



US00RE41530E

(19) **United States**
(12) **Reissued Patent**
Wood, Jr.

(10) **Patent Number:** **US RE41,530 E**
(45) **Date of Reissued Patent:** ***Aug. 17, 2010**

(54) **METHOD AND APPARATUS TO SELECT RADIO FREQUENCY IDENTIFICATION DEVICES IN ACCORDANCE WITH AN ARBITRATION SCHEME**

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(*) Notice: This patent is subject to a terminal disclaimer.

(21) Appl. No.: **10/693,696**

(22) Filed: **Oct. 23, 2003**

Related U.S. Patent Documents

Reissue of:

(64) Patent No.: **6,307,847**
Issued: **Oct. 23, 2001**
Appl. No.: **09/617,390**
Filed: **Jul. 17, 2000**

U.S. Applications:

(63) Continuation of application No. 09/026,043, filed on Feb. 19, 1998, now Pat. No. 6,118,789.

(51) **Int. Cl.**
H04W 4/00 (2006.01)

(52) **U.S. Cl.** **370/329; 370/346; 370/347**

(58) **Field of Classification Search** **370/329, 370/346, 347, 460, 408, 230, 437, 441, 442, 370/449, 458, 463, 342, 447, 445, 448, 432, 370/461, 475**

See application file for complete search history.

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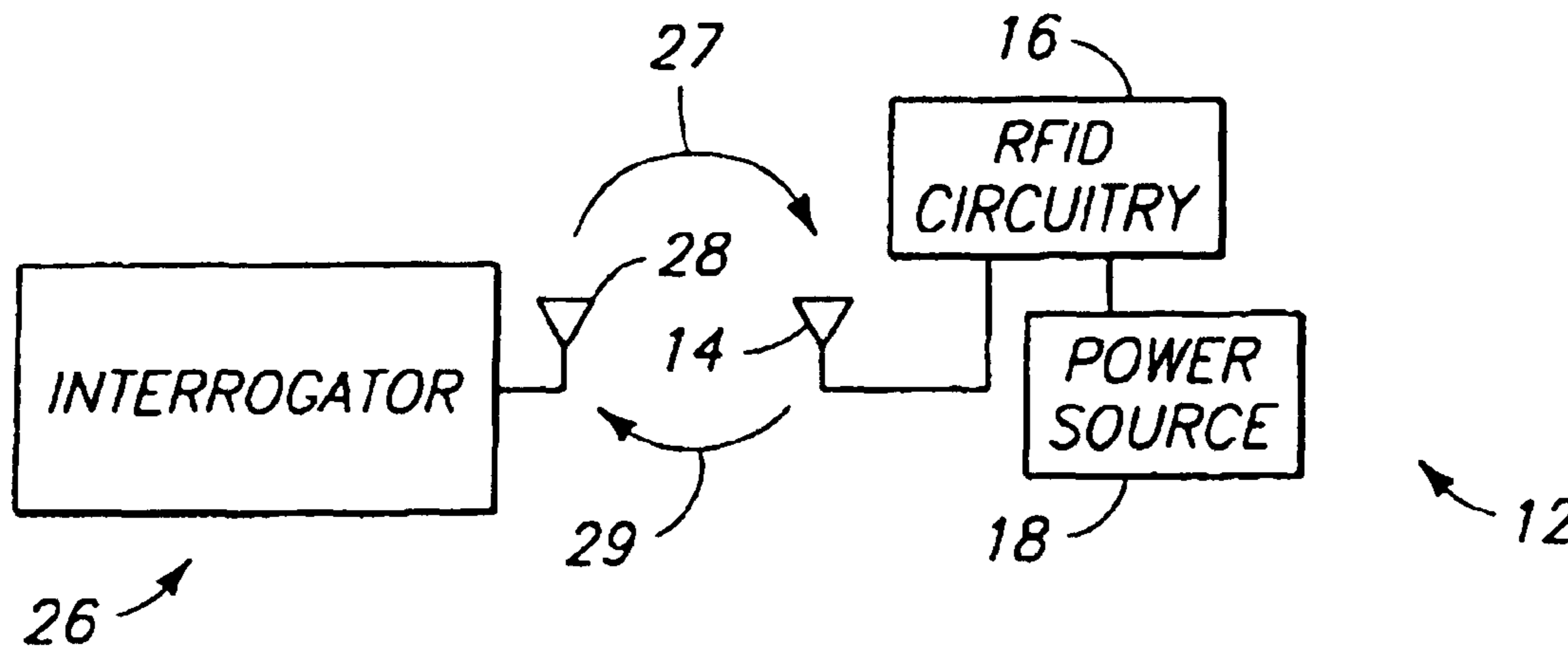
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(57) **ABSTRACT**

[A method of establishing wireless communications between an interrogator and individual ones of multiple wireless identification devices, the method comprising utilizing a tree search method to establish communications without collision between the interrogator and individual ones of the multiple wireless identification devices, a search tree being defined for the tree search method, the tree having multiple levels respectively representing subgroups of the multiple wireless identification devices, the method further comprising starting the tree search at a selectable level of the search tree. A communications system comprising an interrogator, and a plurality of wireless identification devices configured to communicate with the interrogator in a wireless fashion, the respective wireless identification devices having a unique identification number, the interrogator being configured to employ a tree search technique to determine the unique identification numbers of the different wireless identification devices so as to be able to establish communications between the interrogator and individual ones of the multiple wireless identification devices without collision by multiple wireless identification devices attempting to respond to the interrogator at the same time, wherein the interrogator is configured to start the tree search at a selectable level of the search tree.] *RFID devices are selected by an interrogator. The interrogator sends a signal to a plurality of RFID devices. The signal indicates a bit string and a memory range comprising multiple bit locations. RFID devices compare the bits stored in their respective memory ranges to the bit string to determine which of the RFID devices are selected.*

76 Claims, 3 Drawing Sheets



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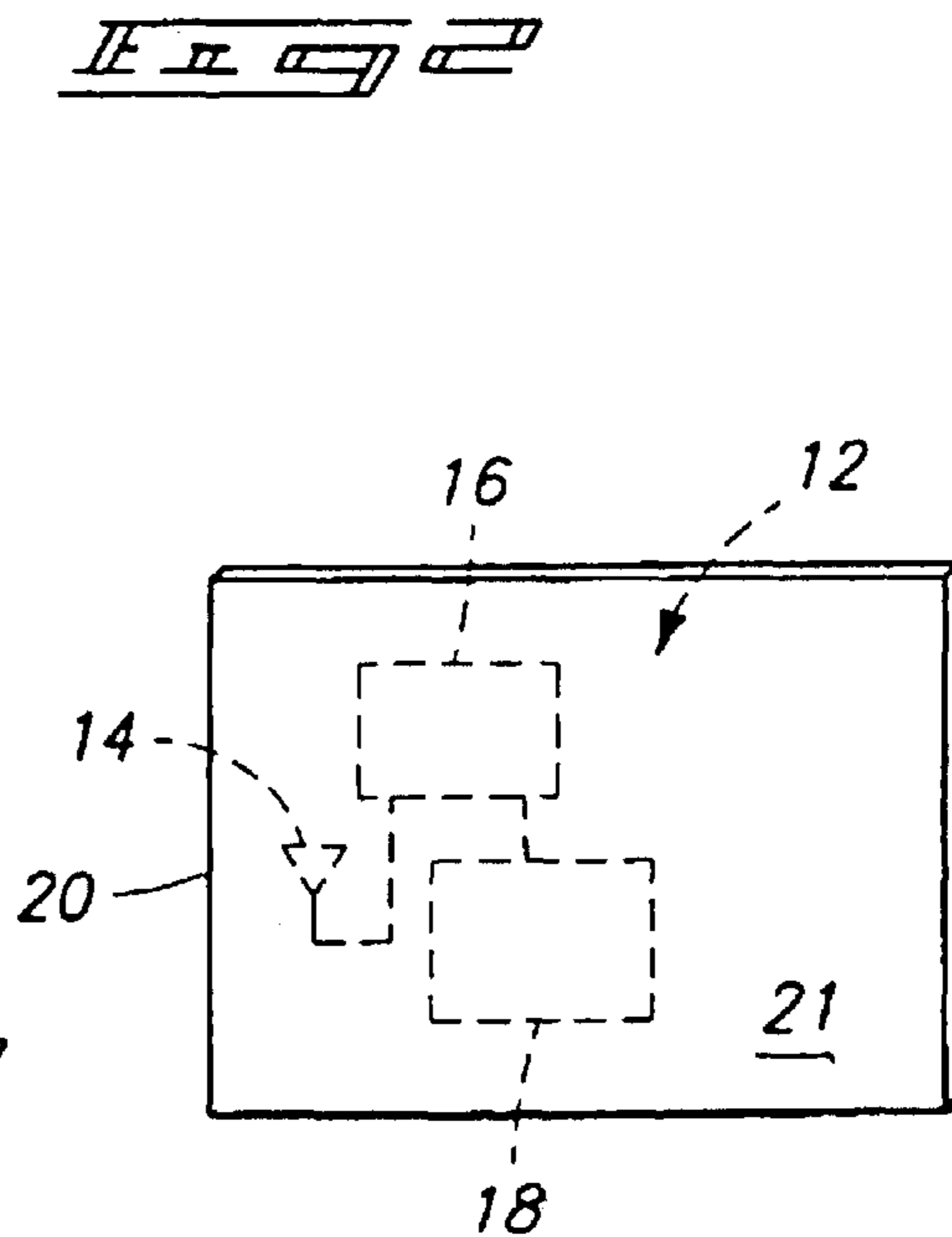
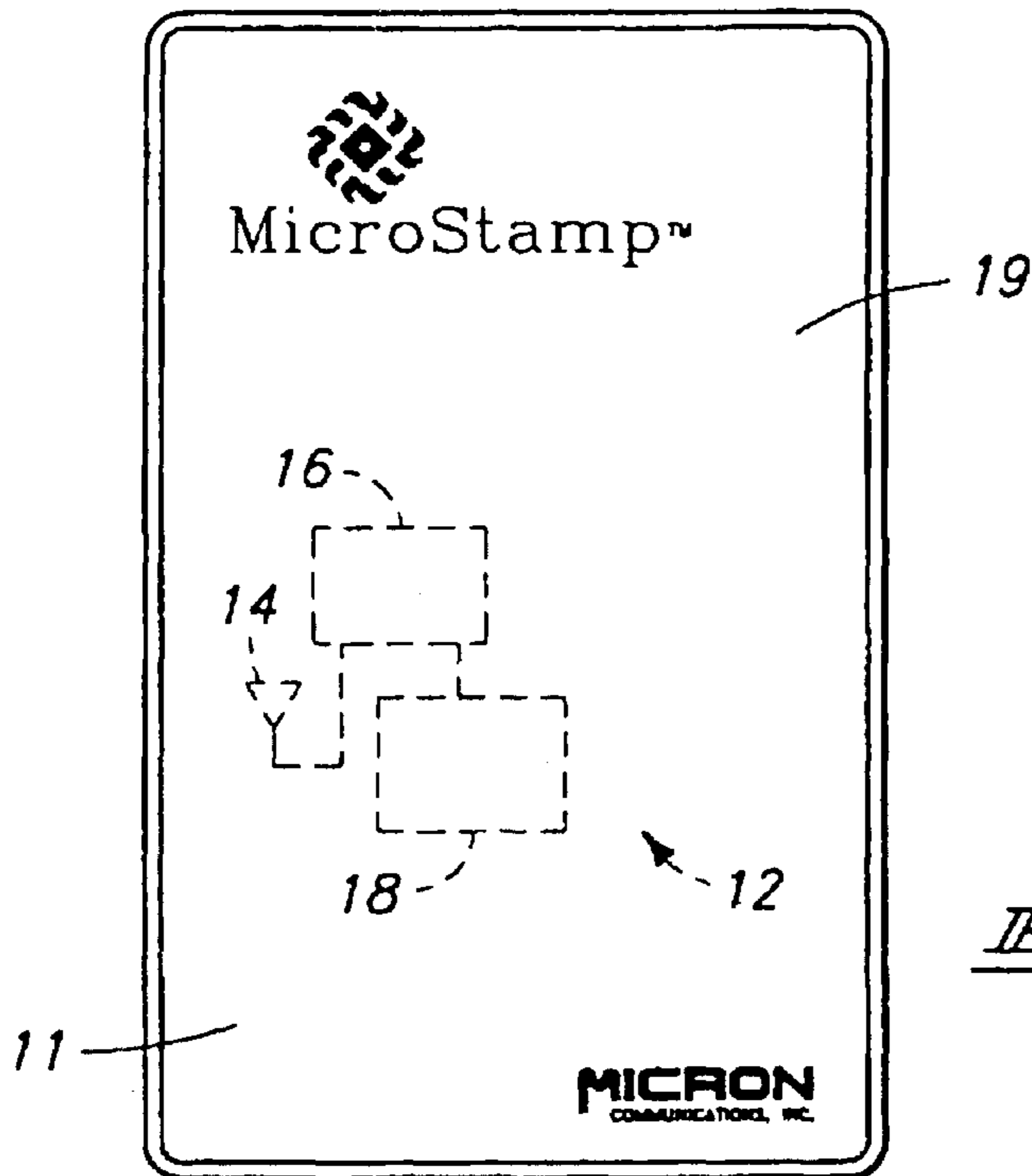
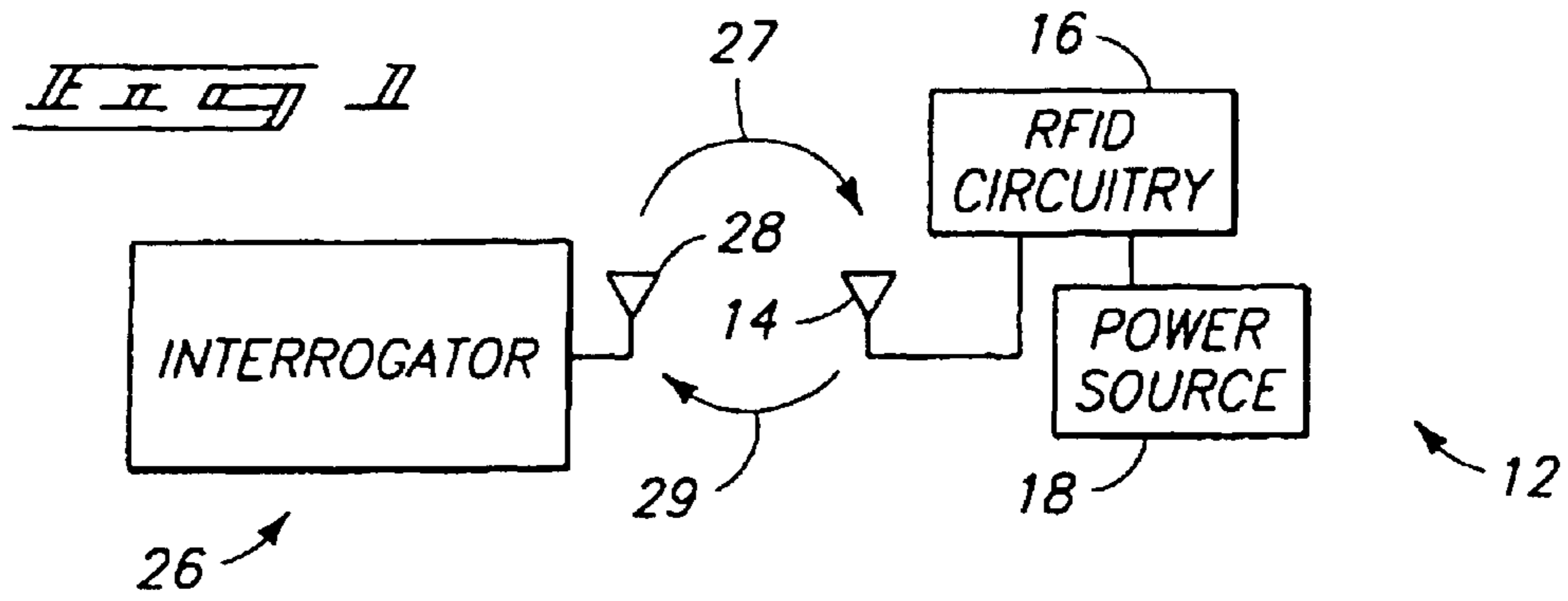
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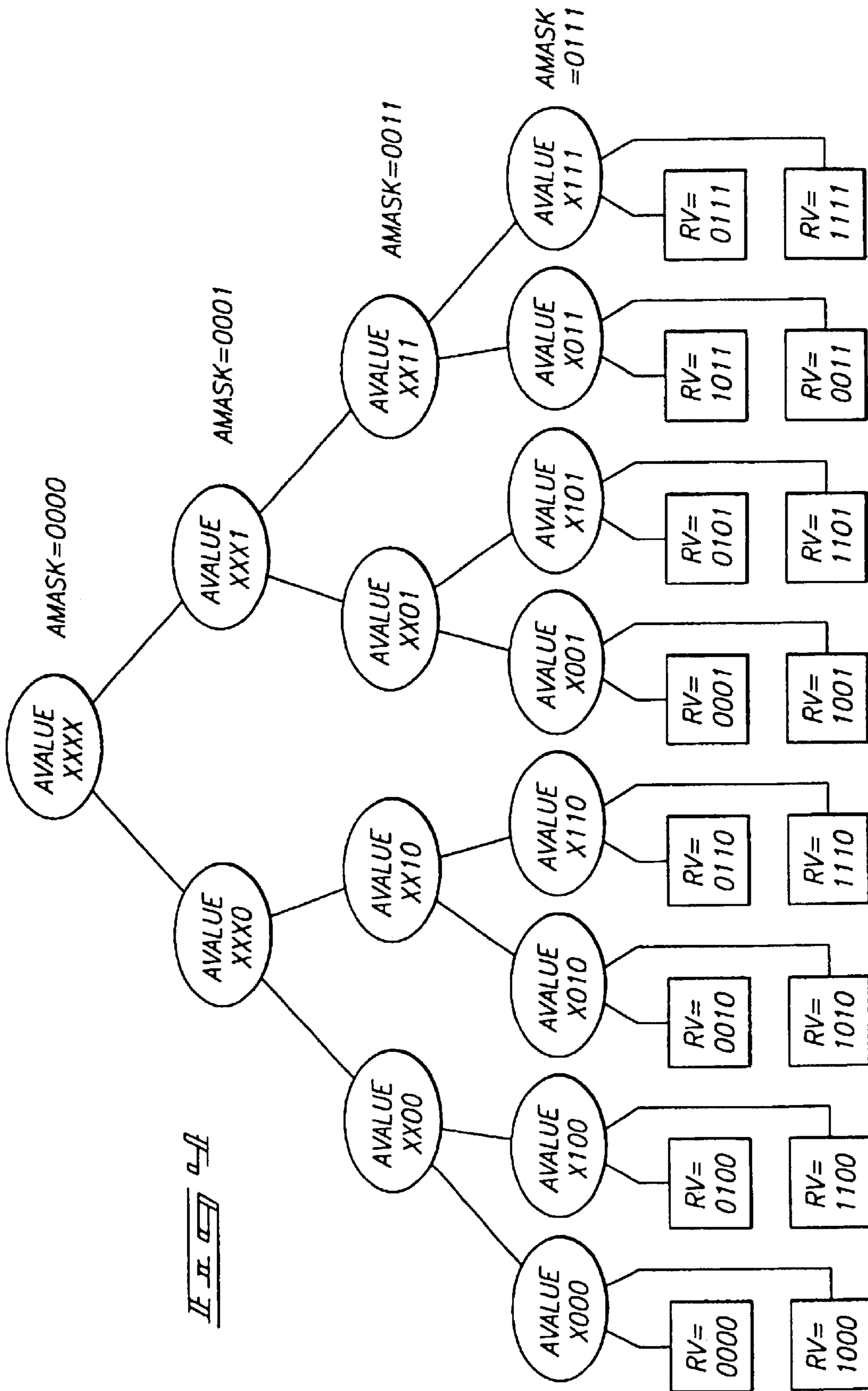
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Il n. a. 11

**METHOD AND APPARATUS TO SELECT
RADIO FREQUENCY IDENTIFICATION
DEVICES IN ACCORDANCE WITH AN
ARBITRATION SCHEME**

Matter enclosed in heavy brackets [] appears in the original patent but forms no part of this reissue specification; matter printed in italics indicates the additions made by reissue.

CROSS REFERENCE TO RELATED
[APPLICATION] APPLICATIONS

[This] *More than one reissue application has been filed for the reissue of U.S. Pat. No. 6,307,847, including the present reissue application Ser. No. 10/693,696, filed Oct. 23, 2003, a continuation reissue application Ser. No. 11/859,360, filed Sep. 21, 2007, a continuation reissue application Ser. No. 11/859,364, filed Sep. 21, 2007, and a continuation reissue application Ser. No. 12/493,542, filed Jun. 29, 2009. The present application is a reissue application of U.S. Pat. No. 6,307,847, issued from U.S. patent application Ser. No. 09/617,390, filed Jul. 17, 2000, and titled "Method of Addressing Messages and Communications System," which is a [Continuation] continuation application of U.S. patent application Ser. No. 09/026,043, filed Feb. 19, 1998, and titled "Method of Addressing Messages and Communications System," now U.S. Pat. No. 6,118,789, each of which is incorporated by reference.*

TECHNICAL FIELD

This invention relates to communications protocols and to digital data communications. Still more particularly, the invention relates to data communications protocols in mediums such as radio communication or the like. The invention also relates to radio frequency identification devices for inventory control, object monitoring, determining the existence, location or movement of objects, or for remote automated payment.

BACKGROUND OF THE INVENTION

Communications protocols are used in various applications. For example, communications protocols can be used in electronic identification systems. As large numbers of objects are moved in inventory, product manufacturing, and merchandising operations, there is a continuous challenge to accurately monitor the location and flow of objects. Additionally, there is a continuing goal to interrogate the location of objects in an inexpensive and streamlined manner. One way of tracking objects is with an electronic identification system.

One presently available electronic identification system utilizes a magnetic coupling system. In some cases, an identification device may be provided with a unique identification code in order to distinguish between a number of different devices. Typically, the devices are entirely passive (have no power supply), which results in a small and portable package. However, such identification systems are only capable of operation over a relatively short range, limited by the size of a magnetic field used to supply power to the devices and to communicate with the devices.

Another wireless electronic identification system utilizes a large active transponder device affixed to an object to be monitored which receives a signal from an interrogator. The device receives the signal, then generates and transmits a responsive signal. The interrogation signal and the respon-

sive signal are typically radio-frequency (RF) signals produced by an RF transmitter circuit. Because active devices have their own power sources, and do not need to be in close proximity to an interrogator or reader to receive power via magnetic coupling. Therefore, active transponder devices tend to be more suitable for applications requiring tracking of a tagged device that may not be in close proximity to an interrogator. For example, active transponder devices tend to be more suitable for inventory control or tracking.

Electronic identification systems can also be used for remote payment. For example, when a radio frequency identification device passes an interrogator at a toll booth, the toll booth can determine an identity of the radio frequency identification device, and thus of the owner of the device, and debit an account held by the owner for payment of toll or can receive a credit card number against which the toll can be charged. Similarly, remote payment is possible for a variety of other goods or services.

A communication system typically includes two transponders: a commander station or interrogator, and a responder station or transponder device which replies to the interrogator.

If the interrogator has prior knowledge of the identification number of a device which the interrogator is looking for, it can specify that a response is requested only from the device with that identification number. Sometimes, such information is not available. For example, there are occasions where the interrogator is attempting to determine which of multiple devices are within communication range.

When the interrogator sends a message to a transponder device requesting a reply, there is a possibility that multiple transponder devices will attempt to respond simultaneously, causing a collision, and thus causing an erroneous message to be received by the interrogator. For example, if the interrogator sends out a command requesting that all devices within a communication range identify themselves, and gets a large number of simultaneous replies, the interrogator may not be able to interpret any of these replies. Thus, arbitration schemes are employed to permit communications free of collisions.

In one arbitration scheme or system, described in commonly assigned U.S. Pat. Nos. 5,627,544; 5,583,850; 5,500,650; and 5,365,551, all to Snodgrass et al. and all incorporated herein by reference, the interrogator sends a command causing each device of a potentially large number of responding devices to select a random number from a known range and use it as that device's arbitration number. By transmitting requests for identification to various subsets of the full range of arbitration numbers, and checking for an error-free response, the interrogator determines the arbitration number of every responder station capable of communicating at the same time. Therefore, the interrogator is able to conduct subsequent uninterrupted communication with devices, one at a time, by addressing only one device.

Another arbitration scheme is referred to as the Aloha or slotted Aloha scheme. This scheme is discussed in various references relating to communications, such as *Digital Communications: Fundamentals and Applications*, Bernard Sklar, published January 1988 by Prentice Hall. In this type of scheme, a device will respond to an interrogator using one of many time domain slots selected randomly by the device. A problem with the Aloha scheme is that if there are many devices, or potentially many devices in the field (i.e. in communications range, capable of responding) then there must be many available slots or many collisions will occur. Having many available slots slows down replies. If the magni-

tude of the number of devices in a field is unknown, then many slots are needed. This results in the system slowing down significantly because the reply time equals the number of slots multiplied by the time period required for one reply.

An electronic identification system which can be used as a radio frequency identification device, arbitration schemes, and various applications for such devices are described in detail in commonly assigned U.S. patent application Ser. No. 08/705,043, filed Aug. 29, 1996, [and] *now U.S. Pat. No. 6,130,602, which is incorporated herein by reference.*

SUMMARY OF THE INVENTION

The invention provides a wireless identification device configured to provide a signal to identify the device in response to an interrogation signal.

One aspect of the invention provides a method of establishing wireless communications between an interrogator and individual ones of multiple wireless identification devices. The method comprises utilizing a tree search method to establish communications without collision between the interrogator and individual ones of the multiple wireless identification devices. A search tree is defined for the tree search method. The tree has multiple levels respectively representing subgroups of the multiple wireless identification devices. The method further comprising starting the tree search at a selectable level of the search tree. In one aspect of the invention, the method further comprises determining the maximum possible number of wireless identification devices that could communicate with the interrogator, and selecting a level of the search tree based on the determined maximum possible number of wireless identification devices that could communicate with the interrogator. In another aspect of the invention, the method further comprises starting the tree search at a level determined by taking the base two logarithm of the determined maximum possible number, wherein the level of the tree containing all subgroups is considered level zero, and lower levels are numbered consecutively.

Another aspect of the invention provides a communications system comprising an interrogator, and a plurality of wireless identification devices configured to communicate with the interrogator in a wireless fashion. The respective wireless identification devices have a unique identification number. The interrogator is configured to employ a tree search technique to determine the unique identification numbers of the different wireless identification devices so as to be able to establish communications between the interrogator and individual ones of the multiple wireless identification devices without collision by multiple wireless identification devices attempting to respond to the interrogator at the same time. The interrogator is configured to start the tree search at a selectable level of the search tree.

One aspect of the invention provides a radio frequency identification device comprising an integrated circuit including a receiver, a transmitter, and a microprocessor. In one embodiment, the integrated circuit is a monolithic single die single metal layer integrated circuit including the receiver, the transmitter, and the microprocessor. The device of this embodiment includes an active transponder, instead of a transponder which relies on magnetic coupling for power, and therefore has a much greater range.

In another aspect, an interrogator sends a signal to a plurality of RFID devices. The signal provides a bit string and indicates a memory range comprising multiple bit locations. RFID devices compare the bits stored in their respective memory ranges to the bit string to determine which of the RFID devices are chosen.

BRIEF DESCRIPTION OF THE DRAWINGS

Preferred embodiments of the invention are described below with reference to the following accompanying drawings.

FIG. 1 is a high level circuit schematic showing an interrogator and a radio frequency identification device embodying the invention.

FIG. 2 is a front view of a housing, in the form of a badge or card, supporting the circuit of FIG. 1 according to one embodiment the invention.

FIG. 3 is a front view of a housing supporting the circuit of FIG. 1 according to another embodiment of the invention.

FIG. 4 is a diagram illustrating a tree splitting sort method for establishing communication with a radio frequency identification device in a field of a plurality of such devices.

FIG. 5, is a diagram illustrating a modified tree splitting sort method for establishing communication with a radio frequency identification device in a field of a plurality of such devices.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENTS

This disclosure of the invention is submitted in furtherance of the constitutional purposes of the U.S. Patent Laws "to promote the progress of science and useful arts" (Article 1, Section 8).

FIG. 1 illustrates a wireless identification device 12 in accordance with one embodiment of the invention. In the illustrated embodiment, the wireless identification device is a radio frequency data communication device 12, and includes RFID circuitry 16. The device 12 further includes at least one antenna 14 connected to the circuitry 16 for wireless or radio frequency transmission and reception by the circuitry 16. In the illustrated embodiment, the RFID circuitry is defined by an integrated circuit as described in the above-incorporated patent application Ser. No. 08/705,043, filed Aug. 29, 1996, *now U.S. Pat. No. 6,130,602*. Other embodiments are possible. A power source or supply 18 is connected to the integrated circuit 16 to supply power to the integrated circuit 16. In one embodiment, the power source 18 comprises a battery.

The device 12 transmits and receives radio frequency communications to and from an interrogator 26. An exemplary interrogator is described in commonly assigned U.S. patent application Ser. No. 08/907,689, filed Aug. 8, 1997[and], *now U.S. Pat. No. 6,289,209, which is incorporated herein by reference.* Preferably, the interrogator 26 includes an antenna 28, as well as dedicated transmitting (e.g., *modulator*) and receiving circuitry, similar to that implemented on the integrated circuit 16.

Generally, the interrogator 26 transmits an interrogation signal or command 27 via the antenna 28. The device 12 receives the incoming interrogation signal via its antenna 14. Upon receiving the signal 27, the device 12 responds by generating and transmitting a responsive signal or reply 29. The responsive signal 29 typically includes information that uniquely identifies, or labels the particular device 12 that is transmitting, so as to identify any object or person with which the device 12 is associated.

Although only one device 12 is shown in FIG. 1, typically there will be multiple devices 12 that correspond with the interrogator 26, and the particular devices 12 that are in communication with the interrogator 26 will typically change over time. In the illustrated embodiment in FIG. 1, there is no communication between multiple devices 12. Instead, the

devices **12** respectively communicate with the interrogator **26**. Multiple devices **12** can be used in the same field of an interrogator **26** (i.e., within communications range of an interrogator **26**).

The radio frequency data communication device **12** can be included in any appropriate housing or packaging. Various methods of manufacturing housings are described in commonly assigned U.S. patent application Ser. No. 08/800,037, filed Feb. 13, 1997, [and] *now U.S. Pat. No. 5,988,510, which is incorporated herein by reference.*

FIG. **2** shows but one embodiment in the form of a card or badge **19** including a housing **11** of plastic or other suitable material supporting the device **12** and the power supply **18**. In one embodiment, the front face of the badge has visual identification features such as graphics, text, information found on identification or credit cards, etc.

FIG. **3** illustrates but one alternative housing supporting the device **12**. More particularly, FIG. **3** shows a miniature housing **20** encasing the device **12** and power supply **18** to define a tag which can be supported by an object (e.g., hung from an object, affixed to an object, etc.). Although two particular types of housings have been disclosed, the device **12** can be included in any appropriate housing.

If the power supply **18** is a battery, the battery can take any suitable form. Preferably, the battery type will be selected depending on weight, size, and life requirements for a particular application. In one embodiment, the battery **18** is a thin profile button-type cell forming a small, thin energy cell more commonly utilized in watches and small electronic devices requiring a thin profile. A conventional button-type cell has a pair of electrodes, an anode formed by one face and a cathode formed by an opposite face. In an alternative embodiment, the power source **18** comprises a series connected pair of button type cells. Instead of using a battery, any suitable power source can be employed.

The circuitry **16** further includes a backscatter transmitter and is configured to provide a responsive signal to the interrogator **26** by radio frequency. More particularly, the circuitry **16** includes a transmitter, a receiver, and memory such as is described in U.S. patent application Ser. No. 08/705, 043, *now U.S. Pat. No. 6,130,602.*

Radio frequency identification has emerged as a viable and affordable alternative to tagging or labeling small to large quantities of items. The interrogator **26** communicates with the devices **12** via an electromagnetic link, such as via an RF link (e.g., at microwave frequencies, in one embodiment), so all transmissions by the interrogator **26** are heard simultaneously by all devices **12** within range.

If the interrogator **26** sends out a command requesting that all devices **12** within range identify themselves, and gets a large number of simultaneous replies, the interrogator **26** may not be able to interpret any of these replies. Therefore, arbitration schemes are provided.

If the interrogator **26** has prior knowledge of the identification number of a device **12** which the interrogator **26** is looking for, it can specify that a response is requested only from the device **12** with that identification number. To target a command at a specific device **12**, (i.e., to initiate point-on-point communication), the interrogator **26** must send a number, identifying a specific device **12** along with the command. At start-up, or in a new or changing environment, these identification numbers are not known by the interrogator **26**. Therefore, the interrogator **26** must identify all devices **12** in the field (within communication range) such as by determining the identification numbers of the devices **12** in the field. After this is accomplished, point-to-point communication can proceed as desired by the interrogator **26**.

Generally speaking, RFID systems are a type of multi-access communication system. The distance between the interrogator **26** and devices **12** within the field is typically fairly short (e.g., several meters), so packet transmission time is determined primarily by packet size and baud rate. Propagation delays are negligible. In such systems, there is a potential for a large number of transmitting devices **12** and there is a need for the interrogator **26** to work in a changing environment, where different devices **12** are swapped in and out frequently (e.g., as inventory is added or removed). In such systems, the inventors have determined that the use of random access methods work effectively for contention resolution (i.e., for dealing with collisions between devices **12** attempting to respond to the interrogator **26** at the same time).

RFID systems have some characteristics that are different from other communications systems. For example, one characteristic of the illustrated RFID systems is that the devices **12** never communicate without being prompted by the interrogator **26**. This is in contrast to typical multiaccess systems where the transmitting units operate more independently. In addition, contention for the communication medium is short lived as compared to the ongoing nature of the problem in other multiaccess systems. For example, in a RFID system, after the devices **12** have been identified, the interrogator can communicate with them in a point-to-point fashion. Thus, arbitration in a RFID system is a transient rather than steady-state phenomenon. Further, the capability of a device **12** is limited by practical restrictions on size, power, and cost. The lifetime of a device **12** can often be measured in terms of number of transmissions before battery power is lost. Therefore, one of the most important measures of system performance in RFID arbitration is total time required to arbitrate a set of devices **12**. Another measure is power consumed by the devices **12** during the process. This is in contrast to the measures of throughput and packet delay in other types of multiaccess systems.

FIG. **4** illustrates one arbitration scheme that can be employed for communication between the interrogator and devices **12**. Generally, the interrogator **26** sends a command causing each device **12** of a potentially large number of responding devices **12** to select a random number from a known range and use it as that device's arbitration number. By transmitting requests for identification to various subsets of the full range of arbitration numbers, and checking for an error-free response, the interrogator **26** determines the arbitration number of every responder station capable of communicating at the same time. Therefore, the interrogator **26** is able to conduct subsequent uninterrupted communication with devices **12**, one at a time, by addressing only one device **12**.

Three variables are used: an arbitration value (AVALUE), an arbitration mask (AMASK), and a random value ID (RV). The interrogator sends an Identify command (IdentifyCmnd) causing each device of a potentially large number of responding devices to select a random number from a known range and use it as that device's arbitration number. The interrogator sends an arbitration value (AVALUE) and an arbitration mask (AMASK) to a set of devices **12**. The receiving devices **12** evaluate the following equation: $(AMASK \& AVALUE) == (AMASK \& RV)$ wherein "&" is a bitwise AND function, and wherein "=" is an equality function. If the equation evaluates to "1" (TRUE), then the device **12** will reply. If the equation evaluates to "0" (FALSE), then the device **12** will not reply. By performing this in a structured manner, with the number of bits in the arbitration mask being increased by one each

time, eventually a device **12** will respond with no collisions. Thus, a binary search tree methodology is employed.

An example using actual numbers will now be provided using only four bits, for simplicity, reference being made to FIG. 4. In one embodiment, sixteen bits are used for AVALUE and AMASK. Other numbers of bits can also be employed depending, for example, on the number of devices **12** expected to be encountered in a particular application, on desired cost points, etc.

Assume, for this example, that there are two devices **12** in the field, one with a random value (RV) of 1100 (binary), and another with a random value (RV) of 1010 (binary). The interrogator is trying to establish communications without collisions being caused by the two devices **12** attempting to communicate at the same time.

The interrogator sets AVALUE to 0000 (or "don't care" for all bits, as indicated by the character "X" in FIG. 4) and AMASK to 0000. The interrogator transmits a command to all devices **12** requesting that they identify themselves. Each of the devices **12** evaluate $(AMASK \& AVALUE) = (AMASK \& RV)$ using the random value RV that the respective devices **12** selected. If the equation evaluates to "1" (TRUE), then the device **12** will reply. If the equation evaluates to "0" (FALSE), then the device **12** will not reply. In the first level of the illustrated tree, AMASK is 0000 and anything bitwise ANDed with all zeros results in all zeros, so both the devices **12** in the field respond, and there is a collision.

Next, the interrogator sets AMASK to 0001 and AVALUE to 0000 and transmits an identify command. Both devices **12** in the field have a zero for their least significant bit, and $(AMASK \& AVALUE) = (AMASK \& RV)$ will be true for both devices **12**. For the device **12** with a random value of 1100, the left side of the equation is evaluated as follows $(0001 \& 0000) = 0000$. The right side is evaluated as $(0001 \& 1100) = 0000$. The left side equals the right side, so the equation is true for the device **12** with the random value of 1100. For the device **12** with a random value of 1010, the left side of the equation is evaluated as $(0001 \& 0000) = 0000$. The right side is evaluated as $(0001 \& 1010) = 0000$. The left side equals the right side, so the equation is true for the device **12** with the random value of 1010. Because the equation is true for both devices **12** in the field, both devices **12** in the field respond, and there is another collision.

Recursively, the interrogator next sets AMASK to 0011 with AVALUE still at 0000 and transmits an Identify command. $(AMASK \& AVALUE) = (AMASK \& RV)$ is evaluated for both devices **12**. For the device **12** with a random value of 1100, the left side of the equation is evaluated as follows $(0011 \& 0000) = 0000$. The right side is evaluated as $(0011 \& 1100) = 0000$. The left side equals the right side, so the equation is true for the device **12** with the random value of 1100, so this device **12** responds. For the device **12** with a random value of 1010, the left side of the equation is evaluated as $(0011 \& 0000) = 0000$. The right side is evaluated as $(0011 \& 1010) = 0010$. The left side does not equal the right side, so the equation is false for the device **12** with the random value of 1010, and this device **12** does not respond. Therefore, there is no collision, and the interrogator can determine the identity (e.g., an identification number) for the device **12** that does respond.

De-recursion takes place, and the devices **12** to the right for the same AMASK level are accessed when AVALUE is set at 0010, and AMASK is set to 0011.

The device **12** with the random value of 1010 receives a command and evaluates the equation $(AMASK \& AVALUE)$

$= (AMASK \& RV)$. The left side of the equation is evaluated as $(0011 \& 0010) = 0010$. The right side of the equation is evaluated as $(0011 \& 1010) = 0010$. The right side equals the left side, so the equation is true for the device **12** with the random value of 1010. Because there are no other devices **12** in the subtree, a good reply is returned by the device **12** with the random value of 1010. There is no collision, and the interrogator **26** can determine the identity (e.g., an identification number) for the device **12** that does respond.

By recursion, what is meant is that a function makes a call to itself. In other words, the function calls itself within the body of the function. After the called function returns, de-recursion takes place and execution continues at the place just after the function call: i.e. at the beginning of the statement after the function call.

For instance, consider a function that has four statements (numbered **1,2,3,4**) in it, and the second statement is a recursive call. Assume that the fourth statement is a return statement. The first time through the loop (iteration **1**) the function executes the statement **2** and (because it is a recursive call) calls itself causing iteration **2** to occur. When iteration **2** gets to statement **2**, it calls itself making iteration **3**. During execution in iteration **3** of statement **1**, assume that the function does a return. The information that was saved on the stack from iteration **2** is loaded and the function resumes execution at statement **3** (in iteration **2**), followed by the execution of statement **4** which is also a return statement. Since there are no more statements in the function, the function de-recurses to iteration **1**. Iteration **1**, had previously recursively called itself in statement **2**. Therefore, it now executes statement **3** (in iteration **1**). Following that it executes a return at statement **4**. Recursion is known in the art.

Consider the following code which can be used to implement operation of the method shown in FIG. 4 and described above.

```

40 Arbitrate(AMASK, AVALUE)
   {
   collision=IdentifyCmnd(AMASK,AVALUE)
   if (collision) then
       {
45             /* recursive call for left side */
               Arbitrate((AMASK>>1)+1, AVALUE)
               /* recursive call for right side */
               Arbitrate((AMASK>>1)+1, AVALUE+(AMASK+1))
       } /* endif */
   } /* return */

```

The symbol "<<" represents a bitwise left shift. "<<" means shift left by one place. Thus, 0001<<1 would be 0010. Note, however, that AMASK is originally called with a value of zero, and 0000<<1 is still 0000. Therefore, for the first recursive call, $AMASK = (AMASK \ll 1) + 1$. So for the first recursive call, the value of AMASK is $0000 + 0001 = 0001$. For the second call, $AMASK = (0001 \ll 1) + 1 = 0010 + 1 = 0011$. For the third recursive call, $AMASK = (0011 \ll 1) + 1 = 0110 + 1 = 0111$.

The routine generates values for AMASK and AVALUE to be used by the interrogator in an identify command "IdentifyCmnd." Note that the routine calls itself if there is a collision. De-recursion occurs when there is no collision. AVALUE and AMASK would have values such as the following assuming collisions take place all the way down to the bottom of the tree.

AVALUE	AMASK
0000	0000
0000	0001
0000	0011
0000	0111
0000	1111*
1000	1111*
0100	0111
0100	1111*
1100	1111*

This sequence of AMASK, AVALUE binary numbers assumes that there are collisions all the way down to the bottom of the tree, at which point the Identify command sent by the interrogator is finally successful so that no collision occurs. Rows in the table for which the interrogator is successful in receiving a reply without collision are marked with the symbol “*”. Note that if the Identify command was successful at, for example, the third line in the table then the interrogator would stop going down that branch of the tree and start down another, so the sequence would be as shown in the following table.

AVALUE	AMASK
0000	0000
0000	0001
0000	0011*
0010	0011
...	...

This method is referred to as a splitting method. It works by splitting groups of colliding devices **12** into subsets that are resolved in turn. The splitting method can also be viewed as a type of tree search. Each split moves the method one level deeper in the tree.

Either depth-first or breadth-first traversals of the tree can be employed. Depth first traversals are performed by using recursion, as is employed in the code listed above. Breadth-first traversals are accomplished by using a queue instead of recursion. The following is an example of code for performing a breadth-first traversal.

```

Arbitrate(AMASK, AVALUE)
{
  enqueue(0,0)
  while (queue != empty)
    (AMASK, AVALUE) = dequeue()
    collision=IdentifyCmnd(AMASK,AVALUE)
    if (collision) then
      {
        TEMP = AMASK+1
        NEW_AMASK = (AMASK>>1)+1
        enqueue(NEW_AMASK, AVALUE)
        enqueue(NEW_AMASK, AVALUE+TEMP)
      } /* endif */
    endwhile
  } /* return */

```

The symbol “!=” means not equal to. AVALUE and AMASK would have values such as those indicated in the following table for such code.

AVALUE	AMASK
0000	0000
0000	0001
0001	0001
0000	0011
0010	0011
0001	0011
0011	0011
0000	0111
0100	0111
...	...

Rows in the table for which the interrogator is successful in receiving a reply without collision are marked with the symbol “*”.

FIG. 5 illustrates an embodiment wherein the interrogator **26** starts the tree search at a selectable level of the search tree. The search tree has a plurality of nodes **51, 52, 53, 54** etc. at respective levels. The size of subgroups of random values decrease in size by half with each node descended. The upper bound of the number of devices **12** in the field (the maximum possible number of devices that could communicate with the interrogator) is determined, and the tree search method is started at a level **32, 34, 36, 38, or 40** in the tree depending on the determined upper bound. In one embodiment, the maximum number of devices **12** potentially capable of responding to the interrogator is determined manually and input into the interrogator **26** via an input device such as a keyboard, graphical user interface, mouse, or other interface. The level of the search tree on which to start the tree search is selected based on the determined maximum possible number of wireless identification devices that could communicate with the interrogator.

The tree search is started at a level determined by taking the base two logarithm of the determined maximum possible number. More particularly, the tree search is started at a level determined by taking the base two logarithm of the power of two nearest the determined maximum possible number of devices **12**. The level of the tree containing all subgroups of random values is considered level zero (see FIG. 5), and lower levels are numbered **1, 2, 3, 4,** etc. consecutively.

By determining the upper bound of the number of devices **12** in the field, and starting the tree search at an appropriate level, the number of collisions is reduced, the battery life of the devices **12** is increased, and arbitration time is reduced.

For example, for the search tree shown in FIG. 5, if it is known that there are seven devices **12** in the field, starting at node **51** (level **0**) results in a collision. Starting at level **1** (nodes **52** and **53**) also results in a collision. The same is true for nodes **54, 55, 56,** and **57** in level **2**. If there are seven devices **12** in the field, the nearest power of two to seven is the level at which the tree search should be started. $\log_2 8=3$, so the tree search should be started at level **3** if there are seven devices **12** in the field.

AVALUE and AMASK would have values such as the following assuming collisions take place from level **3** all the way down to the bottom of the tree.

AVALUE	AMASK
0000	0111
0000	1111*

-continued

AVALUE	AMASK
1000	1111*
0100	0111
0100	1111*
1100	1111*

Rows in the table for which the interrogator is successful in receiving a reply without collision are marked with the symbol “*”.

In operation, the interrogator transmits a command requesting devices **12** having random values RV within a specified group of random values to respond, the specified group being chosen in response to the determined maximum number. Devices **12** receiving the command respectively determine if their chosen random values fall within the specified group and, if so, send a reply to the interrogator. The interrogator determines if a collision occurred between devices that sent a reply and, if so, creates a new, smaller, specified group, descending in the tree, as described above in connection with FIG. 4.

Another arbitration method that can be employed is referred to as the “Aloha” method. In the Aloha method, every time a device **12** is involved in a collision, it waits a random period of time before retransmitting. This method can be improved by dividing time into equally sized slots and forcing transmissions to be aligned with one of these slots. This is referred to as “slotted Aloha.” In operation, the interrogator asks all devices **12** in the field to transmit their identification numbers in the next time slot. If the response is garbled, the interrogator informs the devices **12** that a collision has occurred, and the slotted Aloha scheme is put into action. This means that each device **12** in the field responds within an arbitrary slot determined by a randomly selected value. In other words, in each successive time slot, the devices **12** decide to transmit their identification number with a certain probability.

The Aloha method is based on a system operated by the University of Hawaii. In 1971, the University of Hawaii began operation of a system named Aloha. A communication satellite was used to interconnect several university computers by use of a random access protocol. The system operates as follows. Users or devices transmit at any time they desire. After transmitting, a user listens for an acknowledgment from the receiver or interrogator. Transmissions from different users will sometimes overlap in time (collide), causing reception errors in the data in each of the contending messages. The errors are detected by the receiver, and the receiver sends a negative acknowledgment to the users. When a negative acknowledgment is received, the messages are retransmitted by the colliding users after a random delay. If the colliding users attempted to retransmit without the random delay, they would collide again. If the user does not receive either an acknowledgment or a negative acknowledgment within a certain amount of time, the user “times out” and retransmits the message.

There is a scheme known as slotted Aloha which improves the Aloha scheme by requiring a small amount of coordination among stations. In the slotted Aloha scheme, a sequence of coordination pulses is broadcast to all stations (devices). As is the case with the pure Aloha scheme, packet lengths are constant. Messages are required to be sent in a slot time between synchronization pulses, and can be started only at the beginning of a time slot. This reduces the rate of colli-

sions because only messages transmitted in the same slot can interfere with one another. The retransmission mode of the pure Aloha scheme is modified for slotted Aloha such that if a negative acknowledgment occurs, the device retransmits after a random delay of an integer number of slot times.

Aloha methods are described in [a] commonly assigned patent application [naming Clifton W. Wood, Jr. as an inventor, U.S. patent application] Ser. No. 09/026,248, filed Feb. 19, 1998, [titled “Method of Addressing Messages and Communications System,” filed concurrently herewith, and], now U.S. Pat. No. 6,275,476, which is incorporated herein by reference.

In one alternative embodiment, an Aloha method (such as the method described in the commonly assigned patent application mentioned above) is combined with determining the upper bound on a set of devices and starting at a level in the tree depending on the determined upper bound, such as by combining an Aloha method with the method shown and described in connection with FIG. 5. For example, in one embodiment, devices **12** sending a reply to the interrogator **26** do so within a randomly selected time slot of a number of slots.

In another embodiment, levels of the search tree are skipped. Skipping levels in the tree, after a collision caused by multiple devices **12** responding, reduces the number of subsequent collisions without adding significantly to the number of no replies. In real-time systems, it is desirable to have quick arbitration sessions on a set of devices **12** whose unique identification numbers are unknown. Level skipping reduces the number of collisions, both reducing arbitration time and conserving battery life on a set of devices **12**. In one embodiment, every other level is skipped. In alternative embodiments, more than one level is skipped each time.

The trade off that must be considered in determining how many (if any) levels to skip with each decent down the tree is as follows. Skipping levels reduces the number of collisions, thus saving battery power in the devices **12**. Skipping deeper (skipping more than one level) further reduces the number of collisions. The more levels that are skipped, the greater the reduction in collisions. However, skipping levels results in longer search times because the number of queries (Identify commands) increases. The more levels that are skipped, the longer the search times. Skipping just one level has an almost negligible effect on search time, but drastically reduces the number of collisions. If more than one level is skipped, search time increases substantially. Skipping every other level drastically reduces the number of collisions and saves battery power without significantly increasing the number of queries.

Level skipping methods are described in a commonly assigned patent application 09/026,045 naming Clifton W. Wood, Jr. and Don Hush as inventors, titled “Method of Addressing Messages, Method of Establishing Wireless Communications, and Communications Systems,” filed concurrently herewith, now U.S. Pat. No. 6,072,801, and incorporated herein by reference.

In one alternative embodiment, a level skipping method is combined with determining the upper bound on a set of devices and starting at a level in the tree depending on the determined upper bound, such as by combining a level skipping method with the method shown and described in connection with FIG. 5.

In yet another alternative embodiment, both a level skipping method and an Aloha method (as described in the commonly assigned applications described above) are combined with the method shown and described in connection with FIG. 5.

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In compliance with the statute, the invention has been described in language more or less specific as to structural and methodical features. It is to be understood, however, that the invention is not limited to the specific features shown and described, since the means herein disclosed comprise preferred forms of putting the invention into effect. The invention is, therefore, claimed in any of its forms or modifications within the proper scope of the appended claims appropriately interpreted in accordance with the doctrine of equivalents.

What is claimed is:

1. A method of establishing wireless communications between an interrogator and individual ones of multiple wireless identification devices, the wireless identification devices having respective identification numbers and being addressable by specifying identification numbers with any one of multiple possible degrees of precision, the method comprising utilizing a tree search in an arbitration scheme to determine a degree of precision necessary to establish one-to-one communications between the interrogator and individual ones of the multiple wireless identification devices, a search tree being defined for the tree search method, the tree having multiple selectable levels respectively representing subgroups of the multiple wireless identification devices, the level at which a tree search starts being variable the method further comprising starting the tree search at any [selectable level of the search tree] *selectable level other than the top level of the search tree.*

2. A method in accordance with claim 1 and further comprising determining the maximum possible number of wireless identification devices that could communicate with the interrogator, and selecting a level of the search tree based of the determined maximum possible number of wireless identification devices that could communicate with the interrogator.

3. A method in accordance with claim 2 and further comprising starting the tree search at a level determined by taking the base two logarithm of the determined maximum possible number, wherein the level of the tree containing all subgroups is considered level zero, and lower levels are numbered consecutively.

4. A method in accordance with claim 2 and further comprising starting the tree search at a level determined by taking the base two logarithm of the determined maximum possible number, wherein the level of the tree containing all subgroups is considered level zero, and lower levels are numbered consecutively, and wherein the maximum number of devices in a subgroup in one level is half of the maximum number of devices in the next higher level.

5. A method in accordance with claim 2 and further comprising starting the tree search at a level determined by taking the base two logarithm of the power of two nearest the determined maximum possible number, wherein the level of the tree containing all subgroups is considered level zero, and lower levels are numbered consecutively, and wherein the maximum number of devices in a subgroup in one level is half of the maximum number of devices in the next higher level.

6. A method in accordance with claim 1 wherein the wireless identification device comprises an integrated circuit including a receiver, a modulator, and a microprocessor in communication with the receiver and modulator.

7. A method of addressing messages from an interrogator to a selected one or more of a number of communications device, the method comprising:

establishing for respective devices unique identification numbers respectively having a first predetermined number of bits;

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establishing a second predetermined number of bits to be used for random values;

causing the devices to select random values, wherein respective devices choose random values independently of random values selected by the other devices;

determining the maximum number of devices potentially capable of responding to the interrogator;

transmitting a command from the interrogator requesting devices having random values within a specified group of random values to respond, by using a subset of the second predetermined number of bits, the specified group being chosen in response to the determined maximum number;

receiving the command at multiple devices, devices receiving the command respectively determining if the random value chosen by the device falls within the specified group and, if so, sending a reply to the interrogator; and

determining using the interrogator if a collision occurred between devices that sent a reply and, if so, creating a new, smaller, specified group.

8. A method of addressing messages from an interrogator to a selected one or more of a number of communications devices in accordance with claim 7 wherein sending a reply to the interrogator comprises transmitting the unique identification number of the device sending the reply.

9. A method of addressing messages from an interrogator to a selected one or more of a number of communications devices in accordance with claim 7 wherein sending a reply to the interrogator comprises transmitting the random value of the device sending the reply.

10. A method of addressing messages from an interrogator to a selected one or more of a number of communications devices in accordance with claim 7 wherein sending a reply to the interrogator comprises transmitting both the random value of the device sending the reply and the unique identification number of the device sending the reply.

11. A method of addressing messages from an interrogator to a selected one or more of a number of communications devices in accordance with claim 7 wherein, after receiving a reply without collision from a device, the interrogator sends a command individually addressed to that device.

12. A method of addressing messages from an interrogator to a selected one or more of a number of communications devices, the method comprising:

causing the devices to select random values for use as arbitration numbers, wherein respective devices choose random values independently of random values selected by the other devices, the devices being addressable by specifying arbitration numbers with any one of multiple possible degrees of precision;

transmitting a command from the interrogator requesting devices having random values within a specified group of a plurality of possible groups of random values to respond, the specified group being less than the entire set of random values, the plurality of possible groups being organized in a binary tree defined by a plurality of nodes at respective levels, wherein the size of groups of random values decrease in size by half with each node descended, wherein the specified group is below a node on the tree selected based on the maximum number of devices capable of communicating with the interrogator;

receiving the command at multiple devices, devices receiving the command respectively determining if the random value chosen by the device falls within the

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specified group and, if so, sending a reply to the interrogator; and, if not, not sending a reply; and

determining using the interrogator if a collision occurred between devices that sent a reply and, if so, creating a new, smaller, specified group by descending in the tree.

13. A method of addressing messages from an interrogator to a selected one or more of a number of communications devices in accordance with claim 12 and further including establishing a predetermined number of bits to be used for the random values.

14. A method of addressing messages from an interrogator to a selected one or more of a number of communications devices in accordance with claim 13 wherein the predetermined number of bits to be used for the random values comprises an integer multiple of eight.

15. A method of addressing messages from an interrogator to a selected one or more of a number of communications devices in accordance with claim 13 wherein devices sending a reply to the interrogator do so within a randomly selected time slot of a number of slots.

16. A method of addressing messages from an interrogator to a selected one or more of a number of RFID devices, the method comprising:

establishing for respective devices a predetermined number of bits to be used for random values, the predetermined number being a multiple of sixteen;

causing the devices to select random values, wherein respective devices choose random values independently of random values selected by the other devices;

transmitting a command from the interrogator requesting devices having random values within a specified group of a plurality of possible groups of random values to respond, the specified group being equal to or less than the entire set of random values, the plurality of possible groups being organized in a binary tree defined by a plurality of nodes at respective levels, wherein the maximum size of groups of random values decrease in size by half with each node descended, wherein the specified group is below a node on a level of the tree selected based on the maximum number of devices known to be capable of communicating with the interrogator;

receiving the command at multiple devices, devices receiving the command respectively determining if the random value chosen by the device falls within the specified group and, only if so, sending a reply to the interrogator, wherein sending a reply to the interrogator comprises transmitting both the random value of the device sending the reply and [the unique] a unique identification number of the device sending the reply;

using the interrogator to determine if a collision occurred between devices that sent a reply and, if so, creating a new, smaller, specified group using a level of the tree different from the level used in the interrogator transmitting, the interrogator transmitting a command requesting devices having random values within the new specified group of random values to respond; and if a reply without collision is received from a device, the interrogator subsequently sending a command individually addressed to that device.

17. A method of addressing messages from an interrogator to a selected one or more of a number of RFID devices in accordance with claim 16 and further comprising determining the maximum possible number of wireless identification devices that could communicate with the interrogator.

18. A method of addressing messages from an interrogator to a selected one or more of a number of RFID devices in

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accordance with [claim 16] claim 17 wherein selecting the level of the tree comprises taking the base two logarithm of the determined maximum possible number, wherein a level of the tree containing all subgroups is considered level zero, and lower levels are numbered consecutively.

19. A method of addressing messages from an interrogator to a selected one or more of a number of RFID devices in accordance with [claim 16] claim 17 wherein selecting the level of the tree comprises taking the base two logarithm of the determined maximum possible number, wherein a level of the tree containing all subgroups is considered level zero, and lower levels are numbered consecutively, and wherein the maximum number of devices in a subgroup in one level is half of the maximum number of devices in the next higher level.

20. A method of addressing messages from an interrogator to a selected one or more of a number of RFID devices in accordance with [claim 16] claim 17 wherein selecting the level of the tree comprises taking the base two logarithm of the power of two nearest the determined maximum possible number, wherein the level of the tree containing all subgroups is considered level zero, and lower levels are numbered consecutively, and wherein the maximum number of devices in a subgroup in one level is half of the maximum number of devices in the next higher level.

21. A method of addressing messages from an interrogator to a selected one or more of a number of RFID devices in accordance with claim 16 wherein the wireless identification device comprises an integrated circuit including a receiver, a modulator, and a microprocessor in communication with the receiver and modulator.

22. A method of addressing messages from an interrogator to a selected one or more of a number of RFID devices in accordance with claim 16 and further comprising, after the interrogator transmits a command requesting devices having random values within the new specified group of random values to respond, determining, using devices receiving the command, if their chosen random values fall within the new smaller specified group and, if so, sending a reply to the interrogator.

23. A method of addressing messages from an interrogator to a selected one or more of a number of RFID devices in accordance with claim 22 and further comprising, after the interrogator transmits a command requesting devices having random values within the new specified group of random values to respond, determining if a collision occurred between devices that sent a reply and, if so, creating a new specified group and repeating the transmitting of the command requesting devices having random values within a specified group of random values to respond using different specified groups until all of the devices within communications range are identified.

24. A communications system comprising an interrogator, and a plurality of wireless identification devices configured to communicate with the interrogator in a wireless fashion, the wireless identification devices having respective identification numbers, the interrogator being configured to employ a tree search in a search tree having multiple selectable levels, to determine the identification numbers of the [different] wireless identification devices with sufficient precision so as to be able to establish one-on-one communications between the interrogator and individual ones of the [multiple] wireless identification devices, wherein the interrogator is configured to start the tree search at any [selectable level of the search tree] selectable level other than the top level of the search tree.

25. A communications system in accordance with claim 24 wherein the tree search is a binary tree search.

26. A communications system in accordance with claim 24 wherein the wireless identification device comprises an integrated circuit including a receiver, a modulator, and a microprocessor in communication with the receiver and modulator.

27. A system comprising:

an interrogator;

a number of communications devices capable of wireless communications with the interrogator;

means for establishing a predetermined number of bits to be used as random numbers, and for causing respective devices to select random numbers respectively having the predetermined number of bits;

means for inputting a predetermined number indicative of the maximum number of devices possibly capable of communicating with the [receiver] interrogator;

means for causing the interrogator to transmit a command requesting devices having random values within a specified group of random values to respond, the specified group being chosen in response to the inputted predetermined number;

means for causing devices receiving the command to determine if their chosen random values fall within the specified group and, if so, send a reply to the interrogator; and

means for causing the interrogator to determine if a collision occurred between devices that sent a reply and, if so, create a new, smaller, specified group.

28. A system in accordance with claim 27 wherein sending a reply to the interrogator comprises transmitting the random value of the device sending the reply.

29. A system in accordance with claim 27 wherein the interrogator further includes means for, after receiving a reply without collision from a device, sending a command individually addressed to that device.

30. A system comprising:

an interrogator configured to communicate to a selected one or more of a number of communications devices;

a plurality of communications devices;

the devices being configured to select random values, wherein respective devices choose random values independently of random values selected by the other devices, different sized groups of devices being addressable by specifying random values with differing levels of precision;

the interrogator being configured to transmit a command requesting devices having random values within a specified group of a plurality of possible groups of random values to respond, the specified group being less than the entire set of random values, the plurality of possible groups being organized in a binary tree defined by a plurality of nodes at respective levels, wherein the size of groups of random values decrease in size by half with each node descended, wherein the specified group is below a node on the tree selected based on a predetermined maximum number of devices capable of communicating with the interrogator;

devices receiving the command being configured to respectively determine if their chosen random values fall within the specified group and, if so, send a reply to the interrogator; and, if not, not send a reply; and

the interrogator being configured to determine if a collision occurred between devices that sent a reply and, if so, create a new, smaller, specified group by descending in the tree.

31. A system in accordance with claim 30 wherein the random values respectively have a predetermined number of bits.

32. A system in accordance with claim 30 wherein respective devices are configured to store unique identification numbers of a predetermined number of bits.

33. A system in accordance with claim 30 wherein respective devices are configured to store unique identification numbers of sixteen bits.

34. A system comprising:

an interrogator configured to communicate to a selected one or more of a number of RFID devices;

a plurality of RFID devices, respective devices being configured to store unique identification numbers respectively having a first predetermined number of bits, respective devices being further configured to store a second predetermined number of bits to be used for random values, respective devices being configured to select random values independently of random values selected by the other devices;

the interrogator being configured to transmit an identify command requesting a response from devices having random values within a specified group of a plurality of possible groups [or random] of random values, the specified group being less than or equal to the entire set of random values, the plurality of possible groups being organized in a binary tree defined by a plurality of nodes at respective levels, wherein the maximum size of groups of random values decrease in size by half with each node descended, wherein the specified group is below a node on a level of the tree selected based on the maximum number of devices known to be capable of communicating with the interrogator;

devices receiving the command respectively being configured to determine if their chosen random values fall within the specified group and, only if so, send a reply to the interrogator, wherein sending a reply to the interrogator comprises transmitting both the random value of the device sending the reply and the unique identification number of the device sending the reply;

the interrogator being configured to determine if a collision occurred between devices that sent a reply and, if so, create a new, smaller, specified group using a level of the tree different from the level used in previously transmitting an identify command, the interrogator transmitting an identify command requesting devices having random values within the new specified group of random values to respond; and

the interrogator being configured to send a command individually addressed to a device after communicating with a device without a collision.

35. A system in accordance with 34 wherein the interrogator is configured to input and store [the predetermined number] a number representing the specified group.

36. A system in accordance with 34 wherein the devices are configured to respectively determine if their chosen random values fall within a specified group and, if so, send a reply, upon receiving respective identify commands.

37. A system in accordance with claim 36 wherein the interrogator is configured to determine if a collision occurred between devices that sent a reply in response to respective identify commands and, if so, create further new specified groups and repeat the transmitting of the identify command requesting devices having random values within a specified group of random values to respond using different specified groups until all responding devices are identified.

38. A method comprising:
 disposing a plurality of radio frequency identification (RFID) tags in a communication field of an interrogator, each respective tag of the plurality of tags including respective memory storing a respective identification code that identifies a respective object to which each respective tag is affixed;
 sending a select command from the interrogator to the plurality of tags after disposing the plurality of tags in the field and before any of the plurality of tags communicate to the interrogator, the select command including a set of parameters, the set of parameters including a bit string and describing a memory range, the memory range comprising multiple bit locations;
 each respective tag of the plurality of tags receiving the select command and comparing the bit string against the memory range of the respective memory of each respective tag to determine if the respective tag is a member of a population of tags;
 each respective tag of the population picking a respective random value and associating the random value with a respective slot, wherein a sequence in which the population of tags are to reply to the interrogator is determined by each respective slot;
 each respective tag of at least a portion of the population backscattering a respective reply to the interrogator, each respective reply including a respective random number generated by each respective tag, each respective tag replying in accordance with the sequence; and
 sending an acknowledge command from the interrogator in response to the interrogator receiving a respective reply from a respective tag and determining the respective reply to be collision-free.

39. The method of claim 38, further comprising each respective tag of the at least a portion of the population backscattering at least a portion of the respective identification code.

40. The method of claim 39, further comprising the interrogator accessing a tag individually after receiving the random number from the tag, accessing the tag including the interrogator sending the random number to the tag.

41. The method of claim 40, wherein the memory range of the memory of the tag includes at least a portion of the random number.

42. The method of claim 38, wherein each respective random number generated by each respective tag is sixteen bits in length.

43. A method comprising:
 affixing a radio frequency identification (RFID) tag to an object, the tag including tag memory;
 disposing the tag in a communication field of an interrogator;
 sending a first signal from the interrogator to the tag after disposing the tag in the field and before the tag communicates to the interrogator, the first signal including parameters that describe a memory range and a bit string;
 receiving the first signal at the tag, and in response thereto, comparing the bit string against the memory range of the tag memory to determine if the tag is selected, the memory range of the tag memory storing a plurality of bits;
 the tag picking a random value and associating the random value with a slot in accordance with an arbitration scheme for an inventory operation if the tag is determined to be selected;

sending a second signal from the interrogator to the tag;
 backscattering a random number generated by the tag from the tag to the interrogator in accordance with the slot in response to receiving the second signal; and
 sending an acknowledge command from the interrogator to the tag in response to the interrogator receiving the random number.

44. The method of claim 43, further comprising backscattering at least a portion of an identification code from the tag to the interrogator, wherein the identification code is stored in tag memory and identifies the object.

45. The method of claim 44, further comprising the interrogator individually accessing the tag after the interrogator sends the acknowledge command and receives the at least a portion of the identification code, wherein individually accessing the tag includes the interrogator sending an access command to the tag, the access command including the random number.

46. The method of claim 45, wherein the random number is sixteen bits long.

47. The method of claim 43, wherein the plurality of bits includes at least a portion of the random number.

48. A method comprising:
 disposing a radio frequency identification (RFID) tag in a communication field of an interrogator, the tag including tag memory;
 sending a select command from the interrogator to the tag after disposing the tag in the field and before the tag communicates to the interrogator, the select command including parameters that describe a memory range and a bit string;
 receiving the select command at the tag, and in response thereto, comparing the bit string against the memory range of the tag memory to determine if the tag is selected, the memory range of the tag memory storing at least two bits; and
 communicating a random number generated by the tag from the tag to the interrogator in accordance with an arbitration scheme if the tag is determined to be selected.

49. The method of claim 48, wherein the random number is stored in the tag memory.

50. The method of claim 48, wherein the at least two bits include at least a portion of the random number.

51. The method of claim 48, further comprising communicating at least a portion of an identification code from the tag to the interrogator in accordance with the arbitration scheme, wherein the identification code identifies an object to which the tag is affixed.

52. The method of claim 51, wherein the identification code is stored in the tag memory.

53. The method of claim 48, wherein the random number is sixteen bits long.

54. The method of claim 48, further comprising the tag picking a random value and using the random value as a slot, the tag communicating the random number at a time associated with the slot in accordance with the arbitration scheme.

55. The method of claim 54, further comprising sending an acknowledge command from the interrogator to the tag in response to the interrogator receiving the random number.

56. The method of claim 55, further comprising sending a signal from the interrogator to the tag, after sending the select command from the interrogator to the tag and before communicating the random number from the tag to the interrogator, wherein the signal indicates to the tag the time to communicate the random number.

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57. The method of claim 48, further comprising sending a signal from the interrogator to the tag, after sending the select command from the interrogator to the tag and before communicating the random number from the tag to the interrogator, wherein the signal indicates to the tag when to communicate the random number to the interrogator.

58. The method of claim 48, wherein communicating the random number includes backscattering the random number.

59. The method of claim 48, further comprising sending an acknowledge command from the interrogator to the tag in response to the interrogator receiving the random number.

60. The method of claim 59, further comprising communicating at least a portion of an identification code from the tag to the interrogator in accordance with the arbitration scheme, wherein the identification code identifies an object to which the tag is affixed.

61. The method of claim 60, further comprising the interrogator individually accessing the tag after receiving the random number, wherein individually accessing the tag includes the interrogator sending an access command to the tag, the access command including a sixteen bit random number.

62. The method of claim 61, wherein the sixteen bit random number is the random number generated by the tag and communicated from the tag to the interrogator in accordance with the arbitration scheme.

63. A method comprising:

disposing a plurality of radio frequency (RFID) tags in a communication field of an interrogator;

sending a first signal from the interrogator to first and second tags of the plurality of tags after disposing the plurality of tags in the field and before any of the plurality of tags communicate to the interrogator, the first signal including a bit string and indicating a portion of memory, the portion of memory comprising multiple bit storage locations, the first tag having stored therein a first set of bits in bit storage locations corresponding to the portion of memory, and the second tag having stored therein a second set of bits in bit storage locations corresponding to the portion of memory;

the first tag receiving the first signal and comparing the bit string against the first set of bits to determine if the first tag is selected;

the second tag receiving the first signal and comparing the bit string against the second set of bits to determine if the second tag is selected;

the first tag picking a first random value and associating the first random value with a first slot in accordance with an arbitration scheme;

the second tag picking a second random value and associating the second random value with a second slot in accordance with the arbitration scheme;

the first tag backscattering a first identification code that identifies a first object to which the first tag is affixed; and

the second tag backscattering a second identification code that identifies a second object to which the second tag is affixed.

64. The method of claim 63, further comprising:

the first tag backscattering a first random number generated by the first tag;

and the second tag backscattering a second random number generated by the second tag.

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65. The method of claim 64, further comprising:

the interrogator receiving the first random number from the first tag during a period of time associated with the first slot, and, in response thereto, the interrogator sending a first acknowledge signal to acknowledge the first tag; and

the interrogator receiving the second random number from the second tag during a period of time associated with the second slot, and, in response thereto, the interrogator sending a second acknowledge signal to acknowledge the second tag.

66. The method of claim 65, further comprising the interrogator accessing the first tag individually after receiving both the first random number and the first identification code from the first tag, accessing the first tag including the interrogator sending a command that includes a number randomly generated by the first tag that identifies the first tag.

67. The method of claim 66, wherein the number randomly generated by the first tag that identifies the first tag is the first random number, and the first random number is 16 bits in length.

68. The method of claim 67, further comprising sending a second signal from the interrogator after sending the first signal from the interrogator, the first tag backscattering the first identification code in response to receiving the second signal.

69. The method of claim 63, further comprising:

the interrogator sending a first acknowledge signal to acknowledge the first tag;

and the interrogator sending a second acknowledge signal to acknowledge the second tag.

70. A method comprising:

disposing a radio frequency identification (RFID) tag in a communication field of an interrogator;

sending a first command from the interrogator to the tag after disposing the tag in the field and before the tag communicates to the interrogator, the first command including a first set of fields comprising at least two first bit values;

the tag wirelessly receiving the first command;

the tag backscattering a first reply based, at least in part, on whether the two first bit values received from the interrogator match two corresponding bit values stored in the tag, the first reply including a random number generated by the tag;

sending a second command from the interrogator to the tag, the second command including a second set of fields comprising at least two second bit values;

the tag wirelessly receiving the second command; and

the tag backscattering a second reply based, at least in part, on whether the two second bit values received from the interrogator match the two corresponding bit values stored in the tag, the second reply including a random number generated by the tag.

71. The method of claim 70, further comprising backscattering at least a portion of an identification code from the tag to the interrogator, wherein the identification code identifies an object to which the tag is affixed.

72. The method of claim 70, further comprising the tag picking a random value and using the random value as a slot in accordance with an arbitration scheme, the tag backscattering a signal to the interrogator at a time associated with the slot.

73. The method of claim 72, further comprising sending an acknowledge command from the interrogator to the tag.

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74. The method of claim 70, further comprising the interrogator individually accessing the tag, wherein individually accessing the tag includes the interrogator sending an access command to the tag, the access command including a sixteen bit random number.

75. The method of claim 74, further comprising backscattering at least a portion of an identification code from the tag

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to the interrogator, wherein the identification code identifies an object to which the tag is affixed.

76. The method of claim 70, further comprising the interrogator detecting a collision upon receiving the first reply.

* * * * *

UNITED STATES PATENT AND TRADEMARK OFFICE
CERTIFICATE OF CORRECTION

PATENT NO. : RE41,530 E
APPLICATION NO. : 10/693696
DATED : August 17, 2010
INVENTOR(S) : Clifton W. Wood, Jr.

Page 1 of 1

It is certified that error appears in the above-identified patent and that said Letters Patent is hereby corrected as shown below:

TITLE PAGE, COLUMN 1 - SECTION (56) REFERENCES CITED

INSERT THE FOLLOWING

FOREIGN PATENT DOCUMENTS

-- JP 09-054213 2/1997

JP 2002/228809 8/2002 --

Signed and Sealed this
Fourth Day of January, 2011



David J. Kappos
Director of the United States Patent and Trademark Office