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#### Hansen et al.

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# (54) REDUCED LATENCY CONCATENATED REED SOLOMON-CONVOLUTIONAL CODING FOR MIMO WIRELESS LAN

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#### Related U.S. Application Data

- (60) Provisional application No. 60/544,605, filed on Feb. 13, 2004, provisional application No. 60/545,854, filed on Feb. 19, 2004, provisional application No. 60/568,914, filed on May 7, 2004, provisional application No. 60/573,781, filed on May 24, 2004, provisional application No. 60/575,909, filed on Jun. 1, 2004, provisional application No. 60/587,068, filed on Jul. 12, 2004.
- (51) Int. Cl. *H04L 23/02* (2006.01)

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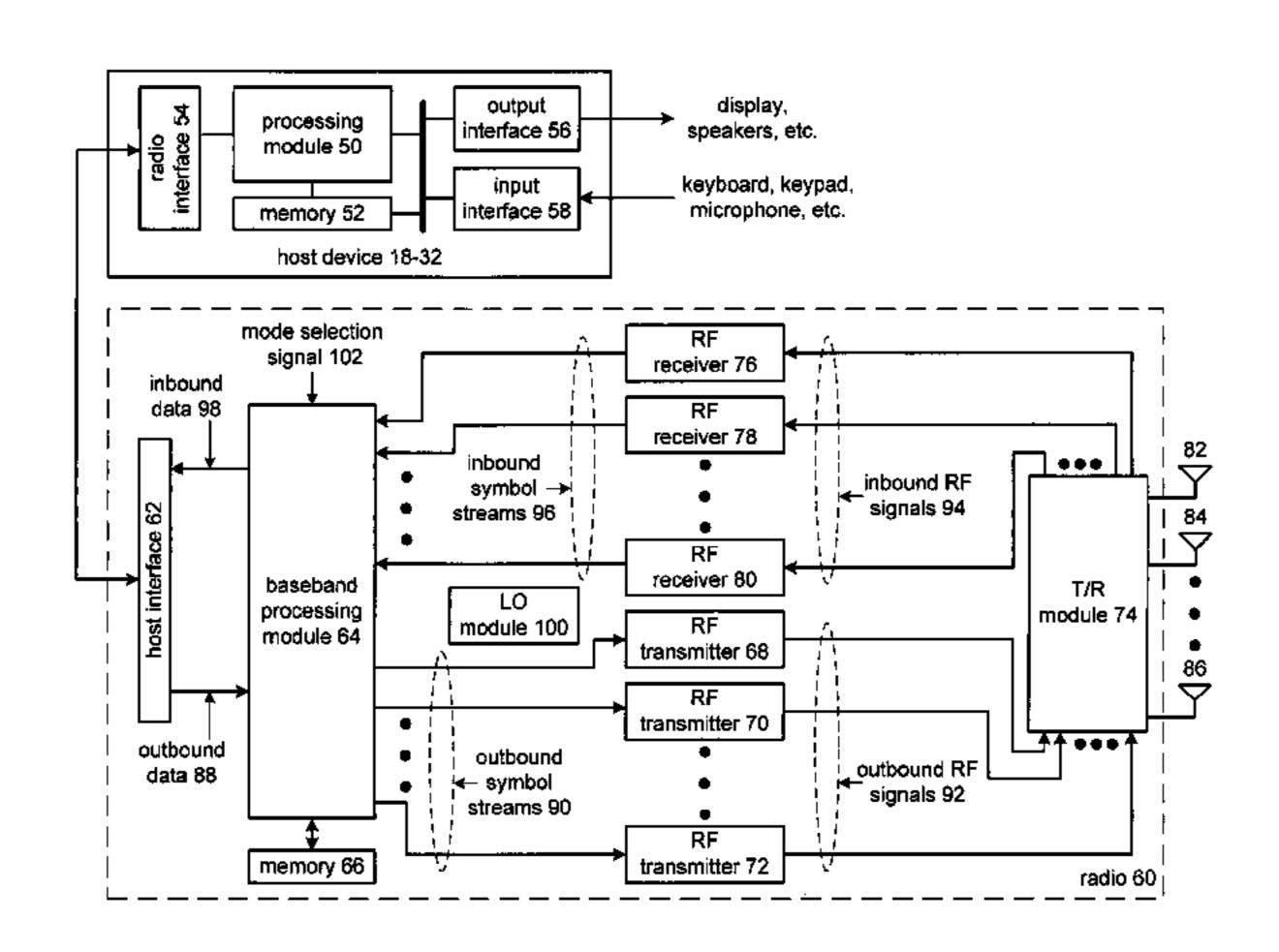
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### (57) ABSTRACT

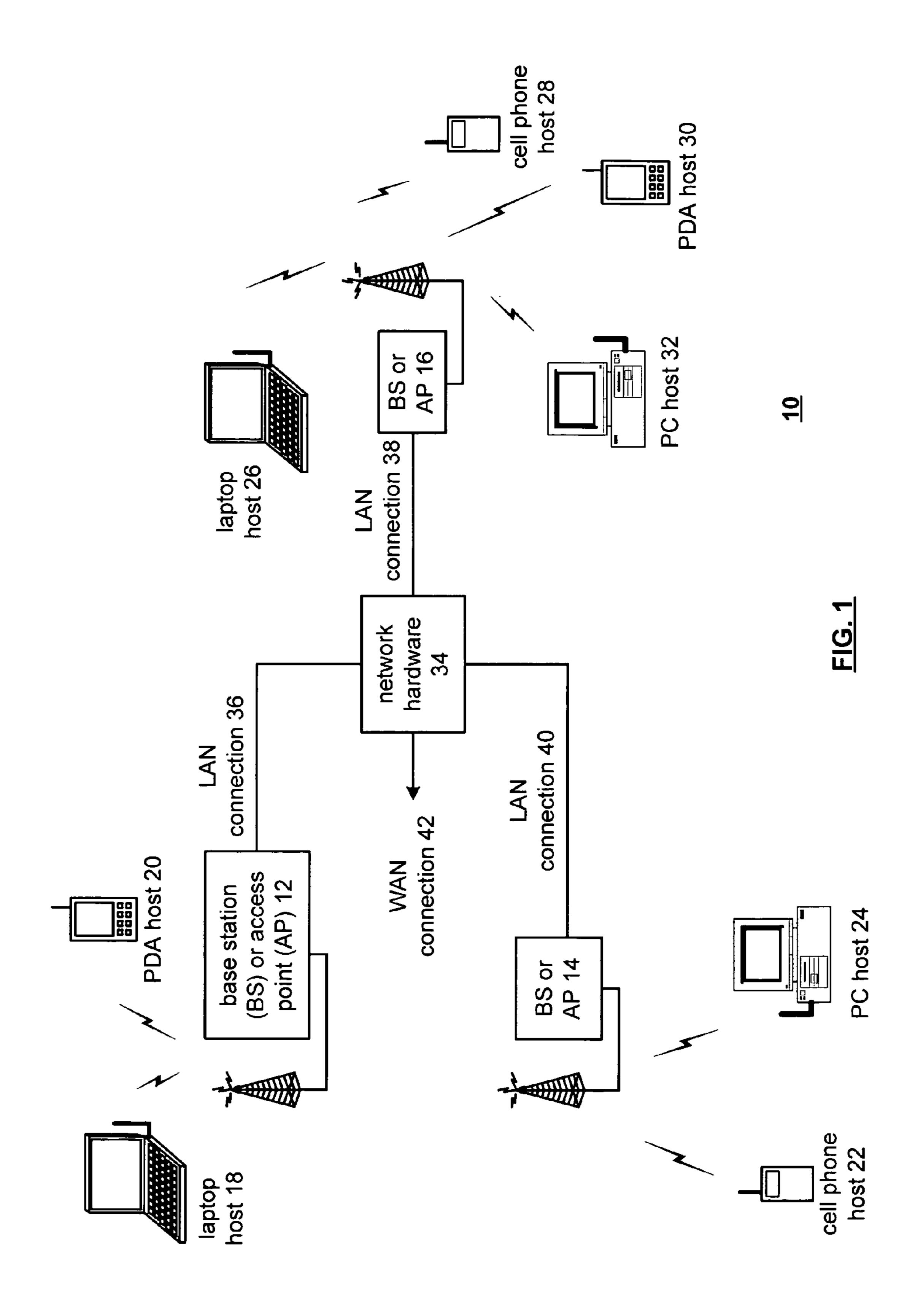
A wireless local area network (WLAN) transmitter includes a baseband processing module and a plurality of radio frequency (RF) transmitters. The baseband processing module operably coupled to scramble data in accordance with a pseudo random sequence to produce scrambled data. The baseband processing module is further operably coupled to interleave, at a word level, the scrambled data to produce interleaved data when the interleaving is enabled. The baseband processing module is further operably coupled to outer Reed-Solomon encode the scrambled data or the interleaved data to produce outer encoded data when the outer Reed-Solomon encoding is enabled. The baseband processing module is further operably coupled to inner puncture convolution encode the outer encoded data or the scrambled data to produce the encoded data. The baseband processing module is further operably coupled to determine a number of transmit streams based on a mode selection signal. The baseband processing module is further operably coupled to convert the encoded data into streams of symbols in accordance with the number of transmit streams and the mode selection signal. The plurality of radio frequency (RF) transmitters, when enabled, converts the streams of symbols into a corresponding number of RF signals.

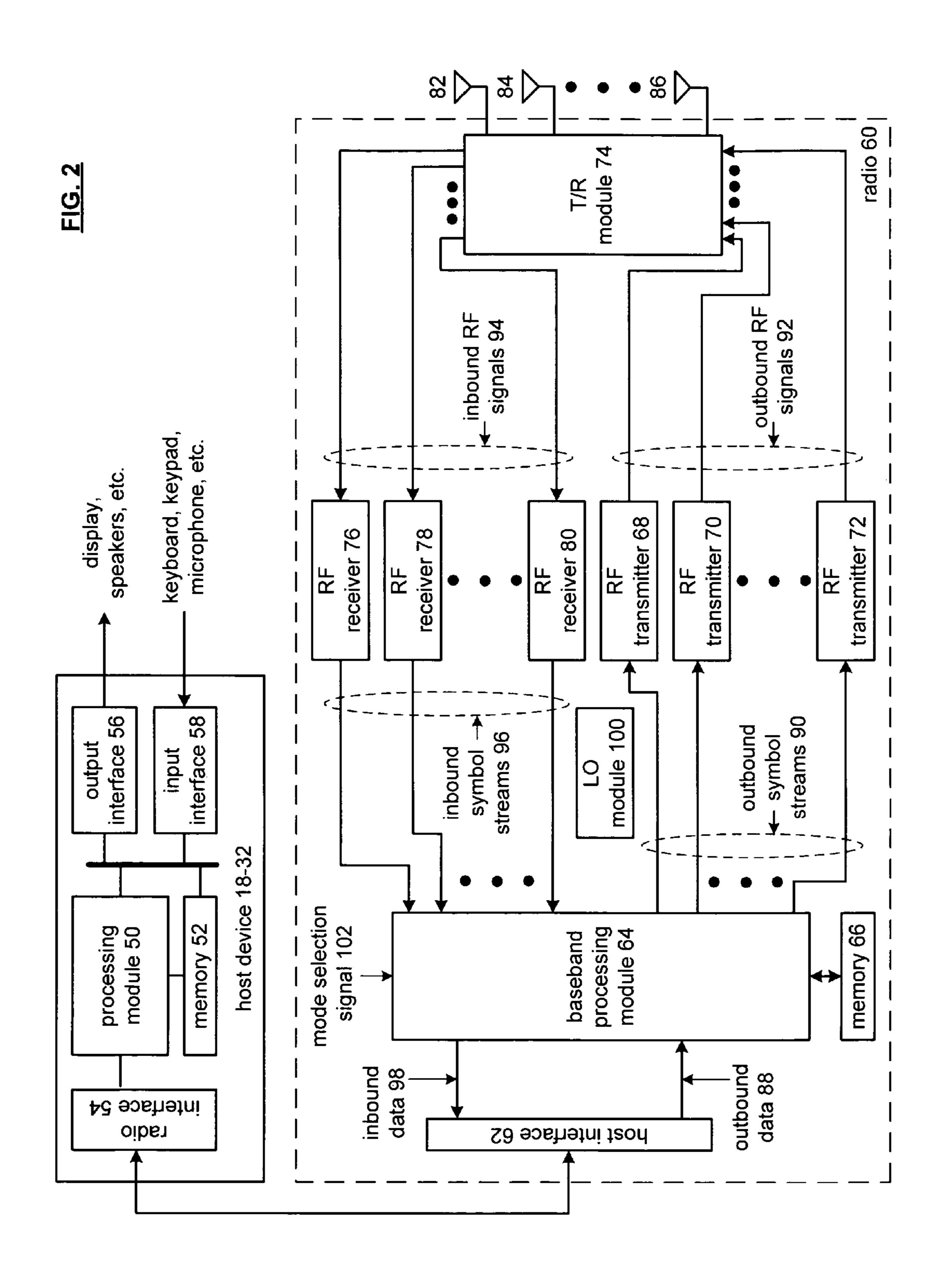
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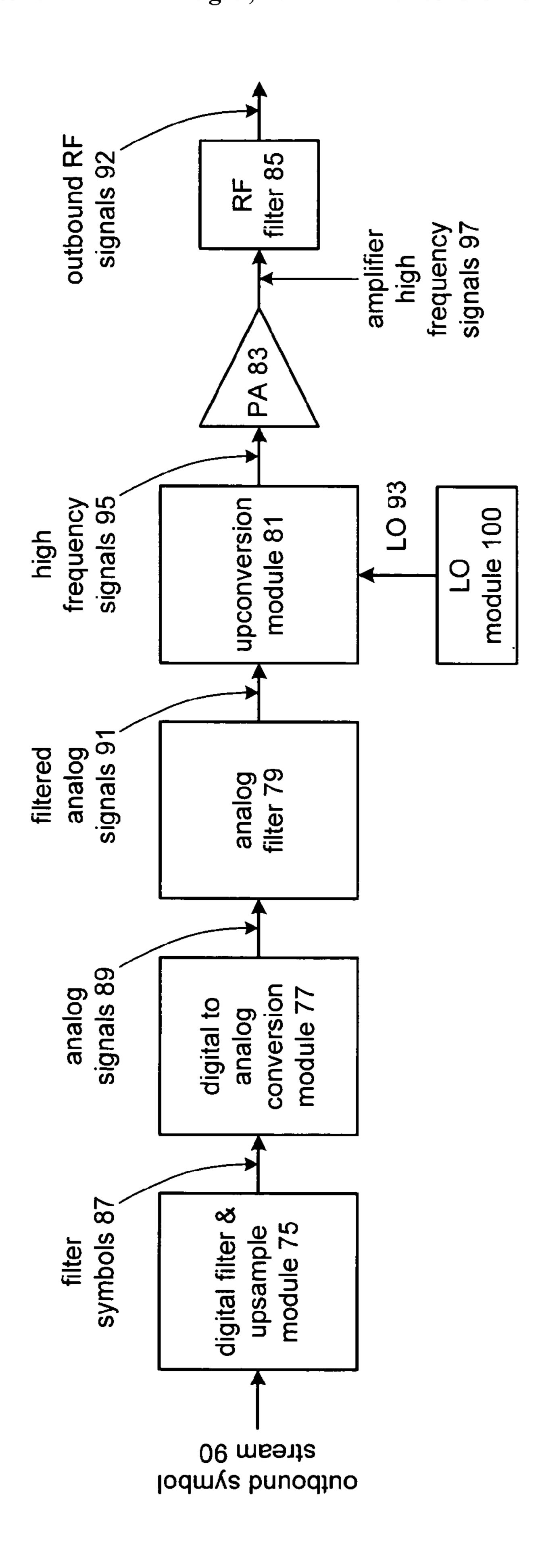


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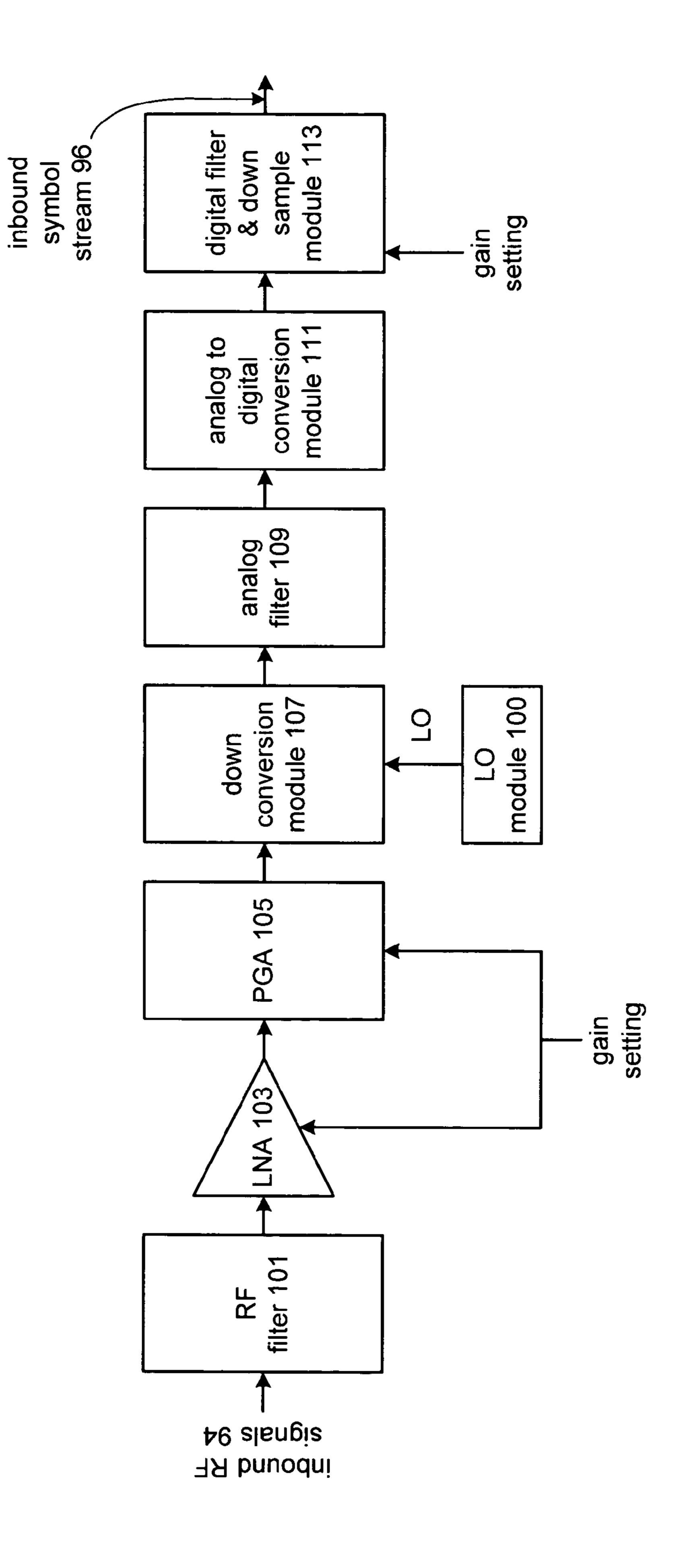
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Fransmitter 68 - 7



RF receiver 76 - 80

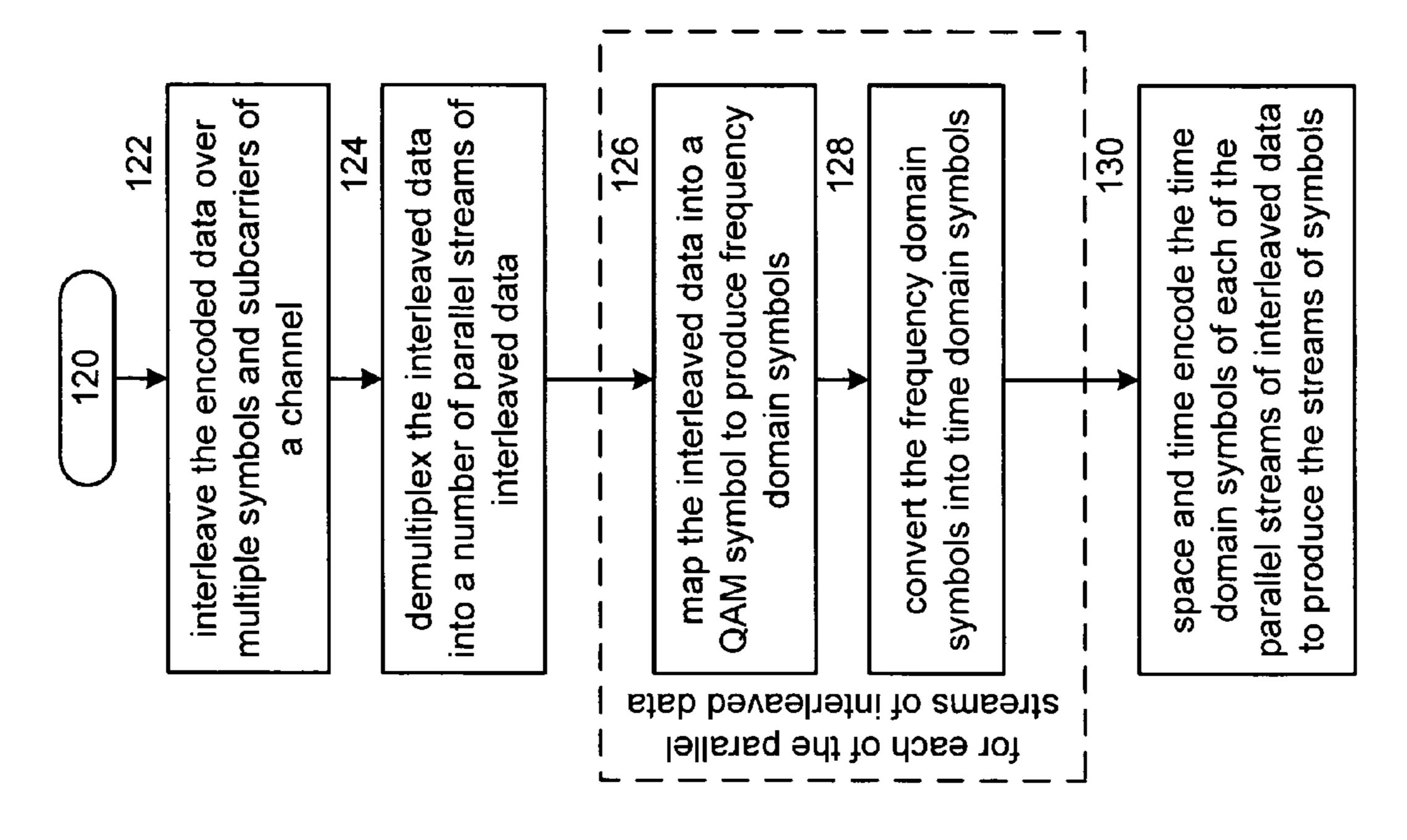


FIG. 6

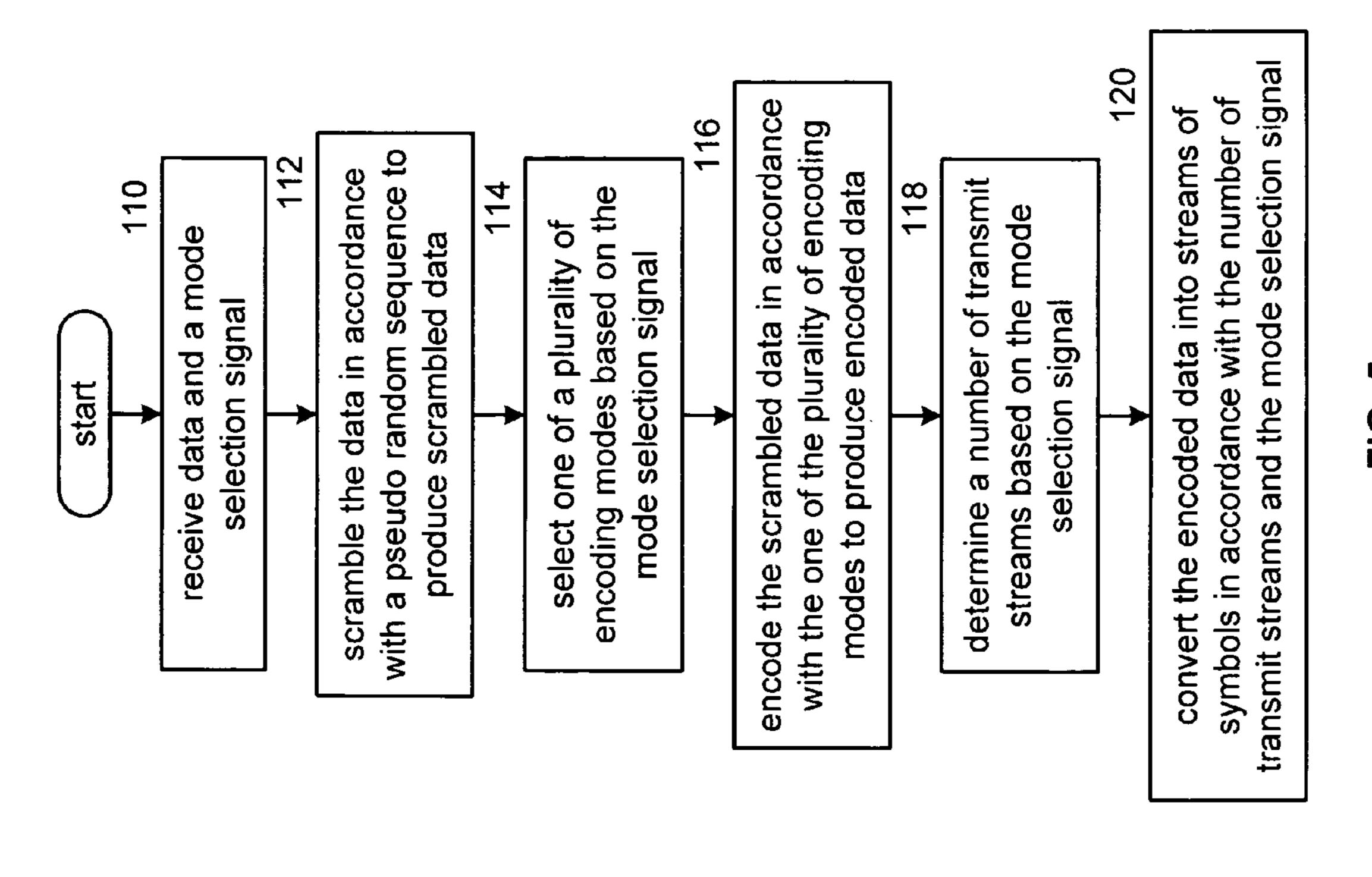


FIG. 5

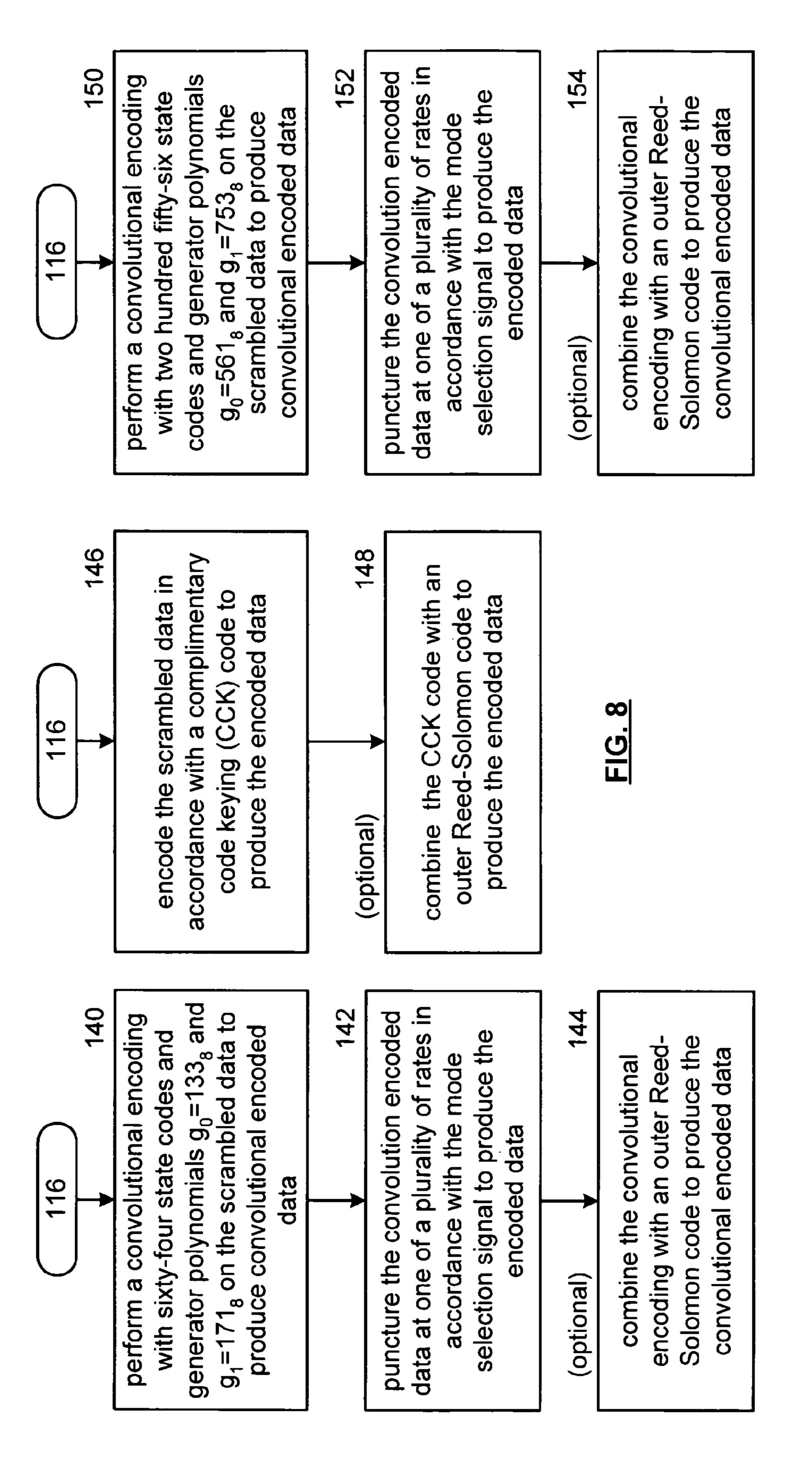
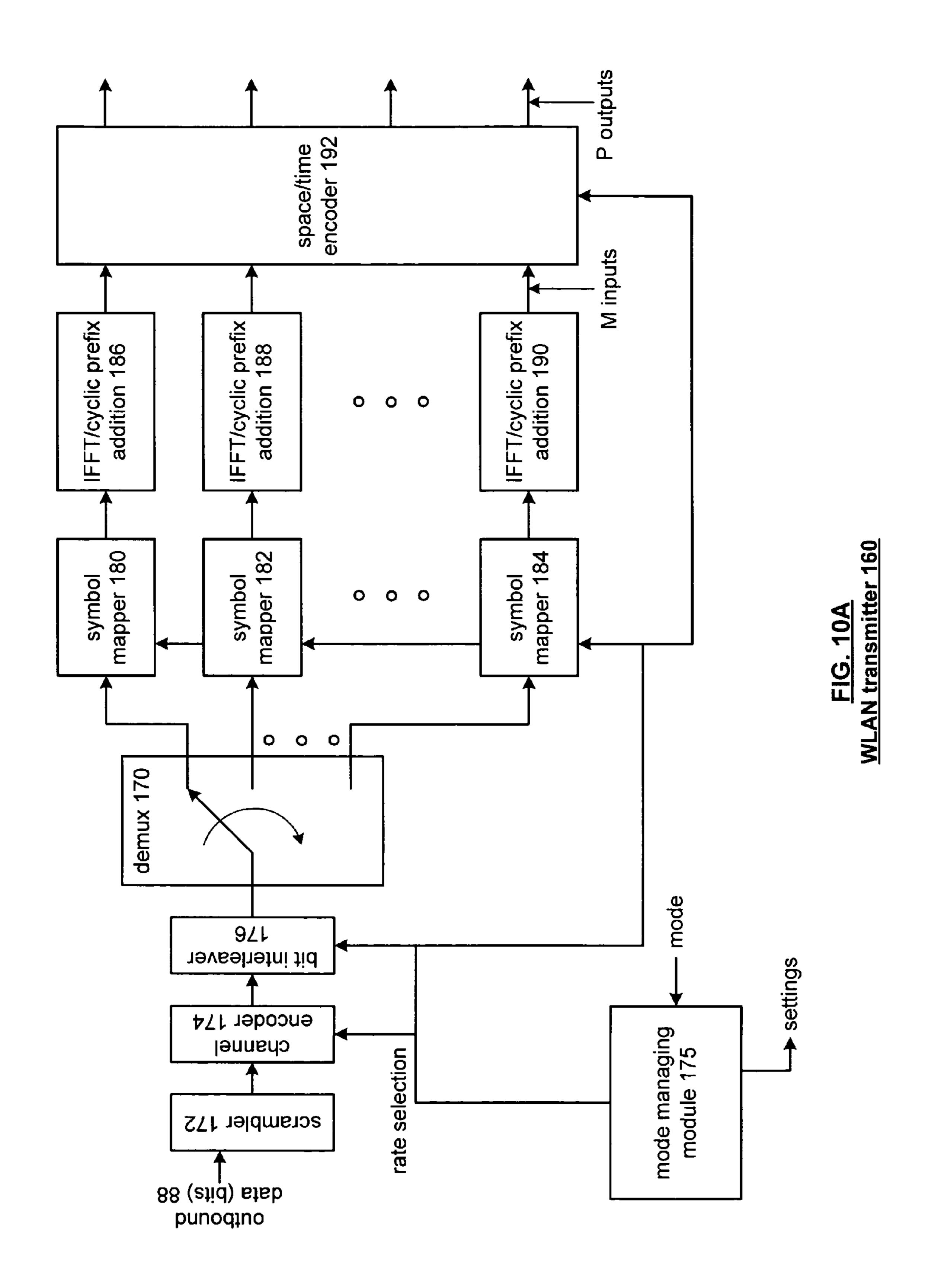
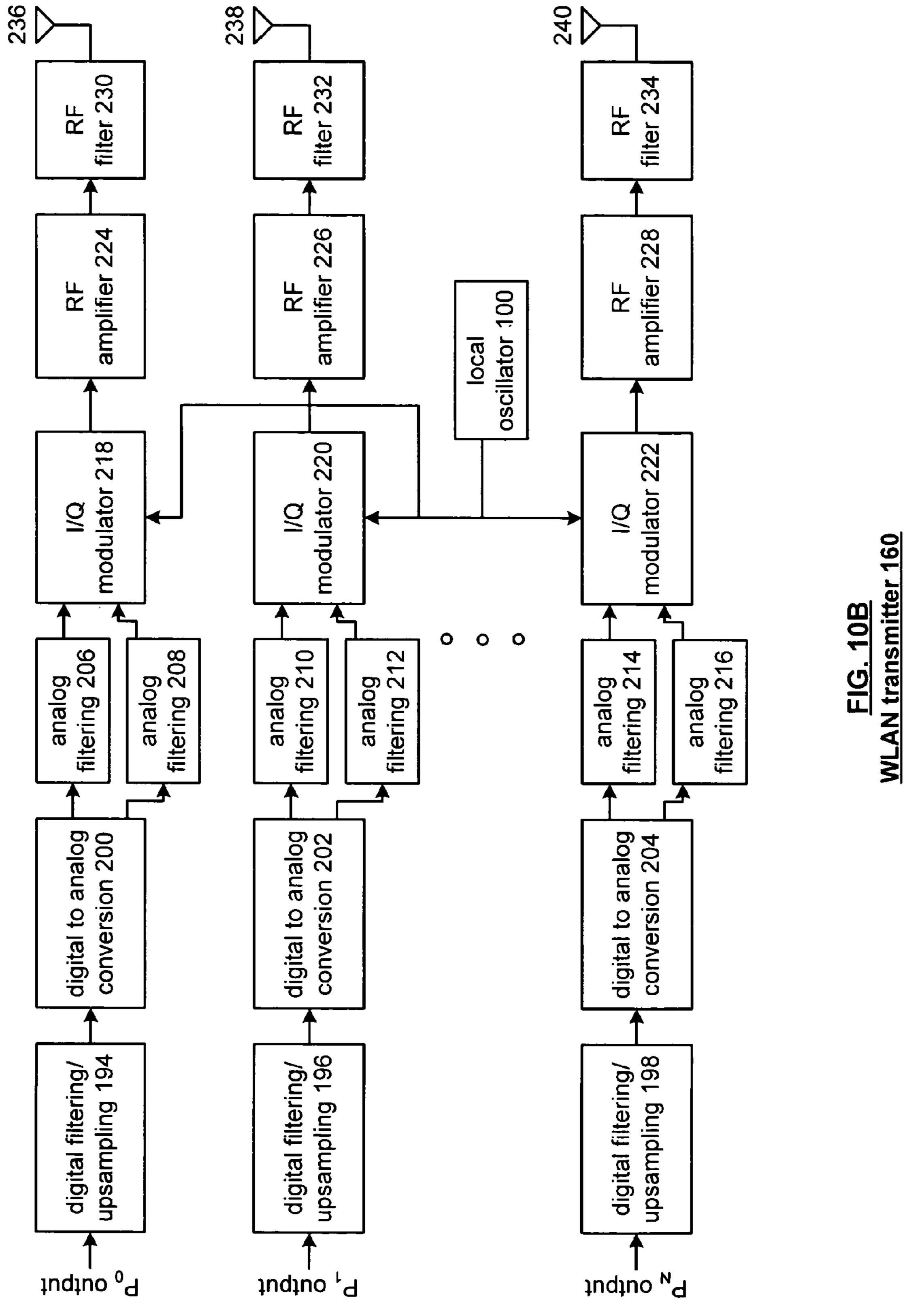
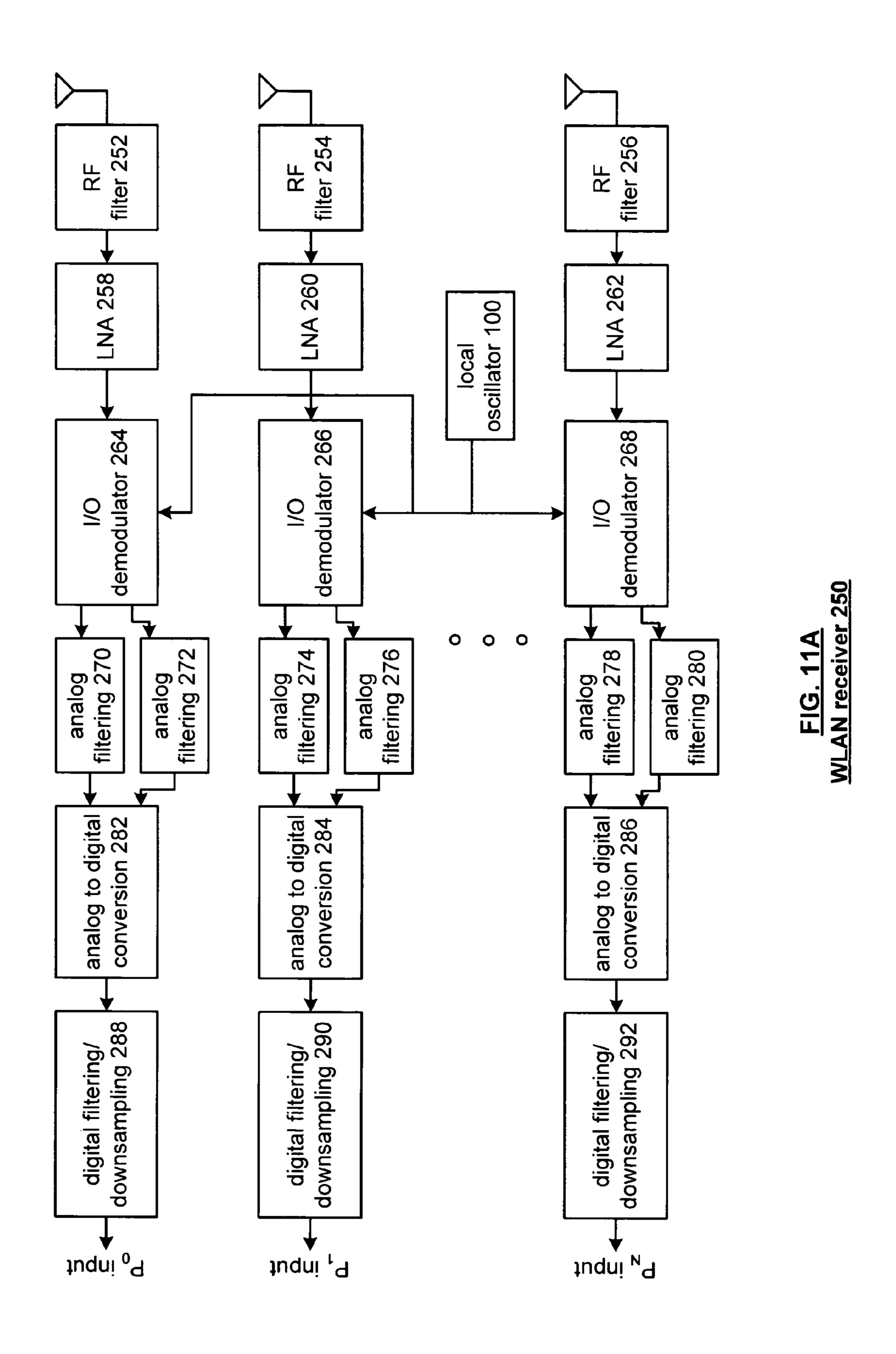
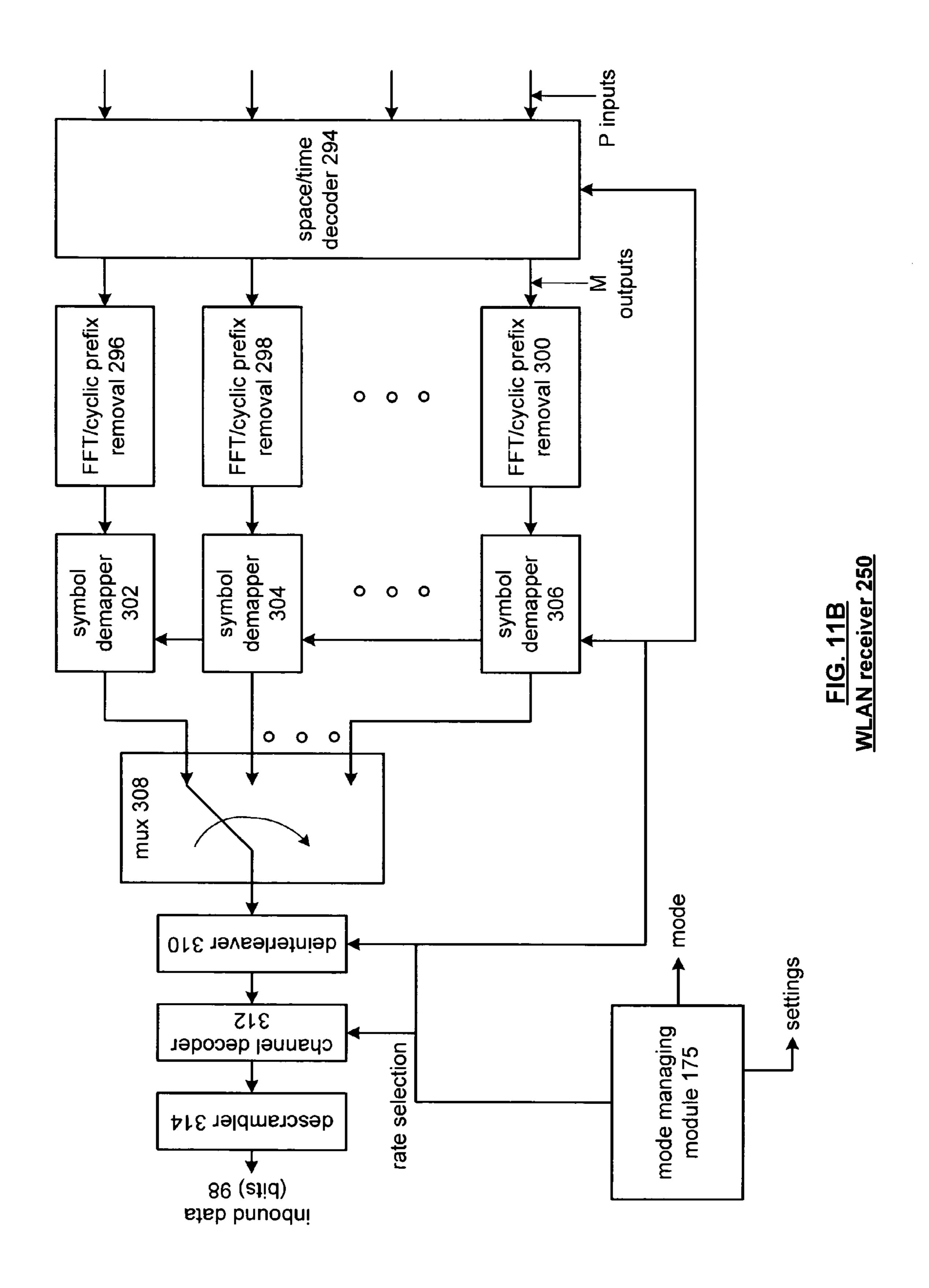


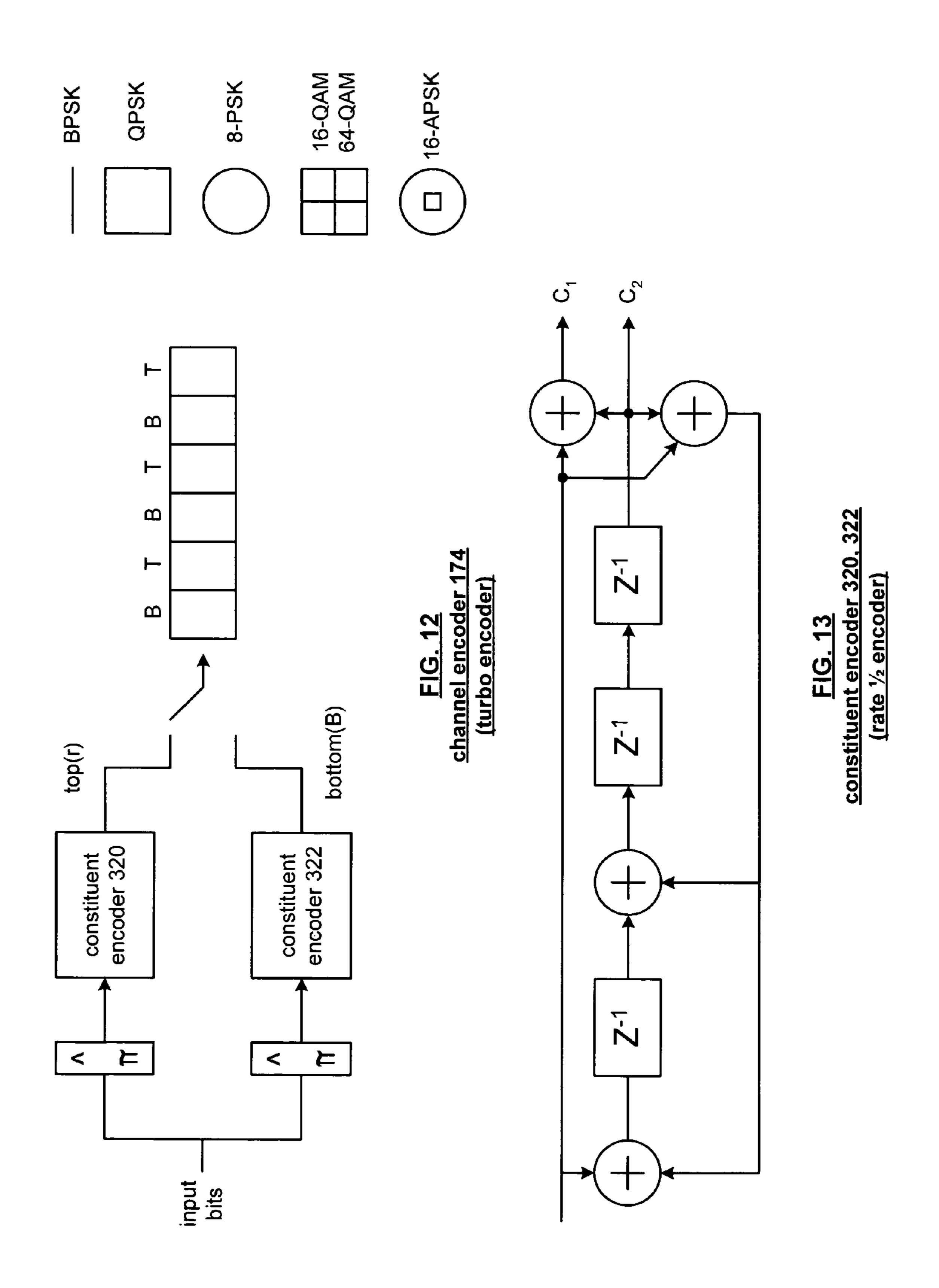
FIG. 9

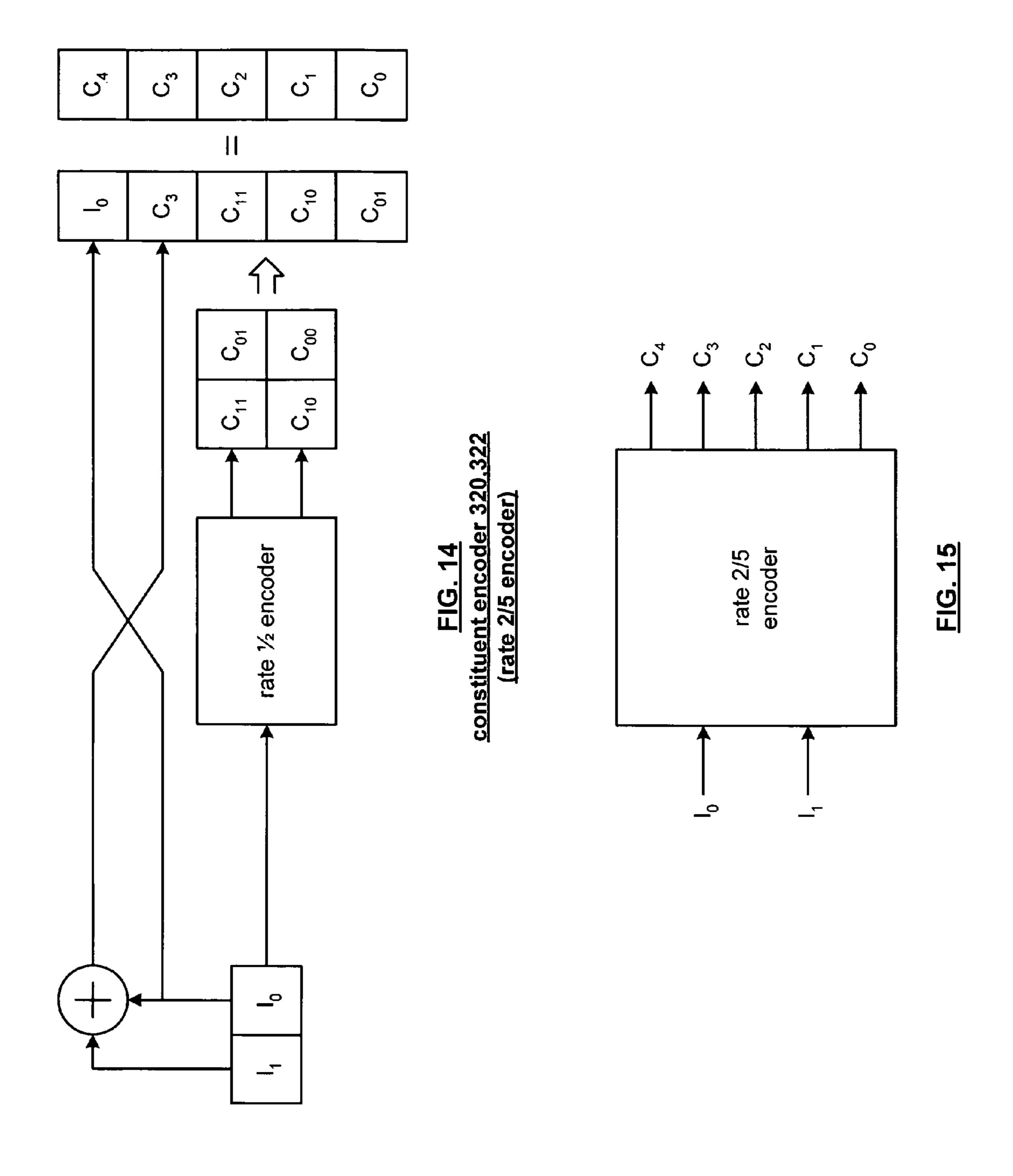


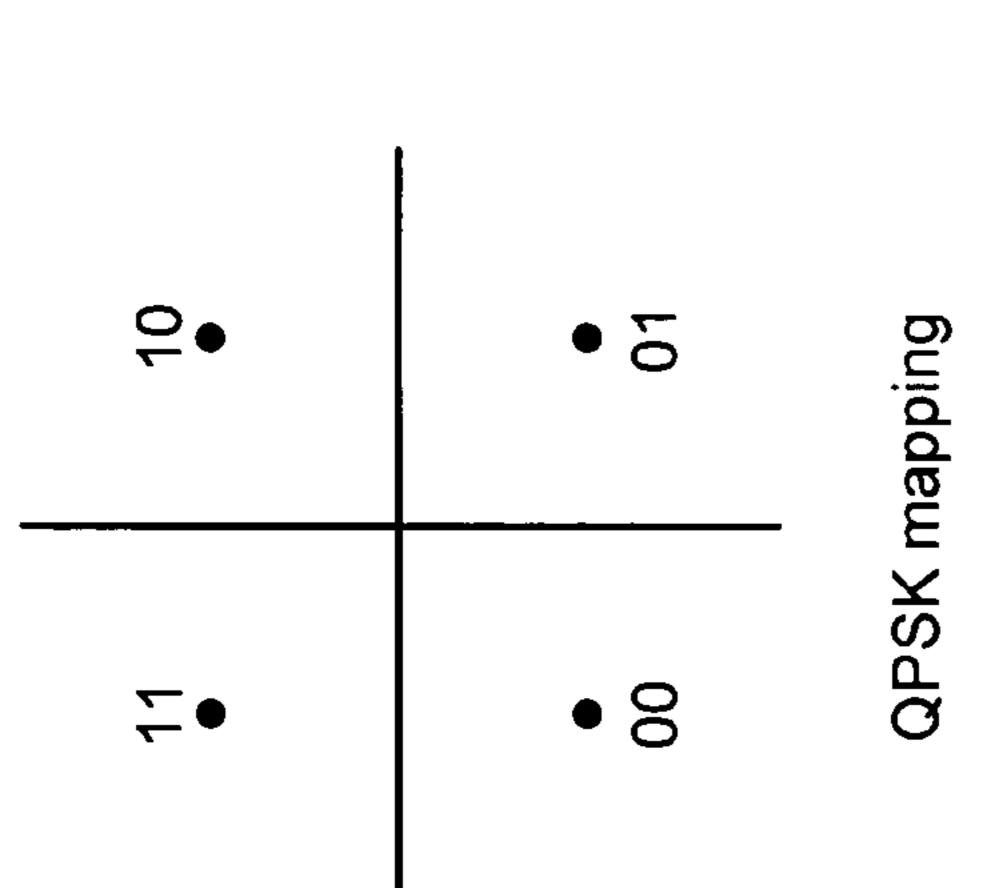




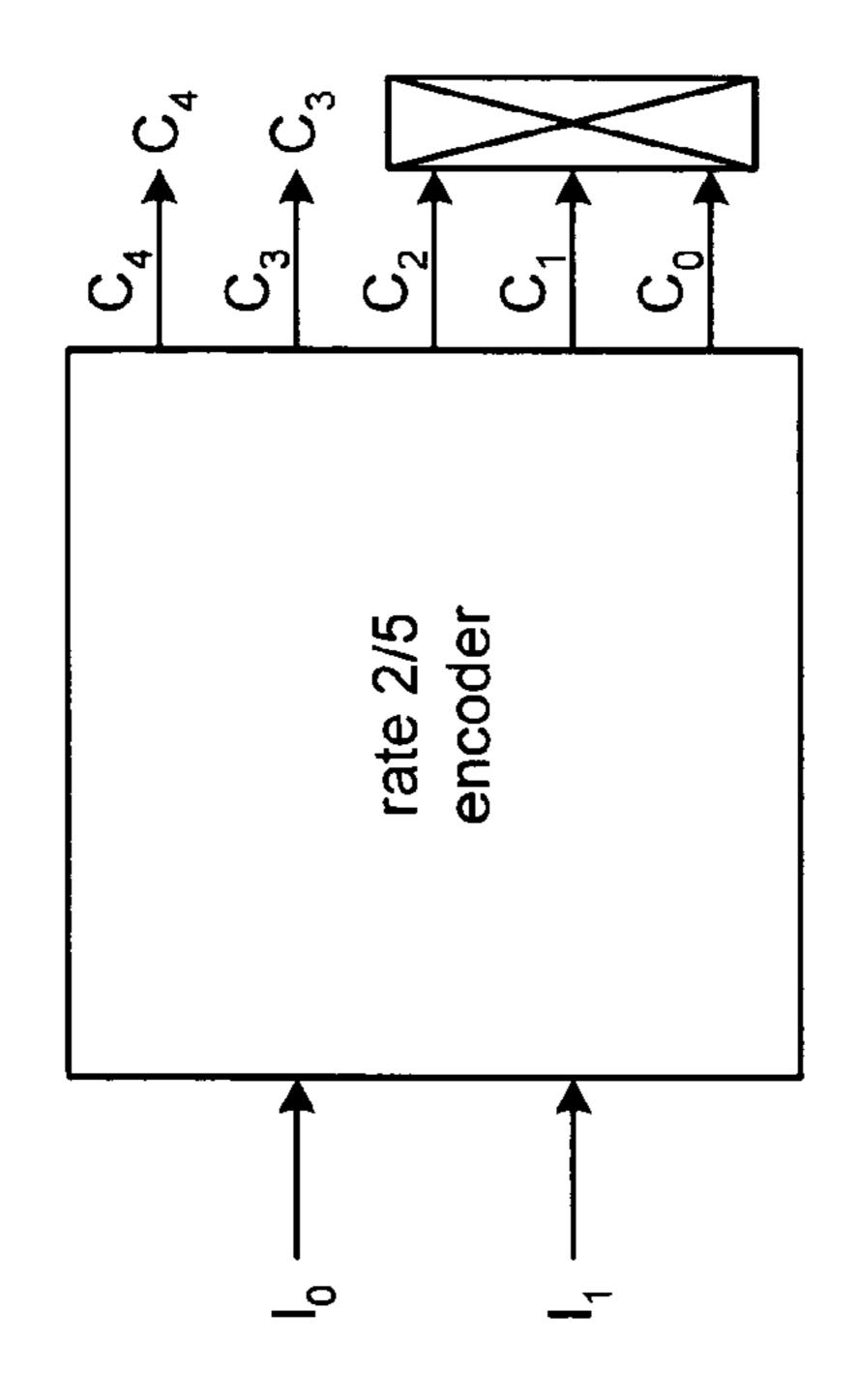




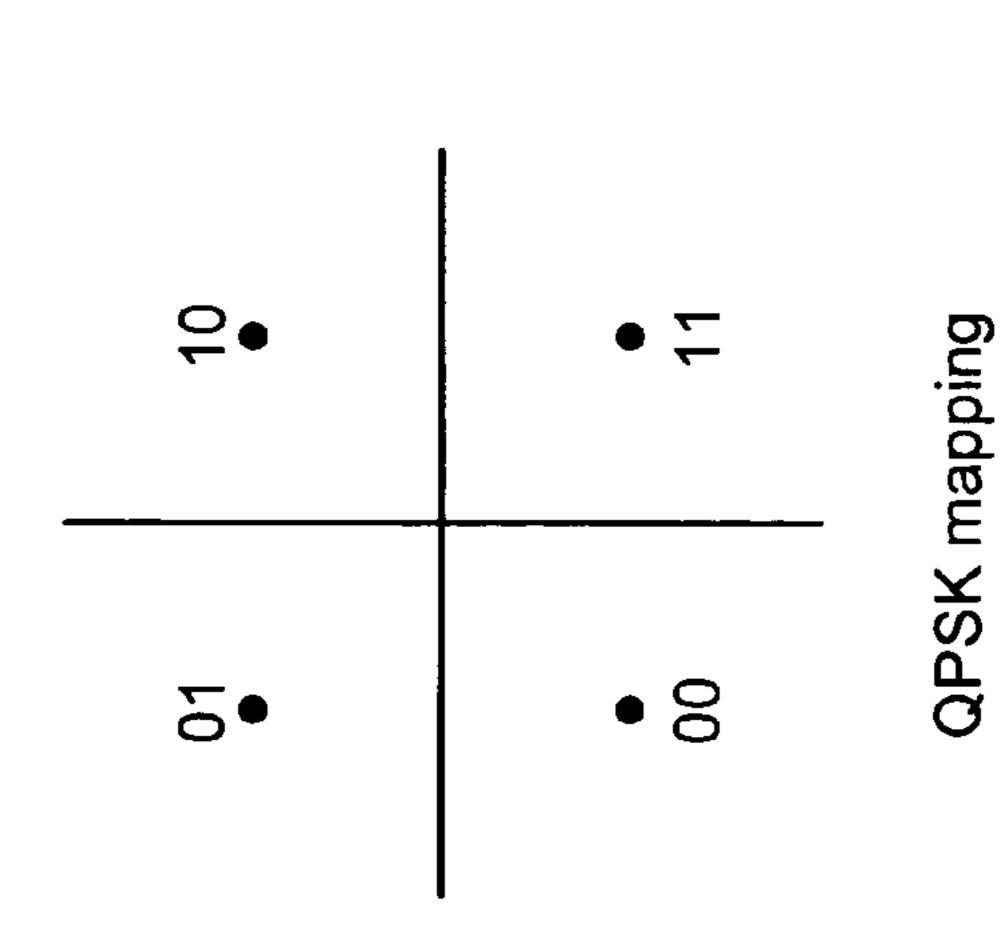


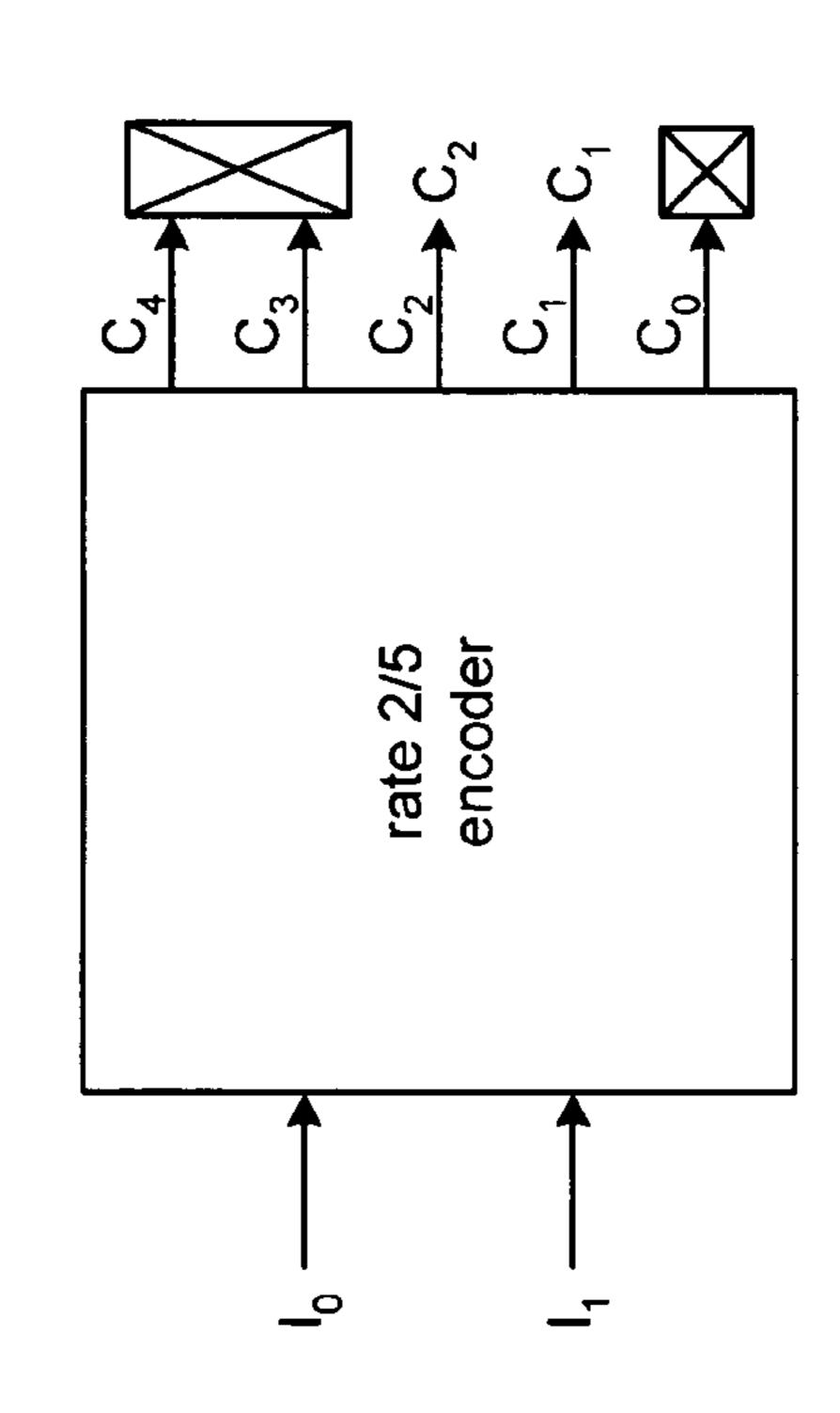


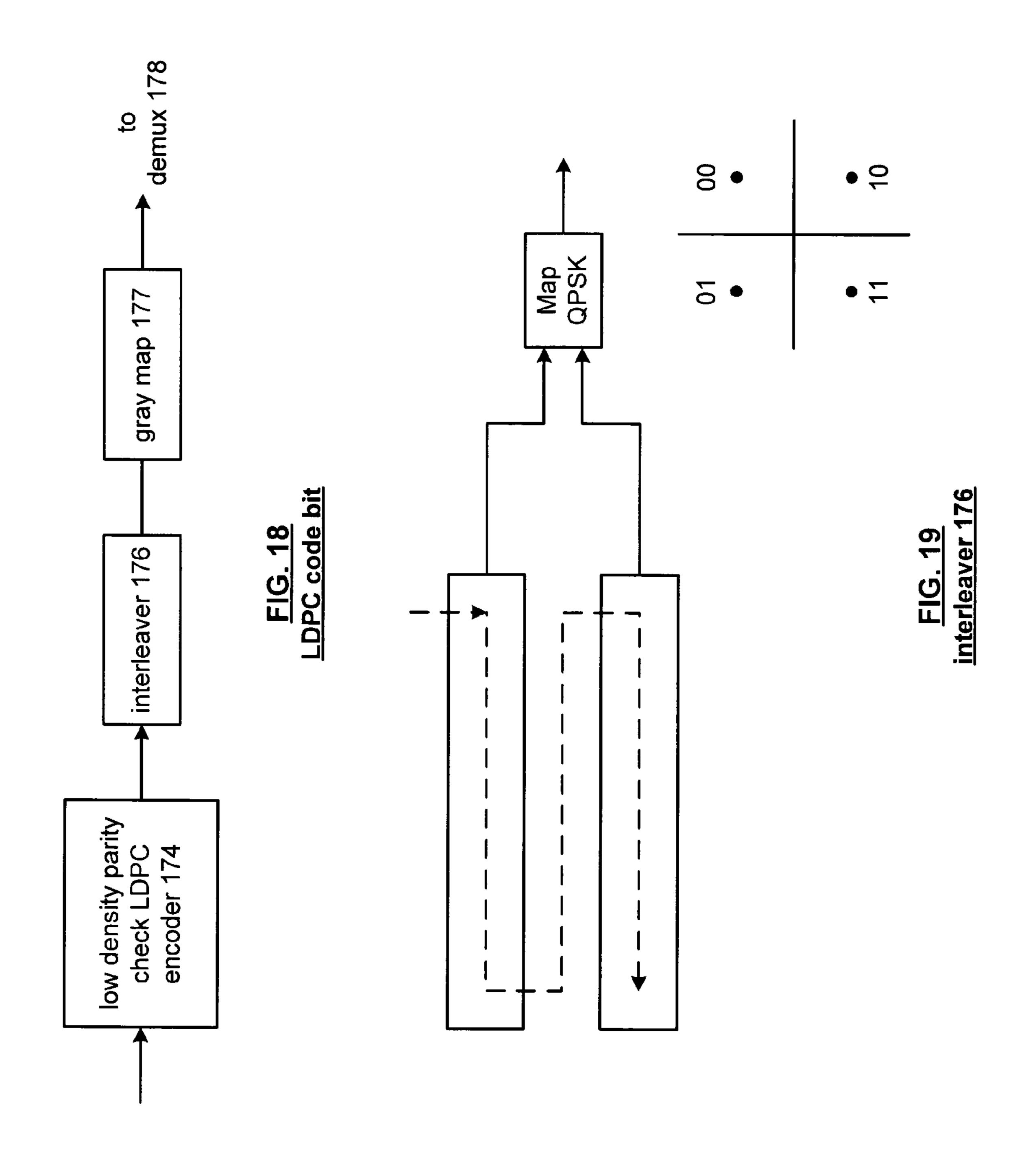
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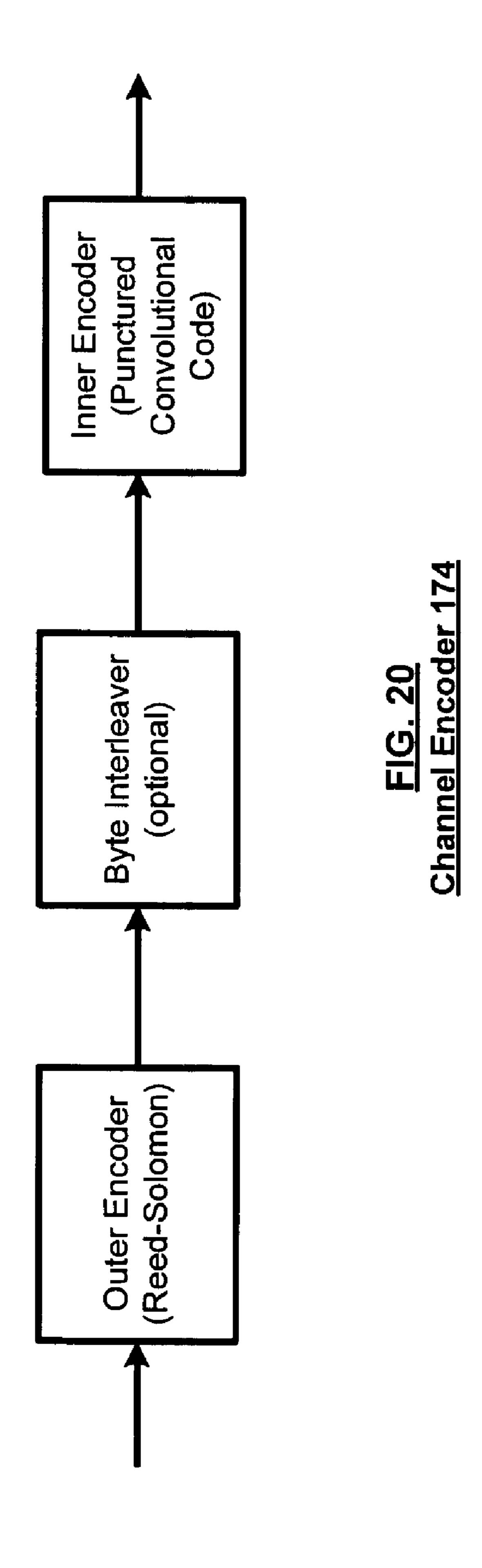


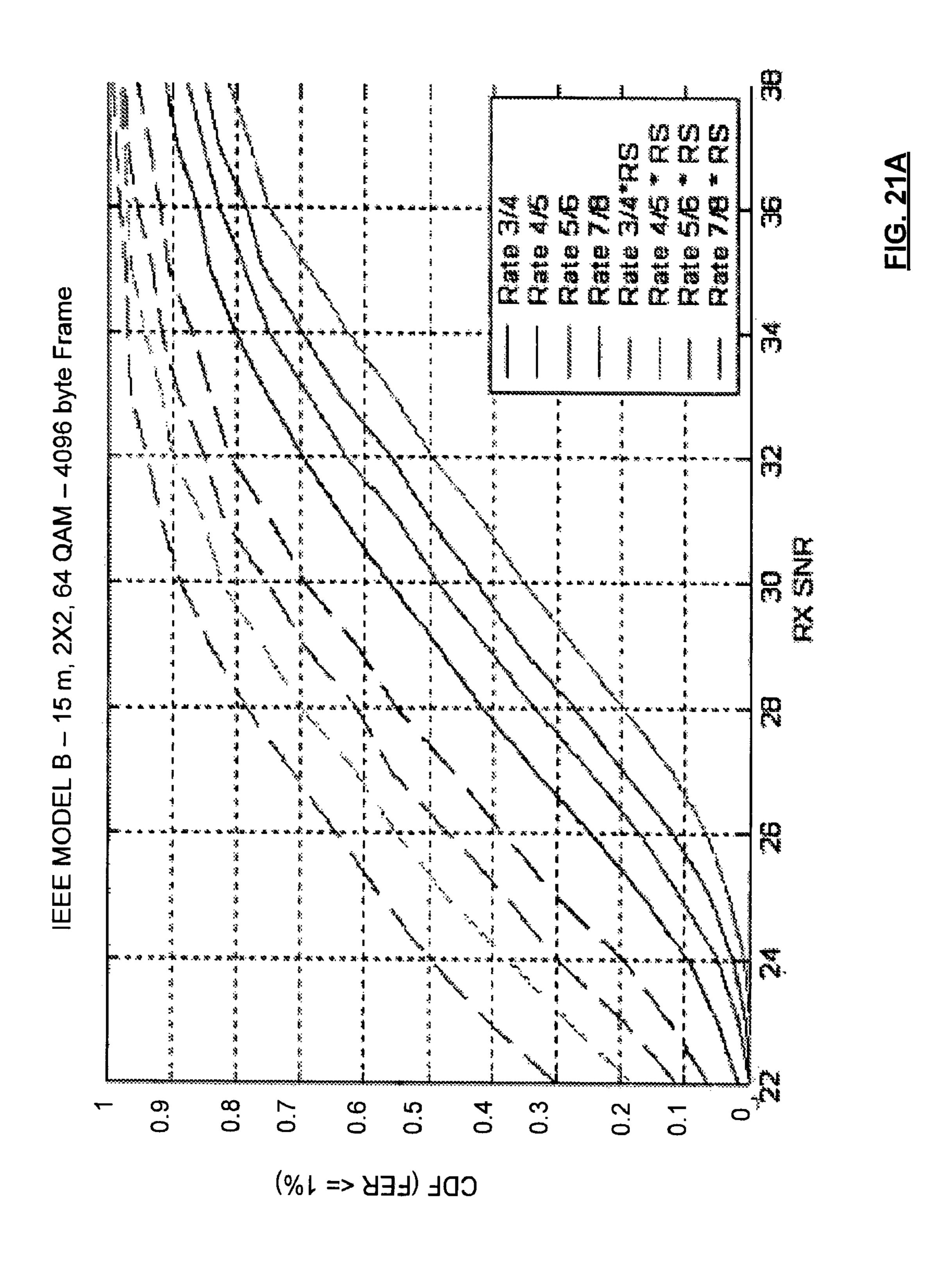
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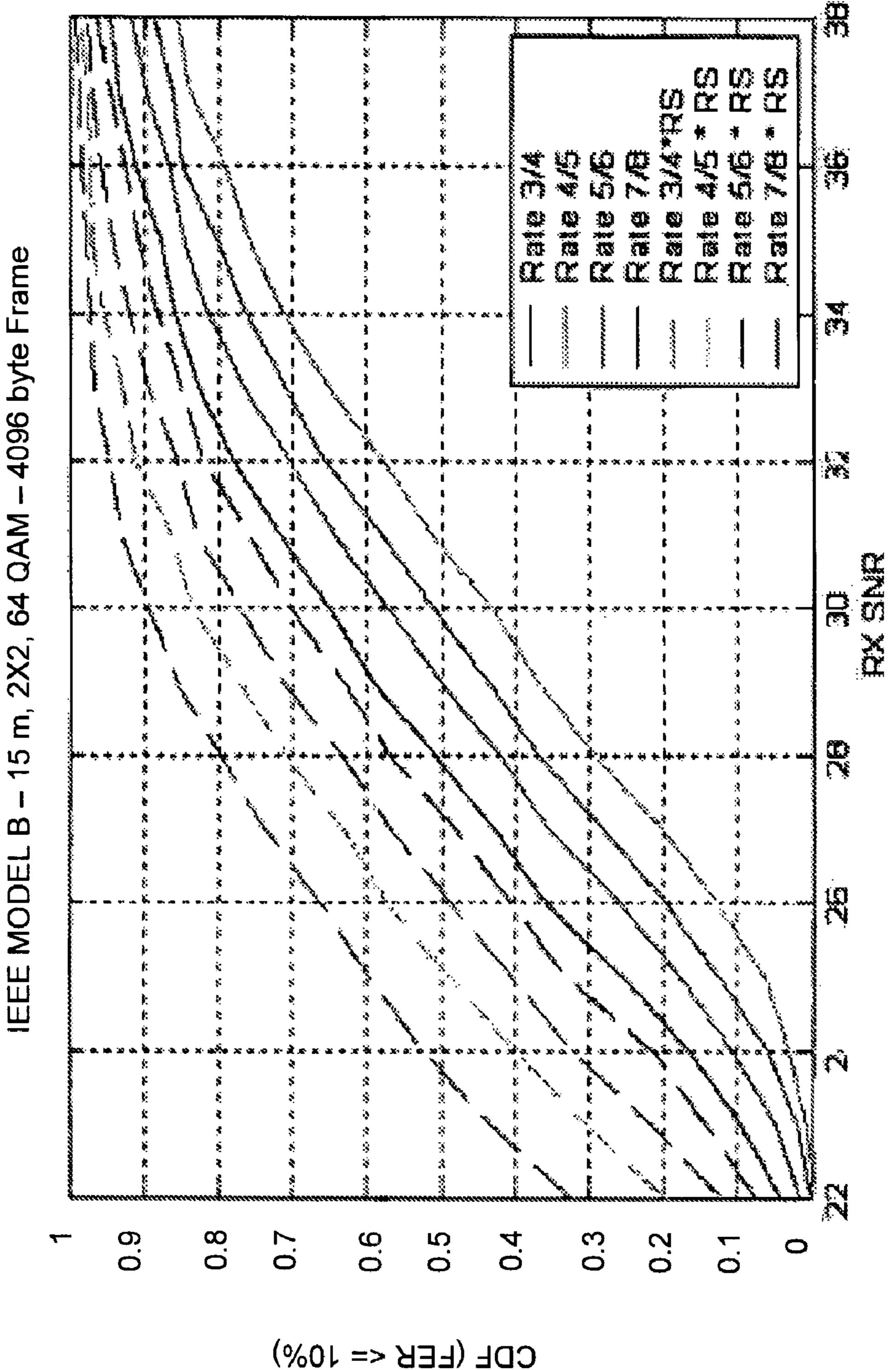


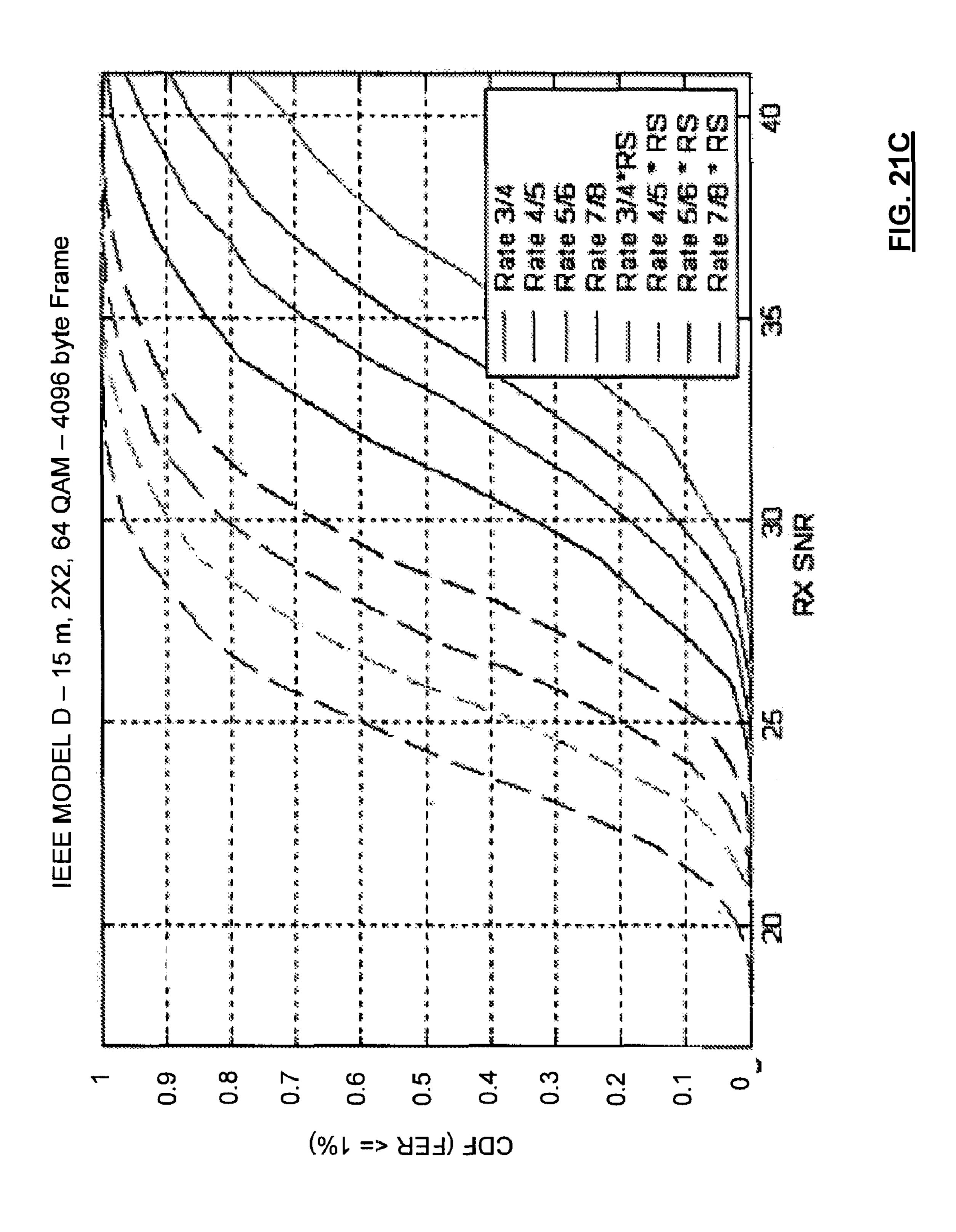


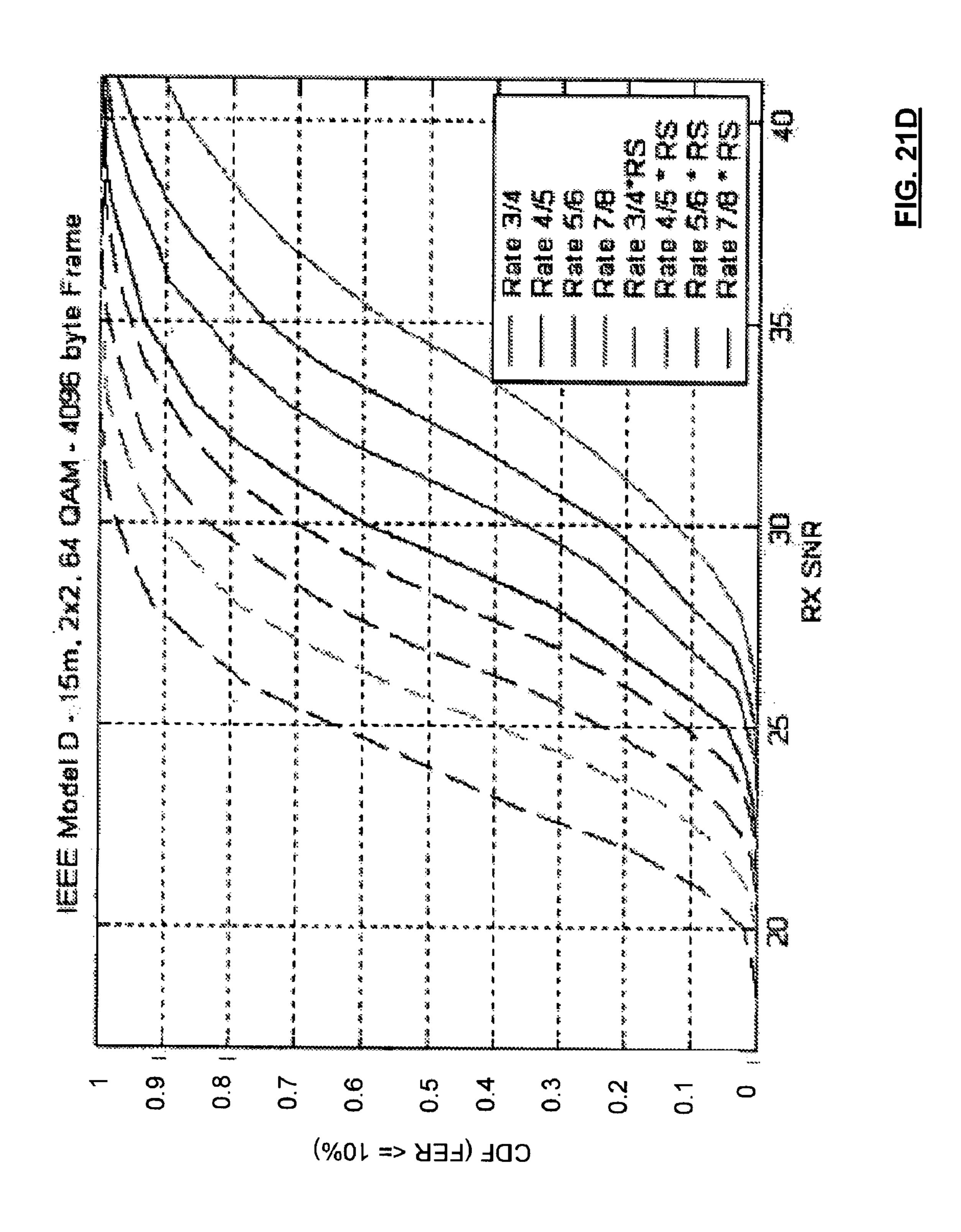


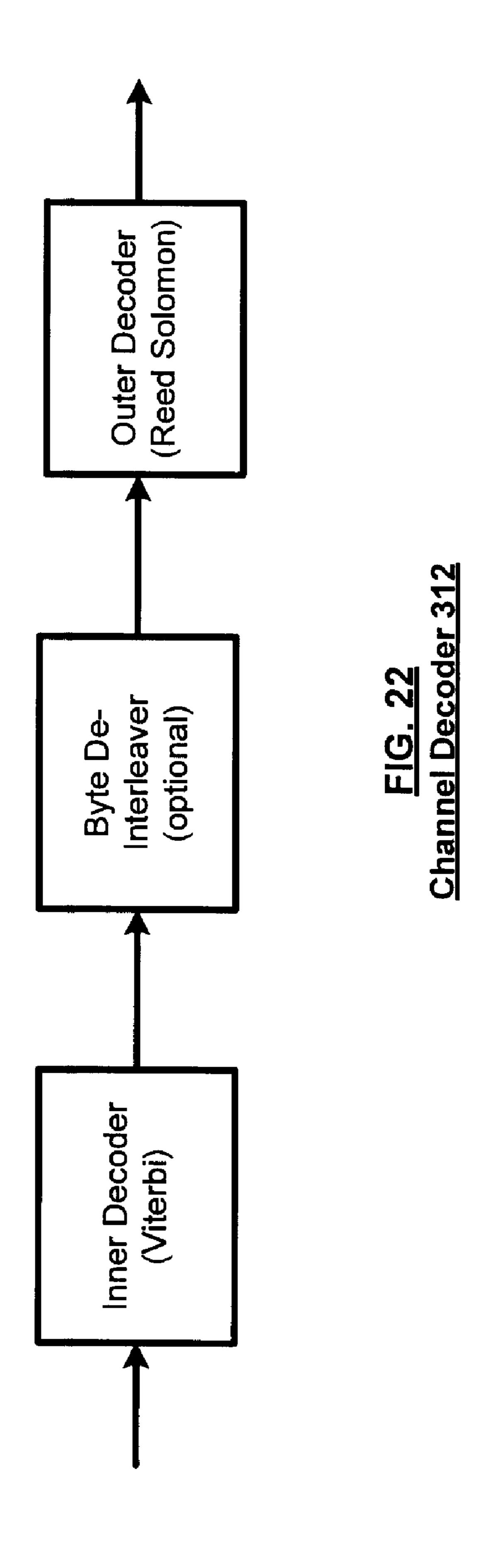


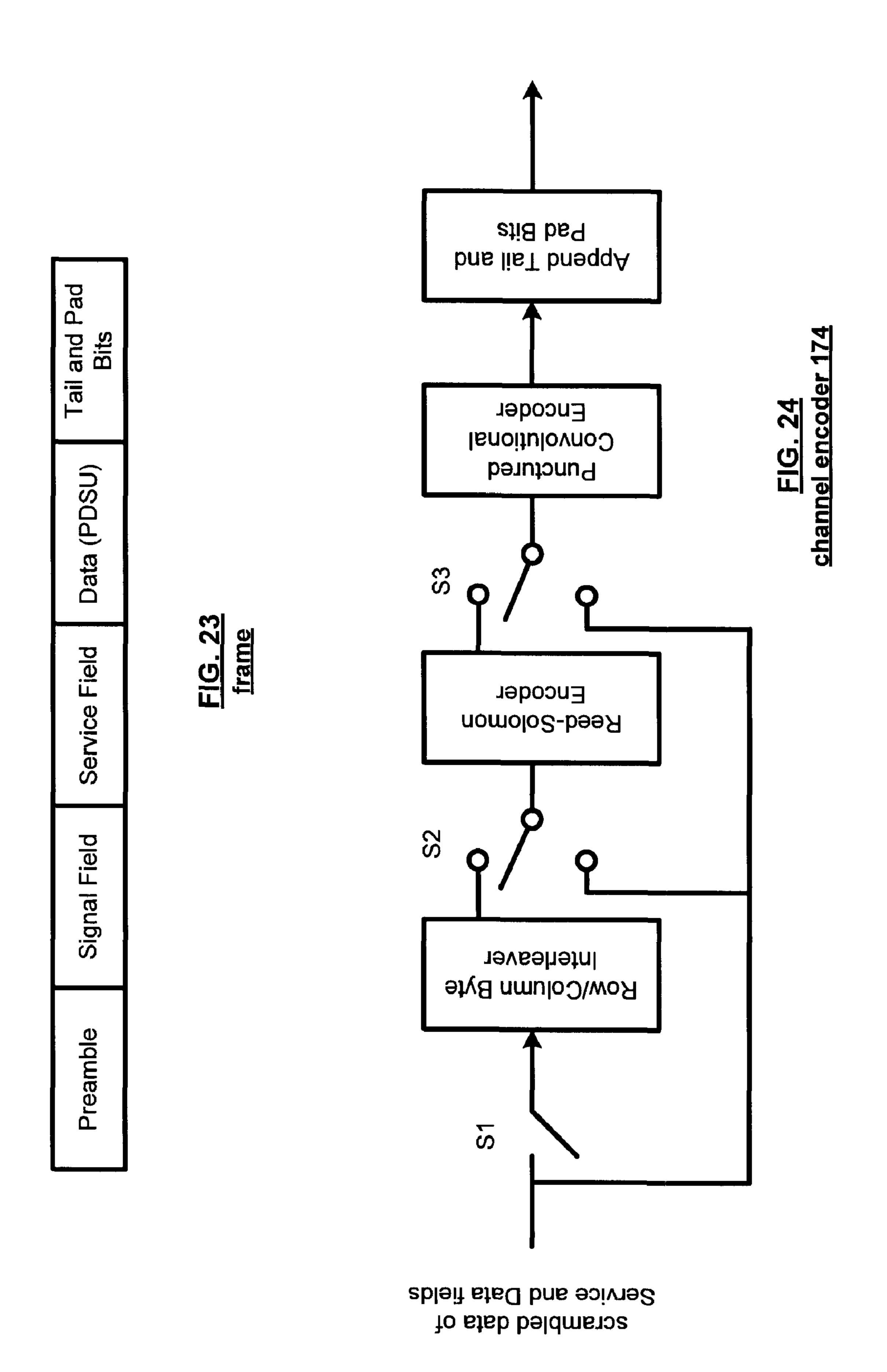


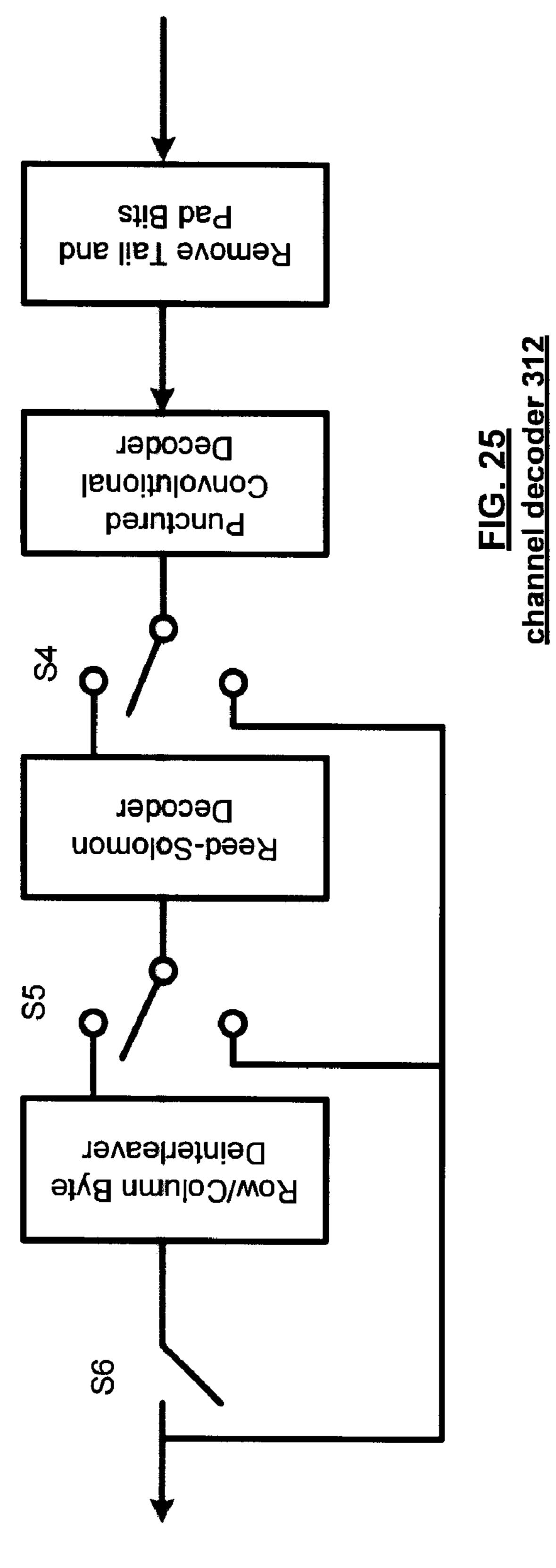








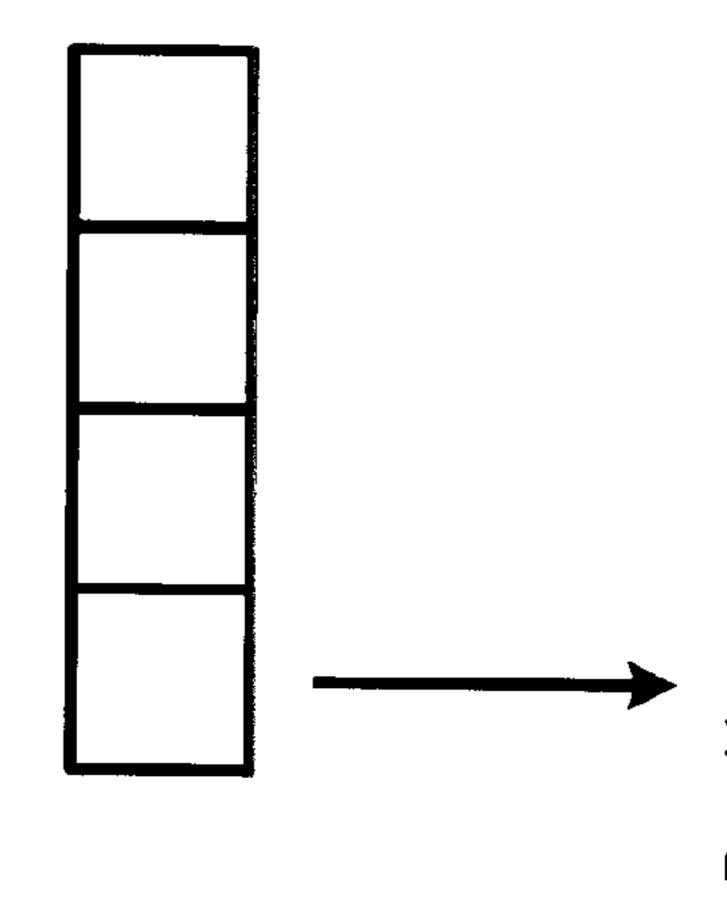




scrambled data of Service and Data fields

Number of columns, N, in inte is equal to the depth of the interleaver. The number of ro interleaver.

(uncoded bytes) in each code in the interleaver is equal ot length k of the data blaocks



Write bytes in row by row

Read bytes out

#### REDUCED LATENCY CONCATENATED REED SOLOMON-CONVOLUTIONAL CODING FOR MIMO WIRELESS LAN

The present U.S. Utility Patent Application claims priority 5 pursuant to 35 U.S.C. §119(e) to the following U.S. Provisional Patent Applications which are hereby incorporated herein by reference in their entirety and made part of the present U.S. Utility Patent Application for all purposes:

- 1. U.S. Provisional Application Ser. No. 60/544,605, 10 entitled "Multiple Protocol Wireless Communications in a WLAN,", filed Feb. 13, 2004, expired.
- 2. U.S. Provisional Application Ser. No. 60/545,854, entitled "WLAN Transmitter Having High Data Throughput,", filed Feb. 19, 2004, expired.
- 3. U.S. Provisional Application Ser. No. 60/568,914, entitled "MIMO Protocol for Wireless Communications,", filed May 7, 2004, expired.
- 4. U.S. Provisional Application Ser. No. 60/573,781, entitled "Encoder and Decoder of a WLAN Transmitter Hav- 20 ing High Data Throughput,", filed May 24, 2004, expired.
- 5. U.S. Provisional Application Ser. No. 60/575,909, entitled "Concatenated Reed Solomon-Convolutional Coding for MIMO Wireless LAN,", filed Jun. 1, 2004, expired.
- 6. U.S. Provisional Application Ser. No. 60/587,068, 25 entitled "Reduced Latency Concatenated Reed Solomon-Convolutional Coding for MIMO Wireless LAN,", filed Jul. 12, 2004, expired.

#### BACKGROUND OF THE INVENTION

#### 1. Technical Field of the Invention

This invention relates generally to wireless communication systems and more particularly to a transmitter transmitsystems.

#### 2. Description of Related Art

Communication systems are known to support wireless and wire lined communications between wireless and/or wire lined communication devices. Such communication systems 40 range from national and/or international cellular telephone systems to the Internet to point-to-point in-home wireless networks. Each type of communication system is constructed, and hence operates, in accordance with one or more communication standards. For instance, wireless communi- 45 cation systems may operate in accordance with one or more standards including, but not limited to, IEEE 802.11, Bluetooth, advanced mobile phone services (AMPS), digital AMPS, global system for mobile communications (GSM), code division multiple access (CDMA), local multi-point 50 distribution systems (LMDS), multi-channel-multi-point distribution systems (MMDS), and/or variations thereof.

Depending on the type of wireless communication system, a wireless communication device, such as a cellular telephone, two-way radio, personal digital assistant (PDA), per- 55 sonal computer (PC), laptop computer, home entertainment equipment, et cetera communicates directly or indirectly with other wireless communication devices. For direct communications (also known as point-to-point communications), the participating wireless communication devices tune their 60 receivers and transmitters to the same channel or channels (e.g., one of the plurality of radio frequency (RF) carriers of the wireless communication system) and communicate over that channel(s). For indirect wireless communications, each wireless communication device communicates directly with 65 an associated base station (e.g., for cellular services) and/or an associated access point (e.g., for an in-home or in-building

wireless network) via an assigned channel. To complete a communication connection between the wireless communication devices, the associated base stations and/or associated access points communicate with each other directly, via a system controller, via the public switch telephone network, via the Internet, and/or via some other wide area network.

For each wireless communication device to participate in wireless communications, it includes a built-in radio transceiver (i.e., receiver and transmitter) or is coupled to an associated radio transceiver (e.g., a station for in-home and/or in-building wireless communication networks, RF modem, etc.). As is known, the receiver is coupled to the antenna and includes a low noise amplifier, one or more intermediate frequency stages, a filtering stage, and a data recovery stage. The low noise amplifier receives inbound RF signals via the antenna and amplifies then. The one or more intermediate frequency stages mix the amplified RF signals with one or more local oscillations to convert the amplified RF signal into baseband signals or intermediate frequency (IF) signals. The filtering stage filters the baseband signals or the IF signals to attenuate unwanted out of band signals to produce filtered signals. The data recovery stage recovers raw data from the filtered signals in accordance with the particular wireless communication standard.

As is also known, the transmitter includes a data modulation stage, one or more intermediate frequency stages, and a power amplifier. The data modulation stage converts raw data into baseband signals in accordance with a particular wireless 30 communication standard. The one or more intermediate frequency stages mix the baseband signals with one or more local oscillations to produce RF signals. The power amplifier amplifies the RF signals prior to transmission via an antenna.

Typically, the transmitter will include one antenna for ting at high data rates with such wireless communication 35 transmitting the RF signals, which are received by a single antenna, or multiple antennas, of a receiver. When the receiver includes two or more antennas, the receiver will select one of them to receive the incoming RF signals. In this instance, the wireless communication between the transmitter and receiver is a single-output-single-input (SOSI) communication, even if the receiver includes multiple antennas that are used as diversity antennas (i.e., selecting one of them to receive the incoming RF signals). For SISO wireless communications, a transceiver includes one transmitter and one receiver. Currently, most wireless local area networks (WLAN) that are IEEE 802.11, 802.11a, 802.11b, or 802.11g employ SISO wireless communications.

> Other types of wireless communications include singleinput-multiple-output (SIMO), multiple-input-single-output (MISO), and multiple-input-multiple-output (MIMO). In a SIMO wireless communication, a single transmitter processes data into radio frequency signals that are transmitted to a receiver. The receiver includes two or more antennas and two or more receiver paths. Each of the antennas receives the RF signals and provides them to a corresponding receiver path (e.g., LNA, down conversion module, filters, and ADCs). Each of the receiver paths processes the received RF signals to produce digital signals, which are combined and then processed to recapture the transmitted data.

> For a multiple-input-single-output (MISO) wireless communication, the transmitter includes two or more transmission paths (e.g., digital to analog converter, filters, up-conversion module, and a power amplifier) that each converts a corresponding portion of baseband signals into RF signals, which are transmitted via corresponding antennas to a receiver. The receiver includes a single receiver path that receives the multiple RF signals from the transmitter. In this

instance, the receiver uses beam forming to combine the multiple RF signals into one signal for processing.

For a multiple-input-multiple-output (MIMO) wireless communication, the transmitter and receiver each include multiple paths. In such a communication, the transmitter parallel processes data using a spatial and time encoding function to produce two or more streams of data. The transmitter includes multiple transmission paths to convert each stream of data into multiple RF signals. The receiver receives the multiple RF signals via multiple receiver paths that recapture the streams of data utilizing a spatial and time decoding function. The recaptured streams of data are combined and subsequently processed to recover the original data.

With the various types of wireless communications (e.g., SISO, MISO, SIMO, and MIMO), it would be desirable to use one or more types of wireless communications to enhance data throughput within a WLAN. For example, high data rates can be achieved with MIMO communications in comparison to SISO communications. However, most WLAN include legacy wireless communication devices (i.e., devices that are compliant with an older version of a wireless communication standard. As such, a transmitter capable of MIMO wireless communications should also be backward compatible with legacy devices to function in a majority of existing WLANs.

In particular, it would be desirable to have a transmitter that used code rates greater than <sup>3</sup>/<sub>4</sub> since smaller constellations are theoretically possible to achieve 100 Mbps data rates using a 20 MHz channel. However, as is known for IEEE 802.11a applications, performance of punctured convolutional codes diminishes at higher rates. For example, R=<sup>1</sup>/<sub>2</sub> dmin=10; R=<sup>3</sup>/<sub>4</sub> dmin=5; R=<sup>4</sup>/<sub>5</sub> or <sup>5</sup>/<sub>6</sub> dmin=4; and R=<sup>7</sup>/<sub>8</sub> dmin=3. Alternate high rate codes such as LDPC and turbo codes may have latency, implementation, and/or other issues.

Therefore, a need exists for a WLAN transmitter and receiver that are capable of high data throughput.

#### BRIEF SUMMARY OF THE INVENTION

# Reduced Latency Concatenated Reed Solomon-Convolutional Coding For MIMO Wireless LAN

The reduced latency concatenated Reed-Solomon convo- 45 lutional coding for MIMO WLAN of the present invention substantially meets these needs and others. In one embodiment, a wireless local area network (WLAN) transmitter includes a baseband processing module and a plurality of radio frequency (RF) transmitters. The baseband processing 50 module operably coupled to scramble data in accordance with a pseudo random sequence to produce scrambled data. The baseband processing module is further operably coupled to interleave, at a word level, the scrambled data to produce interleaved data when the interleaving is enabled. The base- 55 band processing module is further operably coupled to outer Reed-Solomon encode the scrambled data or the interleaved data to produce outer encoded data when the outer Reed-Solomon encoding is enabled. The baseband processing module is further operably coupled to inner puncture convolution 60 encode the outer encoded data or the scrambled data to produce the encoded data. The baseband processing module is further operably coupled to determine a number of transmit streams based on a mode selection signal. The baseband processing module is further operably coupled to convert the 65 tion; encoded data into streams of symbols in accordance with the number of transmit streams and the mode selection signal.

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The plurality of radio frequency (RF) transmitters, when enabled, converts the streams of symbols into a corresponding number of RF signals.

In another embodiment, a wireless local area network (WLAN) receiver includes a plurality of RF receivers and a baseband processing module. The plurality of radio frequency (RF) receivers, based on the mode selection signal, converts a plurality of received RF signals into a number of streams of symbols. The baseband processing module is operably coupled to combine the steams of symbols into a single stream of symbols. The baseband processing module is further operably coupled to inner puncture convolution decode the single stream of symbols to produce inner punctured decoded data. The baseband processing module is further operably coupled to outer Reed-Solomon decode, when enabled, the inner punctured decoded data to produce outer decoded data. The baseband processing module is further operably coupled to deinterleave, at a word level, the inner punctured decoded data or the outer decoded data to produce deinterleaved data when the deinterleaving is enabled. The baseband processing module is further operably coupled to descramble the inner punctured decoded data, the outer decoded data, or the deinterleaved data to produce inbound

## BRIEF DESCRIPTION OF THE SEVERAL VIEWS OF THE DRAWINGS

FIG. 1 is a schematic block diagram of a wireless communication system in accordance with the present invention;

FIG. 2 is a schematic block diagram of a wireless communication device in accordance with the present invention;

FIG. 3 is a schematic block diagram of an RF transmitter in accordance with the present invention;

FIG. 4 is a schematic block diagram of an RF receiver in accordance with the present invention;

FIG. 5 is a logic diagram of a method for baseband processing of data in accordance with the present invention;

FIG. 6 is a logic diagram that provides an embodiment of the encoding step of FIG. 5;

FIGS. 7-9 illustrate logic diagrams of various embodiments for encoding the scrambled data in accordance with the present invention;

FIGS. 10A and 10B are a schematic block diagram of a radio transmitter in accordance with the present invention;

FIGS. 11A and 11B are a schematic block diagram of a radio receiver in accordance with the present invention;

FIG. 12 is a schematic block diagram of a channel encoder in accordance with the present invention;

FIG. 13 is a schematic block diagram of a constituent encoder in accordance with the present invention;

FIG. 14 is a schematic block diagram of an alternate embodiment of a constituent encoder in accordance with the present invention;

FIG. 15 is a schematic block diagram of a rate ½ encoder in accordance with the present invention;

FIG. 16 is a schematic block diagram of a puncture encoder in accordance with the present invention;

FIG. 17 is a schematic block diagram of another embodiment of a puncture encoder in accordance with the present invention;

FIG. 18 is a schematic block diagram of a low density parity check encoder in accordance with the present invention;

FIG. **19** is an illustration of an interleaver in accordance with the present invention;

FIG. 20 is a schematic block diagram of a channel encoder in accordance with the present invention;

FIGS. 21A-D are tables illustrating a comparison of frame error rate (FER) performance of the systems over an ensemble of channels generated with the IEEE channel 5 model;

FIG. 22 is a schematic block diagram of a channel decoder in accordance with the present invention;

FIG. 23 is a diagram of a frame in accordance with the present invention;

FIG. 24 is a schematic block diagram of an embodiment of a channel encoder in accordance with the present invention;

FIG. 25 is a schematic block diagram of an embodiment of a channel decoder in accordance with the present invention; and

FIG. **26** is a diagram of an interleaving function in accordance with the present invention.

#### DETAILED DESCRIPTION OF THE INVENTION

FIG. 1 is a schematic block diagram illustrating a communication system 10 that includes a plurality of base stations and/or access points 12-16, a plurality of wireless communication devices 18-32 and a network hardware component 34. The wireless communication devices 18-32 may be laptop 25 host computers 18 and 26, personal digital assistant hosts 20 and 30, personal computer hosts 24 and 32 and/or cellular telephone hosts 22 and 28. The details of the wireless communication devices will be described in greater detail with reference to FIG. 2.

The base stations or access points 12-16 are operably coupled to the network hardware 34 via local area network connections 36, 38 and 40. The network hardware 34, which may be a router, switch, bridge, modem, system controller, et cetera provides a wide area network connection 42 for the 35 communication system 10. Each of the base stations or access points 12-16 has an associated antenna or antenna array to communicate with the wireless communication devices in its area. Typically, the wireless communication devices register with a particular base station or access point 12-14 to receive 40 services from the communication system 10. For direct connections (i.e., point-to-point communications), wireless communication devices communicate directly via an allocated channel.

Typically, base stations are used for cellular telephone 45 systems and like-type systems, while access points are used for in-home or in-building wireless networks. Regardless of the particular type of communication system, each wireless communication device includes a built-in radio and/or is coupled to a radio. The radio includes a highly linear amplifier 50 and/or programmable multi-stage amplifier as disclosed herein to enhance performance, reduce costs, reduce size, and/or enhance broadband applications.

FIG. 2 is a schematic block diagram illustrating a wireless communication device that includes the host device 18-32 and an associated radio 60. For cellular telephone hosts, the radio 60 is a built-in component. For personal digital assistants hosts, laptop hosts, and/or personal computer hosts, the radio 60 may be built-in or an externally coupled component.

As illustrated, the host device 18-32 includes a processing module 50, memory 52, radio interface 54, input interface 58 and output interface 56. The processing module 50 and memory 52 execute the corresponding instructions that are typically done by the host device. For example, for a cellular telephone host device, the processing module 50 performs the corresponding communication functions in accordance with a particular cellular telephone standard.

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The radio interface **54** allows data to be received from and sent to the radio **60**. For data received from the radio **60** (e.g., inbound data), the radio interface **54** provides the data to the processing module **50** for further processing and/or routing to the output interface **56**. The output interface **56** provides connectivity to an output display device such as a display, monitor, speakers, et cetera such that the received data may be displayed. The radio interface **54** also provides data from the processing module **50** to the radio **60**. The processing module **50** may receive the outbound data from an input device such as a keyboard, keypad, microphone, et cetera via the input interface **58** or generate the data itself. For data received via the input interface **58**, the processing-module **50** may perform a corresponding host function on the data and/or route it to the radio **60** via the radio interface **54**.

Radio 60 includes a host interface 62, a baseband processing module 64, memory 66, a plurality of radio frequency (RF) transmitters **68-72**, a transmit/receive (T/R) module **74**, a plurality of antennas 82-86, a plurality of RF receivers 20 **76-80**, and a local oscillation module **100**. The baseband processing module 64, in combination with operational instructions stored in memory 66, execute digital receiver functions and digital transmitter functions, respectively. The digital receiver functions, as will be described in greater detail with reference to FIG. 11B, include, but are not limited to, digital intermediate frequency to baseband conversion, demodulation, constellation demapping, decoding, de-interleaving, fast Fourier transform, cyclic prefix removal, space and time decoding, and/or descrambling. The digital transmitter functions, as will be described in greater detail with reference to FIGS. 5-19, include, but are not limited to, scrambling, encoding, interleaving, constellation mapping, modulation, inverse fast Fourier transform, cyclic prefix addition, space and time encoding, and/or digital baseband to IF conversion. The baseband processing modules **64** may be implemented using one or more processing devices. Such a processing device may be a microprocessor, micro-controller, digital signal processor, microcomputer, central processing unit, field programmable gate array, programmable logic device, state machine, logic circuitry, analog circuitry, digital circuitry, and/or any device that manipulates signals (analog and/or digital) based on operational instructions. The memory 66 may be a single memory device or a plurality of memory devices. Such a memory device may be a read-only memory, random access memory, volatile memory, non-volatile memory, static memory, dynamic memory, flash memory, and/or any device that stores digital information. Note that when the processing module 64 implements one or more of its functions via a state machine, analog circuitry, digital circuitry, and/or logic circuitry, the memory storing the corresponding operational instructions is embedded with the circuitry comprising the state machine, analog circuitry, digital circuitry, and/or logic circuitry.

In operation, the radio 60 receives outbound data 88 from the host device via the host interface 62. The baseband processing module 64 receives the outbound data 88 and, based on a mode selection signal 102, produces one or more outbound symbol streams 90. The mode selection signal 102 will indicate a particular mode as are illustrated in the mode selection tables, which appear at the end of the detailed discussion. For example, the mode selection signal 102, with reference to table 1 may indicate a frequency band of 2.4 GHz, a channel bandwidth of 20 or 22 MHz and a maximum bit rate of 54 megabits-per-second. In this general category, the mode selection signal will further indicate a particular rate ranging from 1 megabit-per-second to 54 megabits-per-second. In addition, the mode selection signal will indicate a particular

type of modulation, which includes, but is not limited to, Barker Code Modulation, BPSK, QPSK, CCK, 16 QAM and/or 64 QAM. As is further illustrated in table 1, a code rate is supplied as well as number of coded bits per subcarrier (NBPSC), coded bits per OFDM symbol (NCBPS), data bits per OFDM symbol (NDBPS), error vector magnitude in decibels (EVM), sensitivity which indicates the maximum receive power required to obtain a target packet error rate (e.g., 10% for IEEE 802.11a), adjacent channel rejection (ACR), and an alternate adjacent channel rejection (AACR).

The mode selection signal may also indicate a particular channelization for the corresponding mode which for the information in table 1 is illustrated in table 2. As shown, table 2 includes a channel number and corresponding center frequency. The mode select signal may further indicate a power 15 spectral density mask value which for table 1 is illustrated in table 3. The mode select signal may alternatively indicate rates within table 4 that has a 5 GHz frequency band, 20 MHz channel bandwidth and a maximum bit rate of 54 megabitsper-second. If this is the particular mode select, the channel- 20 ization is illustrated in table 5. As a further alternative, the mode select signal 102 may indicate a 2.4 GHz frequency band, 20 MHz channels and a maximum bit rate of 192 megabits-per-second as illustrated in table 6. In table 6, a number of antennas may be utilized to achieve the higher 25 bandwidths. In this instance, the mode select would further indicate the number of antennas to be utilized. Table 7 illustrates the channelization for the set-up of table 6. Table 8 illustrates yet another mode option where the frequency band is 2.4 GHz, the channel bandwidth is 20 MHz and the maximum bit rate is 192 megabits-per-second. [Table 8 is 45 GHz] frequency band.] The corresponding table 8 includes various bit rates ranging from 12 megabits-per-second to 216 megabits-per-second utilizing 2-4 antennas and a spatial time encoding rate as indicated. Table 9 illustrates the channeliza- 35 tion for table 8. The mode select signal **102** may further indicate a particular operating mode as illustrated in table 10, which corresponds to a 5 GHz frequency band having 40 MHz frequency band having 40 MHz channels and a maximum bit rate of 486 megabits-per-second. As shown in table 40 10, the bit rate may range from 13.5 megabits-per-second to 486 megabits-per-second utilizing 1-4 antennas and a corresponding spatial time code rate. Table 10 further illustrates a particular modulation scheme code rate and NBPSC values. Table 11 provides the power spectral density mask for table 45 10 and table 12 provides the channelization for table 10.

The baseband processing module **64**, based on the mode selection signal **102** produces the one or more outbound symbol streams **90**, as will be further described with reference to FIGS. **5-9** from the output data **88**. For example, if the mode selection signal **102** indicates that a single transmit antenna is being utilized for the particular mode that has been selected, the baseband processing module **64** will produce a single outbound symbol stream **90**. Alternatively, if the mode select signal indicates 2, 3 or 4 antennas, the baseband processing module **64** will produce 2, 3 or 4 outbound symbol streams **90** corresponding to the number of antennas from the output data **88**.

Depending on the number of outbound streams 90 produced by the baseband module 64, a corresponding number of 60 the RF transmitters 68-72 will be enabled to convert the outbound symbol streams 90 into outbound RF signals 92. The implementation of the RF transmitters 68-72 will be further described with reference to FIG. 3. The transmit/receive module 74 receives the outbound RF signals 92 and 65 provides each outbound RF signal to a corresponding antenna 82-86.

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When the radio 60 is in the receive mode, the transmit/receive module 74 receives one or more inbound RF signals via the antennas 82-86. The T/R module 74 provides the inbound RF signals 94 to one or more RF receivers 76-80. The RF receiver 76-80, which will be described in greater detail with reference to FIG. 4, converts the inbound RF signals 94 into a corresponding number of inbound symbol streams 96. The number of inbound symbol streams 96 will correspond to the particular mode in which the data was received (recall that the mode may be any one of the modes illustrated in tables 1-12). The baseband processing module 60 receives the inbound symbol streams 90 and converts them into inbound data 98, which is provided to the host device 18-32 via the host interface 62.

As one of average skill in the art will appreciate, the wireless communication device of FIG. 2 may be implemented using one or more integrated circuits. For example, the host device may be implemented on one integrated circuit, the baseband processing module 64 and memory 66 may be implemented on a second integrated circuit, and the remaining components of the radio 60, less the antennas 82-86, may be implemented on a third integrated circuit. As an alternate example, the radio 60 may be implemented on a single integrated circuit. As yet another example, the processing module 50 of the host device and the baseband processing module 64 may be a common processing device implemented on a single integrated circuit. Further, the memory **52** and memory **66** may be implemented on a single integrated circuit and/or on the same integrated circuit as the common processing modules of processing module 50 and the baseband processing module **64**.

FIG. 3 is a schematic block diagram of an embodiment of an RF transmitter 68-72. The RF transmitter 68-72 includes a digital filter and up-sampling module 75, a digital-to-analog conversion module 77, an analog filter 79, and up-conversion module 81, a power amplifier 83 and a RF filter 85. The digital filter and up-sampling module 75 receives one of the out-bound symbol streams 90 and digitally filters it and then up-samples the rate of the symbol streams to a desired rate to produce the filtered symbol streams 87. The digital-to-analog conversion module 77 converts the filtered symbols 87 into analog signals 89. The analog signals may include an in-phase component and a quadrature component.

The analog filter 79 filters the analog signals 89 to produce filtered analog signals 91. The up-conversion module 81, which may include a pair of mixers and a filter, mixes the filtered analog signals 91 with a local oscillation 93, which is produced by local oscillation module 100, to produce high frequency signals 95. The frequency of the high frequency signals 95 corresponds to the frequency of the RF signals 92.

The power amplifier 83 amplifies the high frequency signals 95 to produce amplified high frequency signals 97. The RF filter 85, which may be a high frequency band-pass filter, filters the amplified high frequency signals 97 to produce the desired output RF signals 92.

As one of average skill in the art will appreciate, each of the radio frequency transmitters 68-72 will include a similar architecture as illustrated in FIG. 3 and further include a shut-down mechanism such that when the particular radio frequency transmitter is not required, it is disabled in such a manner that it does not produce interfering signals and/or noise.

FIG. 4 is a schematic block diagram of each of the RF receivers 76-80. In this embodiment, each of the RF receivers 76-80 includes an RF filter 101, a low noise amplifier (LNA) 103, a programmable gain amplifier (PGA) 105, a down-conversion module 107, an analog filter 109, an analog-to-

digital conversion module 111 and a digital filter and down-sampling module 113. The RF filter 101, which may be a high frequency band-pass filter, receives the inbound RF signals 94 and filters them to produce filtered inbound RF signals. The low noise amplifier 103 amplifies the filtered inbound RF signals 94 based on a gain setting and provides the amplified signals to the programmable gain amplifier 105. The programmable gain amplifier further amplifies the inbound RF signals 94 before providing them to the down-conversion module 107.

The down-conversion module 107 includes a pair of mixers, a summation module, and a filter to mix the inbound RF signals with a local oscillation (LO) that is provided by the local oscillation module to produce analog baseband signals. The analog filter 109 filters the analog baseband signals and provides them to the analog-to-digital conversion module 111 which converts them into a digital signal. The digital filter and down-sampling module 113 filters the digital signals and then adjusts the sampling rate to produce the inbound symbol stream 96.

FIG. 5 is a logic diagram of a method for converting outbound data 88 into one or more outbound symbol streams 90 by the baseband processing module 64. The process begins at Step 110 where the baseband processing module receives the outbound data 88 and a mode selection signal 102. The mode 25 selection signal may indicate any one of the various modes of operation as indicated in tables 1-12. The process then proceeds to Step 112 where the baseband processing module scrambles the data in accordance with a pseudo random sequence to produce scrambled data. Note that the pseudo 30 random sequence may be generated from a feedback shift register with the generator polynomial of  $S(x)=x^7+x^4+1$ .

The process then proceeds to Step 114 where the baseband processing module selects one of a plurality of encoding modes based on the mode selection signal. The process then 35 proceeds to Step 116 where the baseband processing module encodes the scrambled data in accordance with a selected encoding mode to produce encoded data. The encoding may be done utilizing a parallel concatenated turbo encoding scheme and/or a low density parity check block encoding 40 scheme. Such encoding schemes will be described in greater detail with reference to FIGS. 12-19. Alternatively, the encoding may be done as further described in FIGS. 7-9 which will be described below.

The process then proceeds to Step 118 where the baseband 45 processing module determines a number of transmit streams based on the mode select signal. For example, the mode select signal will select a particular mode which indicates that 1, 2, 3, 4 or more antennas may be utilized for the transmission. Accordingly, the number of transmit streams will correspond 50 to the number of antennas indicated by the mode select signal. The process then proceeds to Step 120 where the baseband processing module converts the encoded data into streams of symbols in accordance with the number of transmit streams in the mode select signal. This step will be described in greater 55 detail with reference to FIG. 6.

FIG. 6 is a logic diagram of a method performed by the baseband processing module to convert the encoded data into streams of symbols in accordance with the number of transmit streams and the mode select signal. Such processing 60 begins at Step 122 where the baseband processing module interleaves the encoded data over multiple symbols and subcarriers of a channel to produce interleaved data. In general, the interleaving process is designed to spread the encoded data over multiple symbols and transmit streams. This allows 65 improved detection and error correction capability at the receiver. In one embodiment, the interleaving process will

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follow the IEEE 802.11(a) or (g) standard for backward compatible modes. For higher performance modes (e.g., IEEE 802.11(n), the interleaving will also be done over multiple transmit paths or streams.

The process then proceeds to Step 124 where the baseband processing module demultiplexes the interleaved data into a number of parallel streams of interleaved data. The number of parallel streams corresponds to the number of transmit streams, which in turn corresponds to the number of antennas indicated by the particular mode being utilized. The process then continues to Steps 126 and 128, where for each of the parallel streams of interleaved data, the baseband processing module maps the interleaved data into a quadrature amplitude modulated (QAM) symbol to produce frequency domain symbols at Step 126. At Step 128, the baseband processing module converts the frequency domain symbols into time domain symbols, which may be done utilizing an inverse fast Fourier transform. The conversion of the frequency domain symbols into the time domain symbols may further include adding a cyclic prefix to allow removal of intersymbol interference at the receiver. Note that the length of the inverse fast Fourier transform and cyclic prefix are defined in the mode tables of tables 1-12. In general, a 64-point inverse fast Fourier transform is employed for 20 MHz channels and 128point inverse fast Fourier transform is employed for 40 MHz channels.

The process then proceeds to Step 130 where the baseband processing module space and time encodes the time domain symbols for each of the parallel streams of interleaved data to produce the streams of symbols. In one embodiment, the space and time encoding may be done by space and time encoding the time domain symbols of the parallel streams of interleaved data into a corresponding number of streams of symbols utilizing an encoding matrix. Alternatively, the space and time encoding may be done by space and time encoding the time domain symbols of M-parallel streams of interleaved data into P-streams of symbols utilizing the encoding matrix, where P=M+1. In one embodiment the encoding matrix may comprise a form of:

$$\begin{bmatrix} C_1 & C_2 & C_3 & \cdots & C_{2M-1} \\ -C_2^* & C_1^* & C_4 & \cdots & C_{2M} \end{bmatrix}$$

where the number of rows of the encoding matrix corresponds to M and the number of columns of the encoding matrix corresponds to P. The particular values of the constants within the encoding matrix may be real or imaginary numbers.

FIG. 7 is a logic diagram of one method that may be utilized by the baseband processing module to encode the scrambled data at Step 116 of FIG. 5. In this method, the process begins at Step 140 where the baseband processing module performs a convolutional encoding with 64 state codes and generator polynomials of  $G_0$ =133<sub>8</sub> and  $G_1$ =171<sub>8</sub> on the scrambled data to produce convolutional encoded data. The process then proceeds to Step 142 where the baseband processing module punctures the convolutional encoded data at one of a plurality of rates in accordance with the mode selection signal to produce the encoded data. Note that the puncture rates may include  $\frac{1}{2}$ ,  $\frac{2}{3}$ rds and/or  $\frac{3}{4}$ , or any rate as specified in tables 1-12. Note that, for a particular, mode, the rate may be selected for backward compatibility with IEEE 802.11(a) and/or IEEE 802.11(g) rate requirements.

The encoding of FIG. 7 may further include an optional Step 144 where the baseband processing module combines the convolutional encoding with an outer Reed Solomon code

to produce the convolutional encoded data. Note that Step 144 would be conducted in parallel with Step 140.

FIG. **8** is a logic diagram of another encoding method that may be utilized by the baseband processing module to encode the scrambled data at Step **116** of FIG. **5**. In this embodiment, 5 the process begins at Step **146** where the baseband processing module encodes the scrambled data in accordance with a complimentary code keying (CCK) code to produce the encoded data. This may be done in accordance with IEEE 802.11(b) specifications and/or IEEE 802.11(g) specifications. The encoding may include an optional Step **148**, which is performed in parallel with Step **146** that combines the CCK code with an outer Reed Solomon code to produce the encoded data.

FIG. 9 is a logic diagram of yet another method for encoding the scrambled data at Step 116, which may be performed by the baseband processing module. In this embodiment, the process begins at Step 150 where the baseband processing module performs a convolutional encoding with 256 state codes and generator polynomials of  $G_0$ =561 $_8$  and  $G_1$ =753 $_8$  on 20 the scrambled data to produce convolutional encoded data. The process then proceeds to Step 152 where the baseband processing module punctures the convolutional encoded data at one of the plurality of rates in accordance with a mode selection signal to produce encoded data. Note that the puncture rate is indicated in the tables 1-12 for the corresponding mode.

The encoding of FIG. 9 may further include the optional Step 154 where the baseband processing module combines the convolutional encoding with an outer Reed Solomon code 30 to produce the convolutional encoded data.

FIGS. 10A and 10B illustrate a schematic block diagram of a multiple transmitter in accordance with the present invention. In FIG. 10A, the baseband processing is shown to include a scrambler 172, channel encoder 174, a bit interleaver 176, demultiplexer 178, a plurality of symbol mappers 180-184, a plurality of inverse fast Fourier transform (IFFT)/cyclic prefix addition modules 186-190 and a space/time encoder 192. The baseband portion of the transmitter may further include a mode manager module 175 that receives the 40 mode selection signal and produces settings for the radio transmitter portion and produces the rate selection for the baseband portion.

In operations, the scrambler 172 adds (in GF2) a pseudo random sequence to the outbound data bits 88 to make the 45 data appear random. A pseudo random sequence may be generated from a feedback shift register with the generator polynomial of  $S(x)=x^7+x^4+1$  to produce scrambled data. The channel encoder 174 receives the scrambled data and generates a new sequence of bits with redundancy. This will enable 50 improved detection at the receiver. The channel encoder 174 may operate in one of a plurality of modes. For example, for backward compatibility with IEEE 802.11(a) and IEEE 802.11(g), the channel encoder has the form of a rate ½ convolutional encoder with 64 states and a generator polyno- 55 mials of  $G_0=133_8$  and  $G_1=171_8$ . The output of the convolutional encoder may be punctured to rates of ½, ¾rds and ¾ according to the specified rate tables (e.g., tables 1-12). For backward compatibility with IEEE 802.11(b) and the CCK modes of IEEE 802.11 (g), the channel encoder has the form 60 of a CCK code as defined in IEEE 802.11 (b). For higher data rates (such as those illustrated in tables 6, 8 and 10), the channel encoder may use the same convolution encoding as described above or it may use a more powerful code, including a convolutional code with more states, a parallel concat- 65 enated (turbo) code and/or a low density parity check (LDPC) block code. Further, any one of these codes may be combined

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with an outer Reed Solomon code. Based on a balancing of performance, backward compatibility and low latency, one or more of these codes may be optimal. Note that the concatenated turbo encoding and low density parity check will be described in greater detail with reference to FIGS. 12-19.

The interleaver 176 receives the encoded data and spreads it over multiple symbols and transmit streams. This allows improved detection and error correction capabilities at the receiver. In one embodiment, the interleaver 176 will follow the IEEE 802.11(a) or (g) standard in the backward compatible modes. For higher performance modes (e.g., such as those illustrated in tables 6, 8 and 10), the interleaver will interleave data over multiple transmit streams. The demultiplexer 178 converts the serial interleave stream from interleaver 176 into M-parallel streams for transmission.

Each symbol mapper **180-184** receives a corresponding one of the M-parallel paths of data from the demultiplexer. Each symbol mapper **180-182** lock maps bit streams to quadrature amplitude modulated QAM symbols (e.g., BPSK, QPSK, 16 QAM, 64 QAM, 256 QAM, et cetera) according to the rate tables (e.g., tables 1-12). For IEEE 802.11 (a) backward compatibility, double gray coding may be used.

The map symbols produced by each of the symbol mappers 180-184 are provided to the IFFT/cyclic prefix addition modules 186-190, which performs frequency domain to time domain conversions and adds a prefix, which allows removal of inter-symbol interference at the receiver. Note that the length of the IFFT and cyclic prefix are defined in the mode tables of tables 1-12. In general, a 64-point IFFT will be used for 20 MHz channels and 128-point IFFT will be used for 40 MHz channels.

The space/time encoder 192 receives the M-parallel paths of time domain symbols and converts them into P-output symbols. In one embodiment, the number of M-input paths will equal the number of P-output paths. In another embodiment, the number of output paths P will equal M+1 paths. For each of the paths, the space/time encoder multiples the input symbols with an encoding matrix that has the form of

$$\begin{bmatrix} C_1 & C_2 & C_3 & \cdots & C_{2M-1} \\ -C_2^* & C_1^* & C_4 & \cdots & C_{2M} \end{bmatrix}$$

Note that the rows of the encoding matrix correspond to the number of input paths and the columns correspond to the number of output paths.

FIG. 10B illustrates the radio portion of the transmitter that includes a plurality of digital filter/up-sampling modules 194-198, digital-to-analog conversion modules 200-204, analog filters 206-216, I/Q modulators 218-222, RF amplifiers 224-228, RF filters 230-234 and antennas 236-240. The P-outputs from the space/time encoder 192 are received by respective digital filtering/up-sampling modules 194-198.

In operation, the number of radio paths that are active correspond to the number of P-outputs. For example, if only one P-output path is generated, only one of the radio transmitter paths will be active. As one of average skill in the art will appreciate, the number of output paths may range from one to any desired number.

The digital filtering/up-sampling modules 194-198 filter the corresponding symbols and adjust the sampling rates to correspond with the desired sampling rates of the digital-to-analog conversion modules 200-204. The digital-to-analog conversion modules 200-204 convert the digital filtered and up-sampled signals into corresponding in-phase and quadrature analog signals. The analog filters 208-214 filter the cor-

responding in-phase and/or quadrature components of the analog signals, and provide the filtered signals to the corresponding I/Q modulators 218-222. The I/Q modulators 218-222 based on a local oscillation, which is produced by a local oscillator 100, up-converts the I/Q signals into radio frequency signals.

The RF amplifiers 224-228 amplify the RF signals which are then subsequently filtered via RF filters 230-234 before being transmitted via antennas 236-240.

FIGS. 11A and 11B illustrate a schematic block diagram of another embodiment of a receiver in accordance with the present invention. FIG. 11A illustrates the analog portion of the receiver which includes a plurality of receiver paths. Each receiver path includes an antenna, RF filters 252-256, low noise amplifiers 258-260, I/Q demodulators 264-268, analog 15 filters 270-280, analog-to-digital converters 282-286 and digital filters and down-sampling modules 288-290.

In operation, the antennas receive inbound RF signals, which are band-pass filtered via the RF filters **252-256**. The corresponding low noise amplifiers **258-260** amplify the filtered signals and provide them to the corresponding I/Q demodulators **264-268**. The I/Q demodulators **264-268**, based on a local oscillation, which is produced by local oscillator **100**, down-converts the RF signals into baseband inphase and quadrature analog signals.

The corresponding analog filters 270-280 filter the in-phase and quadrature analog components, respectively. The analog-to-digital converters 282-286 convert the in-phase and quadrature analog signals into a digital signal. The digital filtering and down-sampling modules 288-290 30 filter the digital signals and adjust the sampling rate to correspond to the rate of the baseband processing, which will be described in FIG. 11B.

FIG. 11B illustrates the baseband processing of a receiver. The baseband processing includes a space/time decoder 294, 35 a plurality of fast Fourier transform (FFT)/cyclic prefix removal modules 296-300, a plurality of symbol demapping modules 302-306, a multiplexer 308, a deinterleaver 310, a channel decoder 312, and a descramble module 314. The baseband processing module may further include a mode 40 managing module 175. The space/time decoding module 294, which performs the inverse function of space/time encoder 192, receives P-inputs from the receiver paths and produce M-output paths. The M-output paths are processed via the FFT/cyclic prefix removal modules 296-300 which perform 45 the inverse function of the IFFT/cyclic prefix addition modules 186-190 to produce frequency domain symbols.

The symbol demapping modules 302-306 convert the frequency domain symbols into data utilizing an inverse process of the symbol mappers 180-184. The multiplexer 308 combines the demapped symbol streams into a single path.

The deinterleaver 310 deinterleaves the single path utilizing an inverse function of the function performed by interleaver 176. The deinterleaved data is then provided to the channel decoder 312 which performs the inverse function of 55 channel encoder 174. The descrambler 314 receives the decoded data and performs the inverse function of scrambler 172 to produce the inbound data 98.

FIG. 12 is a schematic block diagram of channel encoder 174 implemented as a turbo encoder. In this embodiment, the 60 turbo encoder receives input bits, modifies them, processes them via a constituent encoder 320-322 and interleaves them to produce the corresponding encoded output. Depending on the particular symbol mapping (BPSK, QPSK, 8PSK (phase shift keying), 64 QAM, 16 QAM or 16APSK (amplitude 65 phase shift keying), the turbo encoder will function in the same manner to produce the encoded data. For instance, of  $\pi_0$ 

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and  $\pi_1$  are interleaves of MSB (most significant bit) and LSB (least significant bit), respectively, for a block of 2-bit symbols and  $\pi_L^{-1}$ , L=0, are the inverses, then the modified interleave is as follows:

$$\pi'' l(i) = \begin{pmatrix} i : i \mod 2 = 0 \\ \pi - 1(i) : i \mod 2 = 1 \end{pmatrix} \text{ and }$$

$$\pi l(i) = \begin{pmatrix} i : i \mod 2 = 1 \\ \pi(i) : i \mod 2 = 0 \end{pmatrix}$$

FIG. 13 illustrates an embodiment of the constituent encoders 320-322 of FIG. 12 which may be implemented as rate ½ encoders.

FIG. 14 illustrates a schematic block diagram of another embodiment of a constituent encoder 320-322 that utilizes the ½ rate encoder to produce a rate ½ ths encoder. In this embodiment, two consecutive binary inputs are sent to the rate ½ encoder. The output of the rate ½ encoder is produced as shown.

FIG. 15 represents the generally functionality of FIG. 14. The rate <sup>2</sup>/<sub>s</sub>ths encoder may then be utilized as puncture encoders as shown in FIGS. 16 and 17, which have the corresponding QPSK mapping.

FIG. 18 illustrates the channel encoder 174 being implemented as a low density parity check (LDPC) encoder. In this embodiment, the encoder includes a low density parity check encoder 174, an interleaver 176 and a gray mapping module 177. The block length may be 2000 and the information length may be 1600. In this instance, the low density parity check binary matrix  $H=[H_1, H_2]$ , where  $H_1$  is an irregular  $400\times1600$  low density matrix with 1400 columns of weight 3 and 200 columns of weight 7, and all rows of weight 14. More over, the distribution of the 1's is pseudo random in order to suit a hardware embodiment. The matrix  $H_2$  is a  $400\times400$  matrix which provides a long path with no loops in the bipartite graph between redundancy bit node and check node.

$$H2 = \begin{vmatrix} 100 \dots 00 \\ 11_{-} \dots 00 \\ -11 \dots 00 \\ \dots \\ 000 \dots 10 \\ 000 \dots 11 \end{vmatrix}$$

This parity check matrix provides easy encoding. The code has no circle of loops less than 6. The degree distribution of the bipartite graph of the code is listed in the following table. The total number of edges of the graph is 6399.

		number of nodes
	bit node degree (number of edges emitted from a bit node)	
)	1 2	1 399
	3 7	1400 200
	check node degree	200
	15	1
;	16	399

FIG. 19 illustrates a particular interleaving that may be utilized by the encoder of FIG. 18. In this embodiment, the rate of the code may be ½ and the LDPC code is symmetric. As such, the interleaving is as shown.

FIG. 20 is a schematic block diagram of a channel encoder 5 174 that includes an outer encoder, an optional byte interleaver, and an inner encoder. The outer encoded may be a Reed-Solomon encoder and the inner encoder may be a punctured convolutional encoder.

In one embodiment, the channel encoder 174 includes the 10 byte level interleaver, which has a depth sufficient to interleave the entire frame. Since no decoding can be performed until the entire frame has been received, the interleaver may add more latency than is desired at the receiver. To overcome the potential latency issue, a short interleaver that encom- 15 passes only 2 or 3 Reed Solomon codewords may be used. A short interleaver is a compromise between latency and performance. An interleaver is not required at all if a single Reed Solomon codeword can correct the longest typical error burst from Viterbi decoder in the receiver. For the 802.11a code 20 (with polynomial described by octal g0=133 g1=171), punctured to rate ½ with puncture pattern ([1111010; 1000101] i.e. 0 is a deleted bit) the error bursts have length: at dmin (3)—length 3 and 11, at dmin+1 (4)—length up to 43, at dmin+2 (5)—length up to 83. Thus a Reed Solomon code that 25 can correct 83 bit errors will be effective without an interleaver. This corresponds to t=12 (ie. Ceil(83/8)+1).

In one embodiment, the inner encoder is the same convolutional encoder as 802.11a, with puncturings defined as in 802.11a for rates ½ and ¾, with the addition of new options 30 for this code. Option 1 is to change it to a 256 state code. Option 2 is to add new puncturings for rates ½, ⅙, and ⅙. Both options may be combined. These are viable puncturings for this code and are described in the paper: "High-Rate Punctured Convolutional Codes for Soft Decision Viterbi 35 Decoding" by Yutaka Yasuda, Kanshiro Kashiki, and Tasuo Hirata, IEEE Transactions on Communications, Vol. COM-32, No. 3, March 1984, pp. 315-319.

The Reed-Solomon encoder may be designed to use multiple codeword lengths. In one embodiment, the decoder 40 operates over GF(256), using a codeword length n=255, and information sequence length k=239. This will allow correction of up to t=8 byte errors per code word.

For an effective rate 0.8 code, the concatenated coding scheme has a gain of 4 dB or more over the convolutional code 45 alone. This is demonstrated in the tables of FIGS. **21**A-D, which compare the frame error rate (FER) performance of the systems over an ensemble of channels generated with the IEEE channel model.

By adding the outer Reed-Solomon encoder, coding 50 gain—i.e. effective operation at a lower SNR for the same frame error rate—is achieved. Such a Reed-Solomon encoder uses punctured convolutional code that has been used for 802.11a and .11g. Further, since 802.11n will require longer frames (perhaps 4096 bytes) than .11a (1500 bytes typical) 55 and will require lower frame error rates to achieve high throughput, Reed-Solomon coding is known to work well under these conditions. Note that more sophisticated, higher performance, Reed-Solomon decoders can be implemented at higher complexity, if desired.

For example, a comparison of BER and FER Performance with and without outer RS Code demonstrates that performance bound allows comparison over many channels, SNR levels, and low BERs. For this example, assume MMSE spatial processing, ideal bit interleaving for inner code, bit errors are uniformly distributed, ideal byte interleaving for outer code, and BER computed using union bound. From this per-

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spective, Present Cumulative Distribution Function (CDF) of P(FER<10%) and P(FER<1%) as a function of RX SN, with 2 TX, 2 RX MIMO, and per IEEE Channel Models B and D, 4096 byte frames have a gain with RS codes that is typically 1 dB less than 1024 byte frames and Gain of Rate 7% CC with N=255, K=239 RS relative to CC alone has a net code rate of 0.82.

FIG. 22 is a schematic block diagram of a channel decoder 312 that includes an inner decoder, a deinterleaver, and an outer decoder. The inner decoder may be a Viterbi decoder and the outer decoder may be a Reed-Solomon decoder. In one embodiment, the Reed-Solomon decoder is a simpler structure than the decoders for LDPC (low density parity check) and turbo codes.

While it is possible for the Reed-Solomon decoder in the receiver to operate only on decoded bits from the Viterbi decoder, it is desirable if the Viterbi decoder can provide soft decision information to the Reed-Solomon decoder. This information is in the form of erasures, i.e. indications that certain decoder bits are likely to be in error because they have large error metric. Since, a t error correcting code can correct 2t erasures, it is better to operate on erasures, if possible.

FIG. 23 is a schematic block diagram of a frame that may be transmitted from a transmitter of one wireless communication device to a receiver of one or more other wireless communication devices. The frame includes a preamble section, a signal field, a service field, a data section, and a tail and pad bits section. The preamble includes one or more training sequences to facilitate a wireless communication between the wireless communication devices. The signal field includes information pertaining to the length of the frame, the rate of the data within the frame, etc.

In one embodiment, the preamble section and signal field of a frame bypasses the channel encoder of the transmitter and hence the channel decoder of the receiver. The service field and the data section may or may not bypass the channel encoder, or portions thereof, based on the desired latency within the receiver. For example, a number of bytes of the service field and/or data field may be interleaved and outer Reed-Solomon encoded based on the size of the interleaver and/or based on the size of the Reed-Solomon encoder. Once enough data has been received to fill the Reed-Solomon encoder, subsequent bytes bypass the interleaver and/or Reed-Solomon encoder. The amount of data that bypasses the interleaver and/or Reed-Solomon encoder depends on the speed of the encoding process, the size of the frame, the size of the Reed-Solomon encoder, and/or a desired latency of the receiver.

FIG. 24 is a schematic block diagram of an embodiment of the channel encoder 174, which includes a row/column byte interleaver, a Reed-Solomon encoder, a punctured convolutional encoder, an append tail and pad bits module, and a plurality of switching elements. The row/column byte interleaver interleaves the scrambled data of the service field and/or of the data section as shown in FIG. 25 when the switching element S1 is closed.

As shown in FIG. 25 bytes of data are read into the interleaver, which may be a memory device as defined herein, row by row and read out of the interleaver column by column. Note that, in one embodiment, the number of rows is set by the number of bytes to be encoded in a codeword. (e.g., 239 for a [239, 255] Reed-Solomon encoder). The number of columns is equal to the interleaver depth. Longer depths provide more coding gain, but increase latency. While depths of 3 to 5 provide acceptable tradeoffs for 802.1 In applications, other depths may be used.

Returning back to the discussion of FIG. 24, the Reed-Solomon (RS) encoder, which may be a [239, 255] block encoder, receives its input from switching element S2. As such, the Reed-Solomon encoder either encodes the interleaved data outputted by the interleaver or the scrambled 5 input data. In one embodiment, the initial bytes (usually most of them) of a frame are sent through the entire chain (i.e., the switching elements couple the interleaver to the RS encoder, which is coupled to the punctured convolution encoder). Towards the end of the frame the switches are changed to 10 route the bytes so that the interleaver and/or Reed-Solomon encoder are bypassed. The bypass operation allows the receiver to reduce its latency because it does not have to buffer an entire received frame to apply the Reed-Solomon decoder. In practice, it is desirable to break up the received data into 15 blocks of 239 bytes (for the 239, 255) Reed-Solomon encoder. The remaining bytes at the end of the frame that do not fit into a 239 byte block are not encoded by the Reed-Solomon outer encoder.

To balance maximum performance and minimal latency, 20 input data bytes may be broken up into blocks of R\*N bytes, where N is the interleaver depth and R is the codeword size of the RS encoder (e.g., 239). The bytes that fit into these blocks should pass through the interleaver and encoder. The remaining bytes bypass the interleaver but are still encoded into 239 bytes or shortened Reed-Solomon codewords. Alternatively, the remaining bytes may bypass both the interleaver and the RS encoder.

The punctured convolutional encoder is operably coupled to receive its input from switching element S3. As such, the 30 punctured convolutional encoder is either encoded the output of the RS encoder or the scrambled data of the service field and/or the data section of a frame. In one embodiment, the punctured convolutional encoder may be a 64 state rate  $\frac{1}{2}$  code with polynomials  $G_0=133_8$  and  $G_1=171_8$ . The encoder 35 may be punctured to rate  $\frac{3}{4}$  with puncturing pattern of [110; 101], rate  $\frac{4}{5}$  with a puncture pattern of [1 1 0 1 0; 1 0 1 0 1], and rate  $\frac{7}{8}$  with a puncture pattern of [1 1 1 1 0 1 0; 1 0 0 0 1 0 1]. Within the puncture patterns, a 1 indicates a bit is kept, while a zero 40 indicates a bit that is punctured.

In another embodiment, the encoder may be a 256 state rate  $^{1}$ /2 code with polynomials  $G_0$ =561 $_8$  and  $G_1$ =753 $_8$ . The encoder may be punctured to rate  $^{3}$ /4 with a puncturing pattern of [111; 100], rate  $^{4}$ /5 with a puncture pattern of [1 1 0 1; 1 0 45 1 0], rate  $^{5}$ /6 with a puncture pattern of [1 0 1 1 0; 1 1 0 0 1], and rate  $^{7}$ /8 with a puncture pattern of [1 1 0 1 0 1 1; 1 0 1 0 1 0 0]. This longer constraint length code improves coding gain, but increase decoder complexity with respect to the previous embodiment of the punctured convolutional encoder.

FIG. **25** is a schematic block diagram of an embodiment of a channel decoder **312**, which includes a remove tail and pad bits module, a punctured convolutional decoder, a Reed-Solomon decoder, a deinterleaver and a plurality of switching elements **S4-S6**. The punctured convolutional decoder, which stream of symbols from the bit level deinterleaver **310** of FIG. **11B** and decodes it to produce inner punctured decoded data. In one embodiment, the punctured convolutional decoder may use a 64 state rate  $\frac{1}{2}$  code with polynomials  $G_0$ =133 $_8$  and  $G_1$ =171 $_8$ . The 60 encoder may be punctured to rate  $\frac{3}{4}$  with puncturing pattern of [110; 101], rate  $\frac{4}{5}$  with a puncture pattern of [1 1 1; 1 0], and rate  $\frac{5}{6}$  with a puncture pattern of [1 1 0 1 0; 1 0 1 0], and rate  $\frac{7}{8}$  with a puncture pattern of [1 1 1 1 0 1 0; 1 0 0 0 1 0 1].

In another embodiment, the decoder may use a 256 state 65 rate  $\frac{1}{2}$  code with polynomials  $G_0=561_8$  and  $G_1=753_8$ . The decoder may further use a rate  $\frac{3}{4}$  with a puncturing pattern of

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[111; 100], rate \(^4\)s with a puncture pattern of [1 1 0 1; 1 0 1 0], rate \(^5\)/s with a puncture pattern of [1 0 1 1 0; 1 1 0 0 1], and rate \(^7\)/s with a puncture pattern of [1 1 0 1 0 1 1; 1 0 1 0 1 0 0]. This longer constraint length code improves coding gain, but increase decoder complexity with respect to the previous embodiment of the punctured convolutional encoder.

The output of the punctured convolutional decoder is either provided to the outer Reed-Solomon (RS) decoder or is provided as the output of the channel decoder 312. If outer RS decoder receives the inner punctured decoded data it decodes it to produce outer decoded data. Note that the RS decoder will be complementary to the RS encoder of the channel encoder 174.

Switching element S5 either provides the output of the RS encoder to the deinterleaver or as the output of the channel decoder 312. When switching element S5 provides the RS decoder output to the deinterleaver, the deinterleaver deinterleaves, at a word level, the outer decoded data to produce deinterleaved data when the deinterleaving is enabled. The deinterleaver performs a complementary function to that of the interleaver of the channel encoder 174. With reference to FIG. 26, the deinterleaver would write data in on a column by column basis and read it out on a row by row basis.

Returning to the discussion of FIG. 25, switching element S6 provides the output of the deinterleaver as the output of the channel decoder 312, when the deinterleaver is enabled. In such an embodiment, the output of the channel decoder 312 is the inner punctured decoded data, the outer decoded data, or the deinterleaved data.

The switching elements S1-S6 may be any type of device that provides selective coupling such as, but not limited to, transistors, electrical switches, optical switches, and/or mechanical switches. Note that the activation of switching elements S4-S6 will correspond to the activation of switching elements S1-S3 to provide the desired level of performance and receiver latency.

As one of average skill in the art will appreciate, the term "substantially" or "approximately", as may be used herein, provides an industry-accepted tolerance to its corresponding term. Such an industry-accepted tolerance ranges from less than one percent to twenty percent and corresponds to, but is not limited to, component values, integrated circuit process variations, temperature variations, rise and fall times, and/or thermal noise. As one of average skill in the art will further appreciate, the term "operably coupled", as may be used herein, includes direct coupling and indirect coupling via another component, element, circuit, or module where, for indirect coupling, the intervening component, element, circuit, or module does not modify the information of a signal but may adjust its current level, voltage level, and/or power level. As one of average skill in the art will also appreciate, 50 inferred coupling (i.e., where one element is coupled to another element by inference) includes direct and indirect coupling between two elements in the same manner as "operably coupled". As one of average skill in the art will further appreciate, the term "compares favorably", as may be used herein, indicates that a comparison between two or more elements, items, signals, etc., provides a desired relationship. For example, when the desired relationship is that signal 1 has a greater magnitude than signal 2, a favorable comparison may be achieved when the magnitude of signal 1 is greater than that of signal 2 or when the magnitude of signal 2 is less than that of signal 1.

The preceding discussion has presented various embodiments of a multiple input/multiple output transceiver for use in wireless communication systems. As one of average skill in the art will appreciate, other embodiments may be derived from the teaching of the present invention without deviating from the scope of the claims.

### Mode Selection Tables:

TABLE 1

	2.	4 GHz, 2	20/22 MHz	channel	BW, 54 M	bps max	bit rate		
Rate	Modulation	Code Rate	NBPSC	NCBPS	NDBPS	EVM	Sensi- tivity	ACR	AACR
	Barker								
1	BPSK								
	Barker								
2	QPSK								
5.5	CCK								
6	BPSK	0.5	1	48	24	-5	-82	16	32
9	BPSK	0.75	1	48	36	-8	-81	15	31
11	CCK								
12	QPSK	0.5	2	96	48	-10	-79	13	29
18	QPSK	0.75	2	96	72	-13	-77	11	27
24	16-QAM	0.5	4	192	96	-16	-74	8	24
36	16-QAM	0.75	4	192	144	-19	<b>-7</b> 0	4	20
48	64-QAM	0.666	6	288	192	-22	-66	O	16
54	64-QAM	0.75	6	288	216	-25	-65	-1	15

TABLE 2 TABLE 5

Channelizat	tion for Table 1				Channelization f	for Table 4		
Channel	Frequency (MHz)		Channel	Frequency (MHz)	Country	Channel	Frequency (MHz)	Country
1	2412	30		(11222)			(11222)	
2	2417	30	240	4920	Japan			
3	2422		244	<b>494</b> 0	Japan			
4	2427		248	4960	Japan			
5	2432		252	4980	Japan			
6	2437		8	5040	Japan			
7	2442	2.5	12	5060	Japan			
8	2447	35	16	5080	Japan			
9	2452		36	5180	USA/Europe	34	5170	Japan
10	2457		40	5200	USA/Europe	38	5190	Japan
11	2462		44	5220	USA/Europe	42	5210	Japan
12	2467		48	5240	USA/Europe	46	5230	Japan
			52	5260	USA/Europe			-
		40	56	5280	USA/Europe			
			60	5300	USA/Europe			
ТАТ	BLE 3		64	5320	USA/Europe			
	JLE J		100	5500	USA/Europe			
Darrian Charatual Danaita	(DCD) Maal- fau Tal-la 1		104	5520	USA/Europe			
Power Spectral Density	y (PSD) Mask for Table 1		108	5540	USA/Europe			
DCD Magle	1	45	112	5560	USA/Europe			
PSD Mask	175	15	116	5580	USA/Europe			
Frequency Offset	$\mathrm{dBr}$		120	5600	USA/Europe			
-9 MHz to 9 MHz	0		124	5620	USA/Europe			
+/- 11 MHz	-20		128	5640	USA/Europe			
+/- 20  MHz	-28 -28		132	5660	USA/Europe			
+/- 20 MHz and greater		50	136	5680	USA/Europe			
		50	140	<b>57</b> 00	USA/Europe			

TABLE 4

		5 GHz, 2	20 MHz c	hannel BV	V, 54 Mbp	s max b	it rate		
Rate	Modulation	Code Rate	NBPSC	NCBPS	NDBPS	EVM	Sensi- tivity	ACR	AACR
6	BPSK	0.5	1	48	24	-5	-82	16	32
9	BPSK	0.75	1	48	36	-8	-81	15	31
12	QPSK	0.5	2	96	48	-10	-79	13	29
18	QPSK	0.75	2	96	72	-13	-77	11	27
24	16-QAM	0.5	4	192	96	-16	-74	8	24
36	16-QAM	0.75	4	192	144	-19	-70	4	20
48	64-QAM	0.666	6	288	192	-22	-66	0	16
54	64-QAM	0.75	6	288	216	-25	-65	-1	15

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### TABLE 5-continued

TABLE 5-continued

	(	Channelization	for Table 4					(	Channelization	for Table 4		
Channel	Frequency (MHz)	Country	Channel	Frequency (MHz)	Country	5	Channel	Frequency (MHz)	Country	Channel	Frequency (MHz)	Country
149 153	5745 5765	USA USA					161 165	5805 5825	USA USA			
157	5785	USA										

TABLE 6

	2.4 (	GHz, 20 MH:	z channel BW, 1	192 Mbps 1	max bit rat	e.	
Rate	TX Antennas	ST Code Rate	Modulation	Code Rate	NBPSC	NCBPS	NDBPS
12	2	1	BPSK	0.5	1	48	24
24	2	1	QPSK	0.5	2	96	48
48	2	1	16-QAM	0.5	4	192	96
96	2	1	64-QAM	0.666	6	288	192
108	2	1	64-QAM	0.75	6	288	216
18	3	1	BPSK	0.5	1	48	24
36	3	1	QPSK	0.5	2	96	48
72	3	1	16-QAM	0.5	4	192	96
144	3	1	64-QAM	0.666	6	288	192
162	3	1	64-QAM	0.75	6	288	216
24	4	1	BPSK	0.5	1	48	24
48	4	1	QPSK	0.5	2	96	48
96	4	1	16-QAM	0.5	4	192	96
192	4	1	64-QAM	0.666	6	288	192
216	4	1	64-QAM	0.75	6	288	216

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TABLE 7

	Channeli	zation for Table 6	
5	Channel	Frequency (MHz)	
	1	2412	
	2	2417	
	3	2422	
	4	2427	
	5	2432	
	6	2437	
	7	2442	
	8	2447	
	9	2452	
	10	2457	
	11	2462	
)	12	2467	

TABLE 8

	5 G	Hz, 20 MHz	channel BW, 19	92 Mbps n	nax bit rate	;	
Rate	TX Antennas	ST Code Rate	Modulation	Code Rate	NBPSC	NCBPS	NDBPS
12	2	1	BPSK	0.5	1	48	24
24	2	1	QPSK	0.5	2	96	48
48	2	1	16-QAM	0.5	4	192	96
96	2	1	64-QAM	0.666	6	288	192
108	2	1	64-QAM	0.75	6	288	216
18	3	1	BPSK	0.5	1	48	24
36	3	1	QPSK	0.5	2	96	48
72	3	1	16-QAM	0.5	4	192	96
144	3	1	64-QAM	0.666	6	288	192
162	3	1	64-QAM	0.75	6	288	216
24	4	1	BPSK	0.5	1	48	24
48	4	1	QPSK	0.5	2	96	48
96	4	1	16-QAM	0.5	4	192	96
192	4	1	64-QAM	0.666	6	288	192
216	4	1	64-QAM	0.75	6	288	216

TABLE 9

		channelization f	or Table 8		
Channel	Frequency (MHz)	Country	Channel	Frequency (MHz)	Country
240	4920	Japan			
244	<b>494</b> 0	Japan			
248	<b>496</b> 0	Japan			
252	<b>498</b> 0	Japan			
8	5040	Japan			
12	5060	Japan			
16	5080	Japan			
36	5180	USA/Europe	34	5170	Japan
40	5200	USA/Europe	38	5190	Japan
44	5220	USA/Europe	42	5210	Japan
48	5240	USA/Europe	46	5230	Japan
52	5260	USA/Europe			
56	5280	USA/Europe			
60	5300	USA/Europe			
64	5320	USA/Europe			
100	5500	USA/Europe			
104	5520	USA/Europe			
108	5540	USA/Europe			
112	<b>556</b> 0	USA/Europe			
116	5580	USA/Europe			
120	5600	USA/Europe			
124	5620	USA/Europe			
128	<b>564</b> 0	USA/Europe			
132	5660	USA/Europe			
136	5680	USA/Europe			
140	5700	USA/Europe			
149	5745	USA			
153	5765	USA			
	5785	USA			
157					
161 165	5805 5825	USA USA			

TABLE 10

Rate	<b>;</b>	TX Antennas	ST Code Rate	Modulation	Code Rate	NBPSC
13.5	Mbps	1	1	BPSK	0.5	1
27	Mbps	1	1	QPSK	0.5	2
54	Mbps	1	1	16-QAM	0.5	4
108	Mbps	1	1	64-QAM	0.666	6
121.5	Mbps	1	1	64-QAM	0.75	6
27	Mbps	2	1	BPSK	0.5	1
54	Mbps	2	1	QPSK	0.5	2
108	Mbps	2	1	16-QAM	0.5	4
216	Mbps	2	1	64-QAM	0.666	6
243	Mbps	2	1	64-QAM	0.75	6
40.5	Mbps	3	1	BPSK	0.5	1
81	Mbps	3	1	QPSK	0.5	2
162	Mbps	3	1	16-QAM	0.5	4
324	Mbps	3	1	64-QAM	0.666	6
365.5	Mbps	3	1	64-QAM	0.75	6
54	Mbps	4	1	BPSK	0.5	1
108	Mbps	4	1	QPSK	0.5	2
216	Mbps	4	1	16-QAM	0.5	4
432	Mbps	4	1	64-QAM	0.666	6
	Mbps	1	1	64-QAM	0.75	6

	_
PSD Mask	2
Frequency Offset	$\mathrm{dBr}$
-19 MHz to 19 MHz	0
+/- 21 MHz	-20
+/-30  MHz	-28
+/- 40 MHz and greater	<b>-5</b> 0

TABLE 12

,	Channelization for Table 10								
15	Channel	Frequency (MHz)	Country	Channel	Frequency (MHz)	County			
20	242	4930	Japan						
	250	<b>497</b> 0	Japan						
	12	5060	Japan						
	38	5190	USA/Europe	36	5180	Japan			
	46	5230	USA/Europe	44	5520	Japan			
	54	5270	USA/Europe						
25	62	5310	USA/Europe						
	102	5510	USA/Europe						
	110	5550	USA/Europe						
	118	<b>559</b> 0	USA/Europe						
	126	5630	USA/Europe						
	134	5670	USA/Europe						
	151	5755	USA						
	159	5795	USA						

What is claimed is:

- 1. A wireless local area network (WLAN) receiver having high data throughput, the WLAN receiver comprises:
  - a plurality of radio frequency (RF) receivers, wherein, based on a mode selection signal, a number of the plurality of RF receivers are enabled, wherein each of the number of the plurality of RF receivers that are enabled converts a corresponding one of a plurality of received RF signals into a corresponding stream of symbols such that a corresponding number of streams of symbols is produced; and
  - a baseband processing module operably coupled to: combine the corresponding number of streams of symbols into a single stream of symbols;
    - inner puncture convolution decode the single stream of symbols to produce inner punctured decoded data based upon a state rate code of a plurality of state rate codes;
    - outer Reed-Solomon decode, when enabled, the inner punctured decoded data to produce outer decoded data;
    - deinterleave, when enabled, at a word level the outer decoded data to produce deinterleaved data, the deinterleave includes deinterleaving bytes of the outer decoded data in a row column pattern with R rows and C columns, wherein the R rows correspond to R bytes of the inner punctured decoded data being decoded by the outer Reed-Solomon decode,
    - wherein the deinterleave and the outer Reed-Solomon decode are enabled on a frame-by-frame basis for a block of the inner punctured decoded data including R bytes of data of a frame, and disabled, on the frame-by-frame basis, for at least some of remaining inner punctured decoded data of the frame; and
    - descramble the inner punctured decoded data, the outer decoded data, or the deinterleaved data to produce inbound data.

2. The WLAN receiver of claim 1, wherein the outer Reed-Solomon decoding comprises:

performing the outer Reed-Solomon decoding based on GF(256), a codeword length n=255, and information sequence length k=239.

3. The WLAN receiver of claim 1, wherein the inner puncture convolution decoding comprises:

utilizing a 64 state rate  $\frac{1}{2}$  code of the plurality of state rate codes with polynomials  $G_0=133_8$  and  $G1=171_8$ ; and

utilizing a puncture rate 3/4 with a puncture pattern of [110; 101], a puncture rate 4/5 with a puncture pattern of [1111; 1000], a puncture rate 5/6 with a puncture pattern of [11010; 10101], or a puncture rate 7/8 with a puncture pattern of [1 1 1 1 0 1 0; 1 0 0 0 1 0 1].

4. The WLAN receiver of claim 1, wherein the inner puncture convolution decoding comprises:

utilizing a 256 state rate  $\frac{1}{2}$  code of the plurality of state rate codes with polynomials  $G_0=561_8$  and

 $G_1$ =753<sub>8</sub>; and utilizing a puncture rate  $\frac{3}{4}$  with a puncture pattern of [111; 100], a puncture rate  $\frac{4}{5}$  with a puncture pattern of [1 1 0 1; 1 0 1 0], a puncture rate 5/6 with a puncture pattern of [1 0 1 1 0; 1 1 0 0 1], or a puncture rate 7/8 with a puncture pattern of [1 1 0 1 0 1 1; 1 0 1 0 1 0 0].

5. A method in a wireless local area network (WLAN) receiver having high data throughput, the WLAN receiver includes plurality of radio frequency (RF) receivers, wherein each of a number of the plurality of RF receivers are enabled based upon a mode selection signal and convert a corresponding one of a plurality of received RF signals into a corresponding stream of symbols such that a corresponding number of streams of symbols is produced, the method comprising:

combining the corresponding number of streams of symbols into a single stream of symbols;

inner puncture convolution decoding the single stream of symbols to produce inner punctured decoded data based upon a state rate code of a plurality of state rate codes when enabled, outer Reed-Solomon decoding the inner

punctured decoded data to produce outer decoded data;

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when enabled, deinterleaving, at a word level, the outer decoded data to produce deinterleaved data, the deinterleaving including deinterleaving bytes of the outer decoded data in a row column pattern with R rows and C columns, wherein the R rows correspond to R bytes of the inner punctured decoded data being decoded by the outer Reed-Solomon decoding,

wherein the deinterleaving and the outer Reed-Solomon decoding are enabled on a frame-by-frame basis for a block of the inner punctured decoded data including R bytes of data of a frame, and disabled, on the frame-by-frame basis, for at least some of remaining inner punctured decoded data of the frame; and

descrambling either of the inner punctured decoded data, the outer decoded data, or the deinterleaved data to produce inbound data.

6. The method of claim 5, wherein the outer Reed-Solomon decoding comprises:

performing the outer Reed-Solomon decoding based on GF(256), a codeword length n=255, and information sequence length k=239.

7. The method of claim 5, wherein the inner puncture convolution decoding comprises:

utilizing a 64 state rate  $\frac{1}{2}$  code of the plurality of state rate codes with polynomials  $G_0=133_8$  and  $G_1=171_8$ ; and

utilizing a puncture rate <sup>3</sup>/<sub>4</sub> with a puncture pattern of [110; 101], a puncture rate <sup>4</sup>/<sub>5</sub> with a puncture pattern of [1111; 1000], a puncture rate <sup>5</sup>/<sub>6</sub> with a puncture pattern of [11010; 10101], or a puncture rate <sup>7</sup>/<sub>8</sub> with a puncture pattern of [1 1 1 1 0 1 0; 1 0 0 0 1 01].

8. The method of claim 5, wherein the inner puncture convolution decoding comprises:

utilizing a 256 state rate  $\frac{1}{2}$  code of the plurality of state rate codes with polynomials  $G_0=561_8$  and

 $G_1$ =753<sub>8</sub>; and utilizing a puncture rate <sup>3</sup>/<sub>4</sub> with a puncture pattern of [111; 100], a puncture rate <sup>4</sup>/<sub>5</sub> with a puncture pattern of [1 1 0 1; 1 0 1 0], a puncture rate <sup>5</sup>/<sub>6</sub> with a puncture pattern of [1 0 1 1 0; 1 1 0 0 1], or a puncture rate <sup>7</sup>/<sub>8</sub> with a puncture pattern of [1 10101 1; 1 0 1 0 1 00].

\* \* \* \* \*