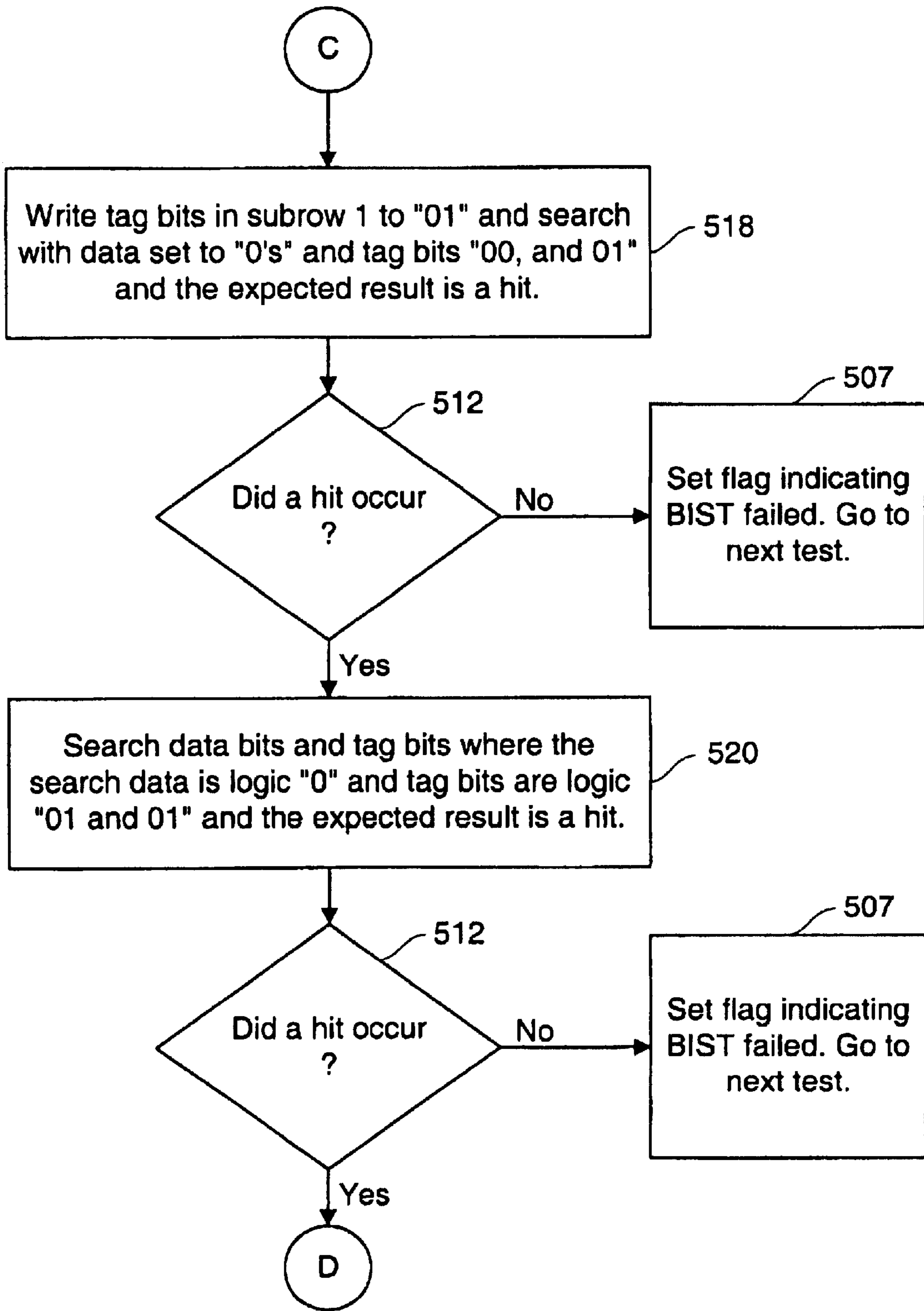


FIG. 5D

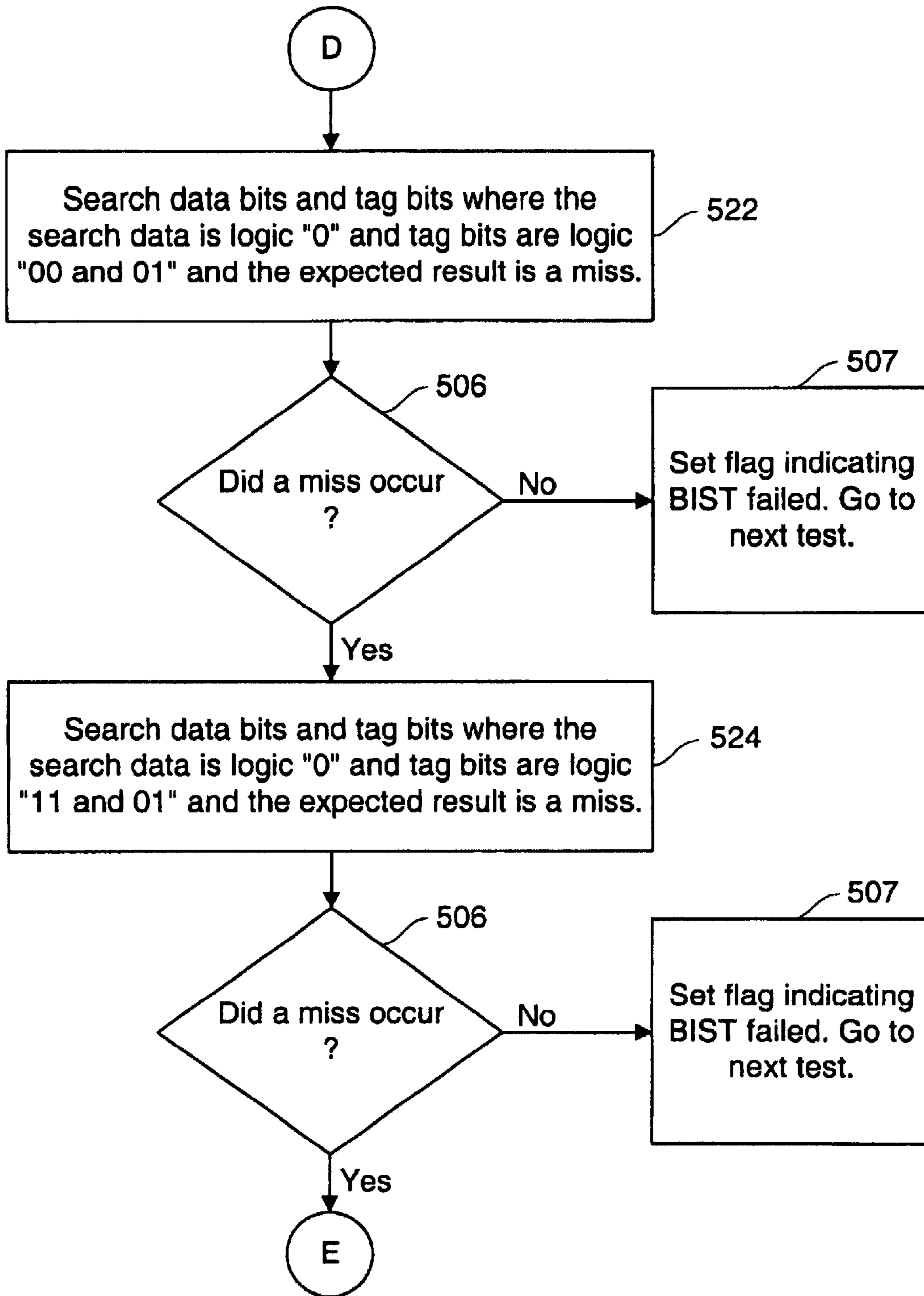


FIG. 5E



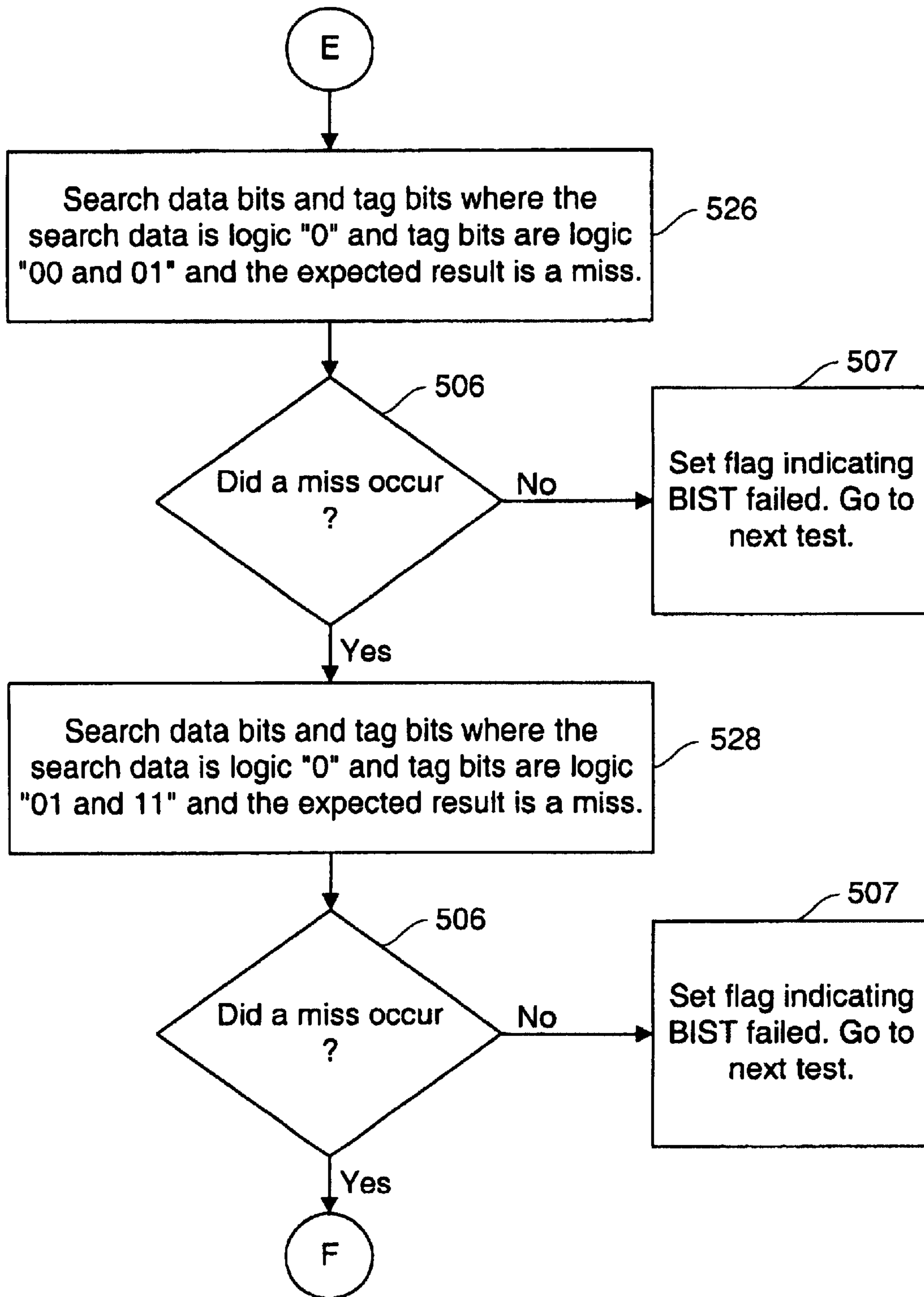


FIG. 5F

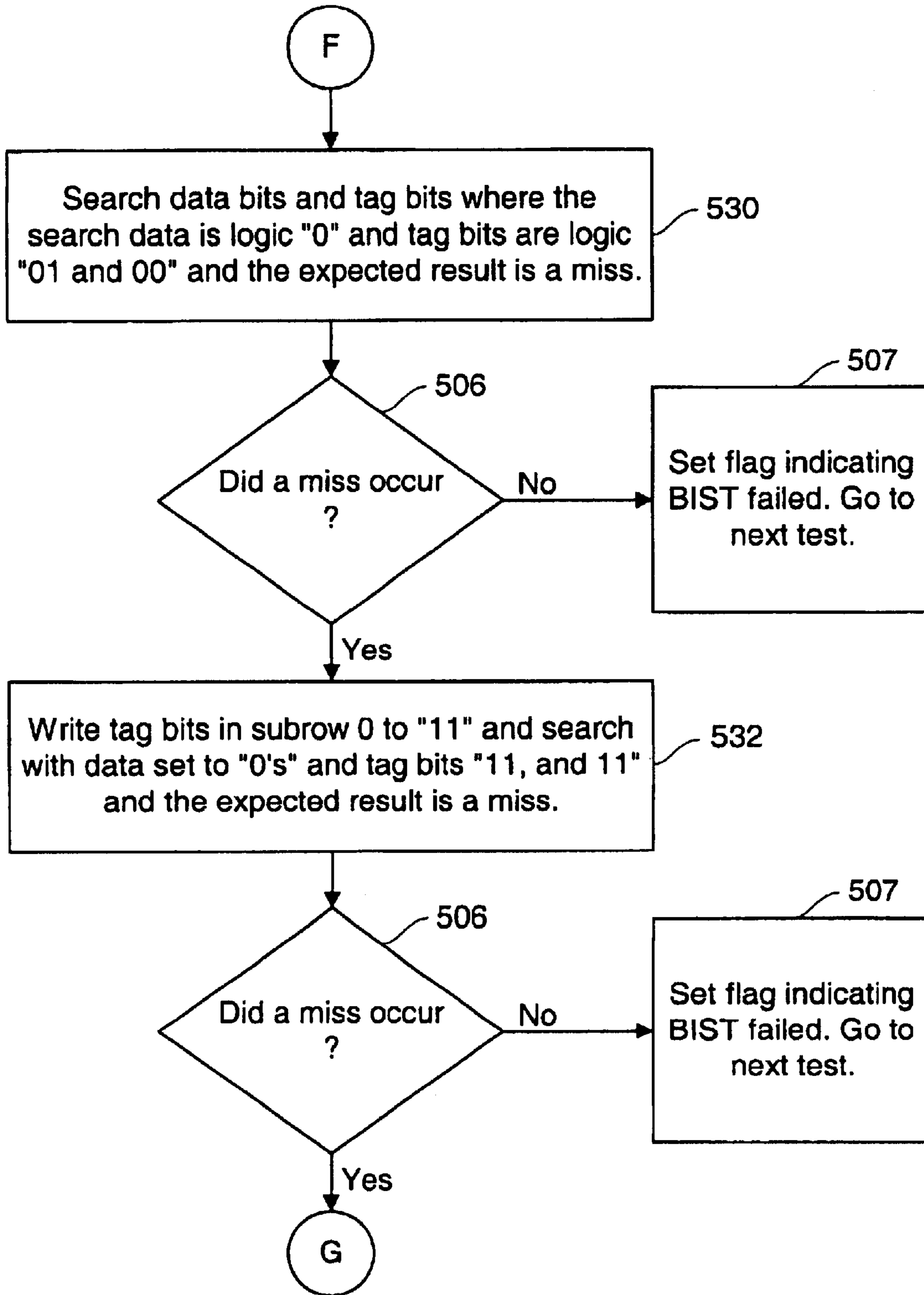


FIG. 5G

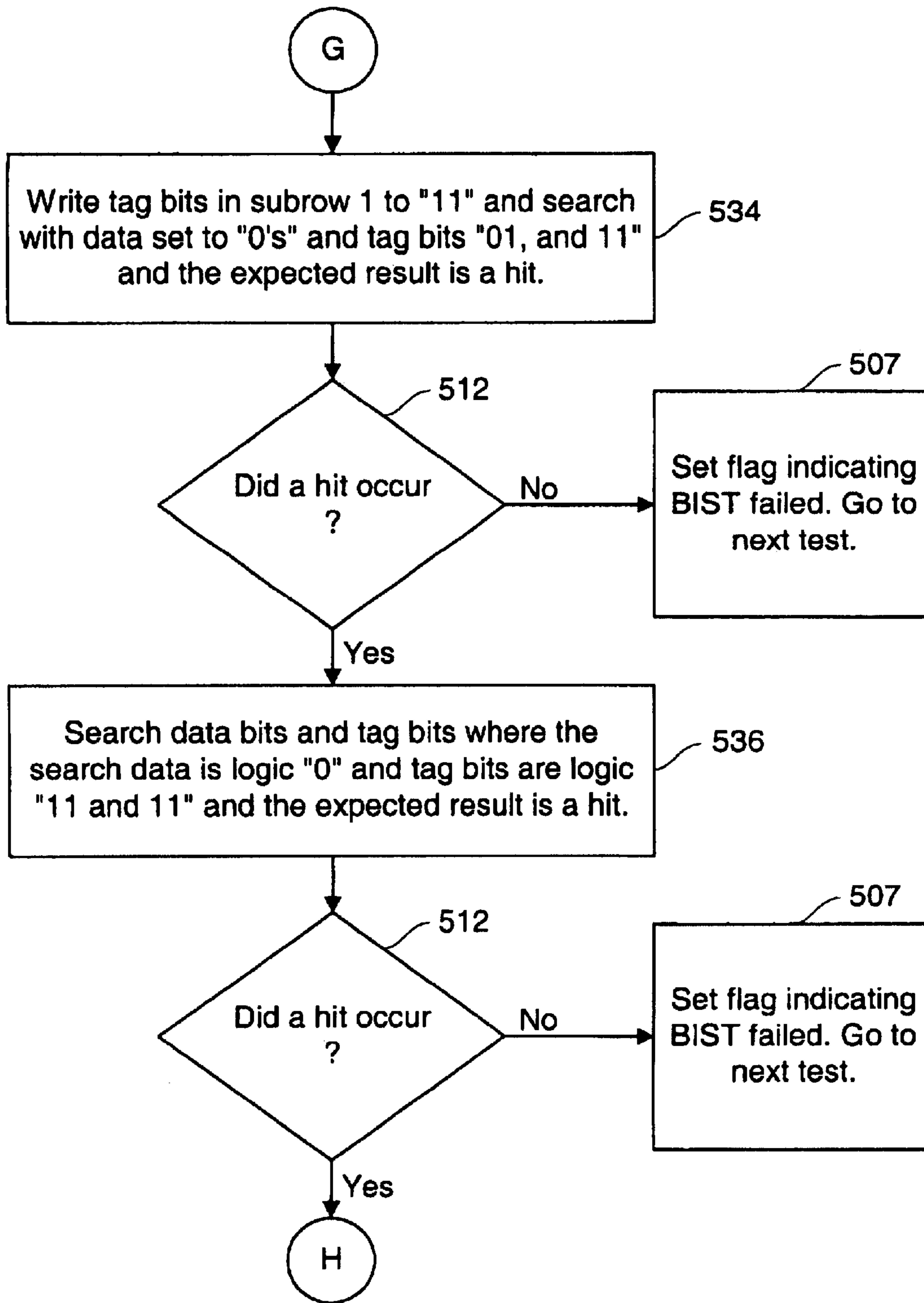


FIG. 5H

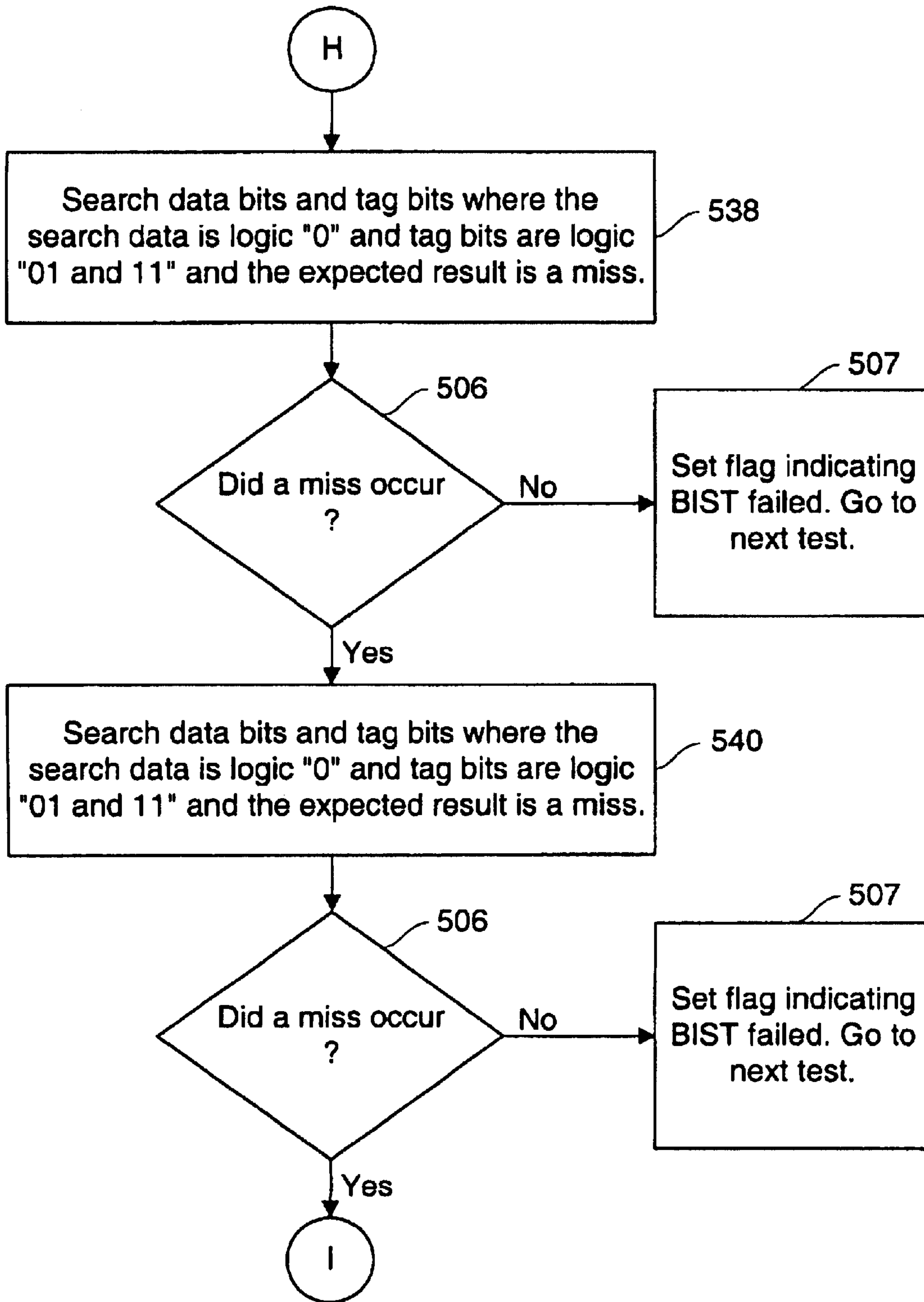


FIG. 5I



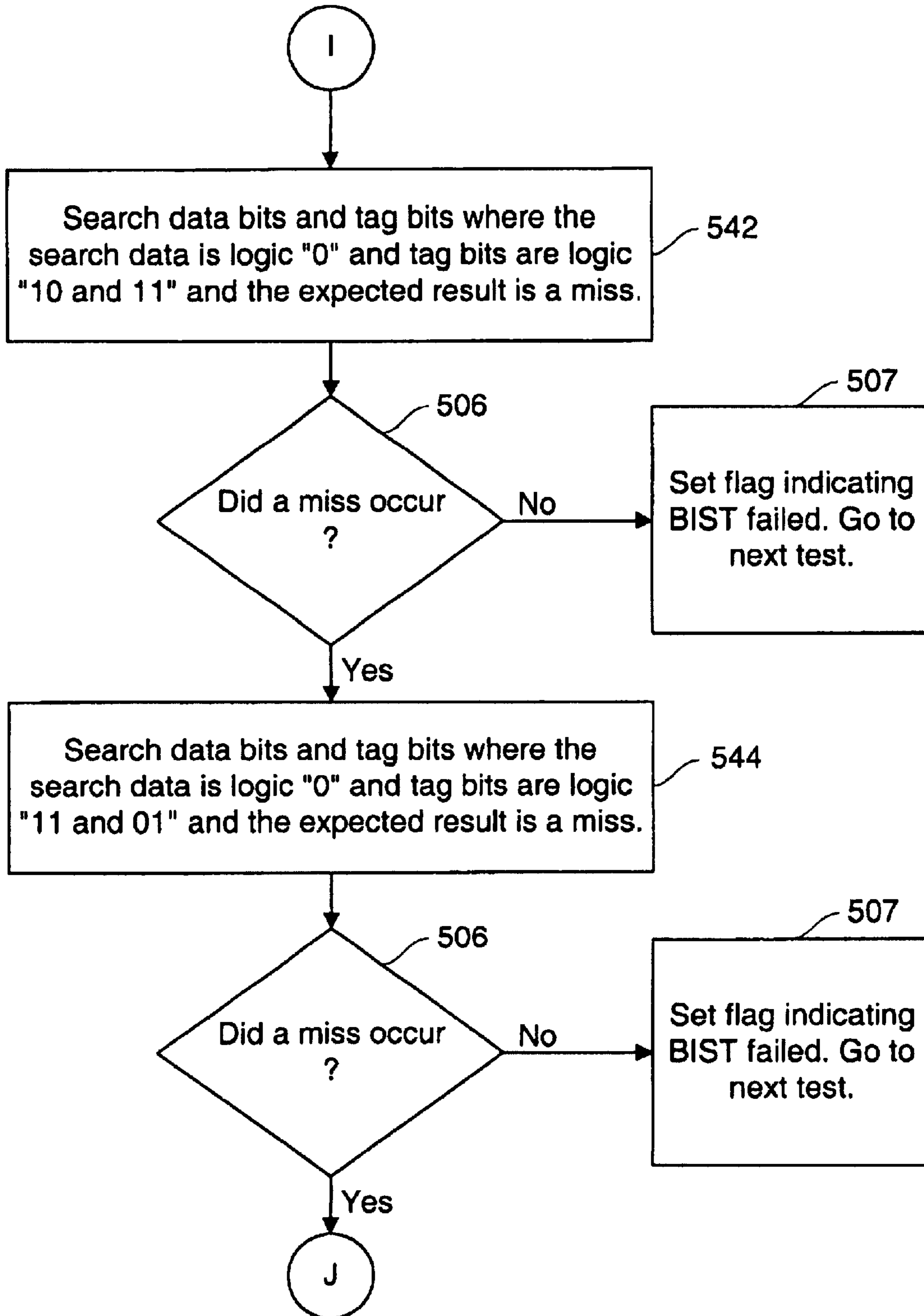


FIG. 5J

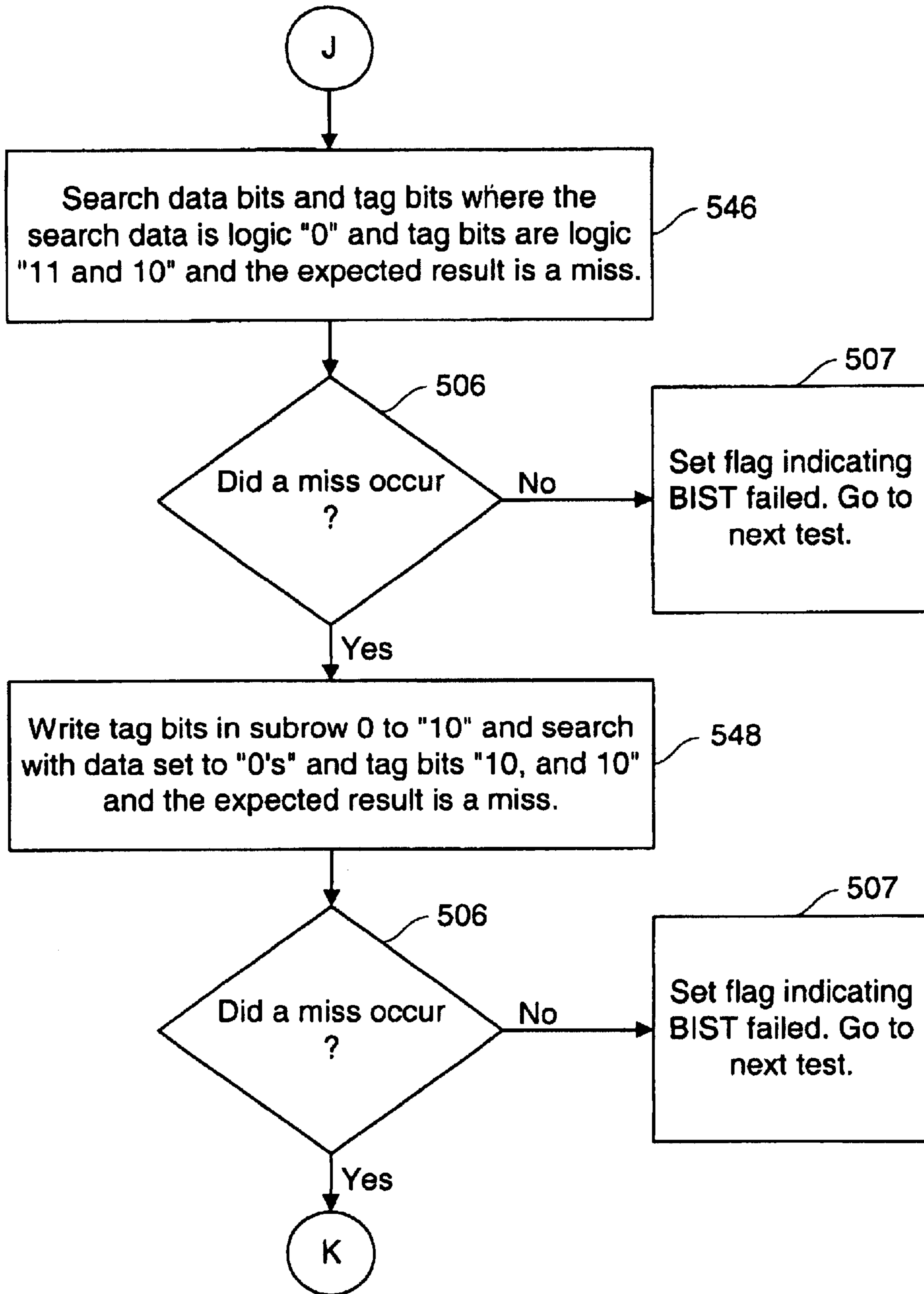


FIG. 5K

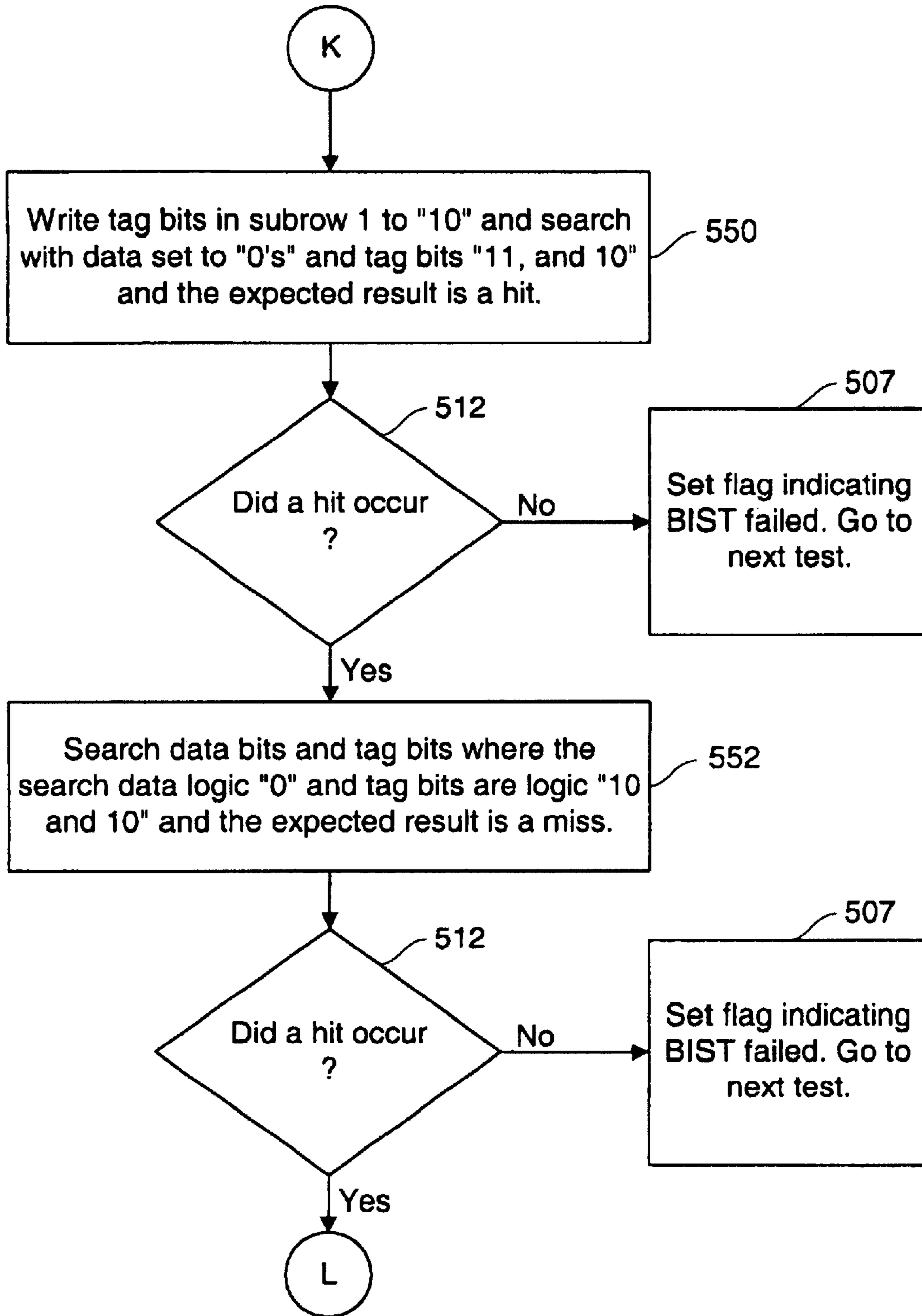


FIG. 5L

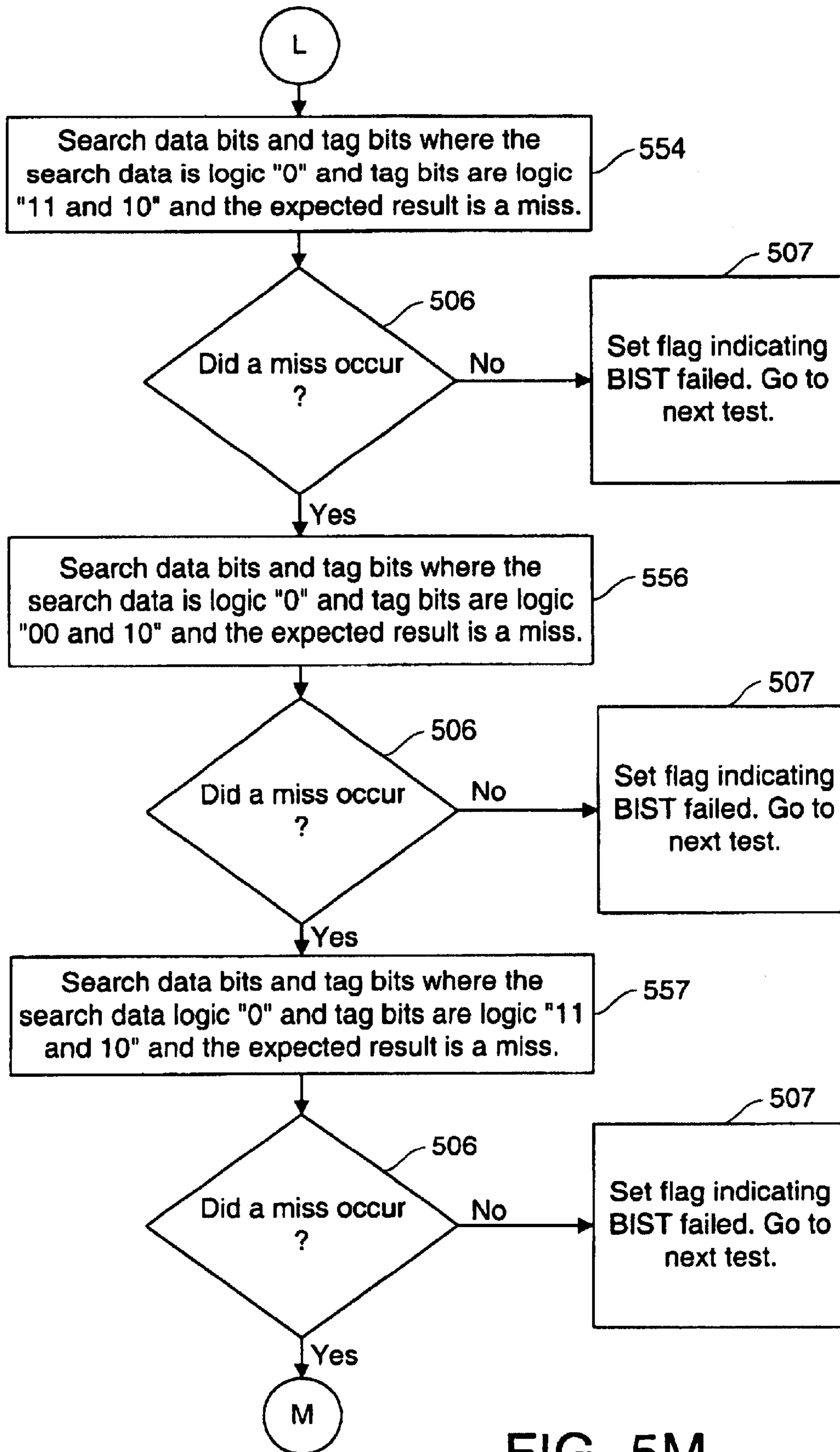


FIG. 5M



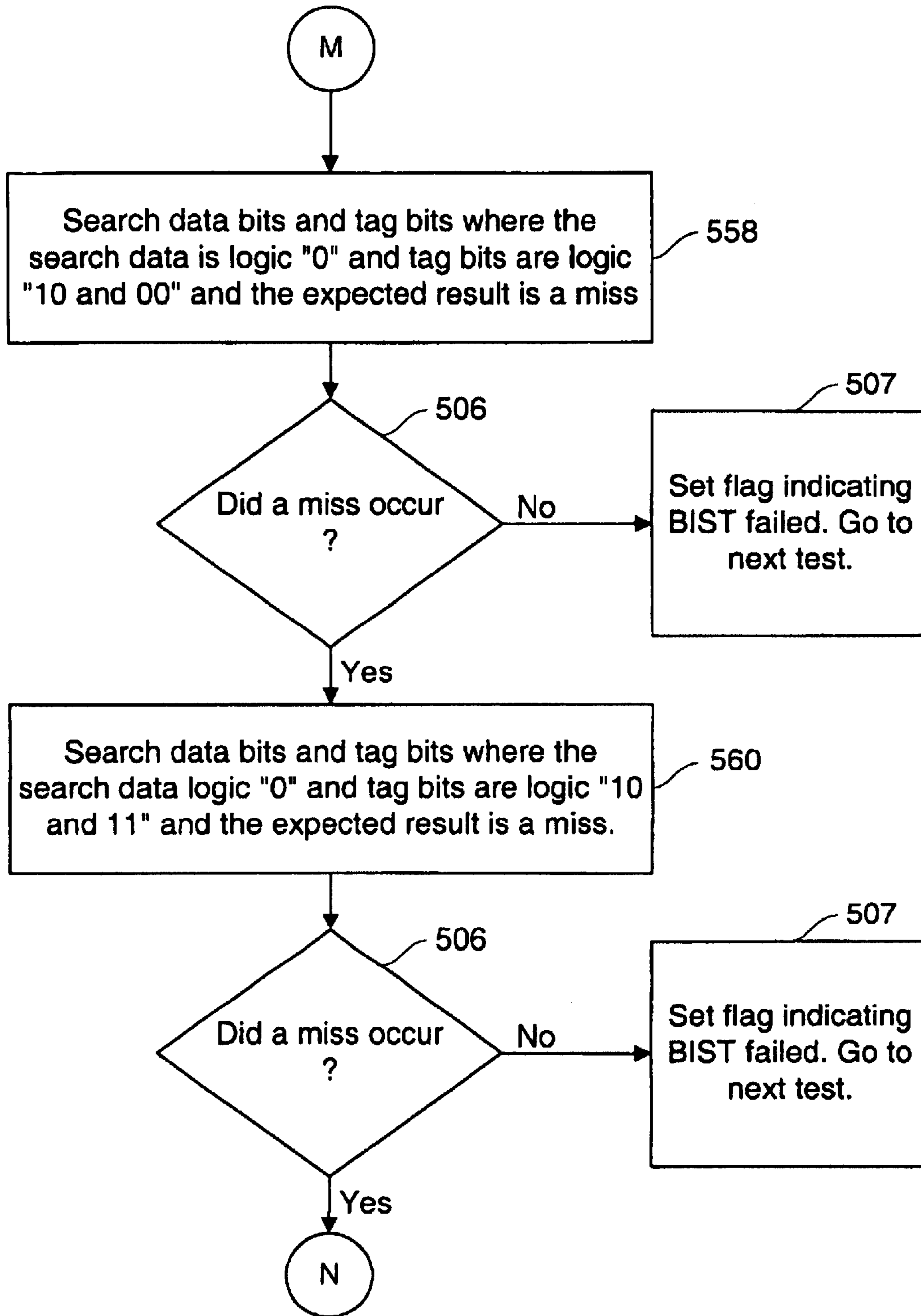


FIG. 5N

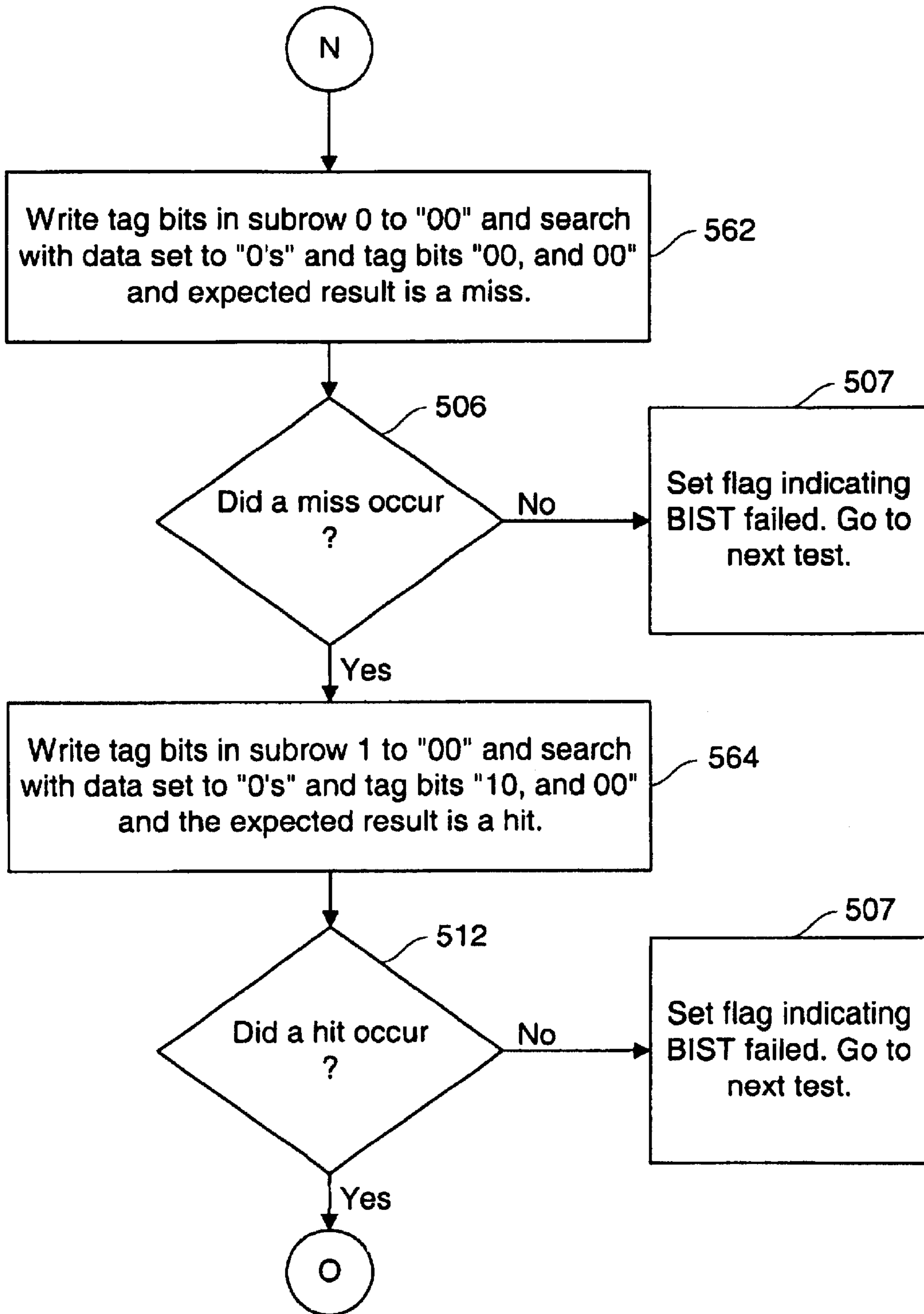


FIG. 50

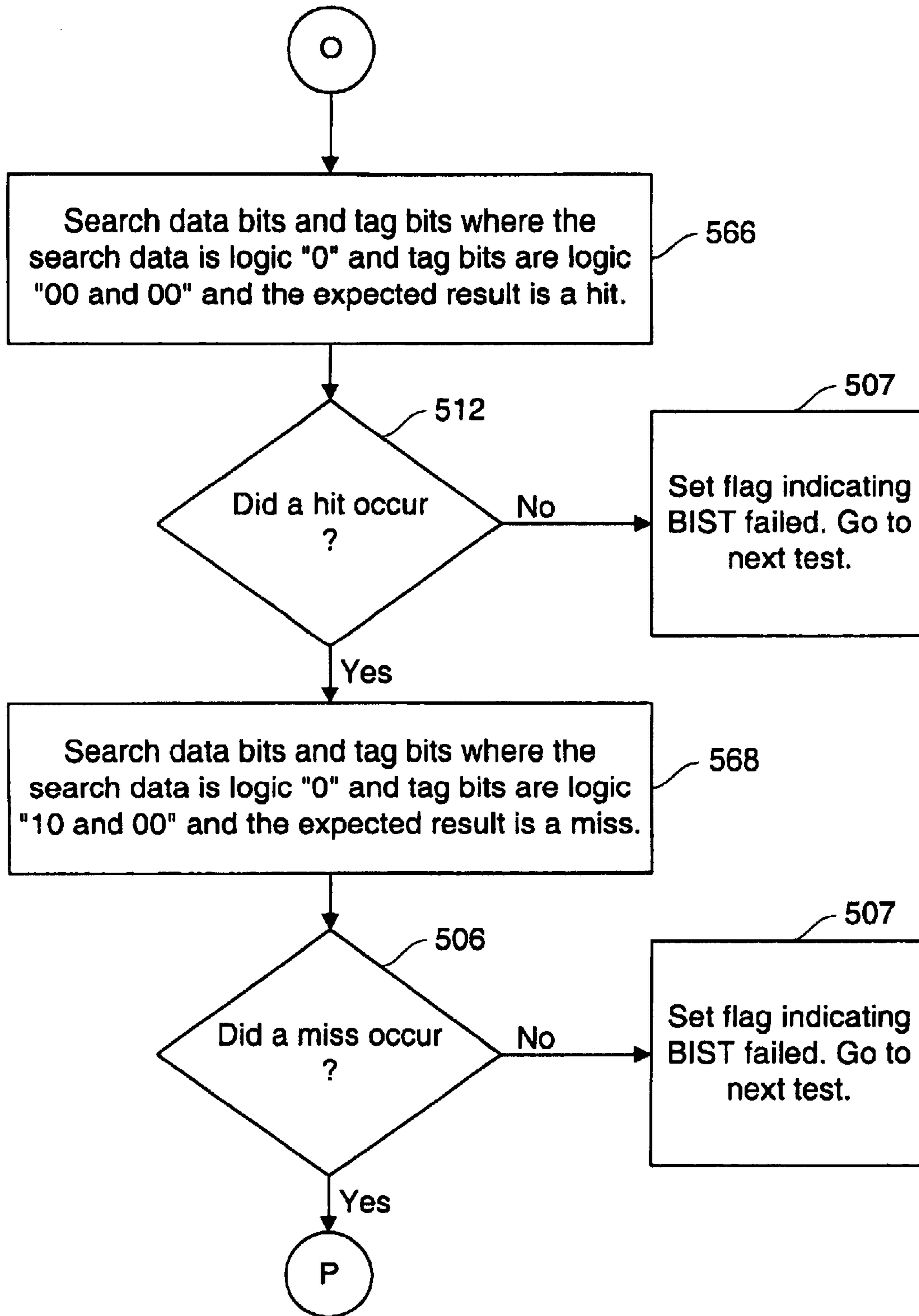


FIG. 5P

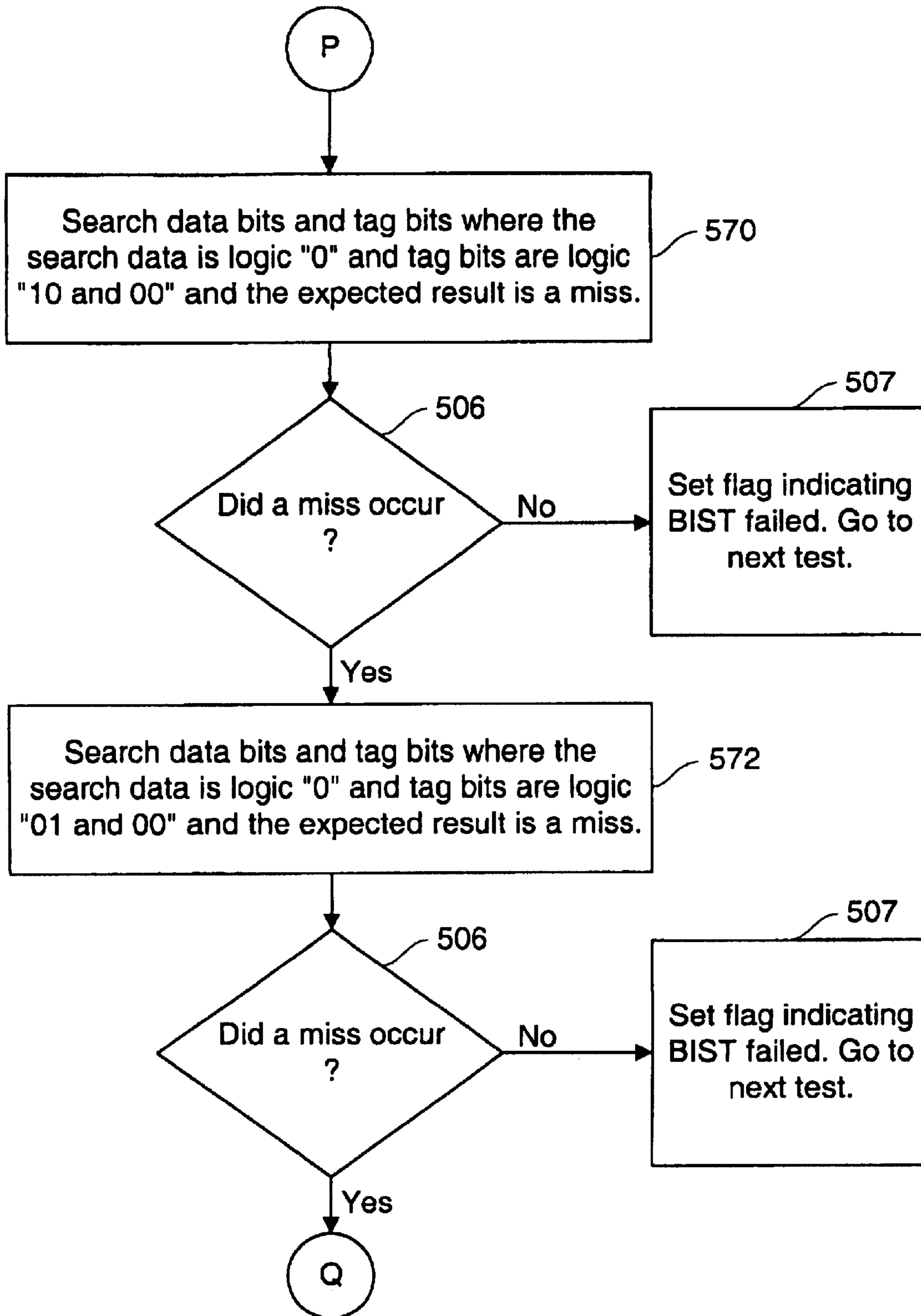


FIG. 5Q



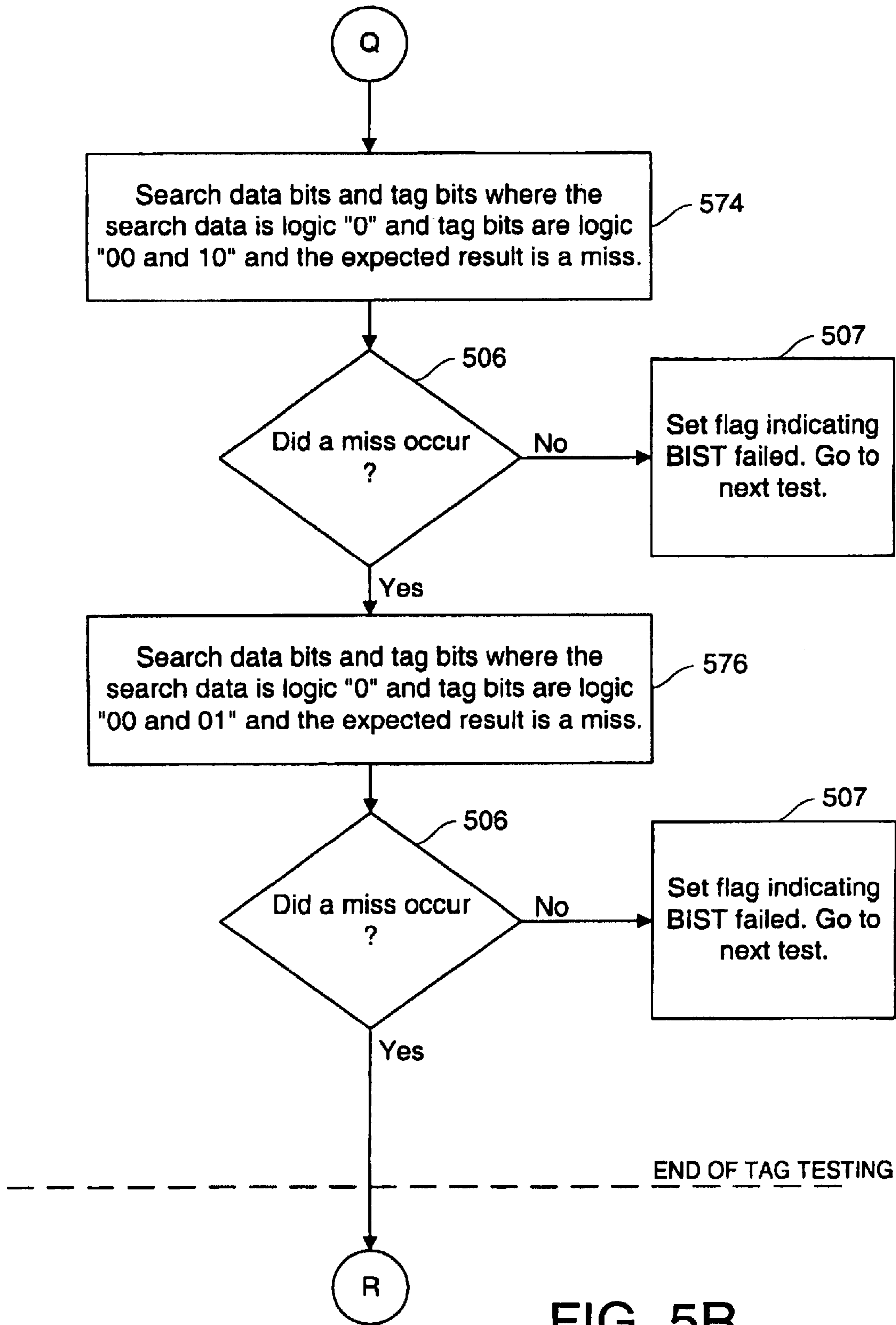


FIG. 5R

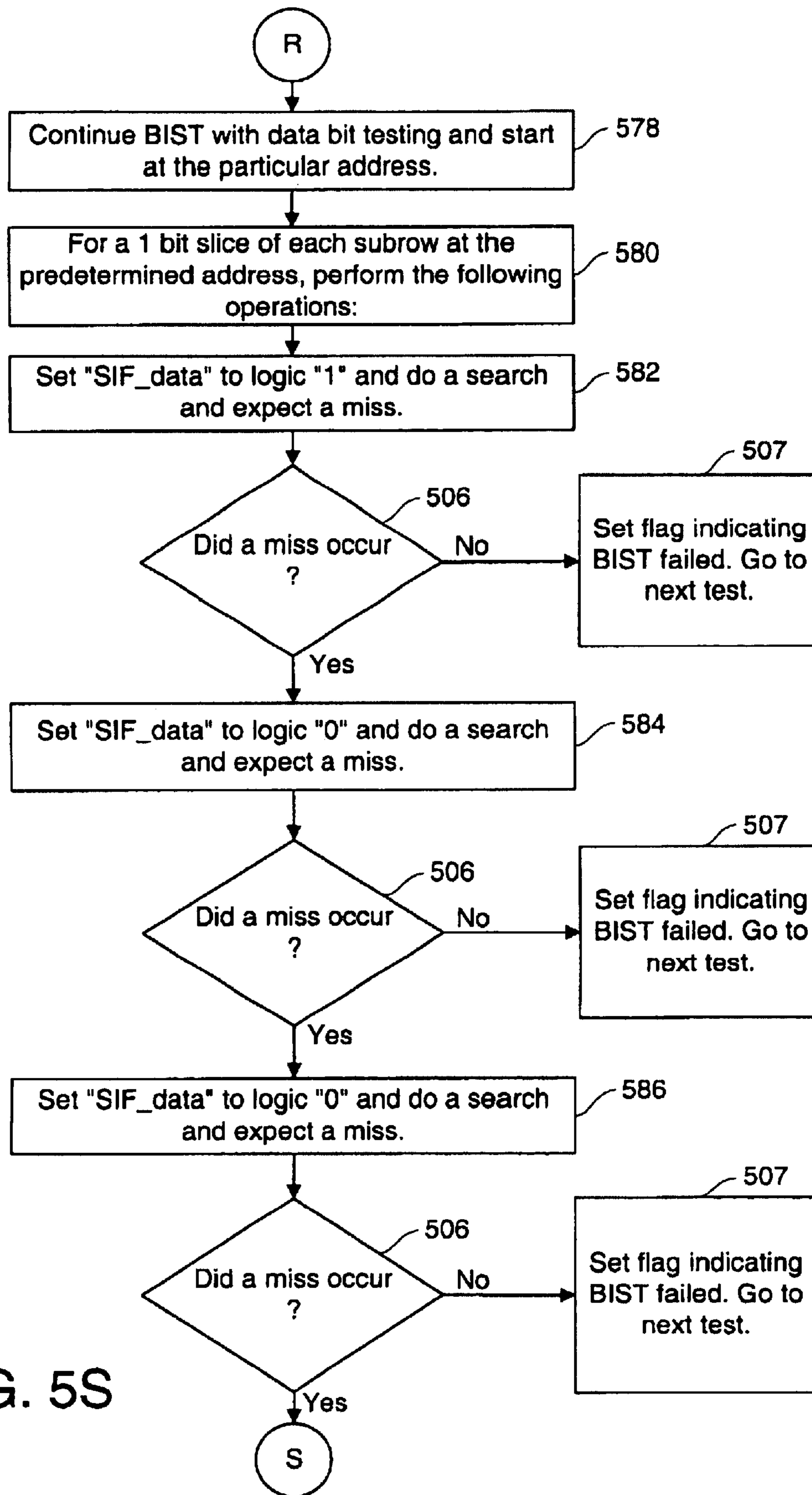


FIG. 5S

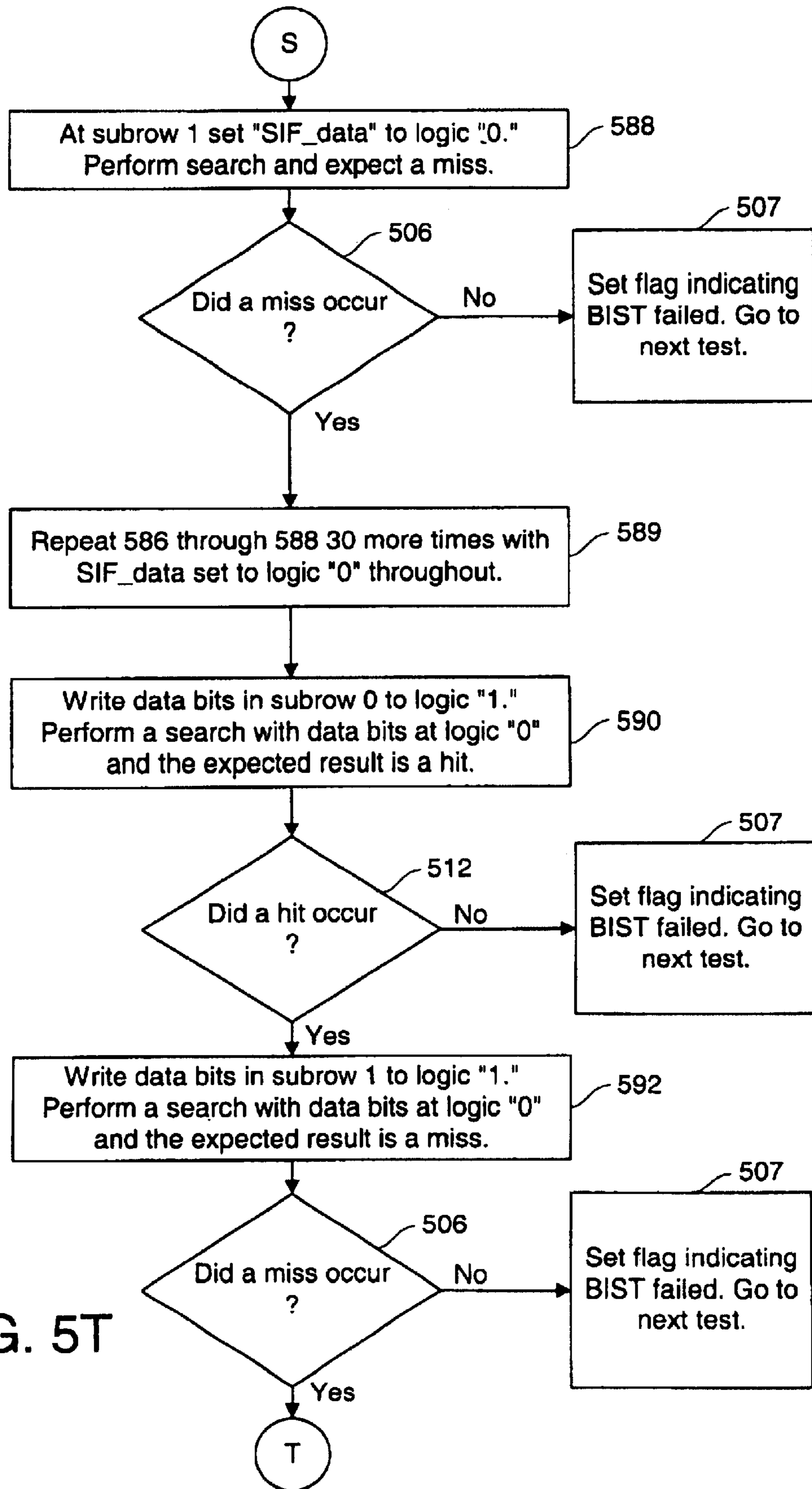


FIG. 5T

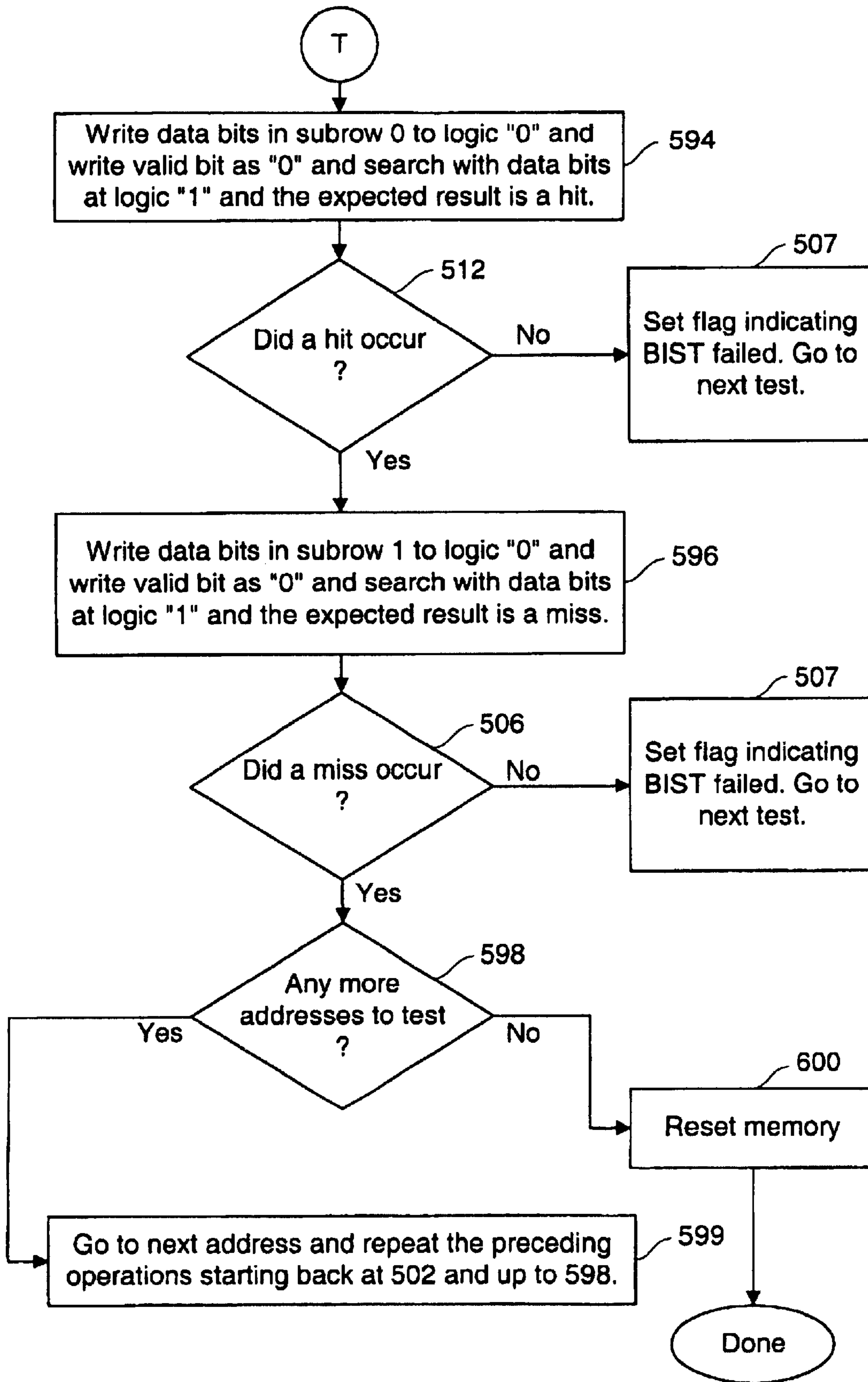


FIG. 5U

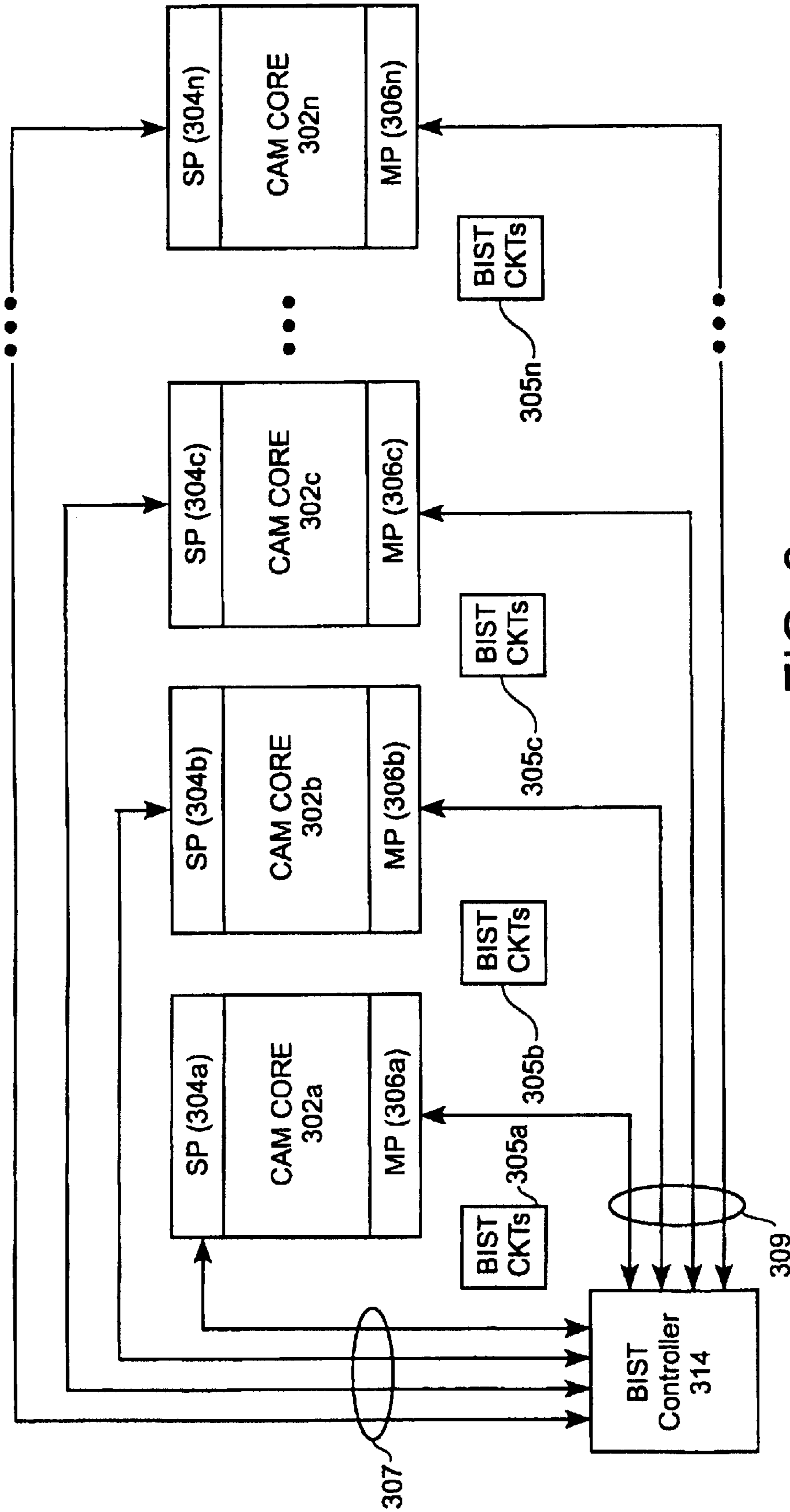


FIG. 6



**METHODS AND CIRCUITRY FOR BUILT-IN  
SELF-TESTING OF CONTENT  
ADDRESSABLE MEMORIES**

**Matter enclosed in heavy brackets [ ] appears in the original patent but forms no part of this reissue specification; matter printed in italics indicates the additions made by reissue.**

CROSS REFERENCE TO RELATED  
APPLICATIONS

**[This application]** *More than one reissue application has been filed for the reissue of U.S. Pat. No. 6,609,222. The reissue applications are application Ser. Nos. 11/208,134 (the present application) filed Aug. 19, 2005, and 11/514,286 (which is a continuation reissue application of the present application) filed Aug. 30, 2006. U.S. Pat. No. 6,609,222 claims priority from U.S. Provisional Patent Application No. 60/153,388 filed Sep. 10, 1999, and entitled "Content Addressable Memory Circuitry." This provisional application is incorporated herein by reference.*

BACKGROUND OF THE INVENTION

1. Field of the Invention

The present invention relates generally to memory circuits, and more particularly to built-in self test (BIST) circuitry for content addressable memory (CAM) circuits.

2. Description of the Related Art

Modern computer systems and computer networks utilize memory devices for storing data and providing fast access to the data stored therein. A content addressable memory (CAM) is a special type of memory device often used for performing fast destination searches for data stored in the CAM. For example, Internet routers often include a CAM for searching the address of specified data. Thus, the use of CAMs allows routers to perform address searches to facilitate more efficient communication between computer systems over computer networks. Besides routers, CAMs are also utilized in other areas where fast-searches are required such as databases, network adapters, image processing, voice recognition applications, etc.

Conventional CAMs typically include a two-dimensional row and column content addressable memory core array of cells. In such an array, each row typically stores an address, pointer, or bit pattern entry. In a read/write configuration, a CAM may perform "read" and "write" operations at specific addresses as is done in conventional random access memories (RAMs). In another configuration, CAMs unlike RAMs, can perform data "search" operations that simultaneously compare a bit pattern of data against an entire list (i.e., column) of prestored entries (i.e., rows).

As with most memory devices, the conventional CAMs may fail to operate properly due to any number of faults. For example, bit cells of the CAM array may become stuck at some value. Another type of fault occurs when bit cells become electrically coupled during write operations. In such cases, when a specific value is written to one bit cell, the other bit cell in the same row or a different row acquires the same value. Another type of fault is a transition fault, which prevents a bit cell from transitioning from one state into another state. Yet another example is when an address decoding fault occurs, thus causing a wrong cell or cells to be addressed, thereby resulting in the access of data in the wrong location.

Still another type of fault occurs when electrical shorts or other defects cause circuit or memory elements to be stuck-

at some value, thereby preventing proper operation of the CAM. Such faults may occur in a CAM itself or its peripheral circuitry. Another fault occurs when CAM circuits operate at a slower speed than expected due to variations in such factors as IC fabrication process, temperature, or voltage. As can be appreciated, such dynamic faults cause undesired delays that can degrade performance of CAM circuits.

To guard against such faults, one approach is to implement a built-in self-test (BIST) technique for detecting some of these types of faults. For example, U.S. Pat. No. 5,107,501 by Yervant Zorian, which is incorporated herein by reference, discloses a BIST technique for testing the data bits in a conventional binary CAM using write and read-match operations. While such method detects faults such as electrical short and open faults and stuck-at faults, it is not well suited for detecting faults in ternary CAMs. Specifically, the bit cells in ternary CAMs are characterized not only by binary states of "0" and "1," but also by a "don't care" state, which denotes that the state of the associated bit cell is not relevant to a particular CAM search. Consequently, not all functional aspects of a ternary CAM can be tested using the teachings of Zorian.

Also, Zorian can only perform writes to set up searches when search operations are not being performed. Consequently, there will be cases in which BIST searching cannot be performed for several cycles until write operations are complete. This, of course, slows down BIST testing and does not allow the BIST testing to be at the functional speed. Furthermore, the BIST technique described by Zorian tests the CAM rows simultaneously. This technique triggers matches in all rows of the CAM and results in very high power consumption. In essence, the Zorian design is impractical for today's larger CAM implementations. Practical CAM devices with large storage capacities are designed to be used only in a manner that does not cause matches in more than some limited number of rows simultaneously.

In view of the foregoing, what is needed is a method for BIST testing of CAM circuitry and associated BIST testing circuitry.

SUMMARY OF THE INVENTION

The present invention fills this need by providing an architecture for performing built-in self-test (BIST) for CAMS. It should be appreciated that the present invention can be implemented in numerous ways, including as a process, an apparatus, a system, a device, or a method. Several embodiments of the present invention are described below.

In one embodiment, a built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core is disclosed. The BIST circuit includes a search port for enabling searches of the CAM core and a maintenance port for enabling addressing of locations of the CAM core. The maintenance port includes writing logic for writing to addresses of the CAM core. The BIST circuit also includes a BIST controller for coordinating BIST testing of the CAM core. The BIST controller is capable of performing a BIST search on the CAM core on every cycle through the search port and performing a BIST write at selected times to the CAM core. Thus, the BIST write is capable of being performed in a same cycle as the BIST search.

In another embodiment, a built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core is disclosed. The BIST circuit includes a search port for enabling searches of the CAM core and a maintenance port for enabling addressing of addresses of the CAM core. The maintenance port further includes writing logic for writing



to addresses of the CAM core. A BIST controller for coordinating BIST testing of the CAM core is also provided. The BIST controller is capable of performing a BIST search and a BIST write on the CAM core at the same time.

In another embodiment, the search port and maintenance port are distributed and are not part of the BIST controller, so as to reduce the number of global wires required to communicate read/write/search data and results to the CAM core.

In yet another embodiment, content addressable memory (CAM) circuitry with BIST testing capabilities is disclosed. The circuitry includes a plurality of CAM cores and a plurality of BIST circuits that are coupled to each of the CAM cores. A single BIST controller is also included and is capable of controlling BIST testing of each of the plurality of CAM cores. Preferably, the single BIST controller is configured to perform BIST searches on each of the plurality of CAM cores during each cycle and is further configured to perform BIST writing during any cycle including a cycle in which the BIST search occurs.

In still a further embodiment, a method for performing built-in self-test (BIST) testing on a content addressable memory (CAM) core is disclosed. The method includes: (a) writing test data to memory addresses in the CAM core; and (b) searching for test data in the CAM core, the searching being continuously performed one cycle after another and the writing of the test data capable of being performed in a same cycle as one or more search performed during the searching. In one aspect of this embodiment, the method can further include an operation of selecting one row of the CAM core to be valid during the searching, such that matches only occur in the one row in one of the cycles. This aspect enables low power BIST testing.

The advantages of the present invention are numerous. Most notably, the BIST testing architecture enables testing of a CAM core separate from the testing of priority encoders (PEs). In one embodiment, separately testing the CAM core is configured to further enable the testing of one row of the CAM core at a time. This capability will reduce the power consumption during BIST testing. A further advantage of the present invention is the ability of the BIST testing architecture (e.g., by way of the BIST controller) to execute uninterrupted searches at each cycle while simultaneously performing writes that set up subsequent searches. For example, the BIST controller is capable simultaneously accessing the search port and the maintenance port of a CAM core. This functionality enables BIST testing of a CAM core at more realistic speeds which resemble realistic CAM operation.

Other advantages of the invention will become apparent from the following detailed description, taken in conjunction with the accompanying drawings, illustrating by way of example the principles of the invention.

#### BRIEF DESCRIPTION OF THE DRAWINGS

The present invention will be readily understood by the following detailed description in conjunction with the accompanying drawings, wherein like reference numerals designate like structural elements.

FIG. 1 shows a block diagram of a CAM core implementing BIST circuitry in accordance with one embodiment of the present invention.

FIG. 2A shows a simplified example of the internal structure of the priority encoder 116.

FIG. 2B shows a detailed view of two subrows (i.e., subrow 0 and subrow 1) that define each ternary word of the CAM core.

FIG. 3 shows a detailed block diagram of a CAM circuit capable of operating BIST operations, in accordance with one embodiment of the invention.

FIG. 4A shows a more detailed diagram of the search port comparator of FIG. 3, in accordance with one embodiment of the present invention.

FIG. 4B shows a more detailed diagram of the BIST maintenance interface of FIG. 3, in accordance with one embodiment of the present invention.

FIG. 4C shows a more detailed diagram of the maintenance port comparator of FIG. 3, in accordance with one embodiment of the present invention.

FIG. 4D shows a more detailed diagram of the BIST search interface (SIF) of FIG. 3, in accordance one embodiment of the present invention.

FIGS. 5A–5U show a flowchart for a method implementing BIST testing for a ternary CAM, in accordance with one embodiment of the present invention.

FIG. 6 shows a block diagram of an alternative embodiment of the present invention, in which one BIST controller coordinates testing of multiple CAM cores.

#### DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENTS

An invention for the built-in self-test (BIST) of a content addressable memory (CAM) core with ternary storage (including a don't care state) is disclosed. In one embodiment, the ternary CAM also contains binary valid and tag bits. The testing of CAM cores in accordance with the present invention is configured to test the CAM cores separate from priority encoders (PE). Separating the testing of the CAM core from the testing of PE enables BIST testing of one row of the CAM core at a time. By doing this, lower power operation can be achieved during BIST testing. A further feature of the present invention is the ability of the BIST testing architecture to execute uninterrupted searches at each cycle while simultaneously performing writes that set up subsequent searches. This functionality enables BIST testing of a CAM core at more realistic speeds which resemble the functional CAM speed. In the following description, numerous specific details are set forth in order to provide a thorough understanding of the present invention. It will be understood, however, to one skilled in the art, that the present invention may be practiced without some or all of these specific details. In other instances, well known process operations have not been described in detail in order not to unnecessarily obscure the present invention.

FIG. 1 shows a block diagram 100 of a CAM core implementing BIST circuitry in accordance with one embodiment of the present invention. In one embodiment, ternary CAMs are preferably used instead of binary CAMs because ternary CAMs speed up routing and address translation, which is important in high speed applications, such as Internet applications. The CAM core 112 is shown in communication with a block decoder 108, a row decoder 110, a search port 114 and a read/write port 120. The CAM core 112 is coupled to a priority encoder 116. The priority encoder 116 is configured to select the highest priority matching result for a CAM search. During BIST operation, a BIST controller 102 is configured to provide inputs to multiplexers 104 and 106, in the form of row addresses and block addresses respectively. The multiplexers 104 and 106 are thus provided with a BIST mode select signal to enable operation in the test mode, as opposed to a functional mode. The search port 114 receives, as inputs, tag bits, search data and valid bits. The read/write port 120 receives, as inputs, write data, tag and valid bits and



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sends, as outputs, read data, tag and valid bits. The functionality of the tag bits, and valid bits will be described in greater detail below.

During a search, whether it be during BIST operation or functional operation, search data is provided by way of the search port **114** and the matches are communicated from the CAM core **112** to the priority encoder **116**. In this example, the CAM core **112** has 4096 rows, and each row is coupled by way of a bus to the priority encoder **116**. Thus, the priority encoder **116** will have 4096 entries with entry **0** having the highest priority and entry **4095** having the lowest priority for a search result. Consequently, the priority encoder **116** is charged with outputting a search result, which represents the match having the highest priority.

For illustration purposes, FIG. 2A shows a simplified example of the internal structure of the priority encoder **116**. Thus, if the priority encoder **116** only received an 8 bit bus instead of the 4096 bit bus, all matching entries are ranked depending on their priority. As shown, entry **3** and entry **6** are both recorded as matches. Of these matches, entry **3** has a high priority than entry **6**. Therefore, the address for entry **3** is designated as the search result.

In this example, the CAM core consists of 4096 rows of 32 ternary data bits **202**, 4 binary tag bits **208**, and 2 binary valid bits **206** each. Since the CAM is ternary, each row of 32 ternary data bits is implemented as 2 sub-rows of 32 binary bits each. FIG. 2B shows a detailed view of the two sub-rows (sub-row **0** and sub-row **1**) that define each 32 bit word of the CAM core **112**. Thus a ternary “don’t care” bit *i* is stored as a logic “11” pattern at bit *i* in the two sub-rows, while a logic “0” and “1” ternary states are stored as “01” and “10” patterns respectively. In addition to the 32 data bits **202**, each sub-row will also include a valid bit **206** and two tag bits **208**. Specifically, sub-row **0** first has a valid bit **206**, followed by a tag bit **208**, data bits **202**, and a tag bit **208**. Sub-row **1** first has a tag bit **208**, followed by data bits **202**, a tag bit **208**, and a valid bit **206**. In operation, both valid bits **206** must be “1” to indicate valid sub-rows before a search operation can find a match. Accordingly, valid bits indicate whether a word has valid data. Tag bits are used to track the width of words. In general, the tag bits may be used to track any information and tag bits are viewed generally as regular data bits. In any event, when BIST testing is performed, the tag bits and valid bits are tested as any other bit is tested in the CAM core.

For read/write operations, the row decoder is designed to address each sub-row individually. Thus, the read/write ports on the CAM have 32 data bits, 2 tag bits, and 1 valid bit. For search operations however, both sub-rows are compared simultaneously against the search port data. Since the search port data is also ternary, it is encoded similarly as above, and thus consists of 64 bits of search data and 4 bits of tag data. Valid bits are not provided since the search is restricted over valid rows only, in this example.

FIG. 3 shows a detailed block diagram **300** of a CAM circuit capable of operating BIST operations, in according to one embodiment of the invention. A CAM core **302** is shown coupled to a search port **304**, a maintenance port **306**, a block decoder **308**, and a row decoder **310**. The search port **304** includes circuitry, such as a multiplexer **320**. Multiplexer **320** accepts search data (e.g., 64 bits of data and 4 tag bits) from a BIST search interface (SIF) **312** and functional search data, and can be controlled to operate in either functional mode or BIST mode. When in BIST mode, the multiplexer **320** selects BIST mode search data. The Maintenance port **306** is shown including a number of multiplexers **322**,

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**324**, and **326**. Each of the multiplexers can be controlled to either operate in the BIST mode or the functional mode. Multiplexer **322** is coupled to the row decoder **310**, the multiplexer **324** is coupled to the block decoder **308**, and the multiplexer **326** is coupled to the read/write port **390** of the CAM core and is configured to introduce data to the CAM core **302** during a writing operation.

The BIST circuitry further includes the BIST search interface (SIF) **312**, a BIST controller **314**, a BIST maintenance interface **316**, the search port comparator **340**, and a maintenance port comparator **318**. The BIST controller **314**, in one embodiment, is a state machine that executes the CAM BIST algorithm of the present invention, and will be described in greater detail with reference to FIGS. 5. The CAM BIST algorithm processed by the BIST controller **314** can be defined by Verilog® code (i.e., register transfer level description), which can then be synthesized into gates. Also provided is an IEEE 1149.1 controller, which is commercially available in Verilog® form. When operating in BIST mode, the BIST search interface (SIF) **312** will provide 64 bits of search data, to the search port **304** by way of a bus **313**. The 64 bits of search data are generated by a shift register in the SIF **312**, which receives a 1 bit data input from the BIST controller **314**. The bus **313** will also be used to communicate 4 tag bits provided by the BIST SIF **312** to the search port. The BIST SIF **312** will also receive control signals controlling the shift register, including a set and a reset signal from the BIST controller **314**. As mentioned earlier, the multiplexer **320** will therefore communicate the 64 bits of search data and 4 bits of tag to the CAM core **302** when the multiplexer **320** operates in BIST mode. The BIST controller **314** is configured to write in different sequences of bits to different locations in the CAM core **302**. The writing performed by the BIST controller **314** is facilitated by the multiplexer **322** and the multiplexer **324** of the maintenance port **306**. In accordance with the present invention, it is possible for the BIST controller **314** to perform writes to the CAM core **302** at the same time searches are performed. That is, searches can be performed in each cycle and the writes can be performed at the same time (e.g., writes can be performed to set up the searches performed during BIST testing). For example, the BIST controller **314** communicates BIST row addresses and BIST block addresses **902** to each of the multiplexers **322** and **324**. At the same time, since the BIST operates on one row at a time, the BIST row addresses and the BIST block addresses **902** are communicated to the search port comparator **340** as the expected row and block addresses for searches.

The BIST controller **314** is also shown receiving signals **952** and **954** from the search port **340**. The signals **952** and **954** represent “hit” (i.e., a CAM row matches the search data) and multiple hit (referred to as “mult” since more than 1 CAM row matches the search data) results obtained from a search. As shown, signals **952** and **954** originate from the search port **304**. During a BIST search, the search port comparator **340** will also be receiving the search addresses by way of bus **904** from the search port **304**. In the search port comparator **340**, the expected search addresses **902** are compared to the search addresses **904** to produce a search result that is communicated by way of bus **906** to the BIST controller **314**. The search port comparator **340** will also compare hit and mults **952** and **954** to expected hits and mults **980** and **982** to produce search results that are also communicated over bus **906**. If the search comparator **340** generates a match between **902** and **904**, between **952** and **980**, or between **954** and **982**, then for each match a logic 1 will be produced (otherwise, a logic 0 will be produced if no match



occurs). In BIST mode, the BIST controller **314** will also be communicating 2 data bits, and 1 valid bit via bus **908** to the BIST maintenance interface **316**. In one embodiment, only 4 types of data are written into the CAM, so 2 data bits are sufficient to express the data. The BIST maintenance inter-  
 5 face **316** is capable of performing well known expansion operations on the data. In this case, the data from the BIST controller is replicated to generate 32 data bits, 2 tag bits, and 1 valid bit over a bus **910** which communicates with the multiplexer **326**. This scheme greatly reduces the number of  
 10 global wires required to communicate BIST data, greatly simplifying chip routing.

The maintenance port comparator **318** is configured to receive 2 bits of expected data, 2 bits of expected tag and 1 bit of expected valid over bus **916**. In one embodiment, there are four types of expected data, and therefore, 2 bits are sufficient to express the expected data. As shown, the main-  
 15 tenance port comparator **318** will receive the output Q for a search from the read/write port **390** over bus **356**. The result of the comparison in the maintenance port comparator **318** (i.e., between **916** and **356**) is then communicated over bus **914** to the BIST controller **314**. If matches are obtained, a logic 1 is generated, otherwise a logic 0 is produced. As shown, the BIST controller **314** is also connected to the  
 20 IEEE 1149.1 controller **360** in order to communicate control and status signals to the external world. The IEEE 1149.1 standard is herein incorporated by reference.

FIG. **4A** shows a more detailed diagram of the search port comparator **340** of FIG. **3**, in accordance with one embodi-  
 25 ment of the present invention. Search port comparator **340** receives search addresses **904**, and row and block addresses through bus **902**, which constitutes the expected search address. The comparator will take the search addresses **904** and the row and block address to produce a search result that is output through bus **906**. In this embodiment, the search  
 30 address, and the row and block addresses are provided over 12 bit buses. As mentioned above, the search result is communicated to the BIST controller **314**. Also performed by the search port comparator **340** is a comparison between the hit and mult **952** and **954** and the expected hit and mult **980** and **982** to produce a result that is also output through bus **960**. If a match is obtained in these comparisons, a logic 1 is pro-  
 35 duced. Of course, the preferred convention of assigning 1's to matches and 0's to non-matches can be changed based on preference and design.

FIG. **4B** shows a more detailed diagram of the BIST main-  
 40 tenance interface **316** of FIG. **3**, in accordance with one embodiment of the present invention. BIST maintenance interface **316** receives 2 data bits, 2 tag bits, and a valid bit as inputs through bus **908** from BIST controller **314**. As mentioned above, the BIST maintenance port **316** is configured to perform an expansion function to produce 32 bits of data. The data bits, the tag bits, and the valid bit are each prefer-  
 45 ably buffered before being communicated as BIST data through bus **910** to maintenance port **306**.

FIG. **4C** shows a more detailed diagram of the mainte-  
 50 nance port comparator **318** of FIG. **3**, in accordance with one embodiment of the present invention. Maintenance port comparator **318** includes three comparators for data bits, tag bits, and valid bits and receives expected data from BIST controller **314** through bus **916** and actual data from read/  
 55 write port **390** through bus **912**, and provides the result of the comparison to BIST controller **314** through bus **914**.

FIG. **4D** shows a more detailed diagram of the BIST  
 60 search interface (SIF) **312** of FIG. **3**, in accordance with one embodiment of the present invention. BIST search interface

**312** is implemented as a 64 bit shift register that receives a 1 bit "SIF\_data," from the BIST controller. The SIF **312** also receives control signals, specifically a reset and set, control signal inputs from BIST controller **314** and provides BIST search data to search port **304** through bus **313**. Once again, the use of a shift register permits BIST search data to be communicated over a single wire plus two control signals, thus significantly reducing the number of BIST related global wires.

FIGS. **5** show a flowchart for a method implementing BIST testing for a ternary CAM in accordance with one embodiment of the present invention. As mentioned above, the BIST testing is controlled by the BIST controller **314**. Therefore, the following method operations represent algo-  
 10 rithm operations carried out by the BIST controller **314**. As will be appreciated by those skilled in the art, the BIST controller **314** of the present invention is capable of simultaneously searching by way of the search port **304** of the CAM core **302** and writing using the maintenance port **306**. This is particularly powerful since conventional designs do not allow for continuous searching during write operations. In contrast, the prior art requires that searches be stopped for one or more cycles while data is written to the CAM core to set up a next search or searches. The powerful architecture of the present invention therefore facilitates high speed BIST testing, which more accurately resembles the functional operation of a CAM. That is, the BIST testing can be carried out at-speed or nearly at speed, thus enabling more realistic testing of CAM cores.

With the foregoing in mind, the method begins at an operation **500** where all cells in the CAM core are initialized to logic 0 by writing successively to all locations in the CAM. That is, during the initialization, the tag bits, data bits, and valid bits are all initialized to zero. Once initialized, the method moves to an operation **502** where the built-in self-  
 35 test (BIST) testing will begin with valid bit testing and start at a particular address. The particular address may be any location in the CAM core. Typically BIST testing begins at one point in the CAM core and will proceed address-by-address until the entire CAM core has been tested using the BIST algorithm. Valid bit testing begins at an operation **504** where logic 1 is written to the valid bit of a subrow 0. Simultaneously, a search is performed where the search data and search tag bits are all logic 0, and the expected result is a miss. As mentioned earlier, searching can be performed at each cycle and writing can be performed at the same time searches are being performed.

In one implementation of a CAM, a write operation takes effect only at the end of the cycle in which the operation is initiated, while a search operation takes effect at the begin-  
 40 ning of the cycle. Thus when a search and a write are initiated in the same cycle, the write data is not included in the current search, but in a search initiated in a subsequent cycle. Moving from operation **504** to operation **506**, it is determined in operation **506** whether a miss occurred, knowing that a miss was expected in operation **504**. If a miss did not occur as expected, the method proceeds to an operation **507** where a flag is set indicating that the BIST test failed.

Once the flag is set, the method proceeds to the next operation and moves to A of FIG. **5B**. At operation **508**, a write of the logic 1 is performed to the valid bit of subrow 1, and simultaneously, a search is performed with the data and tag bits set at logic 0. The expected result will be a miss for operation **508**. In operation **506**, it is determined whether the miss occurred. If the miss did not occur as expected, the flag will be set indicating that the BIST test failed and the method will proceed to the next operation. Alternatively, if a



miss did occur, the method will proceed to operation **510** where a search is performed for data bits and tag bits where the search data and search tag bits are all at logic 0, and the expected result will be a hit.

In response to the search of **510**, the method will proceed to operation **512** where it is determined whether a hit occurred. If a hit does not occur as expected, a flag will be set indicating that the BIST test failed in operation **507**. Otherwise, the method will proceed to B of FIG. **5C**, and operation **511**. In operation **511**, a search of data bits and tag bits where the search data and search tag bits are all at logic 0, is performed and the expected result should be a hit. In operation **512**, if it is determined that a hit occurred, the method will proceed to tag testing and operation **514**. Otherwise, a flag will be set indicating that the BIST test failed in operation **507**. As the method moves to operation **514**, valid bit testing will be done and tag bit testing will begin.

In operation **514**, BIST testing continues with tag testing which starts at the same particular address as the valid bit testing began. Now, the method moves to operation **516** where tag bits are written in subrow 0 to 01, and a search is performed with data set to 0's, and the tag bits set to 01 and 01. The expected result will be a miss. Accordingly, the method will move to operation **506** where it is determined whether a miss occurred. If a miss does not occur, the method will proceed again to operation **507** where the flag is set indicating that the BIST test failed. Otherwise, if the miss did occur, the method proceeds to C of FIG. **5D** and the operation **518**. In operation **518**, tag bits are written in subrow 1 to 01 and a search is performed with data set to 0's, and tag bits set to 00 and 01. The expected result will now be a hit. In operation **512**, it is determined whether a hit occurred. If a hit does not occur, the flag will set indicating that the BIST testing failed. Otherwise, the method will move to an operation **520** where a search of data bits and tag bits is performed where the search data is logic 0, and the tag bits are logic 01 and 01. The expected result for the search should be a hit.

From operation **520**, the method moves to a decision operation **512** where it is determined whether the hit occurred. If the hit does not occur, the flag will be set indicating that the BIST testing failed. Otherwise, the method will move to D of FIG. **5E**, and operation **522**. In operation **522**, a search of data bits and tag bits is performed where the search data is set to logic 0, and the tag bits are set to logic 00 and 01, and the expected result is a miss. In operation **506**, it is determined whether the miss occurred. If the miss did not occur, the method will move to operation **507** where a flag is set indicating that the BIST test failed. Otherwise, if the miss did occur, the method will proceed to operation **524**. In operation **524**, data bits and tag bits are searched where the search data is set to logic 0 and the tag bits are set to logic 11 and 01, and the expected result of the search will be a miss. From operation **524**, the method moves to decision operation **506** where it is determined whether a miss occurred. If a miss did not occur, the method will move to operation **507** where a flag indicating that the BIST test failed to set. Otherwise, the method moves to E of FIG. **5F** and operation **526**. In operation **526**, a search of data bits and tag bits is performed where the search data is logic 0, and the tag bits are logic 00 and 01. The expected result should be a miss.

In operation **506**, it will be determined whether a miss occurred, and if no miss occurred, the flag will be set indicating that the BIST test failed. Otherwise, the method will move to operation **528**. In operation **528**, the search is performed of data bits and tag bits where the search data is logic

0, and the tag bits are logic 01 and 11. The expected result should be a miss. In operation **506**, it is determined whether a miss occurred. If a miss did not occur, the method will move to operation **507** where a flag is set indicating that the BIST testing failed. If a miss did occur, the method will move to F and operation of **530** of FIG. **5G**. In operation **530**, a search is performed of data bits and tag bits where the search data is logic 0, and the tag bits are logic 01 and 00, and the expected result is a miss. If it is determined that a miss did not occur in operation **506**, the method will move to operation **507**, where a flag is set indicating that the BIST test failed. Otherwise, the method will move to operation **532** where tag bits in subrow 0 are written to 11 and the search is performed with data set to 0's and tag bits set to 11, and 11. The expected result will be a miss.

In operation **506**, it will be determined whether the miss occurred. If the miss did not occur, the method again proceeds to operation **507** where a flag is set indicating that the BIST test failed. Otherwise, the method will move to G of FIG. **5H**, and operation **534**. In operation **534**, tag bits are written in subrow 1 to 11, and a search is performed with the data set to 0, and the tag bits set to 01 and 11. The expected result will be a hit. If a hit does not occur in operation **512**, the method will move to operation **507** where a flag is set indicating that the BIST test failed. If a hit does occur, the method proceeds to operation **536** where data bits and tag bits are searched, where the search data is logic 0 and the tag bits are logic 11 and 11, and the expected result is a hit. If a hit does not occur, the method will proceed again to operation **507** where a flag is set indicating that the BIST test failed. Otherwise, the method will move to H of FIG. **5I** and operation **538**. In operation **538**, a search of data bits and tag bits is performed where the search data is logic 0 and the tag bits are logic 01 and 11, and the expected result is a miss. If a miss does not occur in operation **506**, a flag is set indicating that the BIST testing failed in operation **507**. Otherwise, the method will move to operation **540** where a search of data bits and tag bits is performed where the search data is logic 0, and the tag bits are logic 01 and 11. The expected result for this search should be a miss. If a miss does not occur, the flag will be set indicating that the BIST testing failed. Otherwise, the method will move to I of FIG. **5J** and operation **542**.

In operation **542**, a search is performed of data bits and tag bits where the search data is logic 0, and the tag bits are logic 10 and 11. The expected results should be a miss. If a miss does not occur, the flag will be set indicating that the BIST testing failed in operation **507**. Otherwise, the method will move to operation **544** where a search of data bits and tag bits is performed where the search data is logic 0, and the tag bits are logic 11 and 01, and the expected result is a miss. If a miss does not occur, the method will move to operation **507** where a flag is set indicating that BIST test failed. Otherwise, the method will move to J of FIG. **5K** and operation **546**. In operation **546**, a search is performed of data bits and tag bits where the search data is logic 0 and the tag bits are logic 11 and 10, and the expected result is a miss. If a miss does not occur, the method will proceed to operation **507** where a flag is set indicating the BIST testing failed. Otherwise, if a miss does occur, the method will proceed to operation **548** where a writing of tag bits in subrow 0 to 10, and a search is performed with data set to 0's and the tag bits set to 10 and 10. The expected result should be a miss. If a miss does not occur, the flag will be set indicating that the BIST test failed. If miss does occur, the method will move to K of FIG. **5L** and operation **550**.

In operation **550**, tag bits are written in subrow 1 to 10 and a search is performed with data set to 0's, and tag bits set to



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11 and 10, and the expected result is a hit. In operation **512**, it is determined whether a hit occurred. If a hit does not occur, this flag will be set indicating that the BIST test failed. Otherwise, the method will proceed to operation **552** where a search of data bits and tag bits is performed where the search data is logic 0, and the tag bits are logic 10 and 10. The expected result should be a hit. If a miss does not occur in operation **506**, the method will proceed to operation **507** where a flag is set indicating that the BIST test failed. Otherwise, the method will proceed to L of FIG. **5M** in operation **554**.

In operation **554**, a search is performed for data bits and tag bits where the search data is logic 0, and the tag bits are logic 11 and 10, and the expected result is a miss. If a miss does not occur, the flag will be set indicating that the BIST test failed. Otherwise, the method will proceed to operation **556** where a search of data bits and tag bits is performed where the search data is logic 0, and the tag bits are logic 00 and 10. The expected result will be a miss. If a miss does not occur in operation **506**, the flag will be set indicating that the BIST test failed, otherwise, the method will proceed to operation **557** where a search of data bits and tag bits is performed where the search data is logic 0 and the tag bits are logic 11 and 10. The expected result should be a miss. If a miss does not occur in operation **506**, the method will proceed to operation **507** where a flag is set indicating the BIST has failed. Otherwise, the method will proceed to M of FIG. **5N** and operation **558**. In operation **558**, a search is performed of data bits and tag bits where the search data is logic 0, and the tag bits are logic 10 and 00. The expected result should be a miss. If a miss does not occur in operation **506**, the method will proceed to operation **507** where a flag is set indicating that the BIST testing failed. Otherwise, the method will proceed to operation **560** where a search of data bits and tag bits is performed for the search data is logic 0, and the tag bits are logic 10 and 11. The expected result should be a miss. If a miss does not occur, the method will proceed to **507** where a flag is set indicating that the BIST testing failed. Otherwise, the method will proceed to N of FIG. **5O** and operation **562**.

In operation **562**, tag bits are written in subrow 0 to 00, and a search is performed with the data set to 0's and the tag bits set to 00 and 00. The expected result should be a miss. If a miss does not occur, the method will proceed to operation **507** where a flag is set indicating that the BIST test failed. Otherwise, the method will proceed to operation **564**. In operation **564**, tag bits are written in subrow 1 to 00, and a search is performed with the data set to 0's, and the tag bits set to 10 and 00, where the expected result is a hit. If a hit does not occur, the method will proceed to operation **507** where a flag is set indicating that the BIST test failed. Otherwise, the method will proceed to O of FIG. **5P** and operation **566**. In operation **566**, a search is performed of data bits and tag bits where the search data is logic 0 and the tag bits are logic 00 and 00. The expected result is a hit. If it is determined that a hit does not occur in operation **512**, the method will proceed to operation **507** where a flag is set indicating that the BIST test failed.

If a hit did occur, the method proceeds to operation **568** where a search is performed of data bits and tag bits where the search data is logic 0, and the tag bits are logic 10 and 00, and the expected result is a miss. If the result is not a miss, the method will move to operation **507** where the flag is set indicating that the BIST test failed. Otherwise, the method will proceed to P of FIG. **5Q** and operation **570**. In operation **570**, a search will be performed of data bits and tag bits where the search data is logic 0 and the tag bits are logic 10

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and 00. The expected result will be a miss. If a miss does not occur in operation **506**, the method will proceed to an operation **507** where a flag is set indicating that the BIST test failed. If a miss does occur, the method will proceed to operation **572** where a search is performed of data bits and tag where the search data is logic 0, and the tag bits are logic 01 and 00, and the expected result is a miss. If a miss does not occur, the flag will be set indicating that the BIST test failed in operation **507**. Otherwise, the method will proceed to Q of FIG. **5R** and operation **574**.

In operation **574**, a search will be performed of data bits and tag bits where the search data is logic 0, and the tag bits are logic 00 and 10. The expected result will be a miss. If a miss does not occur, the method will proceed to operation **507** where a flag is set indicating that the BIST test failed. Otherwise, the method will proceed to an operation **576** where a search will be performed of data bits and tag bits where the search data is logic 0, and the tag bits are logic 00 and 01. The expected result should be a miss. If a miss does not occur, the flag will be set in operation **507** indicating that BIST test failed. Otherwise, the method will proceed to R of FIG. **5S**. At this point, tag testing will be complete.

The method now moves to R of FIG. **5S** where BIST testing will continue with data bit testing, and the testing will start at the particular address in operations **578**. As mentioned above, the particular address should be the same address at which valid bit testing and tag bit testing began as described above. The method now moves to operation **580** where the following operations are performed for a 1 bit slice of each subrow at the pre-determined address.

In operation **582**, "SIF\_data" is set to logic 1 and a search is performed simultaneously. A miss is expected in operation **582**. The method now moves to operation **506** where it is determined whether a miss occurred. If a miss does not occur, the flag is set indicating that the BIST test failed. Otherwise, the method moves to operation **584** where "SIF\_data" is set to logic 0, and a search is simultaneously performed expecting a miss. If it is determined in operation **506** that a miss did not occur, the method proceeds to operation **507** where a flag is set indicating that the BIST test failed. Otherwise, the method moves to operation **586** where "SIF\_data" is set to logic 0, and a search is performed expecting a miss. If a miss does not occur, the flag will be set indicating that the BIST test failed in operation **507**. If a miss did occur, the method moves to S of FIG. **5P** and operation **588**.

In operation **588**, "SIF\_data" is set to logic 0. At the same time, a search is performed and a miss will be expected. In operation **506**, it is determined whether the miss occurred. If a miss did not occur, a flag will be set indicating that BIST test failed in operation **507**. If a miss did occur, the method will move to operation **589** where the method will repeat operations **586** through **588** "30" more times with SIF\_data set to 0 throughout. This is performed 30 more times since each search word is 64 bits long, and operations **582** through **588** together with operation **589** have the effect of walking a 1 through a background of zeros in the search data. Of course, if the words were of different sizes, this operation would be repeated as many times as needed to operate on all data bits.

Once these operations have been repeated, the method will move to operation **590** where data bits are written in subrow 0 to logic 1. At the same time, a search is performed with the data bits at logic 0, and the expected results being a hit. In operation **512**, it is determined whether a hit occurred. If a hit did not occur, the flag is set indicating that the BIST test failed. Otherwise, the method will move to operation



592 where data bits are written in subrow 1 to logic 1, and a search is simultaneously performed with the data bits at logic 0, and the expected result being a miss. If a miss did not occur, the flag will be set indicating that the BIST test failed. Otherwise, the method will move to P of FIG. 5U and operation 594. In operation 594, data bits are written in subrow 0 to logic 0, and the valid bit is written as 0, and a search is performed with the data bits at logic 1, and the expected result is a hit.

If a hit does not occur, the method will move to operation 507 where the flag is set indicating that the BIST test failed. Otherwise, the method will move to operation 596. In operation 596, data bits are written in subrow 0 to logic 0, and the valid bit is written as 0, and the search is performed with the data bits at logic 1, and the expected result being a miss. If a miss did not occur, the method moves to 507 where the flag is set indicating that the BIST test failed. If a miss did occur, the method moves to operation 598 where it is determined if there are anymore addresses to test. If there are more addresses to test, the method will go to operation 599 where the method proceeds to the next address and repeats the preceding operations starting at 502 and ending at 598. If all of the memory addresses have been tested, then the method will move to operation 600 where the memory is reset and the method will end.

As mentioned above, the algorithm is configured to perform BIST testing on one row at a time. This is enabled by making only the row being tested "valid," thus enabling searches during testing only in the valid rows. Of course, a search will always be performed over the entire CAM core, although, testing for matches will only occur in the valid row. This provides a substantial savings in power, thus making the BIST testing a low power test. Furthermore, the BIST testing executed by the BIST controller is capable of operating searches through the search port at the same time as writes are performed through the maintenance port. This, as mentioned above, improves testing efficiency in that searches can be executed at every cycle, and thus searches need not be stopped to enable writes.

FIG. 6 shows a block diagram of an alternative embodiment of the present invention. In this embodiment, one BIST controller 314 is capable of operating on multiple CAM cores 302. As shown, multiple CAM cores 302a, 302b, 302c and up to 302n each include their respective search ports 304 and maintenance ports 306. Along with replicating the CAM cores 302, the BIST circuits (CKTs) 305 are also replicated. In one embodiment, the BIST circuits 305 include a BIST search interface 312, a search port comparator 340, a BIST maintenance interface 316 and a BIST maintenance comparator 318. For simplicity and for purposes of illustrating the ability of a single BIST controller 314 to operate on any number of CAM cores 302, the signal lines 307 and 309 are shown connected directly to the search ports (SPs) 304 and maintenance ports (MPs) 306. However, in reality, the signals 307 and 309 actually interface with the BIST circuits 305.

In addition to being able to use one BIST controller 314 to execute BIST testing to the CAM cores 302, it is also important to note that the BIST controller 314 can simultaneously perform searches by way of the search ports 304 and writes by way of the maintenance ports 306. The other advantages described above with regard to a single CAM core therefore also apply to embodiments where more than one CAM core is tested using the BIST controller 314.

The present invention may be implemented using any type of integrated circuit logic, state machines, or software driven

computer-implemented operations. By way of example, a hardware description language (HDL) based design and synthesis program may be used to design the silicon-level circuitry necessary to appropriately perform the data and control operations in accordance with one embodiment of the present invention. By way of example, a VHDL® hardware description language available from IEEE of New York, N.Y. may be used to design an appropriate silicon-level layout. Although any suitable design tool may be used, another layout tool may include a hardware description language "Verilog®" tool available from Cadence Design Systems, Inc. of Santa Clara, Calif.

The invention may employ various computer-implemented operations involving data stored in computer systems. These operations are those requiring physical manipulation of physical quantities. Usually, though not necessarily, these quantities take the form of electrical or magnetic signals capable of being stored, transferred, combined, compared, and otherwise manipulated. Further, the manipulation performed are often referred to in terms, such as producing, identifying, determining, or comparing.

Any of the operations described herein that form part of the invention are useful machine operations. The invention also relates to a device or an apparatus for performing these operations. The apparatus may be specially constructed for the required purposes, or it may be a general purpose computer selectively activated or configured by a computer program stored in the computer. In particular, various general purpose machines may be used with computer programs written in accordance with the teachings herein, or it may be more convenient to construct a more specialized apparatus to perform the required operations.

Additionally, the various block diagrams may be embodied in any form which may include, for example, any suitable computer layout, semiconductor substrate, semiconductor chip or chips, printed circuit boards, packaged integrated circuits, or software implementations. Accordingly, those skilled in the art will recognize that the present embodiments are to be considered as illustrative and not restrictive, and the invention is not to be limited to the details given herein, but may be modified within the scope and equivalents of the appended claims.

What is claimed is:

1. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core, comprising:

a search port for enabling searches of the CAM core;

a maintenance port for enabling addressing of locations of the CAM core, the maintenance port further including writing logic for writing to locations of the CAM core;

a BIST controller for coordinating BIST testing of the CAM core, the BIST controller being capable of performing a BIST search on the CAM core on every cycle through the search port and performing a BIST write at selected times to the CAM core, the BIST write is capable of being performed in a same cycle as the BIST search; and

a maintenance port comparator being coupled between the BIST controller and a data output of the maintenance port, the maintenance port *comparator* being configured to compare an expected data generated by the BIST controller with actual data provided from the data output of the maintenance port.

2. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core as recited in claim 1, further comprising:

a search port interface for receiving test data, tag data and control signals from the BIST controller during BIST testing, the search port interface being coupled to the search port.



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3. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core as recited in claim 1, further comprising:

a search port comparator being coupled to the search port and the BIST controller, the search port comparator being configured to,

compare search addresses generated from the search port in response to BIST search and expected addresses generated by the BIST controller and communicated to the maintenance port, and

compare expected hit and multiple hit data generated by the BIST controller with generated hit and multiple hit data output through the search port.

4. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core as recited in claim 1, further comprising:

a BIST maintenance interface for communicating test data to the CAM core.

5. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core, comprising:

a search port for enabling searches of the CAM core;

a maintenance port for enabling addressing of locations of the CAM core, the maintenance port further including writing logic for writing to locations of the CAM core; and

a BIST controller for coordinating BIST testing of the CAM core, the BIST controller being capable of performing a BIST search on the CAM core on every cycle through the search port and performing a BIST write at selected times to the CAM core, the BIST write is capable of being performed in a same cycle as the BIST search.

6. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core as recited in claim 1, wherein the search port includes multiplexer logic for selecting between BIST search data and functional search data.

7. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core as recited in claim 1, wherein the maintenance port includes multiplexer logic for selecting: (a) between BIST row and block addresses and functional mode row and block addresses; and (b) between BIST write data and functional mode write data.

8. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core as recited in claim 1, further comprising an IEEE 1149.1 controller for communicating control signals to and from the BIST controller.

9. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core as recited in claim 1, wherein BIST circuit enables BIST searches on every clock cycle, the searches performed on every cycle enabling at-speed BIST testing.

10. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core as recited in claim 1, further comprising:

a search port comparator;

a search port interface;

a maintenance port comparator; and

a maintenance port interface;

wherein the search port interface, the maintenance port interface, the search port comparator, and the maintenance port comparator are each distributed and have expansion capabilities, which limits a number of global wires required to communicate read/write and search data and results.

11. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core as recited in claim 1,

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wherein the BIST testing only tests one row per cycle and the BIST testing is separate from a priority encoder (PE) BIST testing.

12. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core, comprising:

a search port for enabling searches of the CAM core;

a maintenance port for enabling addressing of locations of the CAM core, the maintenance port further including writing logic for writing to locations of the CAM core;

a BIST controller for coordinating BIST testing of the CAM core, the BIST controller being capable of performing a BIST search on the CAM core on every cycle through the search port and performing a BIST write at selected times to the CAM core, the BIST write is capable of being performed in a same cycle as the BIST search; and

a search port interface that receives only 1 bit of test data from the BIST controller, other bits being generated internally to the search port interface, so as to limit a number of required global wires.

13. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core, comprising:

a search port for enabling searches of the CAM core;

a maintenance port for enabling addressing of locations of the CAM core, the maintenance port further including writing logic for writing to locations of the CAM core;

a BIST controller for coordinating BIST testing of the CAM core, the BIST controller being capable of performing a BIST search on the CAM core on every cycle through the search port and performing a BIST write at selected times to the CAM core, the BIST write is capable of being performed in a same cycle as the BIST search; and

a BIST maintenance interface for communicating test data to the CAM core, the BIST maintenance interface having a capability to expand 2 bits of data from the BIST controller to a required width, the capability to expand being configured to limit a number of needed global wires.

14. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core, comprising:

a search port for enabling searches of the CAM core;

a maintenance port for enabling addressing of locations of the CAM core, the maintenance port further including writing logic for writing to locations of the CAM core; and

a BIST controller for coordinating BIST testing of the CAM core, the BIST controller being capable of performing a BIST search on the CAM core on every cycle through the search port and performing a BIST write at selected times to the CAM core, the BIST write is capable of being performed in a same cycle as the BIST search, wherein BIST testing of the BIST circuit does not generate matches on all rows in the CAM core so as to enable low-power BIST operation.

15. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core, comprising:

a search port for enabling searches of the CAM core;

a maintenance port for enabling addressing of locations of the CAM core, the maintenance port further including writing logic for writing to locations of the CAM core; and

a BIST controller for coordinating BIST testing of the CAM core, the BIST controller being capable of performing a BIST search on the CAM core on every cycle



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through the search port and performing a BIST write at selected times to the CAM core, the BIST write is capable of being performed in a same cycle as the BIST search, wherein the CAM core is a ternary CAM core that is capable of storing three states.

16. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core as recited in claim 15, wherein a word is comprised of two subrows and each of the two subrows include a plurality of binary tag bits and valid bits.

17. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core, comprising:

- a search port for enabling searches of the CAM core;
- a maintenance port for enabling addressing of locations of the CAM core, the maintenance port further including writing logic for writing to locations of the CAM core;
- a BIST controller for coordinating BIST testing of the CAM core, the BIST controller being capable of performing a BIST search and a BIST write on the CAM core at the same time, wherein the BIST search can be performed on every cycle and the BIST write can be performed at any cycle including a cycle in which the BIST search is performed.

18. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core as recited in claim 17, wherein the CAM core is a ternary CAM core that is capable of storing three states.

19. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core as recited in claim 18, wherein a word is comprised of two subrows and each of the two subrows include data bits, and binary tag bits and a valid bit, and the data bits span a 32 bit width, the binary tag bits span 2 bits and the valid bit spans 1 bit.

20. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core as recited in claim 17, further comprising:

- a search port interface for receiving test data, tag data and control signals from the BIST controller during BIST testing, the search port interface being coupled to the search port; and
- a search port comparator being coupled to the search port and the BIST controller, the search port comparator being configured to,
  - compare search addresses generated from the search port in response to BIST search and expected addresses generated by the BIST controller and communicated to the maintenance port, and
  - compare expected hit and multiple hit data generated by the BIST controller with generated hit and multiple hit data output through the search port.

21. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core as recited in claim 17, further comprising:

- a BIST maintenance interface for communicating test data to the CAM core, the BIST maintenance interface having an expansion capability; and
- a maintenance port comparator being coupled between the BIST controller and a data output of the maintenance port, the maintenance port *comparator* being configured to compare an expected data generated by the BIST controller with actual data provided from the data output of the maintenance port.

22. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core as recited in claim 17, wherein the BIST testing only tests one row per cycle and the BIST testing is separate from a priority encoder (PE) BIST testing.

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23. Content addressable memory (CAM) circuitry with BIST testing capabilities, comprising:

- a plurality of CAM cores;
- a plurality of BIST circuits couple to each of the CAM cores; and
- a single BIST controller being capable of controlling BIST testing of each of the plurality of CAM cores.

24. Content addressable memory (CAM) circuitry with BIST testing capabilities as recited in claim 23, wherein the signal BIST controller is configured to perform BIST searches on each of the plurality of CAM cores during each cycle and is further configured to perform BIST writing during any cycle including a cycle in which the BIST search occurs.

25. Content addressable memory (CAM) circuitry with BIST testing capabilities as recited in claim 24, wherein each of the plurality of BIST circuits include:

- (a) a BIST search interface;
- (b) a search port comparator;
- (c) a maintenance port comparator; and
- (d) a BIST maintenance interface.

26. A method for performing built-in self-test (BIST) testing on a content addressable memory (CAM) core, comprising:

- writing test data to memory addresses in the CAM core;
- searching for test data in the CAM core, the searching being continuously performed one cycle after another and the writing of the test data capable of being performed in a same cycle as one or more search performed during the searching, wherein the CAM core is a ternary CAM core.

27. A method for performing built-in self-test (BIST) testing on a content addressable memory (CAM) core as recited in claim 26, further comprising:

- selecting one row of the CAM core to be valid during the searching, such that matches only occur in the one row in one of the cycles.

28. A method for performing built-in self-test (BIST) testing on a content addressable memory (CAM) core as recited in claim 26, wherein the CAM core has a plurality of words, and each of the plurality of words includes a plurality of binary tag bits and a valid bit.

29. A [method for performing] built-in self-test (BIST) circuit for testing [on] a content addressable memory (CAM) core, comprising:

- a search port for enabling searches of the CAM core;
- a maintenance port for enabling addressing of locations of the CAM core, the maintenance port further including writing logic for writing to locations of the CAM core; and
- a BIST controller for coordinating BIST testing of the CAM core, the BIST controller being capable of performing a BIST search and a BIST write on the CAM core at the same time.

30. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core, comprising:

- a maintenance port for enabling addressing of locations of the CAM core, the maintenance port further including writing logic for writing to locations of the CAM core;
- a BIST controller for coordinating BIST testing of the CAM core, the BIST controller being capable of performing a BIST write at selected times to the CAM core, wherein the BIST controller is capable of performing a



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*BIST search on the BIST write on the CAM core at the same time; and*

*a maintenance port comparator being coupled between the BIST controller and a data output of the maintenance port, the maintenance port comparator being configured to compare an expected data generated by the BIST controller with actual data provided from the data output of the maintenance port.*

*31. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core as recited in claim 30, further comprising:*

*a BIST maintenance interface for communicating test data to the CAM core.*

*32. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core, comprising:*

*a maintenance port for enabling addressing of locations of the CAM core, the maintenance port further including writing logic for writing to locations of the CAM core;*

*a BIST controller for coordinating BIST testing of the CAM core, the BIST controller being capable of performing a BIST write at selected times to the CAM core and performing a BIST search on the CAM core, and wherein the BIST controller is capable of performing the BIST search and the BIST write in a same cycle; and*

*a maintenance port comparator being coupled between the BIST controller and a data output of the maintenance port, the maintenance port comparator being configured to compare an expected data generated by the BIST controller with actual data provided from the data output of the maintenance port.*

*33. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core as recited in claim 32, wherein the BIST search is capable of being performed on every cycle.*

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*34. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core as recited in claim 32, further comprising:*

*a search port for enabling searches of the CAM core.*

*35. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core, comprising:*

*a maintenance port for enabling addressing of locations of the CAM core, the maintenance port further including writing logic for writing to locations of the CAM core;*

*a BIST controller for coordinating BIST testing of the CAM core, the BIST controller being capable of performing a BIST write at selected times to the CAM core and performing another operation on the CAM core, the BIST write is capable of being performed in a same cycle as the another operation; and*

*a maintenance port comparator being coupled between the BIST controller and a data output of the maintenance port, the maintenance port comparator being configured to compare an expected data generated by the BIST controller with actual data provided from the data output of the maintenance port.*

*36. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core as recited in claim 35, wherein the another operation is a BIST operation.*

*37. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core as recited in claim 35, wherein the another operation is capable of being performed on the CAM core on every cycle.*

*38. A built-in self-test (BIST) circuit for testing a content addressable memory (CAM) core as recited in claim 35, wherein the maintenance port includes multiplexer logic for selecting:*

*(a) between BIST row and block addresses and functional mode row and block addresses; and*

*(b) between BIST write data and functional mode write data.*

\* \* \* \* \*