

US009100364B2

(12) **United States Patent**
Zuk

(10) **Patent No.:** **US 9,100,364 B2**
(45) **Date of Patent:** ***Aug. 4, 2015**

(54) **INTELLIGENT INTEGRATED NETWORK SECURITY DEVICE**

63/0263 (2013.01); **H04L 63/12** (2013.01);
H04L 63/1416 (2013.01); **H04L 63/1441**
(2013.01); **H04L 69/22** (2013.01); **H04L**
63/0209 (2013.01); **H04L 63/0218** (2013.01)

(71) Applicant: **JUNIPER NETWORKS, INC.**,
Sunnyvale, CA (US)

(58) **Field of Classification Search**

CPC H04L 63/02; H04L 69/22; H04L 63/0263
USPC 726/13
See application file for complete search history.

(72) Inventor: **Nir Zuk**, Palo Alto, CA (US)

(73) Assignee: **Juniper Networks, Inc.**, Sunnyvale, CA
(US)

(*) Notice: Subject to any disclaimer, the term of this
patent is extended or adjusted under 35
U.S.C. 154(b) by 0 days.

This patent is subject to a terminal dis-
claimer.

(56) **References Cited**

U.S. PATENT DOCUMENTS

5,598,410 A 1/1997 Stone
5,606,668 A 2/1997 Shwed

(Continued)

FOREIGN PATENT DOCUMENTS

EP 1 143 660 A2 10/2001
EP 1 427 162 A1 6/2004

(Continued)

OTHER PUBLICATIONS

International Search Report for corresponding PCT application,
PCT/US2004/009607, dated Oct. 22, 2004, 3 pages.

(Continued)

Primary Examiner — Mohammad W Reza

(74) *Attorney, Agent, or Firm* — Harrity & Harrity, LLP

(57) **ABSTRACT**

Methods, computer program products and apparatus for pro-
cessing data packets are described. Methods include receiv-
ing the data packet, examining the data packet, determining a
single flow record associated with the packet and extracting
flow instructions for two or more devices from the single flow
record.

20 Claims, 9 Drawing Sheets

(21) Appl. No.: **14/230,180**

(22) Filed: **Mar. 31, 2014**

(65) **Prior Publication Data**

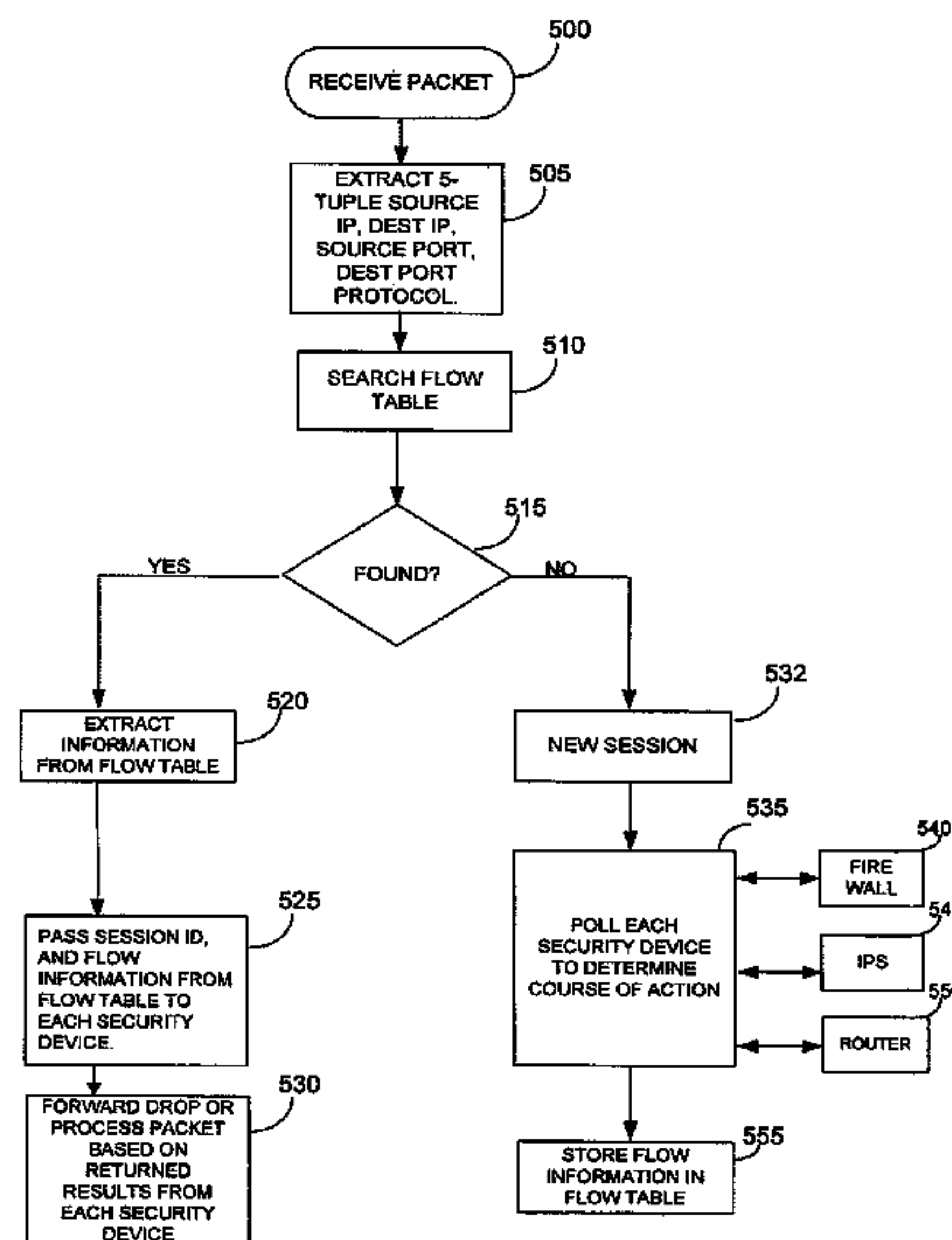
US 2014/0259146 A1 Sep. 11, 2014

Related U.S. Application Data

(63) Continuation of application No. 13/616,067, filed on
Sep. 14, 2012, now Pat. No. 8,726,016, which is a
continuation of application No. 12/575,997, filed on
Oct. 8, 2009, now Pat. No. 8,332,948, which is a
continuation of application No. 10/402,920, filed on
Mar. 28, 2003, now Pat. No. 7,650,634, which is a
continuation-in-part of application No. 10/072,683,
filed on Feb. 8, 2002, now Pat. No. 8,370,936.

(51) **Int. Cl.**
G06F 9/00 (2006.01)
H04L 29/06 (2006.01)

(52) **U.S. Cl.**
CPC **H04L 63/02** (2013.01); **H04L 63/0227**
(2013.01); **H04L 63/0254** (2013.01); **H04L**



(56)

References Cited

U.S. PATENT DOCUMENTS

5,781,550 A 7/1998 Templin et al.
 5,835,726 A 11/1998 Shwed et al.
 5,842,040 A 11/1998 Hughes et al.
 5,909,686 A 6/1999 Muller et al.
 6,006,264 A 12/1999 Colby et al.
 6,049,528 A 4/2000 Hendel et al.
 6,052,788 A 4/2000 Wesinger, Jr. et al.
 6,088,356 A 7/2000 Hendel et al.
 6,098,172 A 8/2000 Coss et al.
 6,119,236 A 9/2000 Shipley
 6,141,749 A 10/2000 Coss et al.
 6,154,775 A 11/2000 Coss et al.
 6,170,012 B1 1/2001 Coss et al.
 6,205,551 B1 3/2001 Grosse
 6,253,321 B1 6/2001 Nikander et al.
 6,275,942 B1 8/2001 Bernhard et al.
 6,279,113 B1 8/2001 Vaidya
 6,301,668 B1 10/2001 Gleichauf et al.
 6,304,975 B1 10/2001 Shipley
 6,311,278 B1 10/2001 Raanan et al.
 6,321,338 B1 11/2001 Porras et al.
 6,370,603 B1 4/2002 Silverman et al.
 6,421,730 B1 7/2002 Narad et al.
 6,449,647 B1 9/2002 Colby et al.
 6,453,345 B2 9/2002 Trcka et al.
 6,466,985 B1 10/2002 Goyal et al.
 6,487,666 B1 11/2002 Shanklin et al.
 6,499,107 B1 12/2002 Gleichauf et al.
 6,590,894 B1 7/2003 Kerr et al.
 6,591,303 B1 7/2003 Hendel et al.
 6,600,744 B1 7/2003 Carr et al.
 6,606,315 B1 8/2003 Menditto et al.
 6,633,560 B1 10/2003 Tiwari et al.
 6,650,641 B1 11/2003 Albert et al.
 6,654,373 B1 11/2003 Maher et al.
 6,704,278 B1 3/2004 Albert et al.
 6,735,169 B1 5/2004 Albert et al.
 6,742,045 B1 5/2004 Jordan et al.
 6,768,738 B1 7/2004 Yazaki et al.
 6,775,692 B1 8/2004 Albert et al.
 6,781,992 B1 8/2004 Rana et al.
 6,788,648 B1 9/2004 Peterson
 6,795,918 B1 9/2004 Trolan
 6,851,061 B1 2/2005 Holland et al.
 6,856,991 B1 2/2005 Srivastava
 6,976,154 B1 12/2005 Dyckerhoff et al.
 6,981,158 B1 12/2005 Sanchez et al.
 7,006,443 B2 2/2006 Storr
 7,032,037 B2 4/2006 Garnett et al.
 7,042,870 B1 5/2006 Albert et al.
 7,051,066 B1 5/2006 Albert et al.
 7,054,930 B1 5/2006 Cheriton
 7,073,196 B1 7/2006 Dowd et al.
 7,076,803 B2 7/2006 Bruton et al.
 7,123,583 B2 10/2006 Hoar et al.
 7,143,438 B1 11/2006 Coss et al.
 7,185,368 B2 2/2007 Copeland, III
 7,346,686 B2 3/2008 Albert et al.
 7,376,085 B2 5/2008 Yazaki et al.
 7,512,980 B2 3/2009 Copeland et al.
 7,535,907 B2 5/2009 Hussain et al.
 7,643,481 B2 1/2010 Kadambi et al.
 7,650,634 B2 1/2010 Zuk
 7,778,254 B2 8/2010 Kadambi et al.
 7,895,431 B2 2/2011 Bouchard et al.
 7,970,886 B1 6/2011 Wetherall et al.
 8,023,413 B2 9/2011 Kadambi et al.
 8,332,948 B2 12/2012 Zuk
 2001/0028650 A1 10/2001 Yoshizawa et al.
 2001/0051864 A1 12/2001 Kerr et al.
 2002/0032797 A1 3/2002 Xu
 2002/0080789 A1 6/2002 Henderson et al.
 2002/0124187 A1 9/2002 Lyle et al.
 2002/0126621 A1 9/2002 Johnson et al.
 2002/0161839 A1 10/2002 Colasurdo et al.

2002/0165956 A1 11/2002 Phaal
 2003/0105976 A1 6/2003 Copeland
 2003/0145225 A1 7/2003 Bruton et al.
 2003/0149887 A1 8/2003 Yadav
 2003/0149888 A1 8/2003 Yadav
 2003/0154399 A1 8/2003 Zuk
 2005/0141503 A1 6/2005 Welfeld
 2005/0163132 A1 7/2005 Mieno et al.
 2005/0210533 A1* 9/2005 Copeland et al. 726/23
 2006/0005231 A1 1/2006 Zuk
 2006/0159019 A1* 7/2006 Buskirk et al. 370/235
 2008/0115204 A1* 5/2008 Ramsey et al. 726/13
 2013/0067561 A1 3/2013 Zuk

FOREIGN PATENT DOCUMENTS

JP 10-107795 4/1998
 JP 11-316677 11/1999
 JP 2000-312225 11/2000
 JP 2001-077857 3/2001
 JP 2001-313640 11/2001
 JP 2002-524891 8/2002
 JP 2003-78549 3/2003
 WO WO 99/67930 12/1999
 WO WO 03/025766 3/2003
 WO WO 03/061238 7/2003

OTHER PUBLICATIONS

Stonesoft, 'StoneBeat Security Cluster White Paper,' Aug. 2000, Finland, pp. 1-9.
 Stonesoft, 'Secure Highly Available Enterprise-A White Paper,' Feb. 2001, Finland, pp. 1-10.
 Stonesoft, 'StoneGate White Paper,' Mar. 2001, Finland, pp. 1-6.
 Stonesoft Corp. 'StoneGate product webpage,' www.stonesoft.com/document/363.html, Mar. 27, 2001 (print date), pp. 1-2.
 Stonesoft Corp. 'Next Level of Network Accessibility' webpage, www.stonesoft.com/document/183.html, Mar. 27, 2001 (print date), p. 1.
 Stonesoft Corp., 'Platforms,' webpage, www.stonesoft.com/document/186.html, Mar. 27, 2001 (print date), p. 1.
 Nokia, 'Technical White Paper: The IP Clustering Power of Nokia VPN—Keeping Customers Connected,' Apr. 2001, pp. 1-13.
 Nokia, 'Nokia VPN Solutions—Nokia VPN CC2500 Gateway,' 2001, product information, pp. 1-2.
 Nokia, 'Nokia VPN Solutions—Nokia VPN CC5200 Gateway,' 2001, product information, pp. 1-2.
 Nokia, 'Nokia VPN Solutions—Nokia VPN CC5205 Gateway,' 2001, product information, pp. 1-2.
 Axelsson, S., "Intrusion Detection Systems: A Survey and Taxonomy," Dept. of Computer Eng., Chalmers Univ. of Technology, Goteborg, Sweden, Mar. 14, 2000, pp. 1-27.
 Avolio, F., "Firewalls and Virtual Private Networks," CSI Firewall Archives, printed Nov. 13, 2001, URL: <http://www.spirit.com/CSI/Papers/fw+vpns.html>, pp. 1-7.
 Bace, R., "An Introduction to Intrusion Detection & Assessment," ICSA Intrusion Detection Systems Consortium White Paper, 1999, URL: <http://www.icsalabs.com/html/communities/ids/whitepaper/Intrusion1.pdf> pp. 1-38.
 Business Wire, Inc., "NetScreen and OneSecure Unite to Deliver Industry's First Total Managed Security Services Platform," San Jose, CA, Feb. 20, 2001, pp. 1-2.
 Business Wire, Inc., "OneSecure Launches the First Co-Managed Security Services Platform," Denver, CO, Jan. 29, 2001, pp. 1-2.
 Carr, Jim, "Intrusion Detection Systems: Back to Front?," Network Magazine, Sep. 5, 2001, URL: <http://www.networkmagazine.com/article/NMG20010823S0007/2>, pp. 1-9.
 Check Point Software Technologies Ltd., Firewall-1® Technical Overview P/N 500326, www.checkpoint.com, Oct. 2000, pp. 1-29.
 Cisco Systems, "Cisco IOS Firewall Intrusion Detection System," Cisco IOS Release 12.0(5)T, 2001, pp. 1-40.
 Cisco Systems, "Cisco IOS Firewall Authentication Proxy," Cisco IOS Release 12.0(5)T, 2001, pp. 1-48.

(56)

References Cited

OTHER PUBLICATIONS

Clark, D., "RFC815-IP Datagram Reassembly Algorithms," Internet RFC/STD/FYI/BCP Archives, <http://www.faqs.org/rfcs/rfc815.html>, Jul. 1982, pp. 1-8.

Copeland, Dr. John A., "Observing Network Traffic-Techniques to Sort Out the Good, the Bad, and the Ugly," PowerPoint Slide Presentation presented to ISSA—Atlanta, Jun. 27, 2001, pp. 1-22.

Denning, Dorothy E., "An Intrusion-Detection Model," IEEE Transactions on Software Engineering, vol. SE-13, No. 2, Feb. 1987, 17 pages.

Farrow, Rik, "An Analysis of Current Firewall Technologies," CSI 1997 Firewalls Matrix, 1998, URL: <http://www.spirit.com/CSI/Papers/farrowpa.htm>, pp. 1-5.

Firewall Product Comparison Table: VelociRaptor, BorderWare Firewall Server and Firewall-1/VPN-1 Gateway, www.spirit.com, printed Nov. 13, 2001, pp. 1-7.

Firewall Product Comparison Table: PIX Firewall, CyberGuard Firewall for UnixWare & CyberGuard Firewall for Windows NT, www.spirit.com, printed Nov. 13, 2001, pp. 1-8.

Firewall Product Comparison Table: CyberGuard Premium Appliance Firewall, InstaGate EX & BizGuardian VPN Firewall, www.spirit.com, printed Nov. 13, 2001, pp. 1-8.

Firewall Product Comparison Table: Server Protector 100, GNAT Box Firewall Software & Lucent Managed Firewall, www.spirit.com, printed Nov. 13, 2001, pp. 1-6.

Firewall Product Comparison Table: Internet Security and Acceleration (ISA) Server 2000, NetBSD/i386 Firewall & Guardian Firewall, www.spirit.com, printed Nov. 13, 2001, pp. 1-7.

Firewall Product Comparison Table: NetScreen-10 and NetScreen-100, CyberwallPLUS & BorderManager, www.spirit.com, printed Nov. 13, 2001, pp. 1-7.

Firewall Product Comparison Table: Gauntlet Firewall, Barricade Classic/XL & Barricade S, www.spirit.com, printed Nov. 13, 2001, pp. 1-8.

Firewall Product Comparison Table: Sidewinder™, SecurePipe Managed Firewall Service & SnapGear, www.spirit.com, printed Nov. 13, 2001, pp. 1-7.

Firewall Product Comparison Table: SonicWALL PRO, Sunscreen Secure Net & WinRoute Pro 4.1, www.spirit.com, printed Nov. 13, 2001, pp. 1-6.

Firewall Product Comparison Table: WatchGuard Technologies, Inc. LiveSecurity System 4.6, www.spirit.com, printed Nov. 13, 2001, pp. 1-4.

Graham, R., "FAQ: Network Intrusion Detection System," www.robertgraham.com/pubs/network-intrusion-detection.html, Ver. 0.8.3, Mar. 21, 2000, pp. 1-43.

Habra, N. et al., "ASAX: Software Architecture and Rule-Based Language for Universal Audit Trail Analysis," Proceedings of the ESORICS '92, European Symposium on Research in Computer Security, Nov. 23-25, 1992, Toulouse, Springer-Verlag, 16 pages.

ICSA Labs, Intrusion Detection System Buyer's Guide, ICSA White Paper, 1999, pp. 1-52.

Jackson, K. et al., "Intrusion Detection System (IDS) Product Survey," Los Alamos National Laboratory, Los Alamos, NM, LA-UR-99-3883 Ver. 2.1, Jun. 25, 1999, pp. 1-103.

Jones, Kyle, "Introduction to Firewalls," IT Audit.org Forum Network Management, vol. 2, May 1, 1999, URL: <http://www.itaudit.org/forum/networkmanagement/f209nm.htm>, pp. 1-5.

Lancope, "The Security Benefits of a Flow-Based Intrusion Detection System," White Paper, date unknown, pp. 1-11.

LapLink, Inc., "Article #178-Introduction to Firewalls," www.laplink.com/support/kb/articie.asp?ID=178, Apr. 24, 2001, pp. 1-3.

McHugh, J. et al., "Defending Yourself: The Role of Intrusion Detection Systems," Software Engineering Institute, IEEE Software Eng., Sep./Oct. 2000, pp. 42-51.

Network ICE Corporation, "Why Firewalls Are Not Enough," at www.networkice.com/products/firewalls.html, 2000, pp. 1-9.

Power, R., et al., "CSI Intrusion Detection System Resource—Five Vendors Answer Some No-Nonsense Questions on IDS," Computer Security Alert #184, Jul. 1998, pp. 1-8.

Power, R., "CSI Roundtable: Experts discuss present and future intrusion detection systems," Computer Security Journal, vol. XIV, #1, URL: <http://www.gocsi.com/roundtable.htm>, 2001, pp. 1-20.

Sample, Char, et al., "Firewall and IDS Shortcomings," SANS Network Security, Monterey, CA, Oct. 2000, pp. 1-13.

Smith, Gary, "A Brief Taxonomy of Firewalls—Great Walls of Fire," SANS Institute's Information Security Reading Room, May 18, 2001, URL: <http://www.sans.org/infosecFAQ/firewall/taxonomy.htm>, pp. 1-21.

Spitzner, Lance, "How Stateful is Stateful Inspection? Understanding the FW-1 State Table," <http://www.enteract.com/~lspitz/fwtable.html>, Nov. 29, 2000, pp. 1-8.

Sundaram, A., "An Introduction to Intrusion Detection," www.acm.org/crossroads/xrds2-4/intrus.html, Jan. 23, 2001, pp. 1-12.

Tyson, Jeff, "How Firewalls Work," <http://www.howstuffworks.com/firewall.htm/printable>, 2001, pp. 1-7.

Xinetica, Ltd., "An Overview of Intrusion Detection Systems," Xinetica White Paper, Nov. 12, 2001 (print date), URL: http://www.xinetica.com/tech_explained/general/ids/wp_ids.html, pp. 1-9.

Zuk, Nir, "Protect Yourself with Firewalls," www.techtv.com, Jul. 12, 2001, URL: <http://www.techtv.com/screensavers/print/0,23102,3325761,00.html>, pp. 1-3.

Zuk, Nir, "How the Code Red Worm Works," www.techtv.com, Sep. 21, 2001, URL: <http://www.techtv.com/screensavers/print/0,23102,3349133,00.html>, pp. 1-2.

Petersen, S., et al., "Web apps pose security threat," ZDNet: Tech Update, Jan. 29, 2001, URL: <http://techupdate.zdnet.com/techupdate/stories/main/0,14179,2679177,00.html>, pp. 1-3.

Lancope, "StealthWatch Provides Early Detection of the Code Red Worm and its Future Variants," www.stealthwatch.com, date unknown, pp. 1-4.

Reavis, J., "Cash and Burn," Jun. 2001, 6 pages.

SOS Corporation, "An Introduction to Firewalls," 1995, URL: <http://www.uclan.ac.uk/facs/destech/compute/staff/haroun/FIREWALS.HTM>, pp. 1-3.

Morgan, Lisa, "Be Afraid, Be Very Afraid," InternetWeek Intrusion Detection Systems, Jan. 3, 2001, pp. 1-6.

Mullins, Robert, "'Cyber war' raises security concerns," Silicon Valley/San Jose Business Journal, May 11, 2001, pp. 1-4.

James P. Anderson Co., "Computer Security Threat Monitoring and Surveillance," Apr. 15, 1980, 56 pages.

Internet Security Systems, Inc., "REALSECURE™, The RealSecure Advantage," 2001, 2 pages.

Chuvakin, A., et al., "Basic Security Checklist for Home and Office Users," SecurityFocus, Nov. 5, 2001, pp. 1-5.

Network Ice, "SMTP WIZ command," 2001, URL: <http://networkice.com/Advice/Intrusions/2001006/default.htm>, pp. 1-2.

Bace, R., et al., "NIST Special Publication on Intrusion Detection Systems," National Institute of Standards and Technology Special Publication, date unknown, pp. 1-51.

G. Navarro: A Partial Deterministic Automaton for Approximate String Matching, 1997, Department of Computer Science, University of Chile, 13 pages.

G. Navarro et al. Improving an Algorithm for Approximate Pattern Matching, 1998, Department of Computer Science, University of Chile, 34 pages.

Network Magazine, vol. 2, No. 2, pp. 116-119 (with English abstract).

Software Design, Nov. 1996, pp. 39-58 (with English abstract).

Julkunen et al., "Enhance Network Security with Dynamic Packet Filter", IEEE (1998), pp. 268-275.

Sharp et al., "Starburst: Building Next-Generation Internet Devices", Bell Labs Technical Journal 6(2), pp. 6-17, 2002.

* cited by examiner

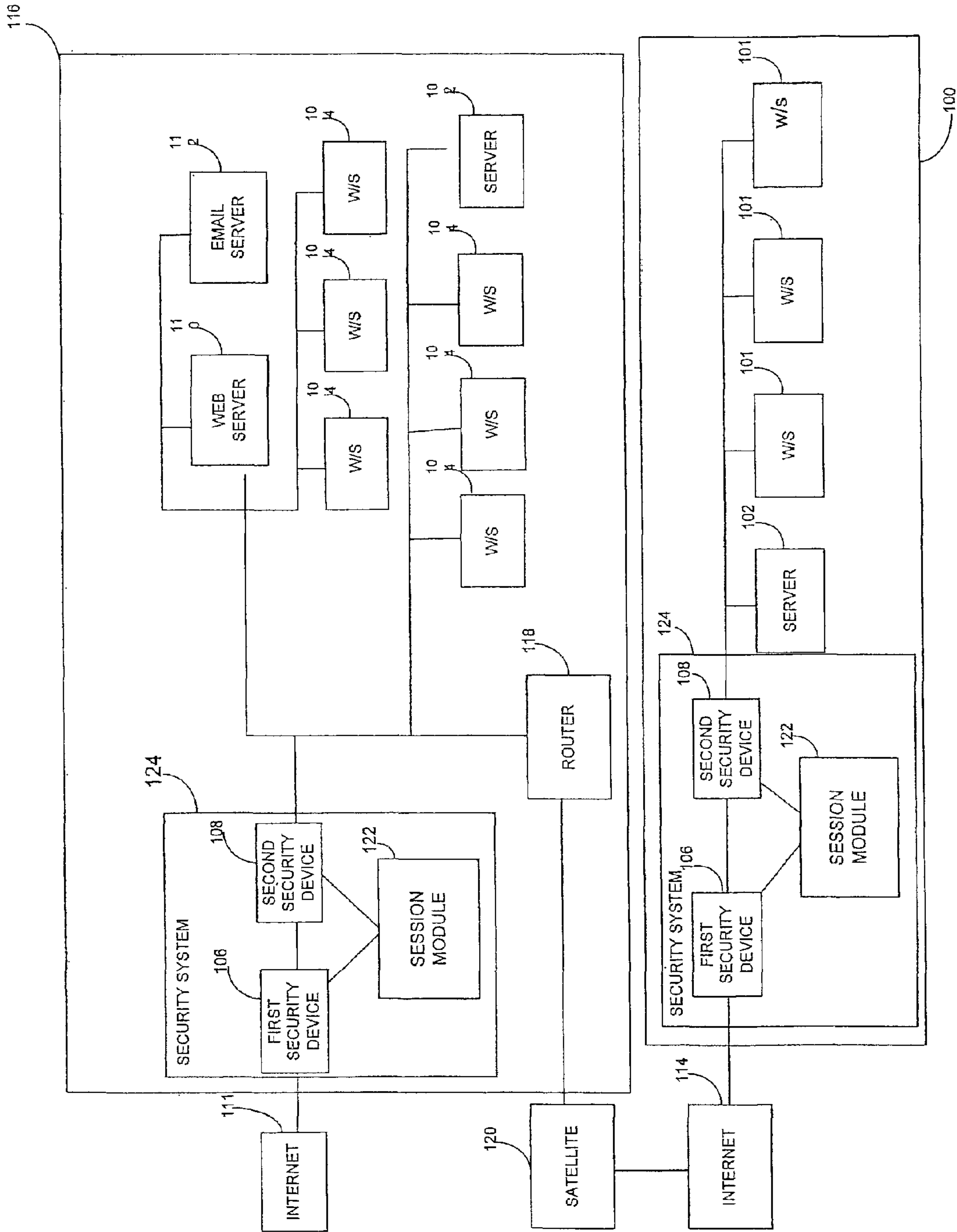


FIG. 1

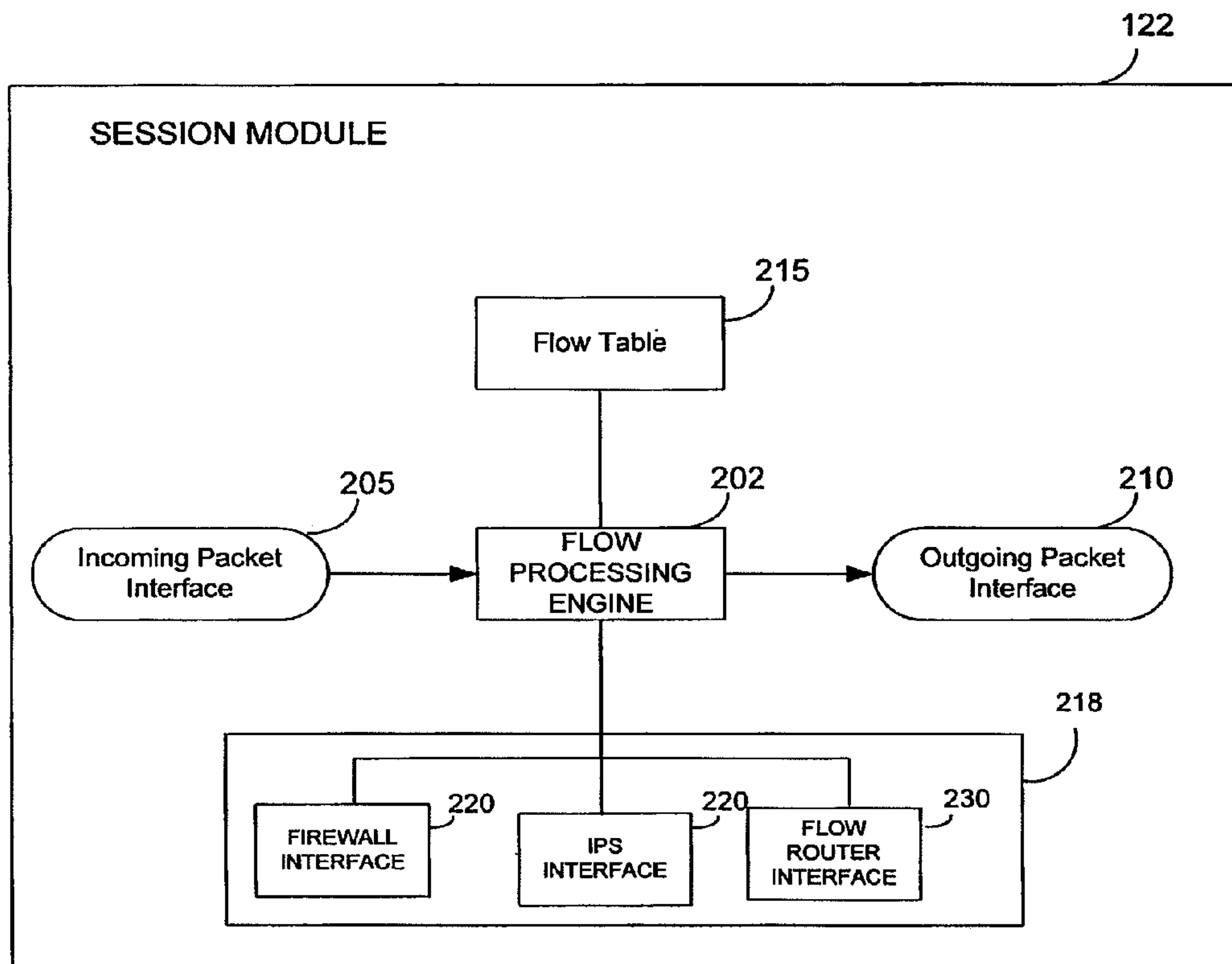


FIG. 2

	305	310	315	320	325
302	KEY	SECURITY DEVICE INFO 1	SECURITY DEVICE INFO 2	SECURITY DEVICE INFO 3	FLOW INFORMATION
302		RECORD #1			
302		RECORD #2			
302		.			
302		RECORD N			

FIG. 3

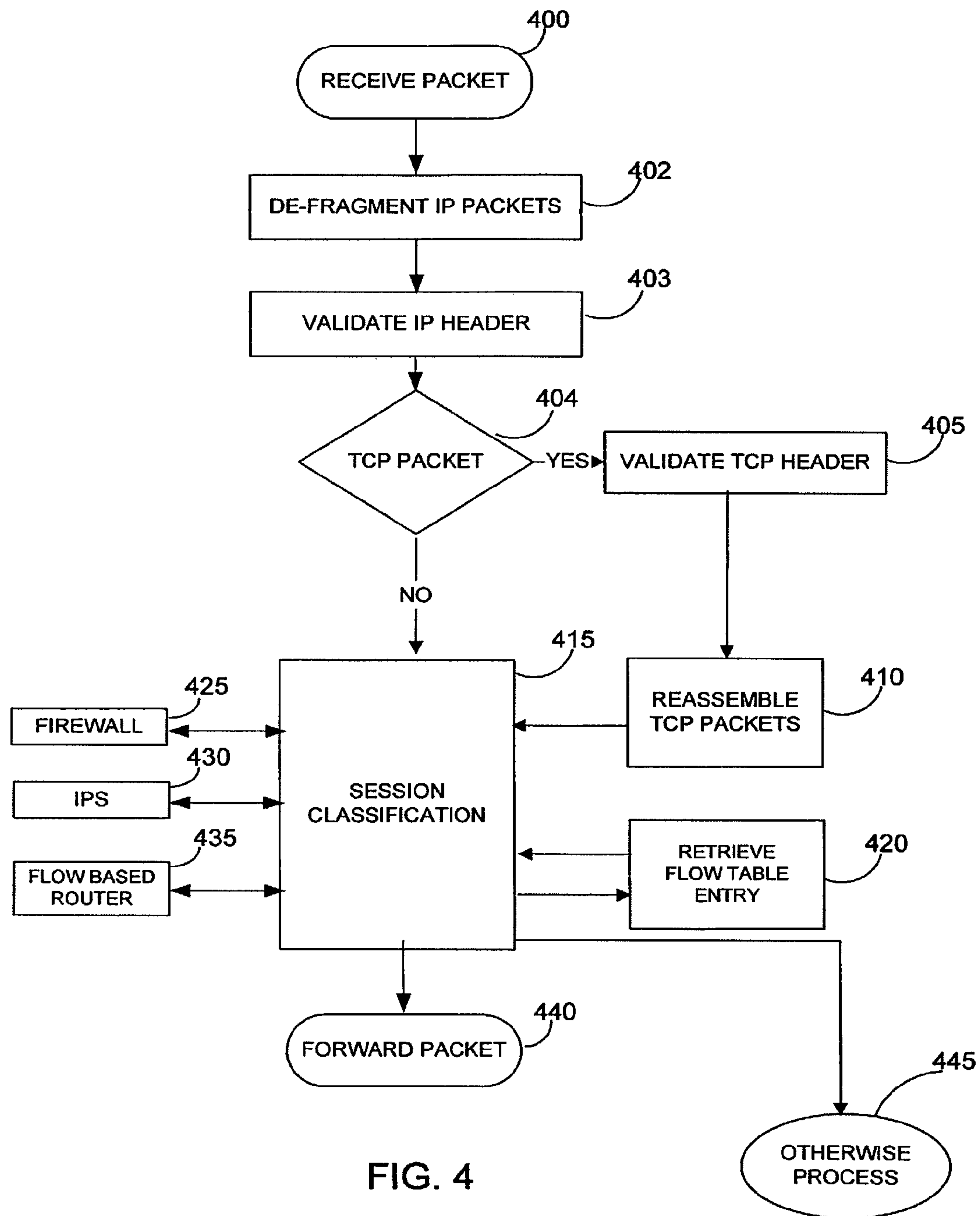


FIG. 4

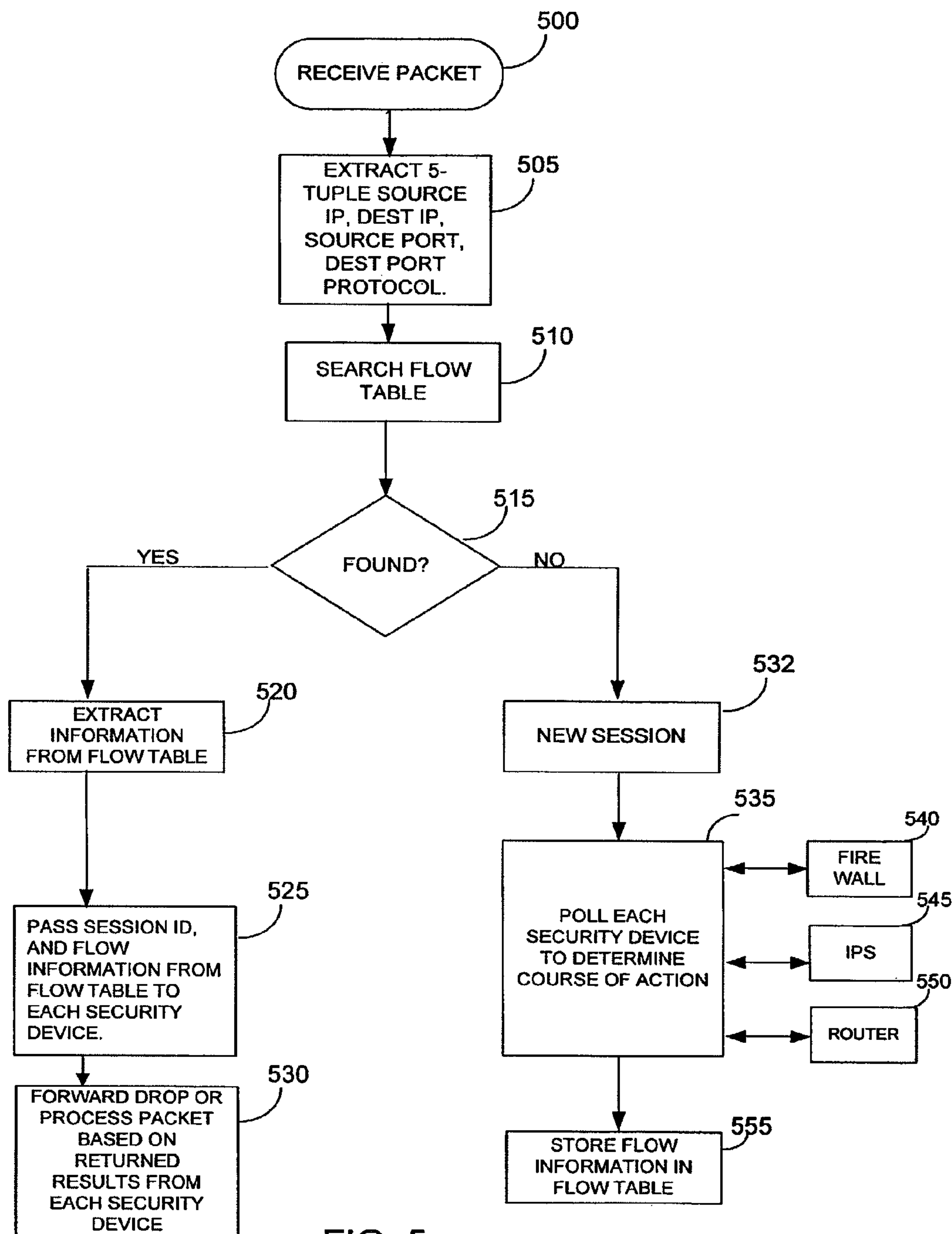


FIG. 5

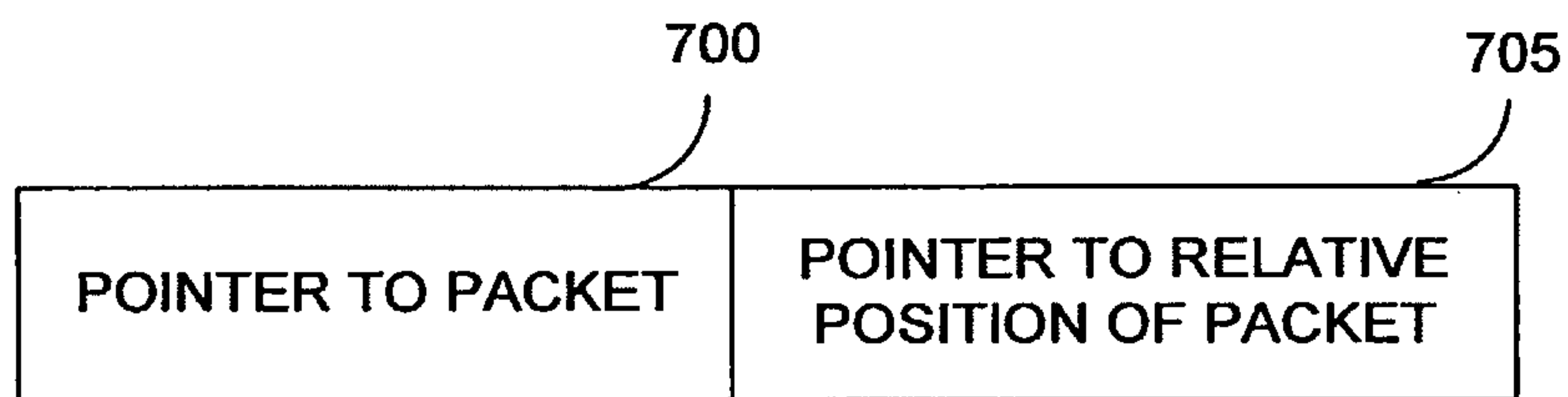


FIG. 6

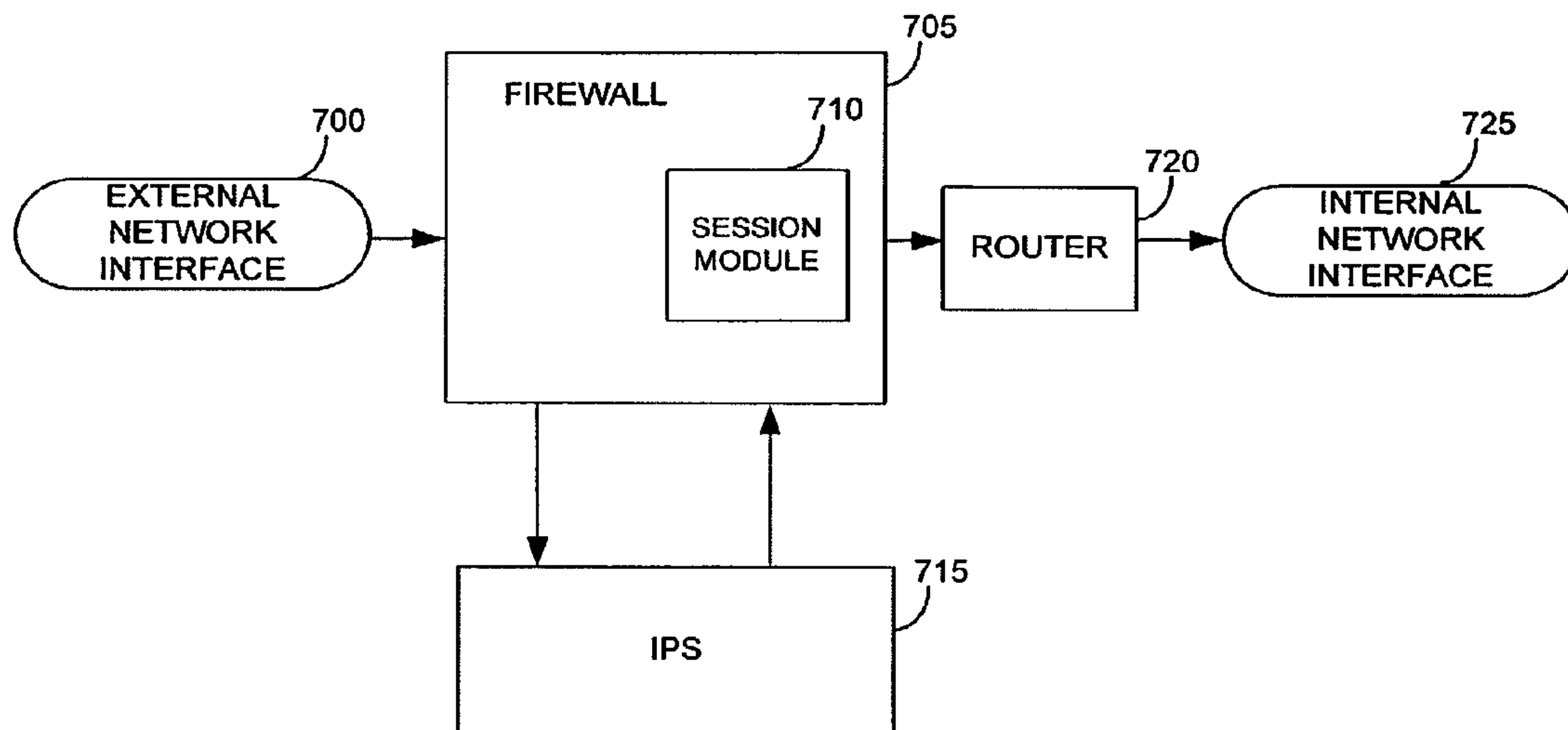


FIG. 7

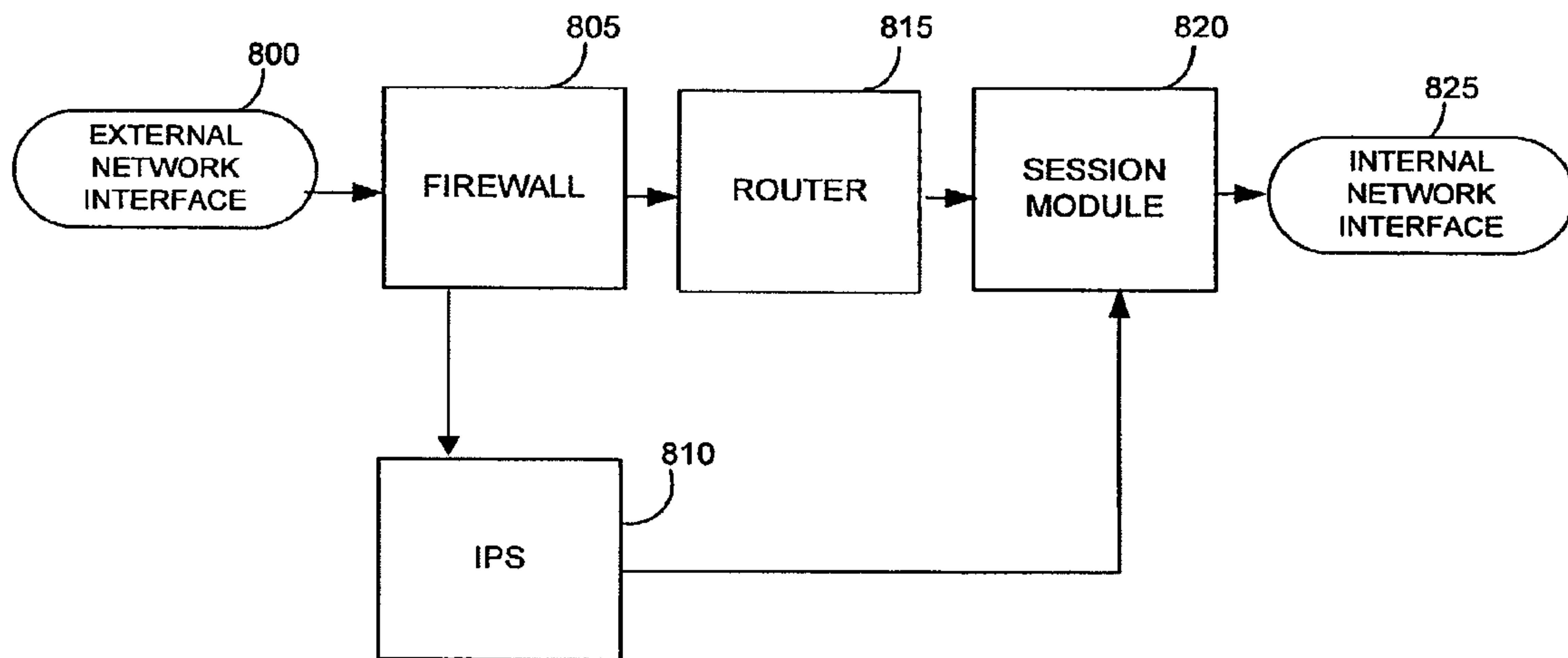


FIG. 8

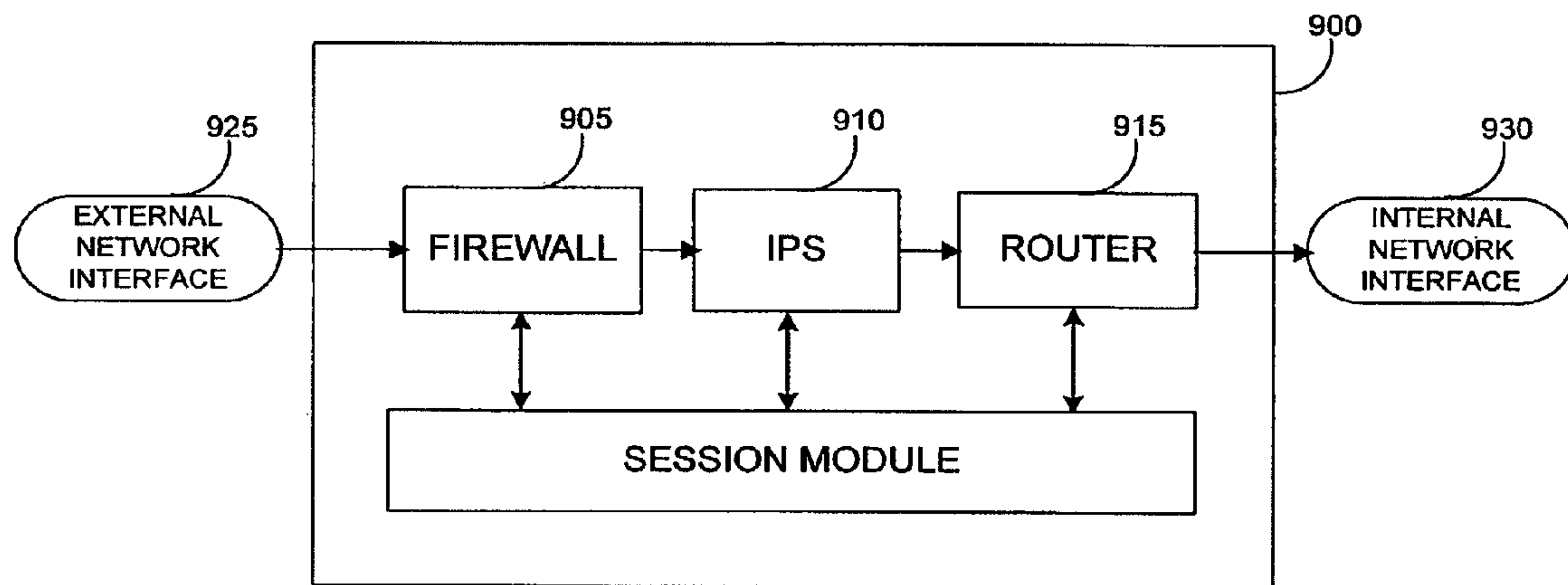


FIG. 9

INTELLIGENT INTEGRATED NETWORK SECURITY DEVICE

CROSS-REFERENCE TO RELATED APPLICATION

This application is a continuation of U.S. application Ser. No. 13/616,067, filed Sep. 14, 2012, which is a continuation of U.S. application Ser. No. 12/575,997, filed Oct. 8, 2009 (now U.S. Pat. No. 8,332,948), which is a continuation of U.S. application Ser. No. 10/402,920, filed Mar. 28, 2003 (now U.S. Pat. No. 7,650,634), which is a continuation-in-part of U.S. application Ser. No. 10/072,683, filed Feb. 8, 2002 (now U.S. Pat. No. 8,370,936). The disclosure disclosures of the prior applications are considered part of (and are incorporated by reference in) the disclosure of this application.

BACKGROUND

The present invention relates to a method for controlling computer network security. Firewalls and intrusion detection systems are devices that are used to protect a computer network from unauthorized or disruptive users. A firewall can be used to secure a local area network from users outside the local area network. A firewall checks, routes, and frequently labels all messages sent to or from users outside the local area network. An intrusion detection system (IDS) can be used to examine information being communicated within a network to recognize suspicious patterns of behavior. Information obtained by the IDS can be used to block unauthorized or disruptive users from accessing the network. An intrusion prevention system (IPS) is an in-line version of an IDS. An IPS can be used to examine information as it is being communicated within a network to recognize suspicious patterns of behavior.

A flow-based router (FBR) can allow network administrators to implement packet forwarding and routing according to network policies defined by a network administrator. FBRs can allow network administrators to implement policies that selectively cause packets to be routed through specific paths in the network. FBRs can also be used to ensure that certain types of packets receive differentiated, preferential service as they are routed. Conventional routers can forward packets to their destination address based on available routing information. Instead of routing solely based on the destination address, FBRs can enable a network administrator to implement routing policies to allow or deny packets based on several other criteria including the application, the protocol, the packet size and the identity of the end system.

A packet filter can operate on the data in the network layer, to defend a trusted network from attack by an untrusted network. Packet filters can operate at the network layer to inspect fields of the TCP/IP header including, the protocol type, the source and destination IP address, and the source and destination port numbers. Disadvantages of packet filters include, speed (i.e., slow) and management in large networks with complex security policies. Packet filters alone may not provide robust protection because packet filters are not aware of the context of the given communication. In addition, packet filters do not inspect the data at the application layer making packet filters vulnerable to attempted security intrusions using the application layer.

A proxy server can operate on values carried in the application layer to insulate a trusted network from an untrusted network. In an application proxy server, two TCP connections are established: one between the packet source and the proxy

server, another between the proxy server and the packet destination. The application proxy server can receive the arriving packets on behalf of the destination server. The application data can be assembled and examined by the proxy server, and a second TCP connection can be opened between the proxy server and the destination server to relay permitted packets to the destination server. Proxy servers can be slow because of the additional protocol stack overhead required to inspect packets at the application layer. Furthermore, because a unique proxy can be required for each application, proxy servers can be complex to implement and difficult to modify for supporting new applications. In addition, because proxy servers only examine application packets proxy servers may not detect an attempted network security intrusion at the TCP or network layers.

SUMMARY

The present invention provides methods and apparatus, including computer program products, for processing data packets and for implementing computer network security.

Advantages of the invention may include one or more of the following features. The technique disclosed can be used to detect an attempted network security intrusion and potentially block the current packet associated with the security intrusion. The disclosed technique can provide robust and efficient network security and includes plural security devices but only one flow table. Network security information is obtained from other network security devices and stored in a single flow record in the flow table. The use of a single flow record to determine whether a packet should be allowed can result in faster response time.

The details of one or more implementations of the invention are set forth in the accompanying drawings and the description below. Other features and advantages of the invention will become apparent from the description, the drawings, and the claims.

DESCRIPTION OF DRAWINGS

FIG. 1 shows a network topology including a session module.

FIG. 2 illustrates a block diagram of the session module.

FIG. 3 shows the structure of a flow table.

FIG. 4 is a flowchart describing the operation of the session module.

FIG. 5 is a flowchart describing session classification.

FIG. 6 shows the quasi-reassembly information generated by the session module.

FIG. 7 shows a network topology where the session module is included in a firewall.

FIG. 8 shows a network topology where the session module operates in series with a firewall, IPS, and router.

FIG. 9 shows a network topology where a session module, firewall, IPS and router are included in a single security device.

Like reference numbers and designations in the various drawings indicate like elements.

DETAILED DESCRIPTION

FIG. 1 shows a network topology including a local area network (LAN) (100), including a server (102), several workstations (W/S) (104), and a security device 124. The security system 124 can include a session module 122 and a plurality of other security devices. In the implementation shown, the security system 124 includes two security devices, a first

security device **106** and a second security device **108**. The LAN **100** is connected to an external network e.g., the Internet (**114**), through the security system **124**. The LAN **100** is also connected to a second LAN (**116**) through a router (**118**), and satellite (**120**). Second LAN **116** includes a web server (**110**), an email server (**112**), a server **102**, several workstations **104** and a security system **124**. The computers, servers and other devices in the LAN are interconnected using a number of data transmission media such as wire, fiber optics, and radio waves. The session module **122** monitors packets being communicated within the network. In one implementation, the first security device **106** can be a firewall and the second security device **108** can be an IPS. The session module **122** can act in conjunction with the first security device **106** and the second security device **108** to facilitate the blocking of packets associated with an attempted network security intrusion.

FIG. **2** shows a block diagram of a session module **122**. The session module **122** includes an incoming packet interface **205** for receiving packets. The received packets are analyzed by a flow processing engine (FPE) **202** to determine if an attempted network security intrusion is in progress. The session module **122** also includes a flow table **215**. The flow table **215** is used to store information regarding flows associated with received packets. The session module **122** also includes interfaces to other security devices on the network. In one implementation, the session module **122** includes a firewall interface **220**, an IPS interface **225**, and a flow router interface **230**. The security device interfaces **220** are used by the session module to obtain information regarding the received packet, and information regarding the flow associated with the packet, in order to determine if the received packet should be allowed or modified. The security device interfaces **218** are also used by the session module **122** to communicate flow information required by the security devices to facilitate processing of the packet.

FIG. **3** illustrates a structure of a flow table **300**. The flow table **300** includes flow records **302** associated with current TCP/IP flows. A TCP/IP flow includes a sequence of data packets communicating information between a source and a destination in one direction. The flow records are indexed using an indexing key **305**. The indexing key **305** is used to store and retrieve the appropriate flow record associated with a received packet. In one implementation, the indexing key **305** can be a hash key and the flow table **300** can be implemented as a hash table. The session module **122** (FIG. **2**) stores instructions for two or more security devices on the network in the same flow record. In one implementation of the session module **122**, instructions for three security devices (i.e. devices **310**, **315**, and **320**) are stored in the flow record **302**. The flow record **302** can store policy information (firewall policy, IPS policy etc., to apply to the flow) as well as other information that is used by the security devices such as encryption parameters, address translation parameters, book-keeping information, and statistics. The flow record **302** can also include flow information **325** required by the session module **122** in order to decide whether the packet should be allowed. Such information can include information required to implement network policies regarding, for example connection time out, time billing, and bandwidth usage. Flows, sessions and flow tables are described in greater detail in co-pending and commonly owned patent application entitled "Multi-Method Gateway-Based Network Security Systems and Methods," and assigned Ser. No. 10/072,683, the contents of which are expressly incorporated herein by reference.

FIG. **4** is a flow diagram describing the operation of the FPE **202** (FIG. **2**). Referring now to FIGS. **2** and **4**, incoming

packets are received by the session module (step **400**). IP packets are de-fragmented (step **402**) and the IP header is validated for each IP packet (step **403**). In the validation step, the IP header associated with a given packet is extracted and the extracted IP header is inspected for fundamental flaws. Thereafter FPE **202** determines if the session is to be allowed (step **415**).

If the packet is a TCP packet (step **404**), the TCP header is validated (step **405**) and the TCP packets are reassembled (step **410**). The validation process includes extracting TCP header data and evaluating the header for fundamental flaws. The quasi-reassembly information developed in step **410** can be communicated by the session module **122** to other security devices to facilitate processing of the packet by the other security devices. Reassembly is described in greater detail below and in "Multi-Method Gateway-Based Network Security Systems and Methods."

In step **415**, FPE **202** performs session classification using the TCP/IP header data associated with a given received packet. The session module **122** can determine if the packet should be allowed based on information obtained regarding the TCP/IP flow associated with the received packet and retrieved from the flow table **420**. In addition, the session module **122** can use information returned from one of the other security devices e.g., the firewall **425**, the IPS **430**, and the flow based router **435**. Further, the session module **122** can also facilitate the operation of the security devices by communicating flow information to a respective device as required by the device to process a given packet. Finally, FPE **202** forwards the packet if the packet should be allowed (step **440**). Otherwise, the packet is otherwise processed at step **445**. Other processing can include logging particular information regarding the packet, holding the packet, modifying and/or dropping the packet. This completes the description of the operation of FPE **202**.

FIG. **5** is a flow diagram showing the steps included in session classification (step **415**). The session classification step receives a packet (step **500**) and extracts information required to determine whether the packet should be allowed. The extracted information can include the source and destination IP addresses, the source and destination port numbers, and the protocol (step **505**). The extracted information can be used to search the flow table (step **510**) in order to determine if the packet is associated with a known session flow. For a known session flow, step **510** will produce a matching flow record in the flow table (step **515**). If a matching flow record is found, the FPE **202** (FIG. **2**) can extract TCP/IP session information for the received packet (step **520**) from the matching flow record. The FPE **202** determines if the received packet should be allowed using the TCP/IP session information obtained during step **520**. More specifically, the FPE **202** extracts information from the matching flow record, and passes the information to the security devices (e.g., communicating the session ID and the TCP/IP session information as well as any other security device specific information from the flow record) (step **525**). Depending on the returned results from the security devices, the FPE **202** can forward, drop, log, store, modify or otherwise process the given packet (step **530**).

If a matching flow record is not found in the flow table during step **515**, the received packet can be associated with a new TCP/IP session (step **532**). For a new TCP/IP session, the FPE **202** can assign a session ID to the new session and the FPE **202** can communicate with the other security devices (e.g. firewall, IPS, flow router) to determine a security policy for packets associated with the new session. For example, the FPE **202** can obtain information from the firewall **540** in order

5

to determine if received packets associated with the new session should be allowed. The FPE 202 can communicate with the IPS 545 in order to determine if the received packet should be blocked because it matches known attack signatures for attempted network security intrusions. The FPE 202 can obtain any network policy associated with the new session from the flow router 550. The FPE 202 can act as an arbiter between the different security devices and use the information obtained from the security devices either individually or in combination to determine if the packets associated with the new TCP/IP session should be allowed. The FPE 202 can use the information obtained from the security devices to create a new flow record and store the new flow record in the flow table (step 555). The new flow record includes the TCP/IP session information for the new session associated with the received packet and any other specific security device information. Thereafter, the FPE 202 can facilitate the processing of received packets associated with a given TCP/IP session as described above in association with FIG. 4 including communicating the session ID, TCP/IP session information and security device specific information to the security devices from a corresponding flow record.

In addition to determining if a received packet is associated with an attempted network security intrusion using the varied security devices, the session module can also perform quasi-reassembly of the received TCP/IP packets as described above in association with FIG. 4. FIG. 6 shows the quasi-reassembly information that can be generated by the session module. The quasi-reassembly information can include a pointer to a location of a given packet 600 in memory and a pointer to information containing the relative position of the packet in a flow 605. In one implementation, an IPS can perform passive TCP/IP reassembly and the pointer to the location of the packet can be used to locate the packet within the IPS. In another implementation, the pointer to information containing the relative position of the packet in the flow can be used to obtain the TCP/IP sequence number included in the TCP/IP header associated with the packet. The quasi-reassembly information can be communicated to the security devices connected to the session module 122 (FIG. 2) as required. The security devices can use the quasi-reassembly information to process the received packet.

The session module can be used in a number of different network topologies. FIG. 7 shows a network topology where a session module 710 is integrated into a firewall 705. The firewall 705 can include an interface to a router 720 and an IPS 715. The firewall 705 receives packets from the external network interface 700. The firewall 705 communicates with the IPS 715 to determine whether the received packet should be blocked based on known attack signatures. If the firewall 705 and IPS 715 determine that the packet should be allowed to pass, the firewall 705 sends the received packet to the router 720. The router 720 forwards the outgoing packet to its intended destination, using the internal network interface 725, based on the network policies stored in the router.

FIG. 8 shows an alternate arrangement for implementing computer network security using a session module. In this arrangement, the session module 820 operates in series with a firewall 805, an IPS 810, and a router 815. Packets received using the external network interface 800 are screened by the firewall 805 before being communicated to the router 815. The firewall 805 also sends information regarding the received packet to the IPS 810. The IPS 810 examines the received packet and informs the session module 820 if the received packet should be blocked based on known attack signatures. The router 815 sends the packet to the session module 820 for further processing. If the session module 820

6

determines that the received packet should be allowed it forwards the received packet to its intended destination using the internal network interface 825.

The invention can be implemented in digital electronic circuitry, or in computer hardware, firmware, software, or in combinations of them. The invention can be implemented as a computer program product, i.e., a computer program tangibly embodied in an information carrier, e.g., in a machine-readable storage device or in a propagated signal, for execution by, or to control the operation of, data processing apparatus, e.g., a programmable processor, a computer, or multiple computers. A computer program can be written in any form of programming language, including compiled or interpreted languages, and it can be deployed in any form, including as a stand-alone program or as a module, component, subroutine, or other unit suitable for use in a computing environment. A computer program can be deployed to be executed on one computer or on multiple computers at one site or distributed across multiple sites and interconnected by a communication network.

Method steps of the invention can be performed by one or more programmable processors executing a computer program to perform functions of the invention by operating on input data and generating output. Method steps can also be performed by, and apparatus of the invention can be implemented as, special purpose logic circuitry, e.g., an FPGA (field programmable gate array) or an ASIC (application-specific integrated circuit).

Processors suitable for the execution of a computer program include, by way of example, both general and special purpose microprocessors, and any one or more processors of any kind of digital computer. Generally, a processor will receive instructions and data from a read-only memory or a random access memory or both. The essential elements of a computer are a processor for executing instructions and one or more memory devices for storing instructions and data. Generally, a computer will also include, or be operatively coupled to receive data from or transfer data to, or both, one or more mass storage devices for storing data, e.g., magnetic, magneto-optical disks, or optical disks. Information carriers suitable for embodying computer program instructions and data include all forms of nonvolatile memory, including by way of example semiconductor memory devices, e.g., EPROM, EEPROM, and flash memory devices; magnetic disks, e.g., internal hard disks or removable disks; magneto-optical disks; and CD-ROM and DVD-ROM disks. The processor and the memory can be supplemented by, or incorporated in special purpose logic circuitry.

The invention can be implemented in a computing system that includes a back-end component, e.g., as a data server, or that includes a middleware component, e.g., an application server, or that includes a front-end component, e.g., a client computer having a graphical user interface or a Web browser through which a user can interact with an implementation of the invention, or any combination of such back-end, middleware, or front-end components. The components of the system can be interconnected by any form or medium of digital data communication, e.g., a communication network. Examples of communication networks include a local area network ("LAN") and a wide area network ("WAN"), e.g., the Internet.

The computing system can include clients and servers. A client and server are generally remote from each other and typically interact through a communication network. The relationship of client and server arises by virtue of computer programs running on the respective computers and having a client-server relationship to each other.

7

This invention has been described in terms of particular embodiments. Nevertheless, it will be understood that various modifications may be made without departing with the spirit and scope of the invention. For instance, the steps of the invention can be performed in a different order and still achieve desirable results. In addition, the session module, IPS, firewall, and router can all be incorporated into a single device such as the configuration shown in FIG. 9. Other configurations of a session module packaged with one or more security devices are also possible. Accordingly, other embodiments are within the scope of the following claims.

What is claimed is:

1. A method comprising:
 - receiving, by one or more processors of a device, a data packet;
 - examining, by the one or more processors, the data packet;
 - determining, by the one or more processors, a single flow record associated with the data packet,
 - determining the single flow record including:
 - determining a packet identifier associated with the data packet, and
 - evaluating a flow table to identify the single flow record using the packet identifier;
 - extracting, by the one or more processors, a session identifier and flow instructions, for two or more security devices, from the single flow record,
 - the session identifier identifying a session associated with the data packet,
 - the two or more security devices being included in the device;
 - sending, by the one or more processors, the flow instructions and the session identifier to respective ones of the two or more security devices to facilitate processing of the data packet;
 - receiving, by the one or more processors and from the two or more security devices, processing results,
 - the processing results being generated by the two or more security devices when processing the data packet based on the flow instructions and the session identifier; and processing, by the one or more processors, the data packet using the processing results.
2. The method of claim 1, where processing the data packet using the processing results includes:
 - dropping the packet; or
 - forwarding the packet toward a destination of the packet.
3. The method of claim 1, where the data packet is included in a flow of packets associated with the session, and where the single flow record further includes:
 - information identifying the flow of packets.
4. The method of claim 3, where sending the flow instructions and the session identifier to the respective ones of the two or more security devices includes:
 - sending, to the respective ones of the two or more security devices, the information identifying the flow of packets and the session identifier for the session.
5. The method of claim 3, where the single flow record further includes device specific information, and where sending the flow instructions and the session identifier to the respective ones of the two or more security devices includes:
 - sending, to the respective ones of the two or more security devices, the information identifying the flow of packets, the session identifier for the session, and the device specific information.

8

6. The method of claim 1, where the two or more security devices include two or more of:
 - an intrusion prevention system,
 - a firewall, or
 - a router.
7. The method of claim 1, where determining the packet identifier includes:
 - extracting, from the data packet, a source address, a destination address, a source port number, a destination port number, and a protocol type associated with the data packet,
 - the packet identifier being based on the source address, the destination address, the source port number, the destination port number, and the protocol type.
8. A non-transitory computer-readable medium storing instructions, the instructions comprising:
 - one or more instructions which, when executed by one or more processors of a device, cause the one or more processors to receive a data packet;
 - one or more instructions which, when executed by the one or more processors, cause the one or more processors to examine the data packet;
 - one or more instructions which, when executed by the one or more processors, cause the one or more processors to determine a single flow record associated with the data packet,
 - the one or more instructions to determine the single flow record including:
 - one or more instructions which, when executed by the one or more processors, cause the one or more processors to determine a packet identifier associated with the data packet, and
 - one or more instructions which, when executed by the one or more processors, cause the one or more processors to evaluate a flow table to identify the single flow record using the packet identifier;
 - one or more instructions which, when executed by the one or more processors, cause the one or more processors to extract a session identifier and flow instructions, for two or more security devices, from the single flow record,
 - the two or more security devices being included in the device;
 - the session identifier identifying a session associated with the data packet;
 - one or more instructions which, when executed by the one or more processors, cause the one or more processors to send the session identifier and the flow instructions to respective ones of the two or more security devices to facilitate processing of the data packet;
 - one or more instructions which, when executed by the one or more processors, cause the one or more processors to receive, from the two or more security devices, processing results,
 - the processing results being generated by the two or more security devices when processing the data packet based on the flow instructions and the session identifier; and
 - one or more instructions which, when executed by the one or more processors, cause the one or more processors to process the data packet using the processing results.
9. The non-transitory computer-readable medium of claim 8, where the one or more instructions to process the data packet using the processing results include:
 - one or more instructions which, when executed by the one or more processors, cause the one or more processors to selectively:

9

log information regarding the data packet based on the processing results,
store the data packet based on the processing results, or
modify the data packet based on the processing results.

10. The non-transitory computer-readable medium of claim **8**, where the instructions further include:

one or more instructions which, when executed by the one or more processors, cause the one or more processors to receive another data packet;

one or more instructions which, when executed by the one or more processors, cause the one or more processors to determine that a single flow record, associated with the other data packet, is not found in the flow table; and

one or more instructions which, when executed by the one or more processors, cause the one or more processors to create and store, in the flow table, a new single flow record associated with the other data packet.

11. The non-transitory computer-readable medium of claim **10**, where the other data packet is associated with another session, and

where the instructions further include:

one or more instructions which, when executed by the one or more processors, cause the one or more processors to communicate with the two or more security devices to determine a security policy for packets associated with the other session;

one or more instructions which, when executed by the one or more processors, cause the one or more processors to assign an identifier to the other session; and

one or more instructions which, when executed by the one or more processors, cause the one or more processors to store, in the new single flow record, the identifier of the other session and the security policy.

12. The non-transitory computer-readable medium of claim **8**, where the one or more instructions to process the data packet using the processing results include:

one or more instructions which, when executed by the one or more processors, cause the one or more processors to selectively:

drop the data packet based on the processing results, or
forward the data packet to a destination of the data packet based on the processing results.

13. The non-transitory computer-readable medium of claim **8**, where each security device, of the two or more security devices, includes a different one of:

an intrusion prevention system,
a firewall, or
a router.

14. The non-transitory computer-readable medium of claim **8**, where the one or more instructions to determine the packet identifier include:

one or more instructions which, when executed by the one or more processors, cause the one or more processors to extract, from the data packet, a source address, a destination address, a source port number, a destination port number, and a protocol type associated with the data packet, and

where the packet identifier is based on the source address, the destination address, the source port number, the destination port number, and the protocol type.

10

15. A system comprising:

a device to:

receive a data packet;

determine a packet identifier associated with the data packet;

search a flow table, using the packet identifier, to identify a single flow record associated with the data packet; extract a session identifier and flow instructions, for two or more security devices, from the single flow record, the two or more security devices being included in the device;

the session identifier identifying a session associated with the data packet;

send the session identifier and the flow instructions to respective ones of the two or more security devices to facilitate processing of the data packet;

receive, from the two or more security devices, processing results,

the processing results being generated by the two or more security devices when processing the data packet based on the flow instructions and the session identifier; and

selectively:

drop the data packet based on the processing results,
or

forward the data packet to a destination of the data packet based on the processing results.

16. The system of claim **15**, where each security device, of the two or more security devices, includes a different one of:

an intrusion prevention system,
a firewall, or
a router.

17. The system of claim **15**, where the data packet is included in a flow of packets associated with the session, and where the single flow record includes:

information identifying the flow of packets.

18. The system of claim **17**, where, when sending the flow instructions, the device is to:

send, to the respective ones of the two or more security devices, the information identifying the flow of packets and the session identifier for the session.

19. The system of claim **15**, where the device is further to: receive another data packet;

determine that a single flow record, associated with the other data packet, is not found in the flow table; and create and store, in the flow table, a new single flow record associated with the other data packet.

20. The system of claim **19**, where the other data packet is associated with another session, and

where the device is further to:

communicate with the two or more security devices to obtain information relating to processing packets associated with the other session;

assign another session identifier to the other session; and store, in the new single flow record, the other session identifier and the information relating to processing the packets associated with the other session.

* * * * *