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(54) **FILTER FOR NETWORK INTRUSION AND VIRUS DETECTION**

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CPC **H04L 63/145** (2013.01); **G06F 7/02** (2013.01); **G06F 21/567** (2013.01); **H04L 63/0245** (2013.01); **H04L 63/1416** (2013.01); **G06F 2207/025** (2013.01)

(58) **Field of Classification Search**

USPC 726/23
See application file for complete search history.

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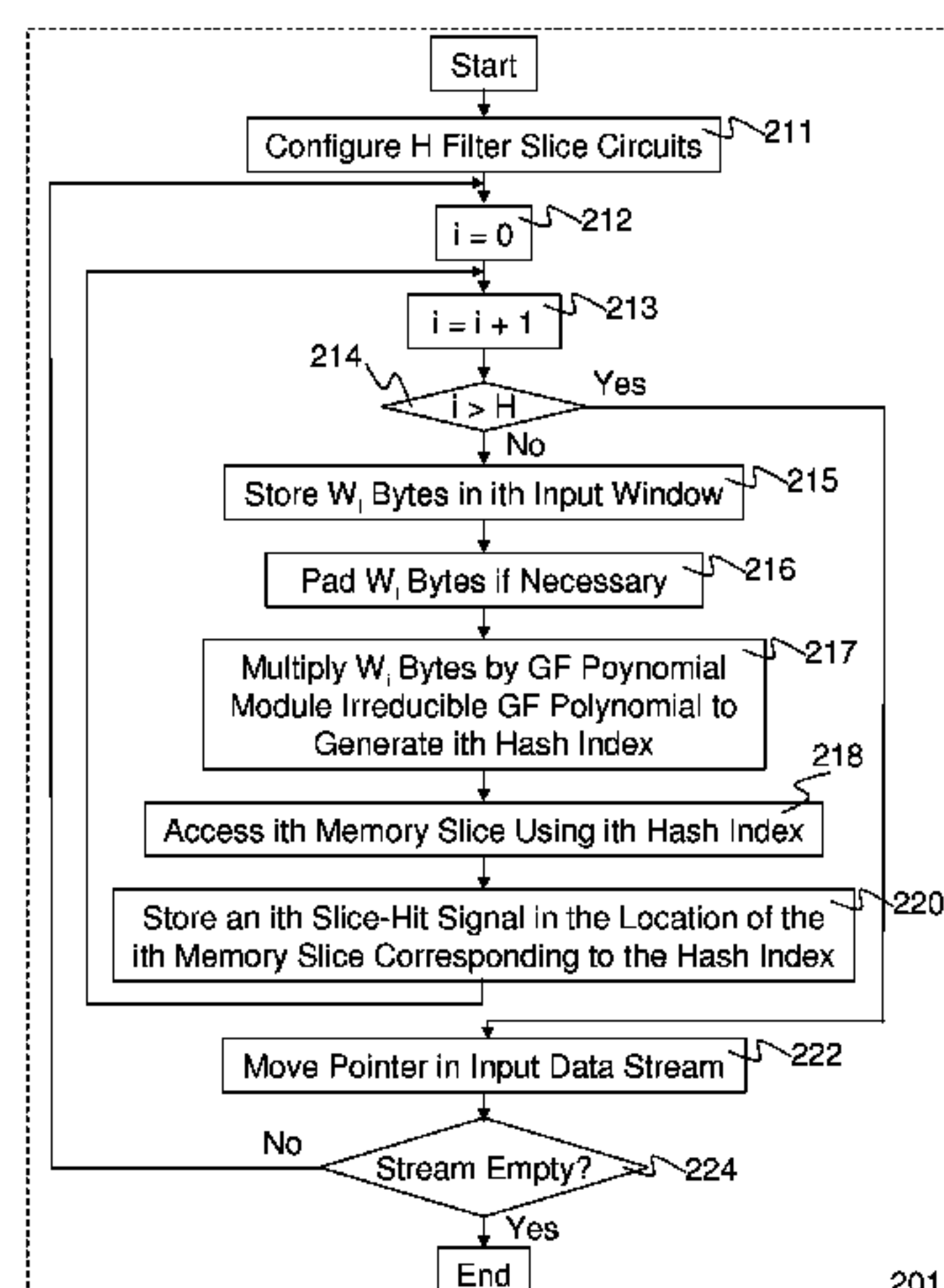
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(57) **ABSTRACT**

Methods and apparatus to perform string matching for network packet inspection are disclosed. In some embodiments there is a set of string matching slice circuits, each slice circuit of the set being configured to perform string matching steps in parallel with other slice circuits. Each slice circuit may include an input window storing some number of bytes of data from an input data stream. The input window of data may be padded if necessary, and then multiplied by a polynomial modulo an irreducible Galois-field polynomial to generate a hash index. A storage location of a memory corresponding to the hash index may be accessed to generate a slice-hit signal of a set of H slice-hit signals. The slice-hit signal may be provided to an AND-OR logic array where the set of H slice-hit signals is logically combined into a match result.

24 Claims, 4 Drawing Sheets



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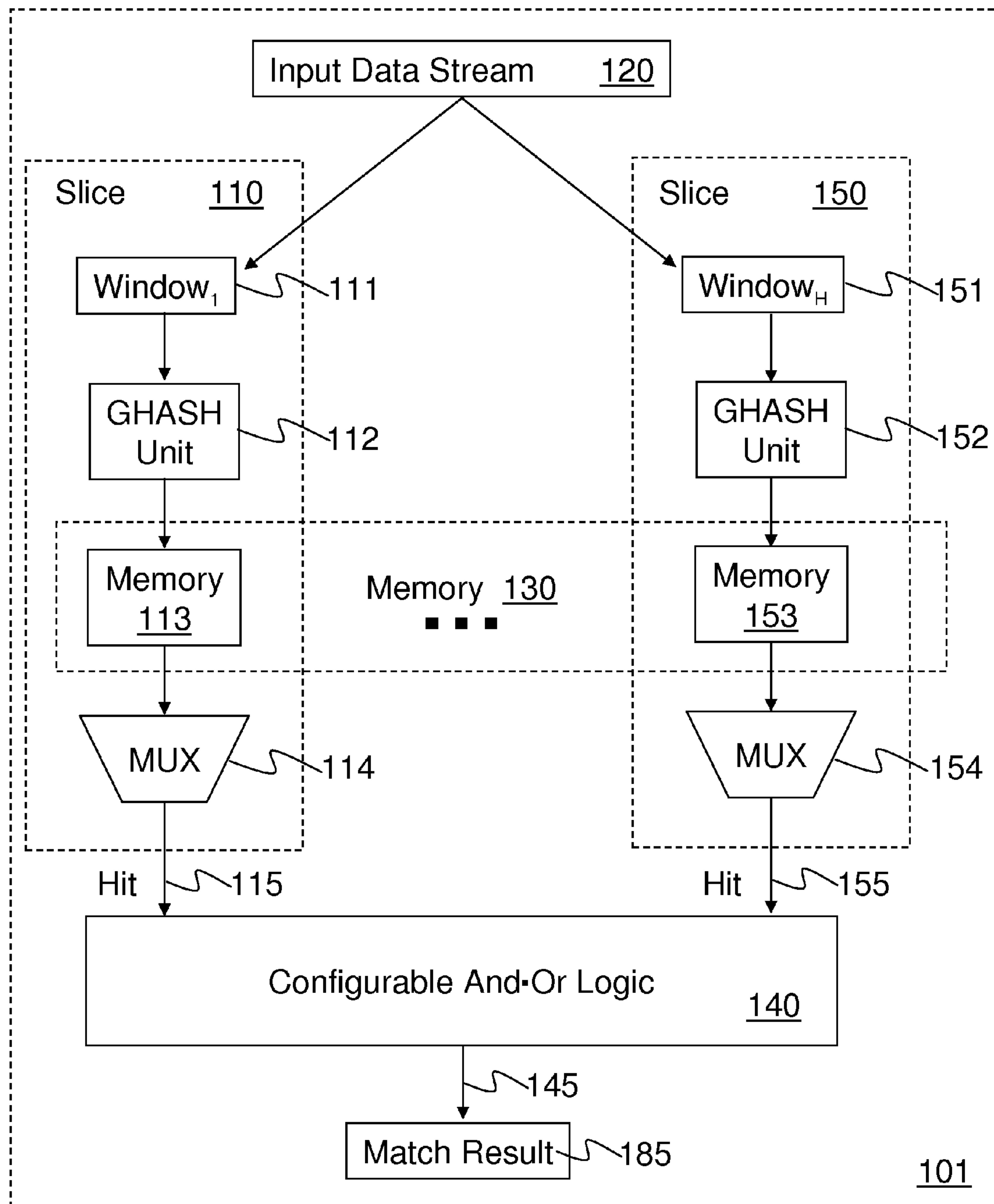


FIG. 1

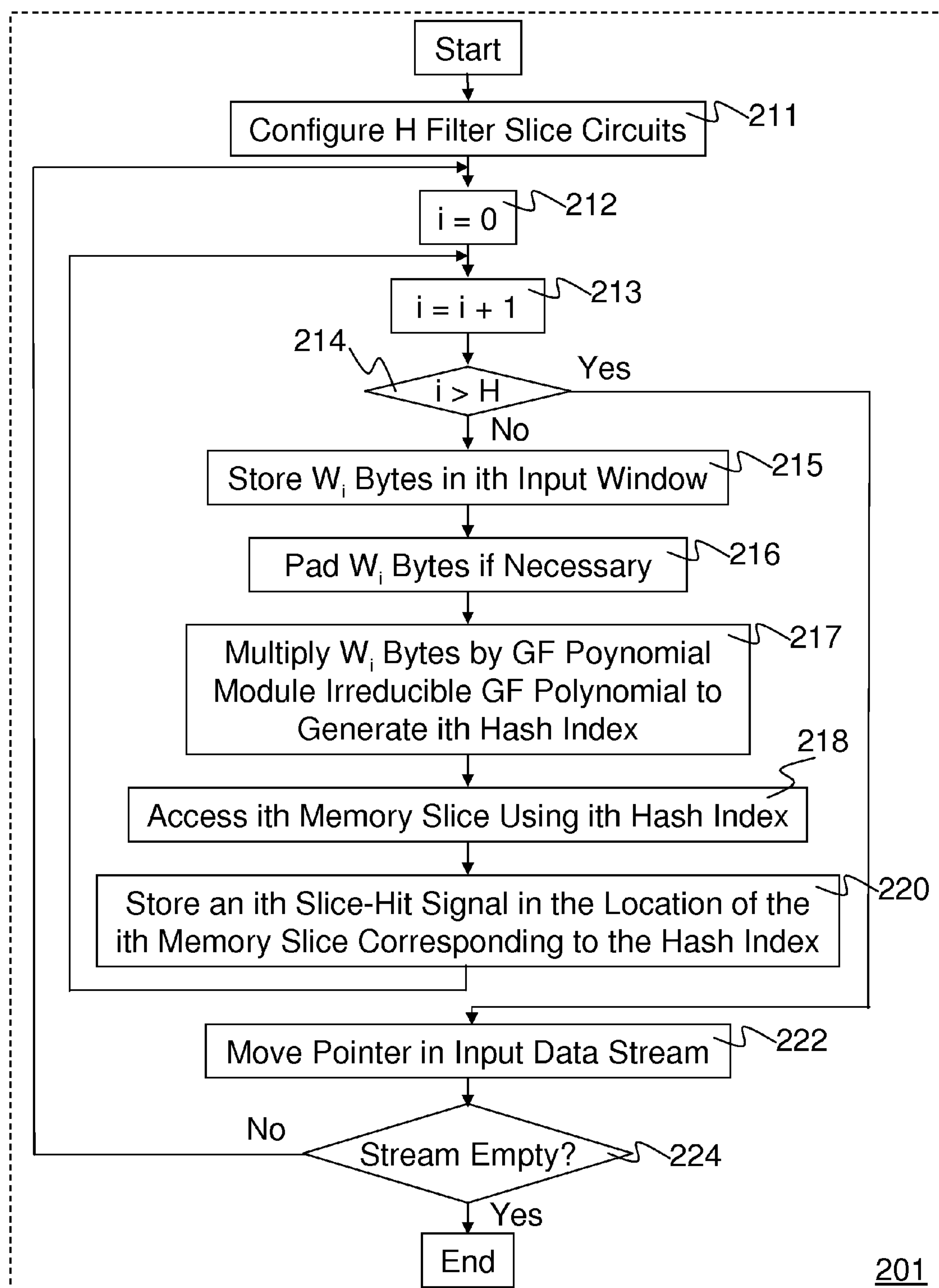


FIG. 2

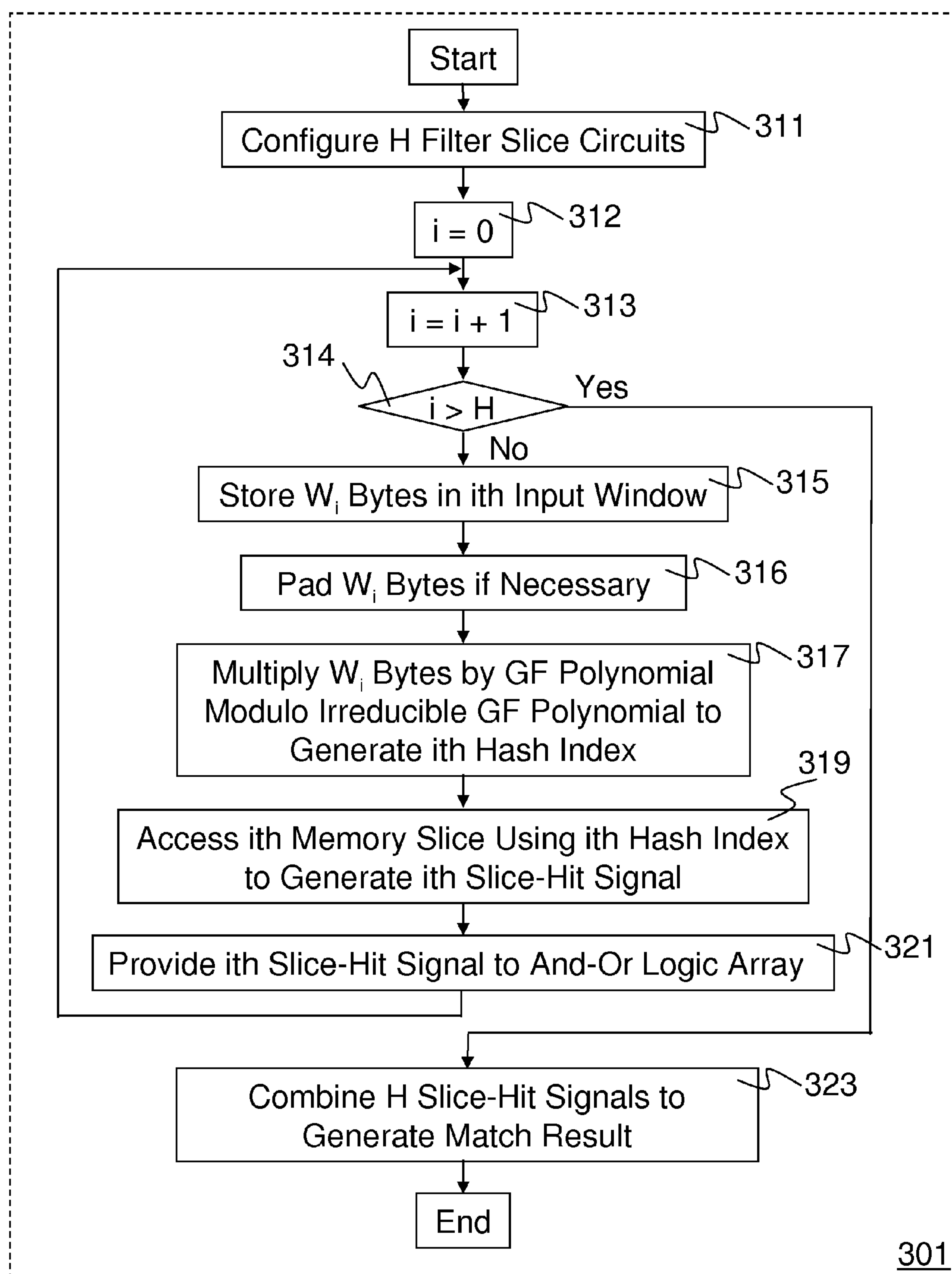


FIG. 3

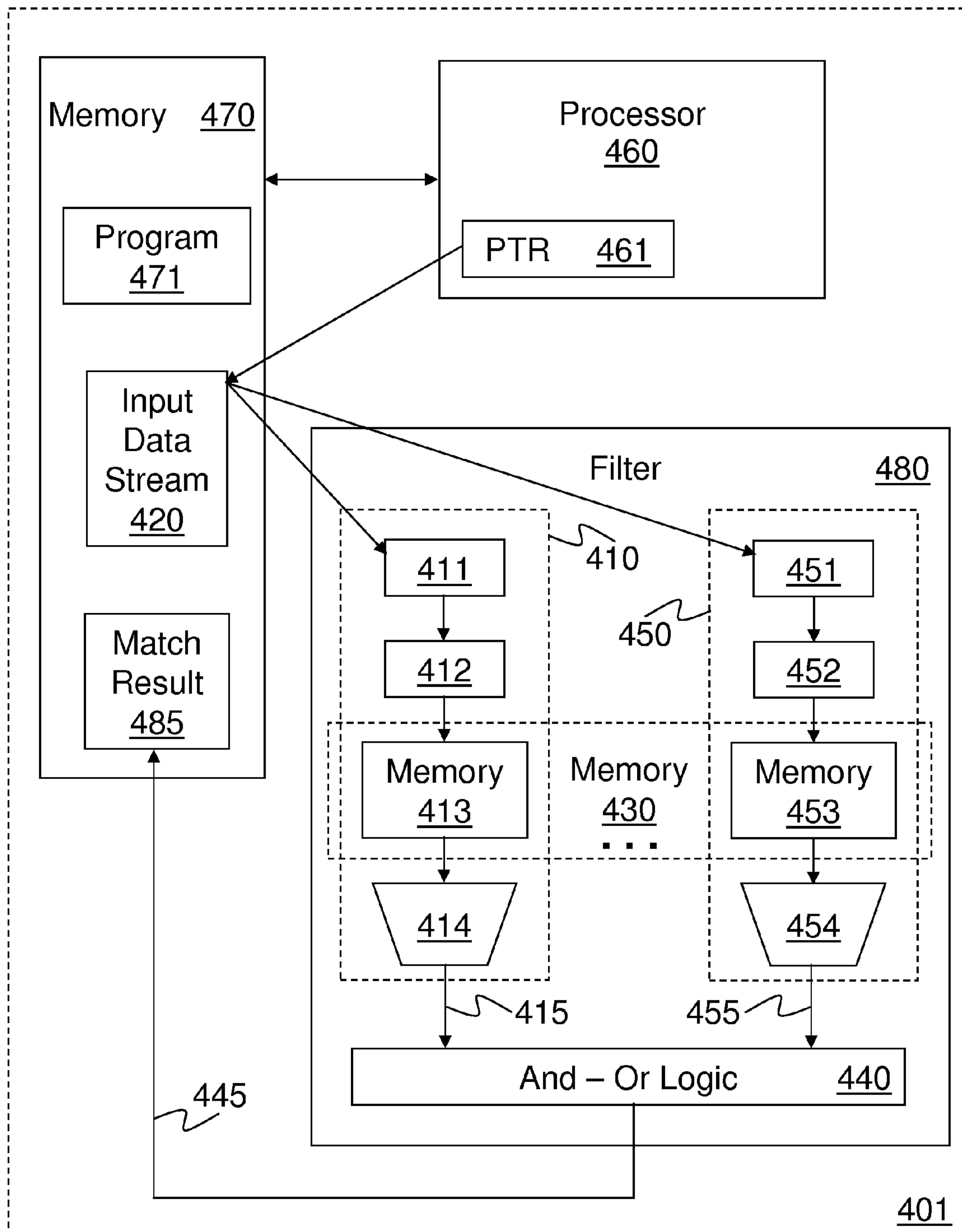


FIG. 4

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FILTER FOR NETWORK INTRUSION AND VIRUS DETECTION

FIELD OF THE DISCLOSURE

This disclosure relates generally to the field of network processing. In particular, the disclosure relates to a novel filter architecture to accelerate string matching in packet inspection for network applications such as intrusion detection/prevention and virus detection.

BACKGROUND OF THE DISCLOSURE

In modern networks, applications such as intrusion detection/prevention and virus detection are important for protecting the networks and/or network users from attacks. In such applications network packets are often inspected to identify problematic packets by finding matches to a known set of data patterns. Matching every byte of an incoming data stream against a large database of patterns (e.g. up to hundreds of thousands) is very compute-intensive. Programs have used techniques such as finite-state machines and filters to find matches to known sets.

A Bloom filter, conceived by Burton H. Bloom in 1970, is a probabilistic structure for determining whether an element is a member of a set. Hashing is performed on the element. Multiple different hash functions are used to generate multiple different hash indices into an array of bits. To add or insert an element into the set, these hash functions are used to index multiple bit locations in the array for the element and these bit locations are then set to one. To query the filter for an arbitrary element the hash functions are used to index multiple bit locations in the array for the element and these bit locations are then checked to see if they are all set to one. If they are not all set to one, the arbitrary element in question is not a member of the set.

Whenever a filter generates a positive outcome for an element, which is not actually a member of the set, the outcome is called a false positive. The Bloom filter will not generate a false negative. It is a goal of any particular filter design, that the probability of false positives is "small." For Bloom filters, after inserting n elements into a set represented by an array of m bits using k different hash functions, the probability of a false positive is $(1 - (1 - 1/m)^{kn})^k$.

Designing a filter for a specific problem may be tedious, and at high data rates it is difficult or impossible for state-of-the-art processors to implement the design at rates even close to line-rate. To achieve rates close to one or more gigabits per second, specialized field-programmable gate array solutions or custom circuits have been proposed.

To date, more generalized reconfigurable architectures to accelerate string matching in packet inspection for network applications such as intrusion detection/prevention and virus detection have not been fully explored.

BRIEF DESCRIPTION OF THE DRAWINGS

The present invention is illustrated by way of example and not limitation in the figures of the accompanying drawings.

FIG. 1 illustrates one embodiment of a filter apparatus to accelerate string matching in packet inspection for network applications such as intrusion detection/prevention and virus detection.

FIG. 2 illustrates a flow diagram for one embodiment of a process to initialize a filter apparatus for string matching in packet inspection.

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FIG. 3 illustrates a flow diagram for one embodiment of a process to utilize a filter apparatus for string matching in packet inspection.

FIG. 4 illustrates one embodiment of a system employing a filter apparatus to accelerate string matching in packet inspection for network applications such as intrusion detection/prevention and virus detection.

DETAILED DESCRIPTION

Methods and apparatus to perform string matching for network packet inspection are disclosed below. In some embodiments, a filter apparatus may be configured as a set of string matching slice circuits, each slice circuit of the set being configured to perform string matching steps in parallel with other slice circuits. Each slice circuit may include an input window storing some number of bytes of data from an input data stream. The input window of data may be padded if necessary, and may be multiplied by a distinct Galois-field polynomial modulo an irreducible Galois-field polynomial to generate a hash index. A storage location of a memory slice corresponding to the hash index may be accessed to generate a slice-hit signal of a plurality of slice-hit signals. The slice-hit signal may be provided to an AND-OR logic array where the plurality of slice-hit signals is logically combined into a match result.

Embodiments of such methods and apparatus represent reconfigurable architectures to accelerate string matching in packet inspection for network applications such as intrusion detection/prevention and virus detection.

In the following description, numerous specific details are set forth. However, it is understood that embodiments of the invention may be practiced without these specific details. In other instances, well-known circuits, structures and techniques have not been shown in detail in order not to obscure the understanding of this description. These and other embodiments of the present invention may be realized in accordance with the following teachings and it should be evident that various modifications and changes may be made in the following teachings without departing from the broader spirit and scope of the invention. The specification and drawings are, accordingly, to be regarded in an illustrative rather than a restrictive sense and the invention measured only in terms of the claims and their equivalents.

FIG. 1 illustrates one embodiment of a filter apparatus **101** to accelerate string matching in packet inspection for network applications such as intrusion detection/prevention and virus detection. Filter apparatus **101** as shown includes an input data stream **120**, which may be in a system memory or may comprise an optional data stream buffer of filter apparatus **101** for storing packed data for inspection and/or a pattern database to initialize filter apparatus **101**. Filter apparatus **101** also includes a set of H (e.g. 1-8) slice circuits **110-150**, each i^{th} slice circuit of the set is configurable for providing an i^{th} slice-hit signal to a configurable AND-OR logic array **140** as one of a set of H slice-hit signals. Slice circuits **110-150**, respectively include input windows **111-151** each configurable to store W_i (e.g. 2-8) bytes of data from input data stream **120**, and Ghash units **112-152** coupled with input windows **111-151** and configurable to receive the W_i bytes of data, to pad the W_i bytes of data if necessary, and to multiply their respective W_i bytes of data by a polynomial modulo an irreducible Galois-field polynomial to generate an index.

It will be appreciated that some embodiments of filter apparatus **101** may use the same irreducible Galois-field polynomial in each of the Ghash units **112-152** with H distinct polynomial multipliers selected at random (each having a

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good mixture of 1's and 0's) to generate H distinct hash indices, thus simplifying the task of generating distinct hash indices for each Ghash unit. It will also be appreciated that in embodiments of filter apparatus **101** where, unlike the Bloom filter, input windows **111-151** are independently configurable to store W_i bytes of data from input data stream **120**, the filter apparatus **101** may be used to solve multiple problems of different sizes (e.g. a 2-byte match, a 3-byte match, a 6-byte match, and an 8-byte match, etc.) at the same time in parallel.

Slice circuits **110-150**, respectively, also include memories **113-153** coupled with the Ghash units **112-152** and configurable to access respective storage locations responsive to their respective indices (e.g. at the addresses specified by some field of bits from respective indices) to each generate an i^{th} slice-hit signal and to provide the an i^{th} slice-hit signal to AND-OR logic array **140** as one of the set of H slice-hit signals **115-155**. Some embodiments of memories **113-153** are configurable from a larger memory **130** to serve as individual memories **113-153** for slice circuits **110-150** respectively. Some alternative embodiments of memories **113-153** may be N-entry (e.g. 1K entries) read/write random-access memories (RAMs) of fixed width (e.g. 64-bits wide) and are configurable to be combined into larger memories (e.g. memory **130**) as necessary (e.g. when a very large set of patterns is required). Slice circuits **110-150** may also include multiplexers **114-154**, respectively, configurable to access respective bit storage locations responsive to portions of their respective indices to generate the i^{th} slice-hit signal and to provide the i^{th} slice-hit signal to AND-OR logic array **140** as one of the set of H slice-hit signals **115-155**.

AND-OR logic array **140** is configurable to receive a set of H slice-hit signals **115-155** and to combine the set of H slice-hit signals **115-155** into a match result **145**, a copy of which may be stored as a match result **185**. Some embodiments of AND-OR logic array **140** may be configurable to perform a simple AND (e.g. as in a Bloom filter) or a simple OR (e.g. as in solving multiple problems of different sizes in parallel) of the set of H slice-hit signals **115-155** to get a match result **145**. Alternative embodiments of AND-OR logic array **140** may be configurable to perform a complex AND-OR of the set of H slice-hit signals **115-155** (e.g. $\text{temp}_k = (\text{AND slice-hit signal}_i \text{ for all } i \text{ in a set } S_k)$ and then the final match result = $(\text{OR temp}_k \text{ for all } k)$) to get a match result **145**. The complex AND-OR of the set of H slice-hit signals **115-155** may be used, for example, in embodiments of filter apparatus **101** to provide multiple Bloom filters in parallel.

It will be appreciated that when a final match result is positive, a verification process may be used to check against false positives. Such verification process may be relatively slower than using filter apparatus **101** and so the configuration of filter apparatus **101** should be carefully made to avoid frequent false positives.

FIG. 2 illustrates a flow diagram for one embodiment of a process **201** to initialize a filter apparatus for string matching in packet inspection. Process **201** and other processes herein disclosed are performed by processing blocks that may comprise dedicated hardware or software or firmware operation codes executable by general purpose machines or by special purpose machines or by a combination of both.

In processing block **211** a set of H slice circuits are configured. In processing block **212**, i is set to zero (0). In processing block **213**, i is incremented. In processing block **214**, i is checked to see if it has exceeded H. It will be appreciated that even though initialization of the H slice circuits is shown as an iterative process **201**, in at least some preferred embodiment of process **201**, the set of H slice circuits are configured to concurrently perform initialization according to processing

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blocks **215-220** of process **201** for use in string matching during network packet inspections. Therefore, for each of the H slice circuits processing blocks **215-220** are executed as follows, before proceeding to processing block **222**.

In processing block **215** W_i bytes of data is stored from an input data stream in an i^{th} input window. In processing block **216** the W_i bytes of data are padded if necessary. Then in processing block **217** the W_i bytes of data are multiplied by a Galois-field polynomial modulo an irreducible Galois-field polynomial to generate an i^{th} hash index. In processing block **218** a storage location of a memory corresponding to the i^{th} hash index is accessed, and in processing block **220** an i^{th} slice-hit signal is stored (i.e. set) in the storage location of the memory corresponding to the i^{th} hash index. When all of the H slice circuits have completed processing blocks **215-220** of process **201**, processing proceeding to processing block **222** where a pointer in the input data stream is moved (e.g. to a new string in the database). Then from processing block **224**, if the data stream is empty processing terminates. Otherwise processing repeats in processing block **212**.

It will be appreciated that the process **201** may be iterated for hundreds to hundreds of thousands of times in order to initialize a filter apparatus for string matching patterns in packet inspection. Thus when the set of H slice circuits are configured to concurrently perform initialization substantial performance improvements may be realized. It will also be appreciated that the process **201** of initializing a filter apparatus (by setting slice-hit signals) may be performed in a manner substantially similar to a process of utilizing a filter apparatus for string matching (by reading the slice-hit signals) in packet inspection. In some embodiments of processing block **222** a pointer into the input data stream may be moved for each i^{th} slice, in such a way as to provide each i^{th} slice with a new compare pattern, whereas in utilizing a filter apparatus for string matching a pointer into the input data stream may be simply incremented.

FIG. 3 illustrates a flow diagram for one embodiment of a process **301** to utilize a filter apparatus for string matching in packet inspection. In processing block **311** a set of H slice circuits are configured. In processing block **312**, i is set to zero (0). In processing block **313**, i is incremented. In processing block **314**, i is checked to see if it has exceeded H. Again, it will be appreciated that even though utilization of the H slice circuits is shown as an iterative process **301**, in at least some preferred embodiment of process **301**, the set of H slice circuits are configured to concurrently perform string matching according to processing blocks **315-321** of process **301** for use during network packet inspections. Therefore, for each of the H slice circuits processing blocks **315-321** are executed as follows, before proceeding to processing block **323**.

In processing block **315** W_i bytes of data is stored from an input data stream in an i^{th} input window. In processing block **316** the W_i bytes of data are padded if necessary. Then in processing block **317** the W_i bytes of data are multiplied by a Galois-field polynomial modulo an irreducible Galois-field polynomial to generate an i^{th} hash index. In processing block **319** a storage location of a memory corresponding to the i^{th} hash index is accessed to generate an i^{th} slice-hit signal of a set of H slice-hit signals. In processing block **321** the i^{th} slice-hit signal is provided to an AND-OR logic array as one of the set of H slice-hit signals. When all of the H slice circuits have completed processing blocks **315-321** of process **301**, processing proceeding to processing block **323** where the AND-OR logic array is configured to receive the set of H slice-hit

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signals and to combine the set of H slice-hit signals into a match result. Then from processing block 323 processing terminates.

It will be appreciated that iterations of process 301 may be configured in accordance with embodiments of filter apparatus 101 to substantially accelerate string matching in packet inspection.

FIG. 4 illustrates one embodiment of a system 401 employing a filter 480 to accelerate string matching in packet inspection for network applications such as intrusion detection/prevention and virus detection.

System 401 includes an input data stream 420, which may be in system memory 470 as shown, or may comprise an optional data stream buffer of filter 480 for storing packed data for inspection and/or a pattern database to initialize filter 480.

Filter 480 includes a set of H slice circuits 410-450, each i^{th} slice circuit of the set is configurable for providing an i^{th} slice-hit signal to a configurable AND-OR logic array 440 as one of a set of H slice-hit signals. Slice circuits 410-450, respectively include input windows 411-451 each configurable to store W_i bytes of data from input data stream 420, and Ghash units 412-452 coupled with input windows 411-451 and configurable to receive the W_i bytes of data, to pad the W_i bytes of data if necessary, and to multiply their respective W_i bytes of data by a polynomial modulo an irreducible Galois-field polynomial to generate an index.

Slice circuits 410-450, respectively, also include memories 413-453 coupled with the Ghash units 412-452 and configurable to access respective storage locations responsive to their respective indices to each generate an i^{th} slice-hit signal and to provide the an i^{th} slice-hit signal to AND-OR logic array 440 as one of the set of H slice-hit signals 415-455. Memories 413-453 may be N-entry read/write RAMs of any fixed width and configurable to be combined into larger memories (e.g. memory 430) as necessary. Alternatively some embodiments of memories 413-453 may be configurable from a larger memory 430. Slice circuits 410-450 may also include multiplexers 414-454, respectively, configurable to access respective bit storage locations responsive to portions of their respective indices to generate the i^{th} slice-hit signal and to provide the i^{th} slice-hit signal to AND-OR logic array 440 as one of the set of H slice-hit signals 415-455. AND-OR logic array 440 may receive the set of H slice-hit signals 415-455 and combine the set of H slice-hit signals 415-455 into a match result 445.

System 401 also includes system processor 460 to executed a program 471 in system memory 470 to accelerate string matching in packet inspection for network applications using filter 480, and to move or increment a pointer 461 into input data stream 420 until a match result 445 is positive (in the case of string matching for packet inspections) or until an end-of-file is reached in the input data stream 420. In some embodiments of system 401, processor 460 may check a copy of match result 445 stored in system memory 470 as a match result 485 when string matching for packet inspections to determine if match result 445 was positive.

The above description is intended to illustrate preferred embodiments of the present invention. From the discussion above it should also be apparent that especially in such an area of technology, where growth is fast and further advancements are not easily foreseen, the invention can may be modified in arrangement and detail by those skilled in the art without departing from the principles of the present invention within the scope of the accompanying claims and their equivalents.

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What is claimed is:

1. A method to perform string matching for network packet inspection, the method comprising:

configuring a set of H slice circuits, each i^{th} slice circuit of the set of H slice circuits being configured to perform the steps of:

independently storing an i^{th} input window of W_i bytes of data from an input data stream;

padding the W_i bytes of data if necessary, and multiplying the W_i bytes of data by a Galois-field polynomial modulo an irreducible Galois-field polynomial combined with a randomly generated polynomial multiplier to generate an i^{th} hash index;

accessing a storage location of a memory corresponding to the i^{th} hash index to generate an i^{th} slice-hit signal of a set of H slice-hit signals; and

providing the i^{th} slice-hit signal to an AND-OR logic array as one of the set of H slice-hit signals; and

configuring the AND-OR logic array to receive the set of H slice-hit signals and to combine the set of H slice-hit signals into a match result.

2. The method of claim 1 wherein configuring each i^{th} slice circuit of the set of H slice circuits to perform the step of providing the i^{th} slice-hit signal to the AND-OR logic array comprises:

storing the i^{th} slice-hit signal in the storage location of the memory corresponding to the i^{th} hash index.

3. The method of claim 2 wherein each i^{th} input window of W_i bytes of data from the input data stream comprises a complete data pattern.

4. The method of claim 2 wherein providing the i^{th} slice-hit signal to the AND-OR logic array comprises:

reading out the i^{th} slice-hit signal, from the storage location of the memory corresponding to the i^{th} hash index, to the AND-OR logic array as the i^{th} one of the set of H slice-hit signals.

5. The method of claim 2 wherein providing the i^{th} slice-hit signal to the AND-OR logic array comprises:

multiplexing the i^{th} slice-hit signal from the storage location of the memory corresponding to the i^{th} hash index, to the AND-OR logic array as the i^{th} one of the set of H slice-hit signals.

6. The method of claim 1, wherein the AND-OR logic array is configured to receive the set of H slice-hit signals and to logically AND the set of H slice-hit signals into a match result.

7. The method of claim 1, wherein the AND-OR logic array is configured to receive the set of H slice-hit signals and to logically OR the set of H slice-hit signals into a match result.

8. The method of claim 1, wherein the AND-OR logic array is configured to receive the set of H slice-hit signals and to logically AND subsets of the set of H slice-hit signals into temporary results, and to logically OR the temporary results into a match result.

9. An apparatus comprising:

an AND-OR logic array configurable to receive a set of H slice-hit signals and to combine the set of H slice-hit signals into a match result; and

a set of H slice circuits, each i^{th} slice circuit of the set comprising:

an input window configurable to independently store W_i bytes of data from an input data stream;

a Ghash unit coupled with the input window and configurable to receive the W_i bytes of data, pad the W_i bytes of data if necessary, and multiply the W_i bytes of data by a Galois-field polynomial modulo an irreduc-

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ible Galois-field polynomial combined with a randomly generated polynomial multiplier to generate an index; and

a memory coupled with the Ghash unit and configurable to access a storage location responsive to the index to generate a slice-hit signal and to provide the slice-hit signal to said AND-OR logic array as one of the set of H slice-hit signals.

10. The apparatus of claim 9 wherein providing the slice-hit signal to the AND-OR logic array comprises:

reading out the slice-hit signal, from the storage location of the memory corresponding to the index of the i^{th} slice circuit, to the AND-OR logic array as the i^{th} one of the set of H slice-hit signals.

11. The apparatus of claim 9 wherein providing the slice-hit signal to the AND-OR logic array comprises:

multiplexing the slice-hit signal, from the storage location of the memory corresponding to the index of the i^{th} slice circuit, to the AND-OR logic array as the i^{th} one of the set of H slice-hit signals.

12. The apparatus of claim 9 wherein the AND-OR logic array is configurable to receive the set of H slice-hit signals and to logically AND the set of H slice-hit signals into a match result.

13. The apparatus of claim 9 wherein the AND-OR logic array is configurable to receive the set of H slice-hit signals and to logically OR the set of H slice-hit signals into a match result.

14. The apparatus of claim 9 wherein the AND-OR logic array is configurable to receive the set of H slice-hit signals and to logically AND subsets of the set of H slice-hit signals into temporary results, and to logically OR the temporary results into a match result.

15. The apparatus of claim 9 wherein the same irreducible Galois-field polynomial is used in each i^{th} slice circuit of the set of H slice circuits.

16. The apparatus of claim 15 wherein each the W_i bytes of data are multiplied by a different distinct Galois-field polynomial in each i^{th} slice circuit of the set of H slice circuits.

17. A packet processing system to perform string matching for network packet inspection, the system comprising:

a system processor;

an AND-OR logic array configurable to receive a set of H slice-hit signals and to combine the set of H slice-hit signals into a match result; and

a set of H slice circuits, each i^{th} slice circuit of the set comprising:

an input window configurable to independently store W_i bytes of data from an input data stream;

a Ghash unit coupled with the input window and configurable to receive the W_i bytes of data, pad the W_i

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bytes of data if necessary, and multiply the W_i bytes of data by a Galois-field polynomial modulo an irreducible Galois-field polynomial combined with a randomly generated polynomial multiplier to generate an index; and

a memory coupled with the Ghash unit and configurable to access a storage location responsive to the index to generate a slice-hit signal and to provide the slice-hit signal to said AND-OR logic array as one of the set of H slice-hit signals; and

a machine readable medium to store executable instructions, such that when said executable instructions are executed by the system processor, the system processor is caused to:

set a pointer to a first character of the input data stream to establish a starting point for the input window of each i^{th} slice circuit, and

increment the pointer until the match result is positive or until an end-of-file is reached in the input data stream.

18. The system of claim 17 wherein the same irreducible Galois-field polynomial is used in each i^{th} slice circuit of the set of H slice circuits.

19. The system of claim 18 wherein each the W_i bytes of data are multiplied by a different distinct Galois-field polynomial in each i^{th} slice circuit of the set of H slice circuits.

20. The system of claim 17 wherein the AND-OR logic array is configurable to receive the set of H slice-hit signals and to logically AND the set of H slice-hit signals into a match result.

21. The system of claim 17 wherein the AND-OR logic array is configurable to receive the set of H slice-hit signals and to logically OR the set of H slice-hit signals into a match result.

22. The system of claim 17 wherein the AND-OR logic array is configurable to receive the set of H slice-hit signals and to logically AND subsets of the set of H slice-hit signals into temporary results, and to logically OR the temporary results into a match result.

23. The system of claim 18 wherein providing the slice-hit signal to the AND-OR logic array comprises:

reading out the slice-hit signal, from the storage location of the memory corresponding to the index of the i^{th} slice circuit, to the AND-OR logic array as the i^{th} one of the set of H slice-hit signals.

24. The system of claim 17 wherein providing the slice-hit signal to the AND-OR logic array comprises:

multiplexing the slice-hit signal, from the storage location of the memory corresponding to the index of the i^{th} slice circuit, to the AND-OR logic array as the i^{th} one of the set of H slice-hit signals.

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