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(54) **METHOD AND SYSTEM FOR SECURING DIGITAL ASSETS USING PROCESS-DRIVEN SECURITY POLICIES**

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See application file for complete search history.

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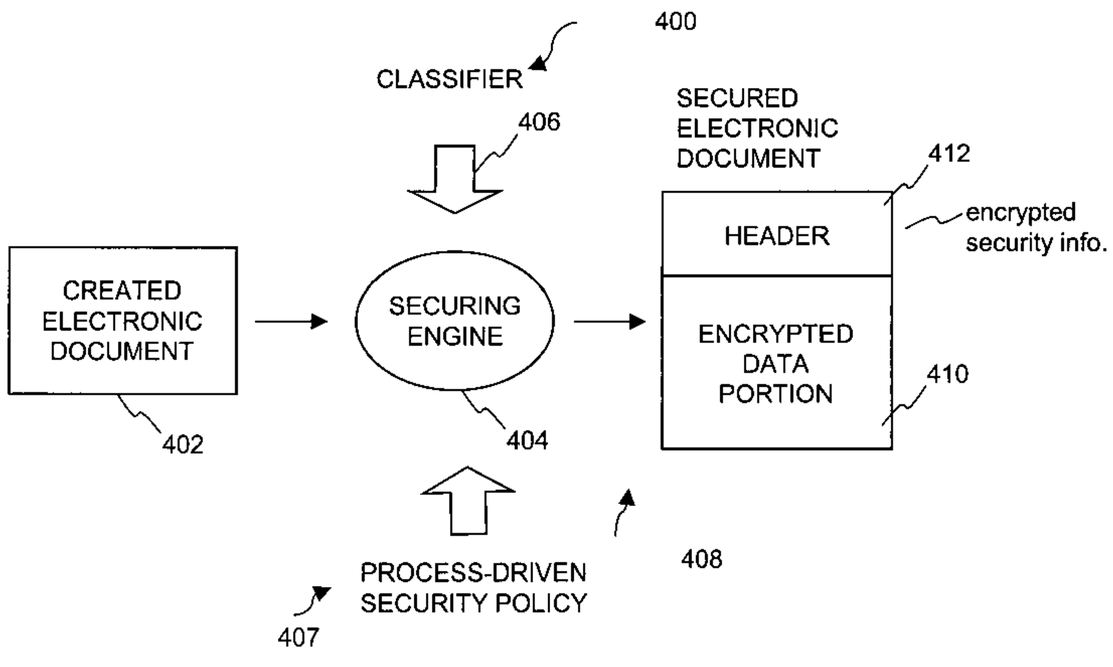
(57) **ABSTRACT**

Techniques for dynamically altering security criteria used in a file security system are disclosed. The security criteria pertains to keys (or ciphers) used by the file security system to encrypt electronic files to be secured or to decrypt electronic files already secured. The security criteria can, among other things, include keys that are required to gain access to electronic files. Here, the keys can be changed automatically as electronic files transition between different states of a process-driven security policy. The dynamic alteration of security criteria enhances the flexibility and robustness of the security system. In other words, access restrictions on electronic files can be dependent on the state of the process-driven security policy.

19 Claims, 13 Drawing Sheets



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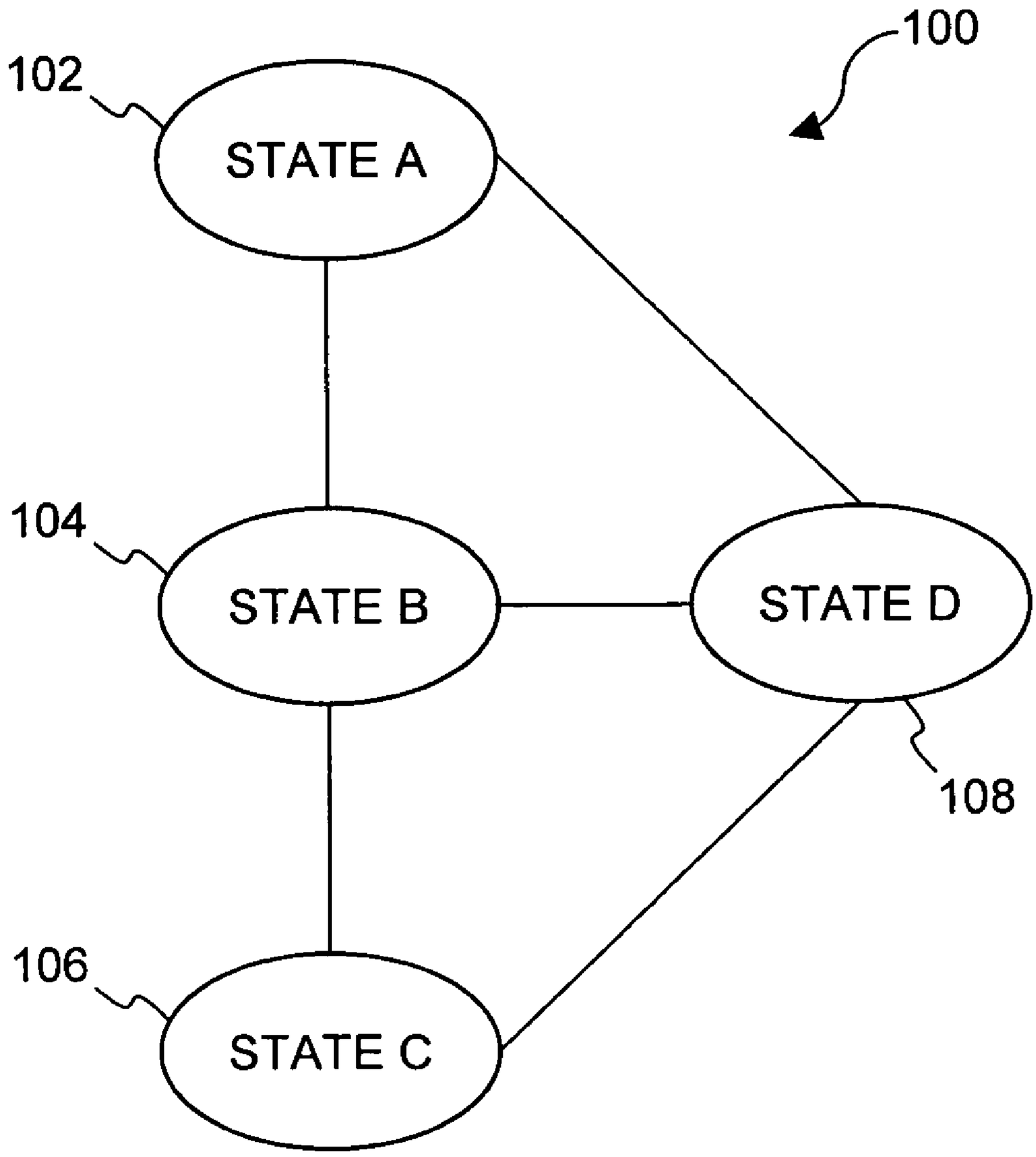


FIG. 1

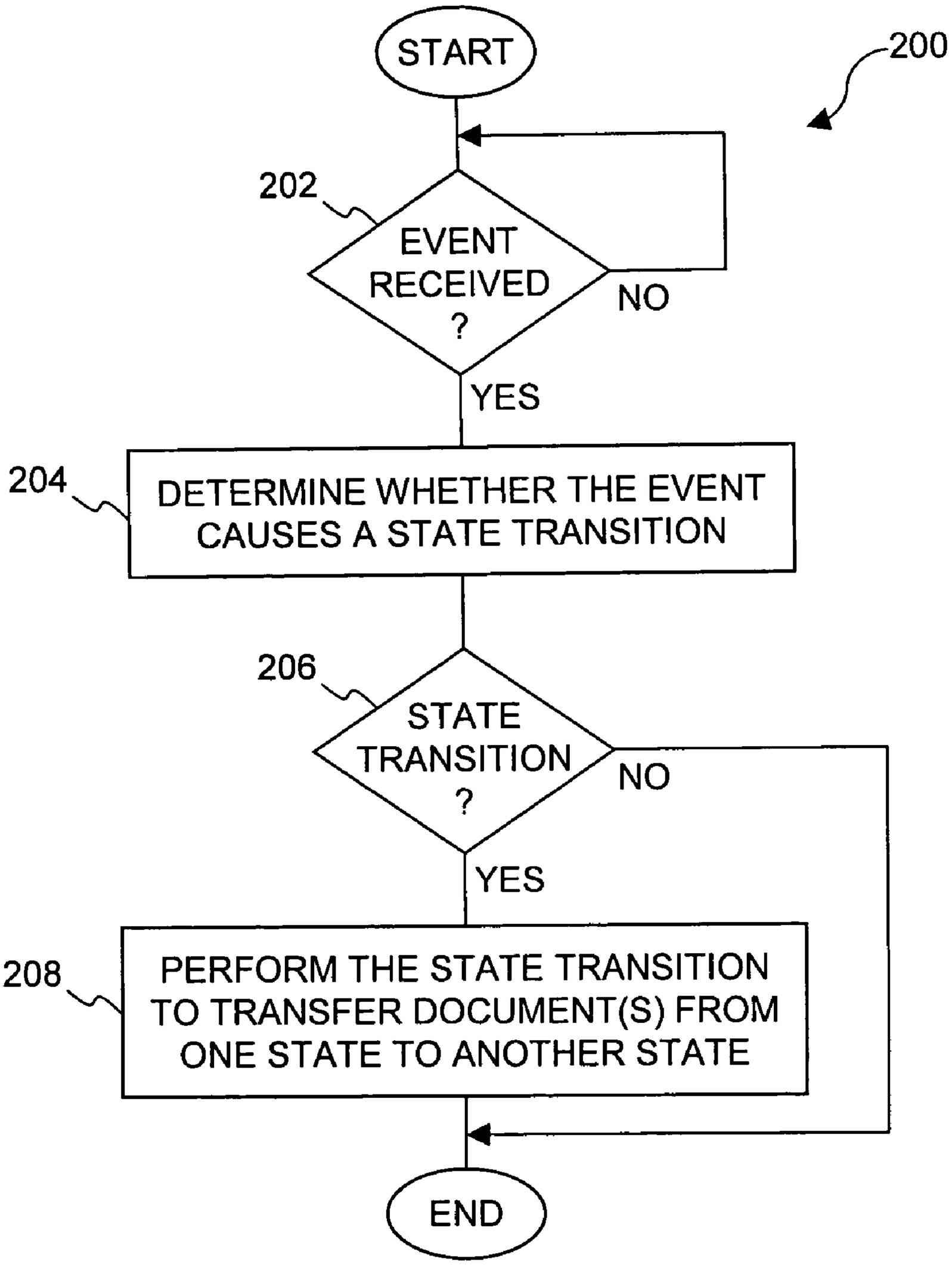

START
200
202
EVENT RECEIVED ?
NO
YES
204
DETERMINE WHETHER THE EVENT CAUSES A STATE TRANSITION
206
STATE TRANSITION ?
NO
YES
208
PERFORM THE STATE TRANSITION TO TRANSFER DOCUMENT(S) FROM ONE STATE TO ANOTHER STATE
END

FIG. 2

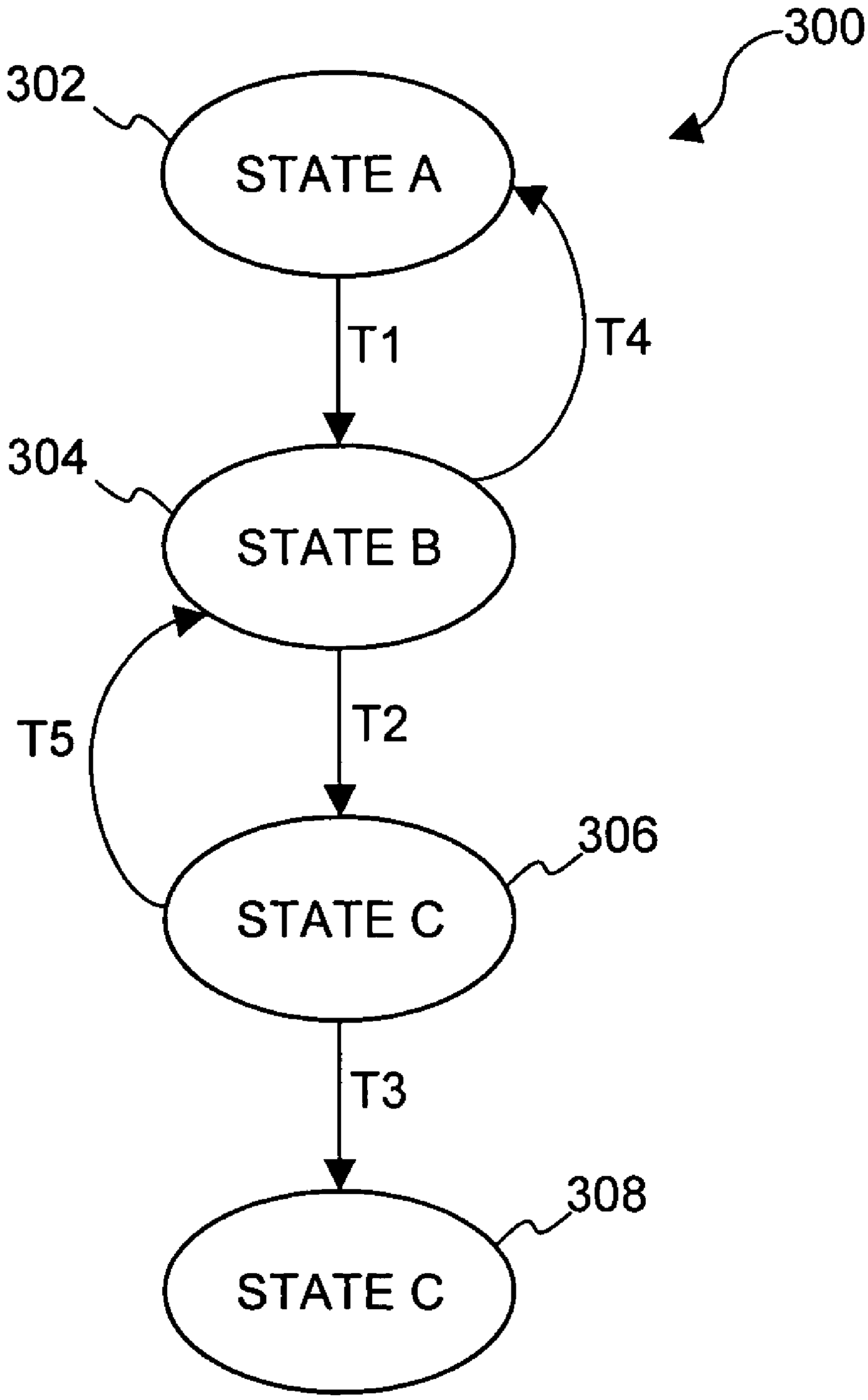


FIG. 3

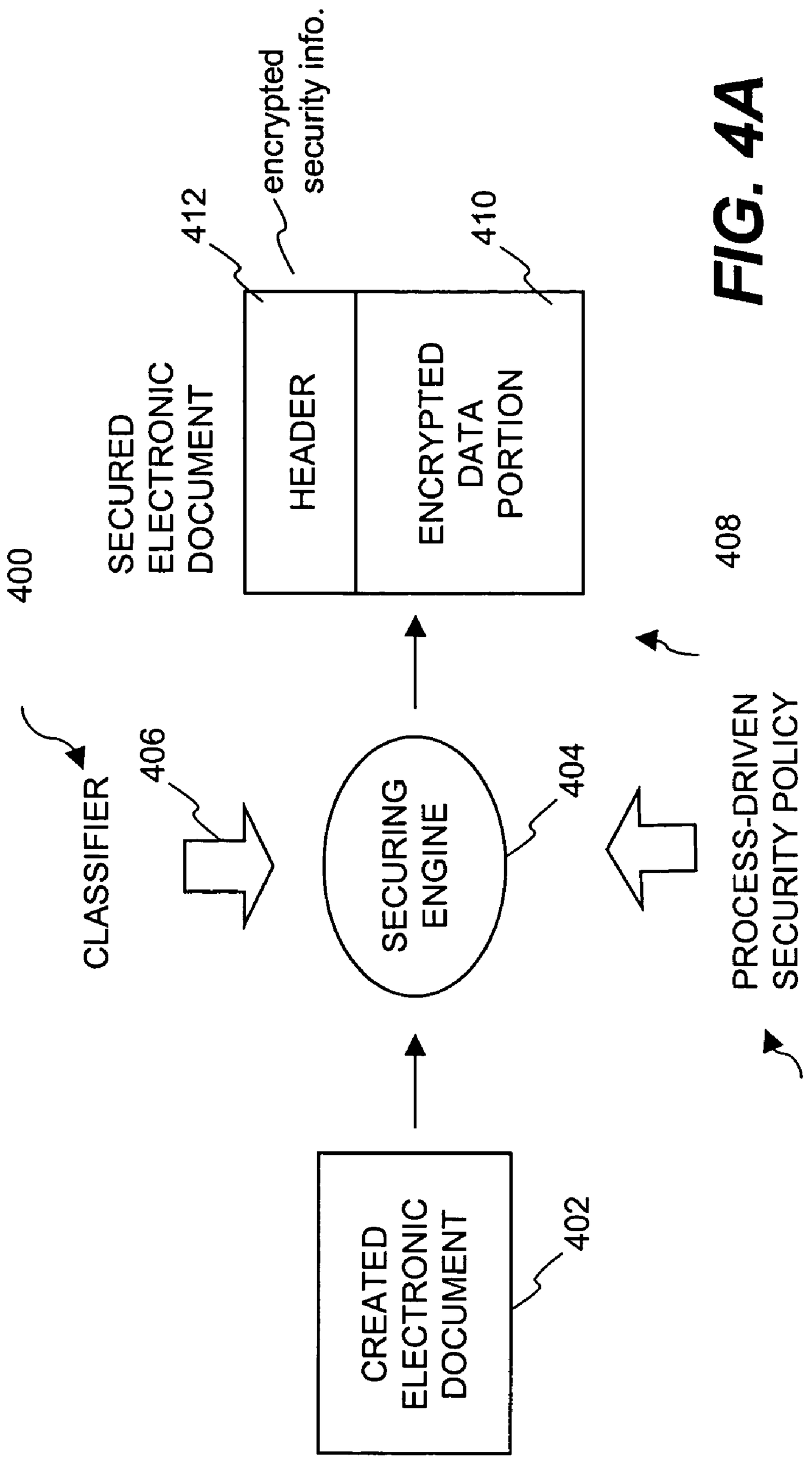

400
CLASSIFIER
406
SECURING ENGINE
404
CREATED ELECTRONIC DOCUMENT
402
SECURED ELECTRONIC DOCUMENT
412
encrypted security info.
HEADER
ENCRYPTED DATA PORTION
410
408
PROCESS-DRIVEN SECURITY POLICY
407

FIG. 4A

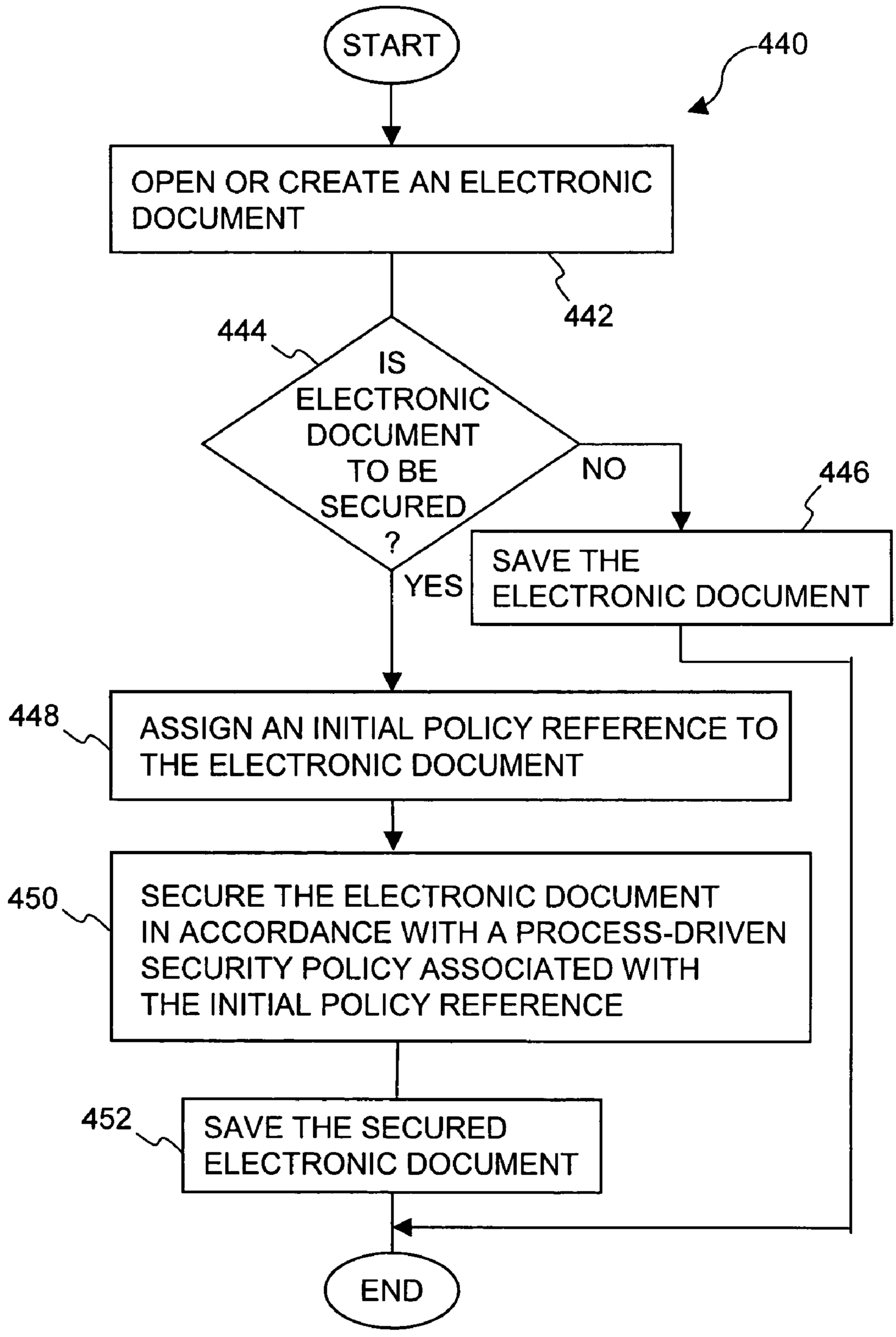

START
440
OPEN OR CREATE AN ELECTRONIC DOCUMENT
442
444
IS ELECTRONIC DOCUMENT TO BE SECURED ?
NO
446
SAVE THE ELECTRONIC DOCUMENT
YES
448
ASSIGN AN INITIAL POLICY REFERENCE TO THE ELECTRONIC DOCUMENT
450
SECURE THE ELECTRONIC DOCUMENT IN ACCORDANCE WITH A PROCESS-DRIVEN SECURITY POLICY ASSOCIATED WITH THE INITIAL POLICY REFERENCE
452
SAVE THE SECURED ELECTRONIC DOCUMENT
END

FIG. 4B

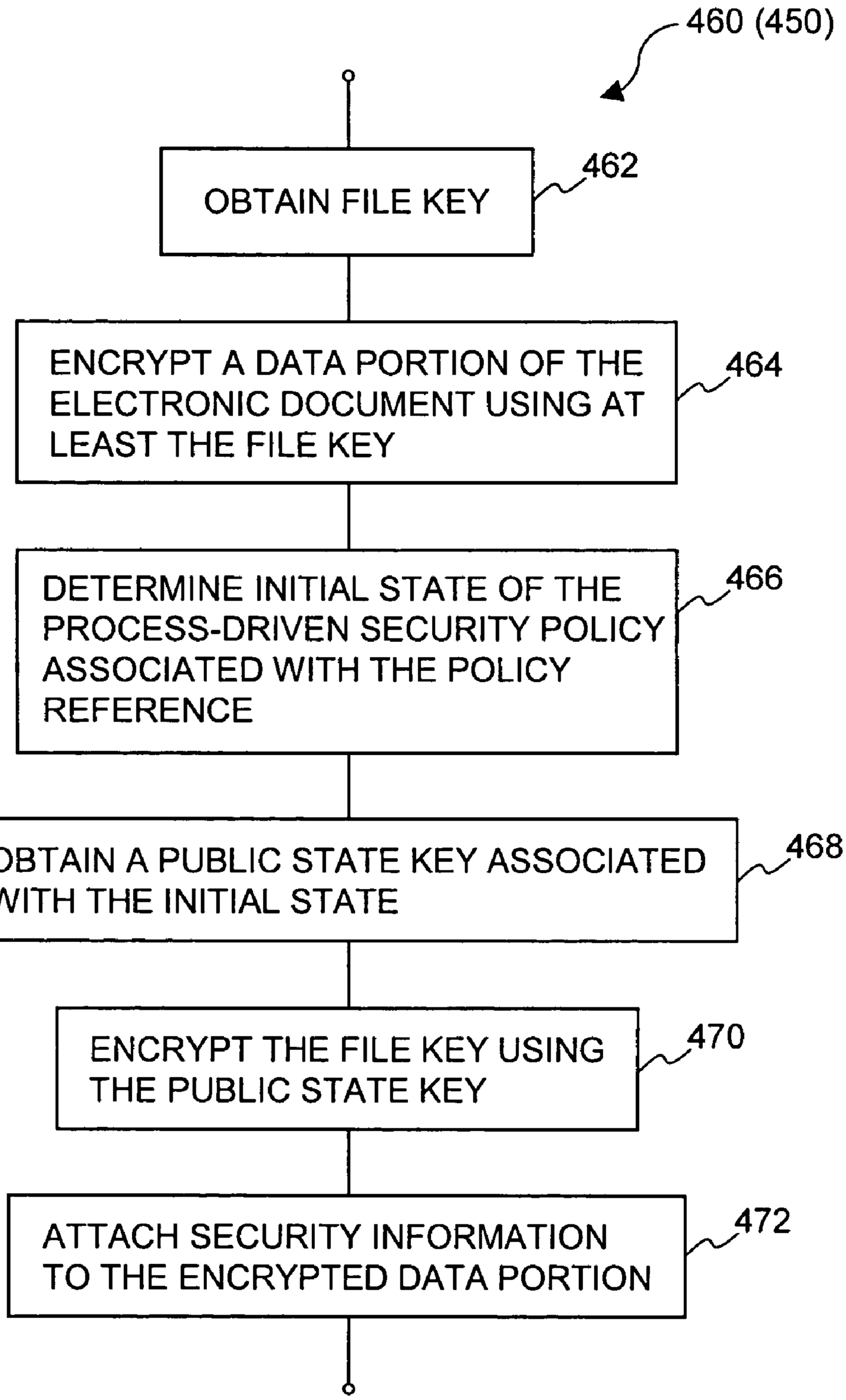

460 (450)
OBTAIN FILE KEY
462
ENCRYPT A DATA PORTION OF THE ELECTRONIC DOCUMENT USING AT LEAST THE FILE KEY
464
DETERMINE INITIAL STATE OF THE PROCESS-DRIVEN SECURITY POLICY ASSOCIATED WITH THE POLICY REFERENCE
466
OBTAIN A PUBLIC STATE KEY ASSOCIATED WITH THE INITIAL STATE
468
ENCRYPT THE FILE KEY USING THE PUBLIC STATE KEY
470
ATTACH SECURITY INFORMATION TO THE ENCRYPTED DATA PORTION
472

FIG. 4C

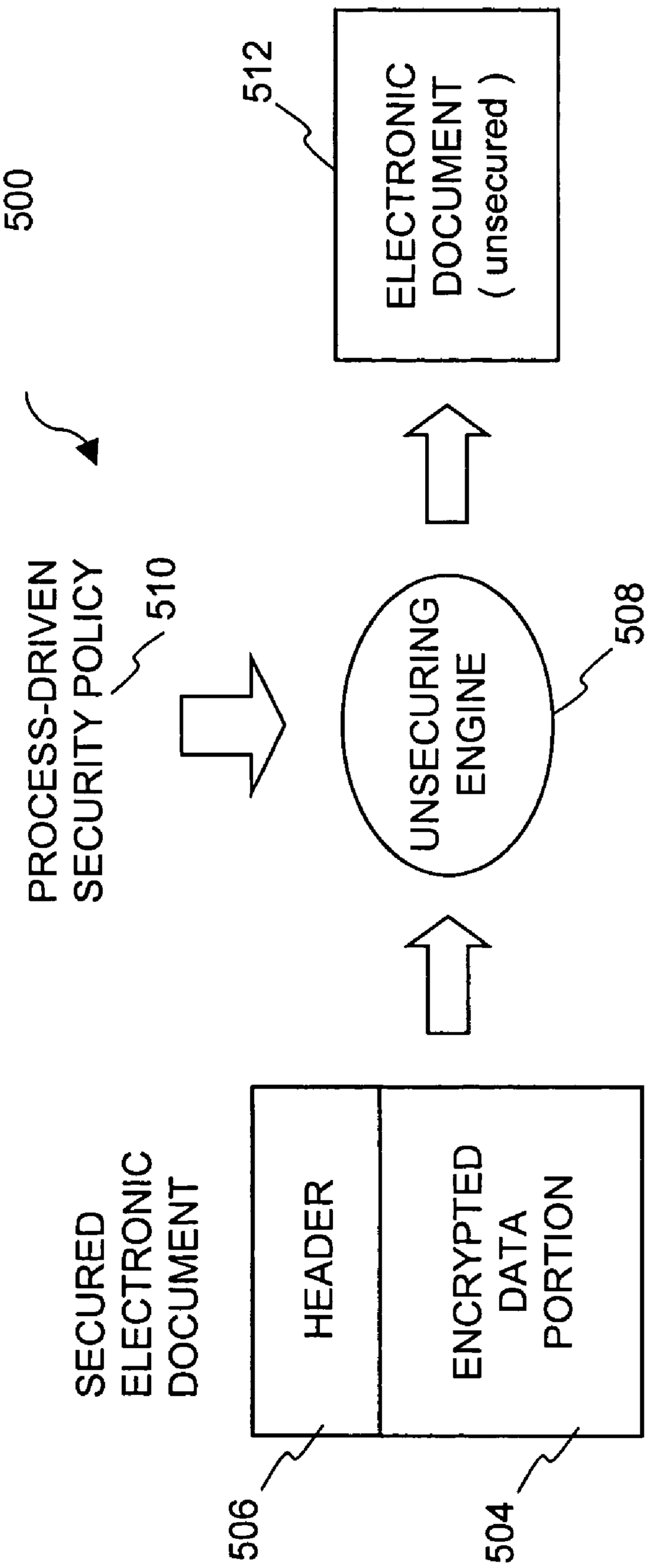

500
512
ELECTRONIC DOCUMENT (unsecured)
PROCESS-DRIVEN SECURITY POLICY
510
UNSECURING ENGINE
508
SECURED ELECTRONIC DOCUMENT
HEADER
ENCRYPTED DATA PORTION
506
504
502

FIG. 5A

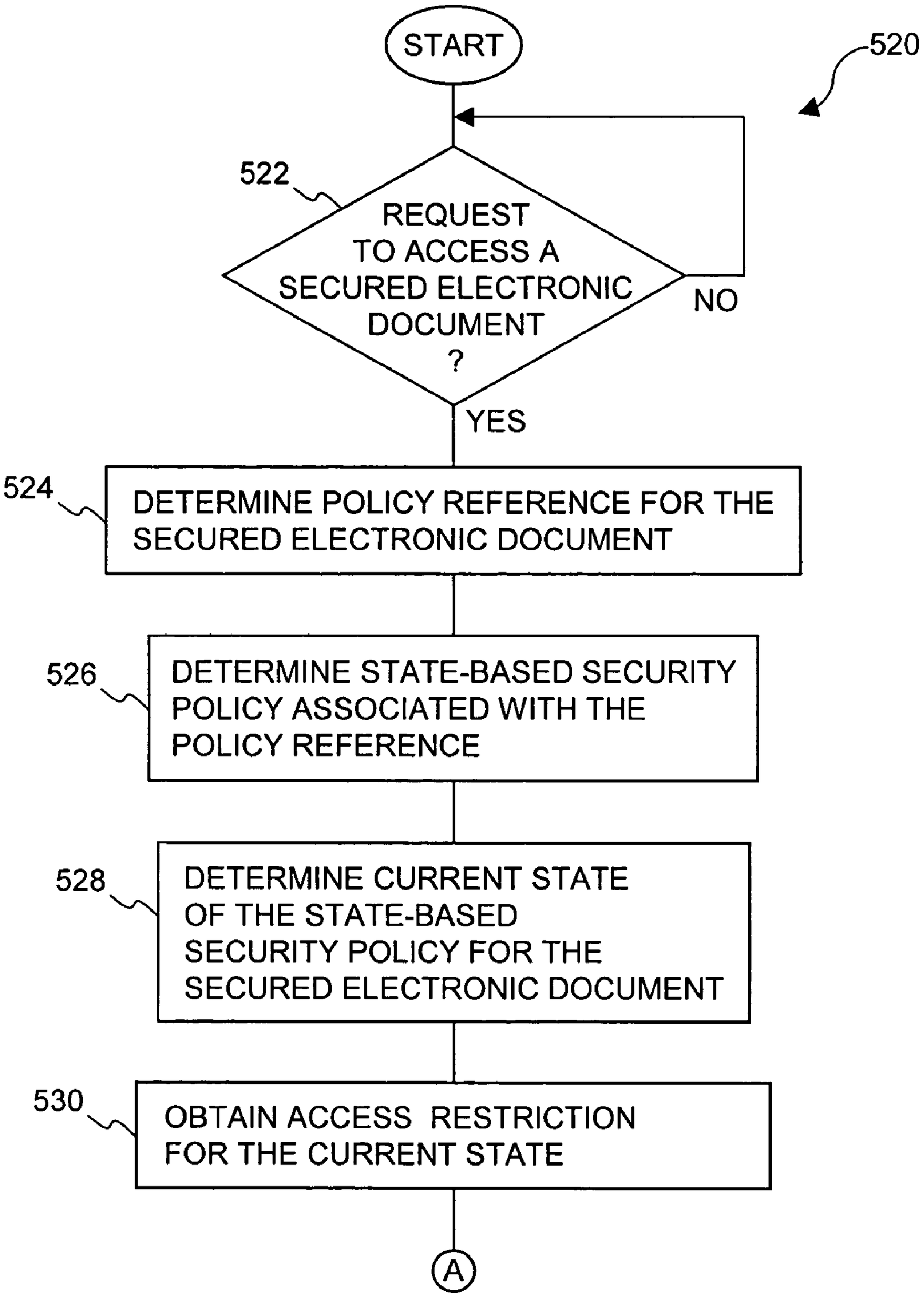

START
520
522
REQUEST TO ACCESS A SECURED ELECTRONIC DOCUMENT ?
NO
YES
524
DETERMINE POLICY REFERENCE FOR THE SECURED ELECTRONIC DOCUMENT
526
DETERMINE STATE-BASED SECURITY POLICY ASSOCIATED WITH THE POLICY REFERENCE
528
DETERMINE CURRENT STATE OF THE STATE-BASED SECURITY POLICY FOR THE SECURED ELECTRONIC DOCUMENT
530
OBTAIN ACCESS RESTRICTION FOR THE CURRENT STATE
A

FIG. 5B

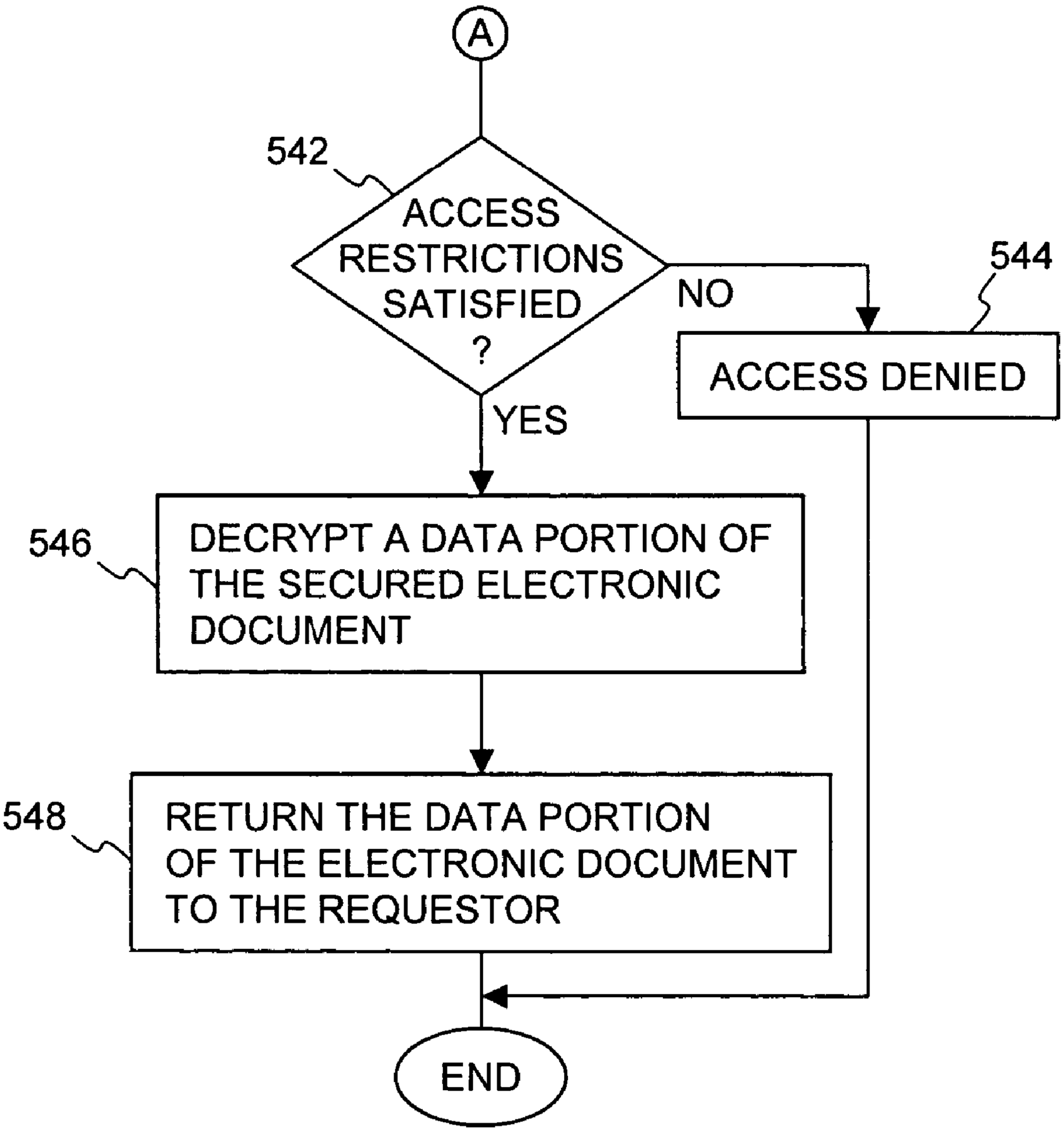

A
542
ACCESS RESTRICTIONS SATISFIED ?
NO
544
ACCESS DENIED
YES
546
DECRYPT A DATA PORTION OF THE SECURED ELECTRONIC DOCUMENT
548
RETURN THE DATA PORTION OF THE ELECTRONIC DOCUMENT TO THE REQUESTOR
END

FIG. 5C

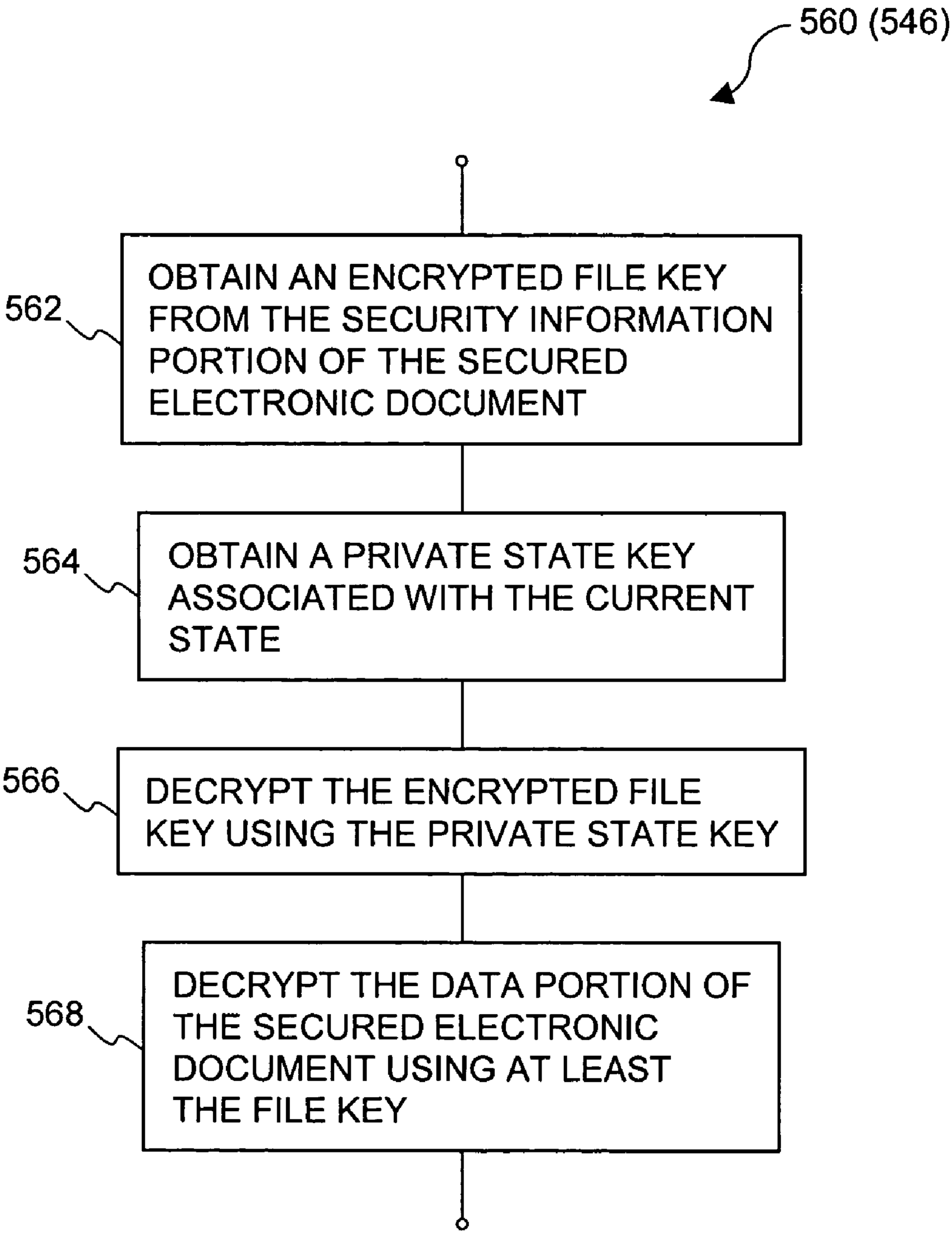

560 (546)
562
OBTAIN AN ENCRYPTED FILE KEY FROM THE SECURITY INFORMATION PORTION OF THE SECURED ELECTRONIC DOCUMENT
564
OBTAIN A PRIVATE STATE KEY ASSOCIATED WITH THE CURRENT STATE
566
DECRYPT THE ENCRYPTED FILE KEY USING THE PRIVATE STATE KEY
568
DECRYPT THE DATA PORTION OF THE SECURED ELECTRONIC DOCUMENT USING AT LEAST THE FILE KEY

FIG. 5D

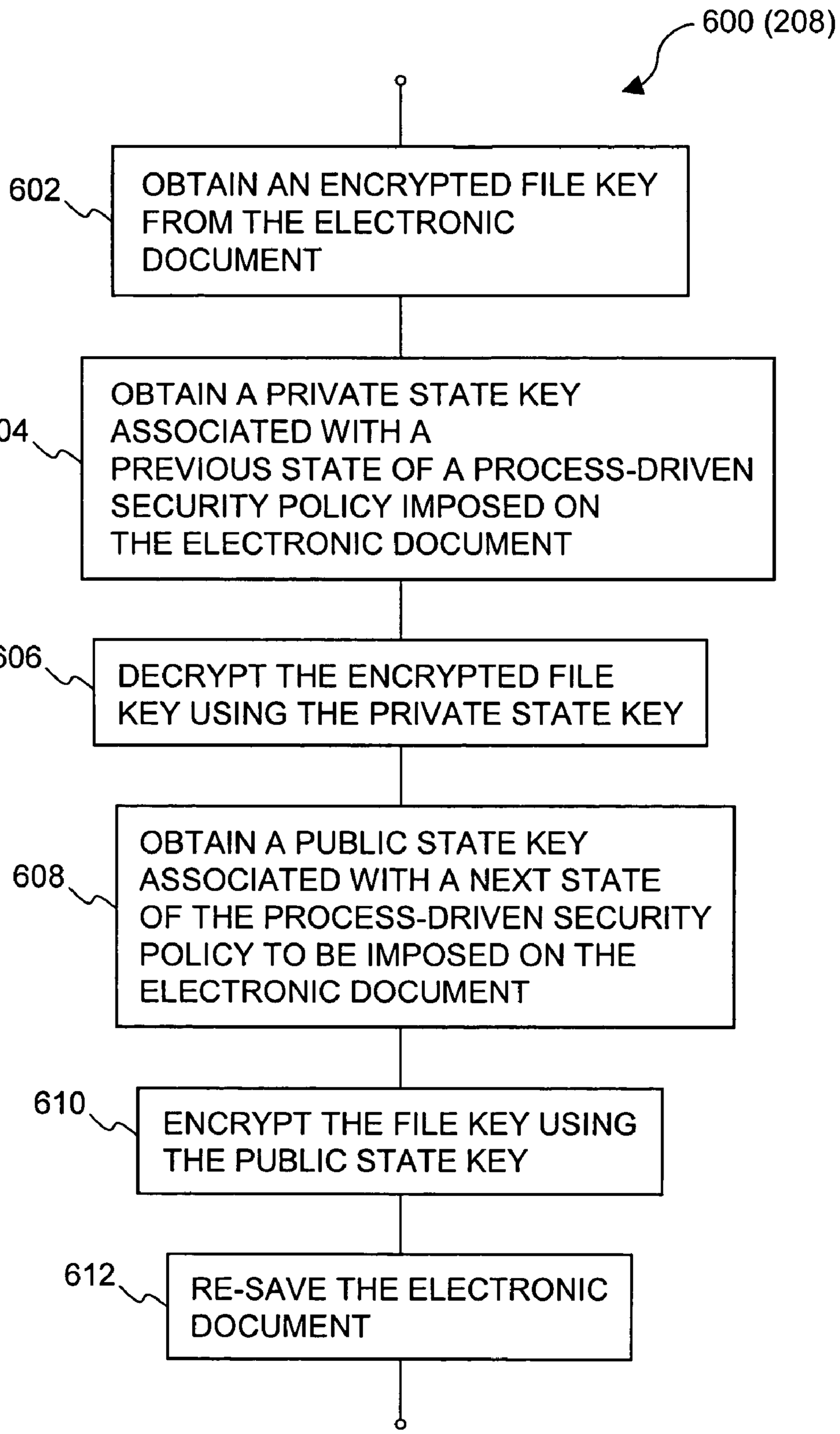

600 (208)
602
OBTAIN AN ENCRYPTED FILE KEY FROM THE ELECTRONIC DOCUMENT
604
OBTAIN A PRIVATE STATE KEY ASSOCIATED WITH A PREVIOUS STATE OF A PROCESS-DRIVEN SECURITY POLICY IMPOSED ON THE ELECTRONIC DOCUMENT
606
DECRYPT THE ENCRYPTED FILE KEY USING THE PRIVATE STATE KEY
608
OBTAIN A PUBLIC STATE KEY ASSOCIATED WITH A NEXT STATE OF THE PROCESS-DRIVEN SECURITY POLICY TO BE IMPOSED ON THE ELECTRONIC DOCUMENT
610
ENCRYPT THE FILE KEY USING THE PUBLIC STATE KEY
612
RE-SAVE THE ELECTRONIC DOCUMENT

FIG. 6

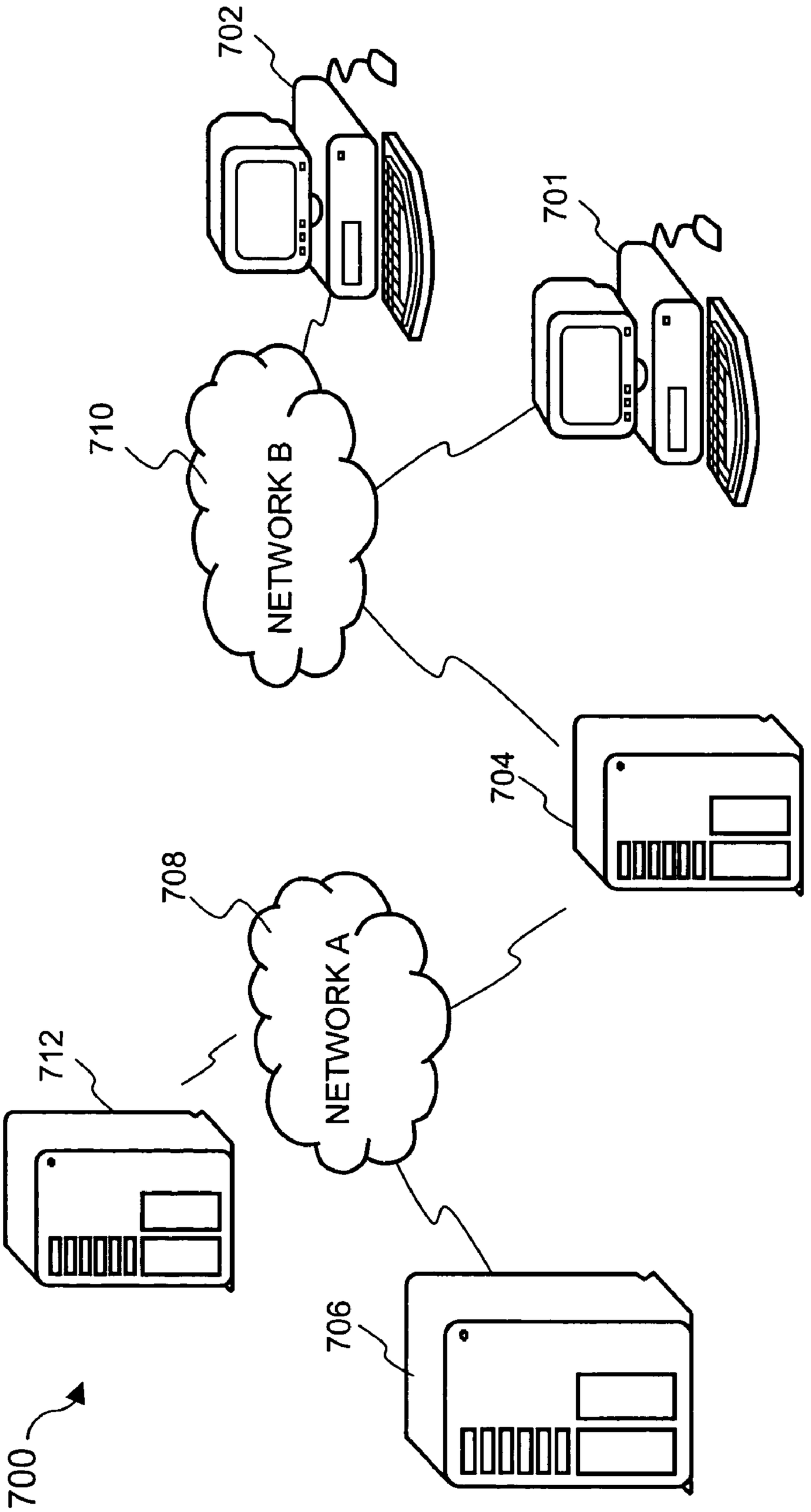


FIG. 7

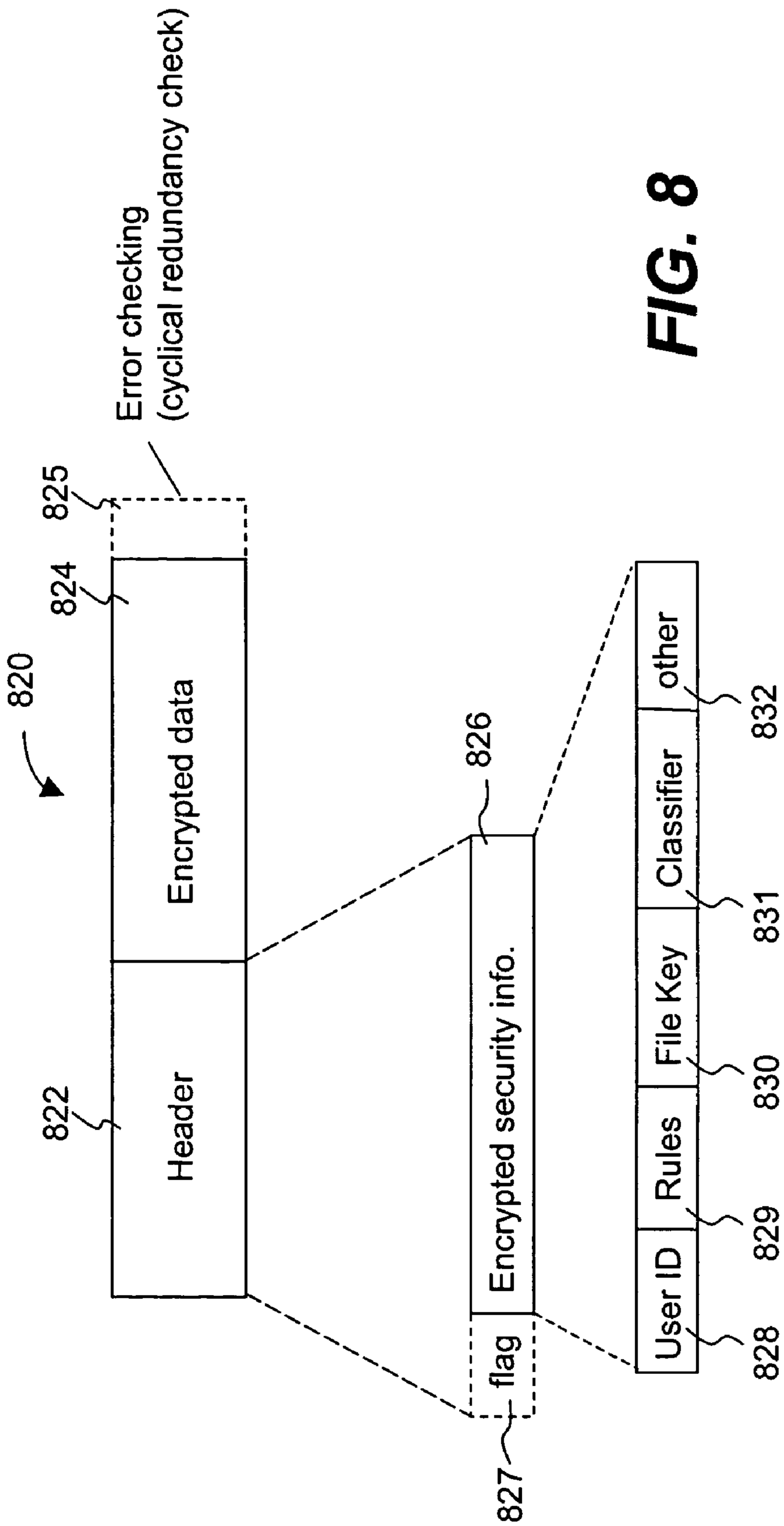

820
822
Header
824
Encrypted data
825
Error checking
(cyclical redundancy check)
826
Encrypted security info.
827
flag
828
User ID
829
Rules
830
File Key
831
Classifier
832
other

FIG. 8

METHOD AND SYSTEM FOR SECURING DIGITAL ASSETS USING PROCESS-DRIVEN SECURITY POLICIES

CROSS-REFERENCE TO RELATED APPLICATIONS

This application is related to: (i) U.S. patent application Ser. No.: 10/676,474, filed concurrently herewith, and entitled "METHOD AND APPARATUS FOR TRANSITIONING BETWEEN STATES OF SECURITY POLICIES USED TO SECURE ELECTRONIC DOCUMENTS," which is hereby incorporated herein by reference; (ii) U.S. patent application Ser. No.: 10/405,587, filed Apr. 1, 2003, and entitled "METHOD AND APPARATUS FOR SECURING DIGITAL ASSETS USING CONTENT TYPE DESIGNATIONS," which is hereby incorporated herein by reference; (iii) U.S. patent application Ser. No.: 10/159,537, filed May 5, 2002, and entitled "METHOD AND APPARATUS FOR SECURING DIGITAL ASSETS," which is hereby incorporated herein by reference; and (iv) U.S. patent application Ser. No.: 10/127,109, filed Apr. 22, 2002, and entitled "EVALUATION OF ACCESS RIGHTS TO SECURED DIGITAL ASSETS," which is hereby incorporated herein by reference.

BACKGROUND OF THE INVENTION

1. Field of the Invention

The present invention relates to security systems for data and, more particularly, to security systems that protect data in an inter/intra enterprise environment.

2. Description of Related Art

The Internet is the fastest growing telecommunications medium in history. This growth and the easy access it affords have significantly enhanced the opportunity to use advanced information technology for both the public and private sectors. It provides unprecedented opportunities for interaction and data sharing among businesses and individuals. However, the advantages provided by the Internet come with a significantly greater element of risk to the confidentiality and integrity of information. The Internet is an open, public and international network of interconnected computers and electronic devices. Without proper security means, an unauthorized person or machine may intercept information traveling across the Internet and even gain access to proprietary information stored in computers that interconnect to the Internet.

There are many efforts in progress aimed at protecting proprietary information traveling across the Internet and controlling access to computers carrying the proprietary information. Cryptography allows people to carry over the confidence found in the physical world to the electronic world, thus allowing people to do business electronically without worries of deceit and deception. Every day millions of people interact electronically, whether it is through e-mail, e-commerce (business conducted over the Internet), ATM machines, or cellular phones. The perpetual increase of information transmitted electronically has led to an increased reliance on cryptography.

One of the ongoing efforts in protecting the proprietary information traveling across the Internet is to use one or more cryptographic techniques to secure a private communication session between two communicating computers on the Internet. The cryptographic techniques provide a way to transmit information across an unsecure communication channel without disclosing the contents of the information to anyone eavesdropping on the communication channel. Using an encryption process in a cryptographic technique, one party can protect the contents of the data in transit from access by an unauthorized third party, yet the intended party can read the encrypted data after using a corresponding decryption process.

A firewall is another security measure that protects the resources of a private network from users of other networks. However, it has been reported that many unauthorized accesses to proprietary information occur from the inside, as opposed to from the outside. An example of someone gaining unauthorized access from the inside is when restricted or proprietary information is accessed by someone within an organization who is not supposed to do so. Due to the open nature of networks, contractual information, customer data, executive communications, product specifications, and a host of other confidential and proprietary intellectual property remain available and vulnerable to improper access and usage by unauthorized users within or outside a supposedly protected perimeter.

Many businesses and organizations have been looking for effective ways to protect their proprietary information. Typically, businesses and organizations have deployed firewalls, Virtual Private Networks (VPNs), and Intrusion Detection Systems (IDS) to provide protection. Unfortunately, these various security means have been proven insufficient to reliably protect proprietary information residing on private networks. For example, depending on passwords to access sensitive documents from within often causes security breaches when the password of a few characters long is leaked or detected. Consequently, various cryptographic means are deployed to provide restricted access to electronic data in security systems.

Various security criteria, such as encryption or decryption keys, are often used to facilitate restricted access to data in security systems. Conventional uses of security criteria provide static assignment of security criteria to electronic resources being secured. However, the assigning of security criteria in a static manner does not permit subsequent alteration of the security criteria under certain conditions. Although an administrator may be able to change the security criteria for an electronic resource that has already been secured, such alteration would be a manual process only available to the administrator. Further, given that an administrator is managing secure electronic resources (e.g., data) for many users, it is not feasible for the administrator to participate in the changing of security criteria for a large volume of electronic resources. Therefore, there is a need to provide more effective ways for security systems to permit security criteria imposed on electronic resources to be changed, thereby altering the security used to protect the electronic resources.

SUMMARY OF THE INVENTION

The invention relates to techniques for dynamically altering security criteria used in a file security system. The security criteria pertains to keys (or ciphers) used by the file security system to encrypt electronic files to be secured, or to decrypt electronic files already secured. The security criteria can, among other things, include keys that are required to gain access to electronic files. Here, the keys can be changed automatically as electronic files transition between different states of a process-driven security policy. The dynamic alteration of security criteria enhances the flexibility and robustness of the security system. In other words, access restrictions on electronic files can be dependent on the state of the process-driven security policy.

According to one aspect of the invention, methods and systems for securing electronic files use process-driven security policies. As an electronic file transitions through a process, access restrictions can automatically change. The process can be defined by a number of states, with each state having different security policies associated therewith. The security policies control which users are permitted to access the electronic files. In one embodiment, the access restrictions are imposed by one or more keys that are required to decrypt electronic files that were previously secured. The process can also be referred to as a workflow, where the workflow has a series of states through which files (documents) can move, where different security policies can be imposed at different states.

Another aspect of the invention is that process-driven security policies are controlled at a server of a file security system. A group of one or more electronic documents are bound together and progress together through states of a process specified by process-driven security policies. The server can automatically and remotely enforce the process-driven security policies on the group of electronic documents.

Still another aspect of the invention is that process-driven security policies are controlled at a client of a file security system. Here, each individual electronic document can be separately and independently bound to process-driven security policies. The process-driven security policies can thus operate at the client with little or no communication with a central server.

The process-driven security policies typically offer persistent states. Each state can specify a different set of users that are permitted access to an electronic document. The states are also independent of the electronic documents themselves.

The invention can be implemented in numerous ways, including as a method, system, device, and computer readable medium. Several embodiments of the invention are discussed below.

As a method for limiting access to electronic documents, one embodiment includes at least the acts of: creating a process-driven security policy having a plurality of states, with each of the states having a different set of access restrictions; associating an identifier to the process-driven security policy; and making the identifier available to certain of users or groups of users.

As a method for imposing access restrictions on electronic documents, one embodiment includes at least the acts of: providing at least one process-driven security policy from a server machine to a client machine, the process-driven security policy having a plurality of states associated therewith; and associating the electronic document with at least one of the states of the process-driven security policy to impose access restrictions on an electronic document, the access restrictions being dependent on the at least one of the states of the process-driven security policy.

As a computer readable medium including at least computer program code for imposing access restrictions on electronic documents, one embodiment of the invention includes at least: computer program code for providing at least one process-driven security policy from a server machine to a client machine, the process-driven security policy having a plurality of states associated therewith; and computer program code for associating the electronic document with at least one of the states of the process-driven security policy to impose access restrictions on an electronic document, the access restrictions being dependent on the at least one of the states of the process-driven security policy.

Other objects, features, and advantages of the present invention will become apparent upon examining the following detailed description of an embodiment thereof, taken in conjunction with the attached drawings.

BRIEF DESCRIPTION OF THE DRAWINGS

These and other features, aspects, and advantages of the invention will become better understood with regard to the following description, appended claims and accompanying drawings, wherein:

FIG. **1** is a diagram of an exemplary process-driven security policy (PDSP) according to one embodiment of the invention.

FIG. **2** is a flow diagram of a transition process according to one embodiment of the invention.

FIG. **3** illustrates a security policy state machine according to one embodiment of the invention.

FIG. **4**A is a diagram of a document securing system according to one embodiment of the invention.

FIG. **4**B is a flow diagram of a document securing process according to one embodiment of the invention.

FIG. **4**C is a detailed flow diagram of an encryption process according to one embodiment of the invention.

FIG. **5**A is a diagram of a document unsecuring system according to one embodiment of the invention.

FIGS. **5**B and **5**C are flow diagrams of a document access process according to one embodiment of the invention.

FIG. **5**D is a flow diagram of a decryption process according to one embodiment of the invention.

FIG. **6** is a flow diagram of a transition process according to one embodiment of the invention.

FIG. **7** shows a basic security system in which the invention may be practiced in accordance with one embodiment thereof.

FIG. **8** shows an exemplary data structure of a secured file that may be used in one embodiment of the invention.

DETAILED DESCRIPTION OF THE INVENTION

The invention relates to techniques for dynamically altering security criteria used in a file security system. The security criteria pertains to keys (or ciphers) used by the file security system to encrypt electronic files to be secured, or to decrypt electronic files already secured. The security criteria can, among other things, include keys that are required to gain access to electronic files. Here, the keys can be changed automatically as electronic files transition between different states of a process-driven security policy. The dynamic alteration of security criteria enhances the flexibility and robustness of the security system. In other words, access restrictions on electronic files can be dependent on the state of the process-driven security policy.

According to one aspect of the invention, methods and systems for securing electronic files use process-driven security policies. As an electronic file transitions through a process, access restrictions can automatically change. The process can be defined by a number of states, with each state having different security policies associated therewith. The security policies control which users are permitted to access the electronic files. In one embodiment, the access restrictions are imposed by one or more keys that are required to decrypt electronic files that were previously secured. The process can also be referred to as a workflow, where the workflow has a series of states through which files (documents) can move, where different security policies can be imposed at different states.

Another aspect of the invention is that process-driven security policies are controlled at a server of a file security system.

A group of one or more electronic documents are bound together and progress together through states of a process specified by process-driven security policies. The server can automatically and remotely enforce the process-driven security policies on the group of electronic documents.

Still another aspect of the invention is that process-driven security policies are controlled at a client of a file security system. Here, each individual electronic document can be separately and independently bound to process-driven security policies. The process-driven security policies can thus operate at the client with little or no communication with a central server.

The process-driven security policies typically offer persistent states. Each state can specify a different set of users that are permitted access to an electronic document. The states are also independent of the electronic documents themselves.

Secured files are files that require one or more keys, passwords, access privileges, etc. to gain access to their content. The security is often provided through encryption and access rules. The files, for example, can pertain to documents, multimedia files, data, executable code, images and text. In general, a secured file can only be accessed by authenticated users with appropriate access rights or privileges. In one embodiment, each secured file is provided with a header portion and a data portion, where the header portion contains, or points to, security information. The security information is used to determine whether access to associated data portions of secured files is permitted.

In one embodiment, security information provided with an electronic document controls restrictive access to a data portion which is encrypted. The security information can employ access rules together with cipher keys (e.g., a file key and various other keys) to ensure that only those users with proper access privileges or rights can access the encrypted data portion.

As used herein, a user may mean a human user, a software agent, a group of users, a member of the group, a device and/or application. Besides a human user who needs to access a secured document, a software application or agent sometimes needs to access secured files in order to proceed. Accordingly, unless specifically stated, the "user" as used herein does not necessarily pertain to a human being.

The invention is related to processes, systems, architectures and software products for providing pervasive security to digital assets (e.g., electronic documents). The invention is particularly suitable in an enterprise environment. In general, pervasive security means that digital assets are secured (i.e., secured data) and can only be accessed by authenticated users with appropriate access rights or privileges. Digital assets may include, but not be limited to, various types of documents, multimedia files, data, executable code, images and texts.

In the following description, numerous specific details are set forth in order to provide a thorough understanding of the invention. However, it will become obvious to those skilled in the art that the invention may be practiced without these specific details. The description and representation herein are the common meanings used by those experienced or skilled in the art to most effectively convey the substance of their work to others skilled in the art. In other instances, well-known methods, procedures, components, and circuitry have not been described in detail to avoid unnecessarily obscuring aspects of the invention.

Reference herein to "one embodiment" or "an embodiment" means that a particular feature, structure, or characteristic described in connection with the embodiment can be included in at least one embodiment of the invention. The appearances of the phrase "in one embodiment" in various places in the specification are not necessarily all referring to the same embodiment, nor are separate or alternative embodiments mutually exclusive of other embodiments. Further, the order of blocks in process flowcharts or diagrams representing one or more embodiments of the invention do not inherently indicate any particular order, nor imply any limitations in the invention.

Embodiments of the invention are discussed herein with reference to FIGS. **1-8**. However, those skilled in the art will readily appreciate that the detailed description given herein with respect to these figures is for explanatory purposes as the invention extends beyond these limited embodiments.

FIG. **1** is a diagram of an exemplary process-driven security policy (PDSP) **100** according to one embodiment of the invention. The process-driven security policy **100** includes a plurality of different states. As shown in FIG. **1**, the process-driven security policy **100** can include state A **102**, state B **104**, state C **106**, and state D **108**. Each of these different states can be associated with one or more access restrictions.

The process-driven security policy **100** is used by a file (document) security system to restrict access to files (documents). As files are placed in different ones of the states of the process-driven security policy **100**, the access restrictions being utilized to secure access to the files typically changes. More particularly, as the files move from state-to-state in accordance with a process, the access restrictions utilized by the files often changes. Indeed, the access restrictions can change automatically based upon the state the file is in, and thus does not necessarily require user or administrator interaction in order to change the access restrictions. Typically, access restrictions will designate which users (or groups of users) are able to access secure documents, whether certain clearance levels are needed, whether off-line access is permitted, and which of various other possible criteria or considerations are utilized. A set of access restrictions for the various states can be referred to as a security policy.

A file can transition between the various states of the process-driven security policy **100** in a controlled manner. Often, the process-driven security policy **100** defines the transitions that are permissible. In one embodiment, the state transitions are event-driven. The events can be either internal to the file security system or external to the file security system. When event-driven, the transitions between states can be automatic and thus do not require user or administrator interaction. However, some events can be triggered or initiated by user or administrator interaction.

As shown in FIG. **1**, a file (document) can transition between the different states **102-108** offered by the process-driven security policy **100**. For example, a file currently in state A **102** can transition to state B **104** or state D **108**, depending upon process-related conditions (e.g., events). Similarly, a file in state D **108**, depending upon process considerations, can transition to state A **102**, state B **104** or state C **106**. Likewise, a file in state B **104** or state C **106** can transition to one or more other states. Additional details on states, security policies and transitions between states are discussed in additional detail below.

FIG. **2** is a flow diagram of a transition process **200** according to one embodiment of the invention. The transition process **200** can be used to transition a document (file) between different states of a process-driven security policy, such as the process-driven security policy shown in FIG. **1**. The transition process **200** is typically deemed process-driven because it is driven by a process. The process is, for example, defined by transition rules. The transition rules typically rely upon events to cause transitions between states. Often user or administra-

tor interaction is not involved when activating transitions. However, the transition process **200** can permit a user or administrator to participate in activating transitions, such as by causing an event to occur which initiates a transition.

The transition process **200** begins with a decision **202** that determines whether an event relevant to the process-driven security policy imposed on a document has been received. Typically, the process-driven security policy is imposed on the document by a file security system. One implementation of a process-driven security policy is a security policy state machine. The process-driven security policy (or security policy state machine) has a plurality of states, and transition rules for transitioning between the various states. In any case, the transition process **200** monitors events that are relevant to the process-driven security policy (or the security policy state machine). When the decision **202** determines that an event has not yet been received, the transition process **200** awaits such an event.

On the other hand, when the decision **202** determines that an event has been received, then the transition process **200** determines **204** whether the event causes a state transition. Here, the rules by which transitions between states occur, i.e., transition rules, can be specified by the process-driven security policy. For example, an administrator for the document security system may have created the process-driven security policy and thus defined its states and its transition rules. Hence, when an event is received, it is evaluated to determine **204** whether the event causes a state transition. When the decision **206** determines that a state transition is to occur, the state transition is performed **208** to transfer one or more documents from one state to another state. Alternatively, when the decision **206** determines that a state transition is not to occur, the block **208** is bypassed so that no state transition is performed. Once the one or more documents transition to the new state, the access restrictions for the new state govern when access to the documents, which are secured, is permitted. Following the block **208** or its being bypassed, the transition process **200** is complete and ends.

FIG. **3** illustrates a security policy state machine **300** according to one embodiment of the invention. As previously noted, a security policy state machine is one implementation of a process-driven security policy. In this exemplary embodiment, the security policy state machine **300** includes four distinct states, namely, a state A ("Draft") **302**, state B ("Final Draft") **304**, state C ("Retain") **306**, and state D ("Delete") **308**. Each of these states has one or more associated access restriction for documents (files) which reside in that state. Further, the permitted transitions between the various states **302-308** are identified by transitions T**1**-T**5**. In particular, a document in the Draft state **302** can follow the transition T**1** to the Final Draft state **304**. A document in the Final Draft state **304** can follow the transition T**2** to the Retain state **306**. A document in the Retain state **306** can follow transition T**3** to the Delete state **308**. Further, a document in the Final Draft state **304** can follow transition T**4** to the Draft state **302**, and a document in the Retain state **306** can follow transition T**5** to the Final Draft state **304**.

A file security system can enforce the security policy state machine **300** on one or more electronic documents. In doing so, the security policy state machine **300** is typically described in a textual manner, such as in a markup language (e.g., XML), pseudo-code, and the like. One representative example of a textual description of the security policy state machine **300** is as follows.

State=DRAFT
Accessors=Finance, unrestricted
Deny off-line access
Grant audit access
State=FINAL DRAFT
Accessors=Finance, restricted; Finance Managers, unrestricted
Deny off-line access
Grant audit access
State=RETAIN
Accessors=All
Allow off-line access
Deny audit access
State=DELETE
Accessors=None

Note that in the Draft state, the users with permission to access the electronic document (referred to as "Accessors") include those users that are members of a Finance group. The access is also unrestricted in this Draft state. Also, in the Draft state, offline access to the electronic document is not permitted, but audit access is permitted. Note, however, in the Final Draft state, those users that are members of the Finance group now only have restricted access. In one embodiment, restricted access means that the data (content) of the document can be accessed but that such data cannot be further disseminated through operations such as cut, paste, print, etc.

Additionally, the security policy state machine **300** transitions between the various states in accordance with transition rules. Typically, the transition rules are triggered by the occurrence of events. The events can be internal or external. The external events can originate from users or from another system (e.g., a document management system). In a specific case of the security policy state machine **300**, a representative description of a transition rule is as follows.

On event ( ), transition from STATE1 to STATE2

Some exemplary transition rules using internal or external events are as follows.

On (time=Sep. 1, 2008), RETAIN to DELETE
On (ExtEvent==docCheckIn), FINAL DRAFT to RETAIN
On (ExtEvent==docFinalize), DRAFT to FINAL DRAFT
On (ExtEvent==docReject), FINAL DRAFT to DRAFT
On (period=event transition day (FINAL DRAFT)+90 days), FINAL DRAFT to RETAIN

Of these exemplary transition rules, the first and last transition rules are triggered by internal events and the others are triggered by external events. For example, the external events can be from a document management system that is separate from the file (document) security system.

FIG. **4**A is a diagram of a document securing system **400** according to one embodiment of the invention. The document securing system **400** is, for example, performed by a computing device, such as client computer **701** or **702** shown in FIG. **7** below.

The document securing system **400** creates or obtains an electronic document **402** that is to be secured. The electronic document **402** is then supplied to a securing engine **404**. The securing engine **404** receives a designation of a classifier **406** to be associated with the electronic document **402**. The classifier **406** refers to an accessor user list, and possibly other forms of access restriction. In one embodiment, the classifier **406** can be a label to a categorization of the electronic document with respect to a plurality of different types of content. Examples of classifiers include: External, Financial, Sales Forecast, Sales Quota, Press Release, Budget, Marketing Presentation, Marketing Planning, Engineering Planning, Engineering Project X, Engineering Specification, and Engineering Design. In addition, the securing engine **404** can receive a process-driven security policy **407** to be used to secure the electronic document **402**. In one embodiment, the process-

driven security policy **407** is chosen from a plurality of process-driven security policies based on the classifier **406**. In another embodiment, the process-driven security policy **407** is made up of states, and each of the states correspond to one of the classifiers **406**.

The securing engine **404** operates to produce a secured electronic document **408**. The secured electronic document **408** includes an encrypted data portion **410** and a header portion **412**. The encrypted data portion **410** is the electronic document **402** after having been encrypted. The encryption can result from the use of one or more keys and encryption algorithms. For stronger security, a hierarchy of encryption may be used. The header portion **412** is also referred to as encrypted security information, because the header portion **412** includes the encrypted security information as at least a substantial component of the header portion **412**. The encrypted security information can include a classifier, access rules and at least one key (e.g., file key, private state key). The access rules and the keys utilized to encrypt the electronic document **402** depend on the state of the associated process-driven security policy **407** which is indicated by the classifier. Initially, the electronic document **402** is encrypted in accordance with an initial state of the process-driven security policy **407**. Typically, one of the states of the process-driven security policy **407** is designated as its initial state.

Hence, if the encrypted security information is able to be decrypted, the file key is able to be retrieved from the header portion **412** and used to decrypt the encrypted data portion **410** of the secured electronic document **408**, as will be discussed in more detail below with respect to FIG. **5**C. However, the encrypted security information in the header portion **412** is often secured through one or multiple layers of encryption, which can use various keys. These various keys are used to encrypt the security information. Typically, these various keys are managed by a server, but made available to client computers so that decryption can be performed locally. In one implementation, the encrypted security information within the header portion **412** can be decrypted if, and only if, the decrypting party has possession of both of the following: a group key (a private key for a group specified in the header), and a state key (a private key for the classifier specified in the header). As previously noted, the classifier is used to determine the state of the process-driven security policy **407**.

Additional details on securing files or documents is provided in U.S. patent application Ser. No.: 10/159,537, filed May 5, 2002, and entitled "METHOD AND APPARATUS FOR SECURING DIGITAL ASSETS," which is hereby incorporated by reference.

FIG. **4**B is a flow diagram of a document securing process **440** according to one embodiment of the invention. The document securing process **440** represents processing performed by a document securing system, such as the document securing system **400** illustrated in FIG. **4**A.

The document securing process **440** initially opens or creates **442** an electronic document. Next, a decision **444** determines whether the electronic document is to be secured. When the decision **444** determines that the electronic document is not to be secured, then the electronic document is saved **446** in the normal course. Here, the electronic document is not secured but simply stored in a conventional fashion.

On the other hand, when the decision **444** determines that the electronic document is to be secured, then an initial policy reference for the electronic document is assigned **448**. In one implementation, the policy reference is a pointer to an accessor user list. A classifier for an electronic document can be assigned in a variety of different ways. In one implementation, a user or creator of the electronic document is able to assign the classifier. For example, the user or creator of the electronic document might interact with a graphical user interface to select a classifier from a list of available classifiers.

After the policy reference is assigned **448**, the electronic document is secured **450** in accordance with a process-driven security policy associated with the policy reference. Here, the electronic document is typically secured in accordance with the initial state of the process-driven security policy. Thereafter, the secured electronic document is saved **452**. Following the operations **452** and **446**, the document securing process **440** is complete and ends. The subsequent transitions to other states of the process-driven security policy is discussed below with reference to FIG. **6**.

FIG. **4**C is a detailed flow diagram of an encryption process **460** according to one embodiment of the invention. The encryption process **460** is, for example, processing suitable for being performed by the block **450** shown in FIG. **4**B in which an electronic document is secured in accordance with a process-driven security policy.

According to the encryption process **460**, a file key is obtained **462**. In one implementation, the file key is a symmetric key used to encrypt and decrypt a data portion of a secured document. After the file key is obtained **462**, the data portion of the electronic document is then encrypted **464** using at least the file key.

In one embodiment, each of the different states of the process-driven security policy would include a different public state key that would be used to encrypt documents being placed into such state. An initial state of the process-driven security policy associated with the policy reference is then determined **466**. Next, a public state key associated with the initial state is obtained **468**. Typically, the public state key is a public key of a public and private cryptography key pair that is to be utilized to encrypt documents associated with the initial state of the process-driven security policy. Once the public state key associated with the initial state has been obtained **468**, the file key is encrypted **470** using the public state key. Thereafter, security information is attached **472** to the encrypted data portion. The security information, for example, can include the policy reference and the encrypted file key. For example, the policy reference can be used as a state indicator to identify the applicable state of the process-driven security policy.

In one embodiment, the policy reference has a key pair associated therewith. The file (document) security system (e.g., server) maintains the current state of the process-driven security policy associated with the policy reference. The public key in this pair is used to encrypt the document and bind it with the process-driven security policy.

In this implementation, the electronic document has at least a data portion and a security information portion. The data portion is encrypted using at least the file key. In one embodiment, the electronic document can be encrypted many times over such that a plurality of different keys are needed to encrypt (and consequently to decrypt) the electronic document. In another embodiment, a key used to encrypt the electronic document can be encrypted many times over after being used to encrypt the electronic document. In other words, although the document securing process **440** refers to encryption of the data portion through use of the file key and then encryption of the file key through use of the public state key, it should be understood that additional keys can be used to directly encrypt the electronic document, or indirectly encrypt the electronic document by encrypting a key used to encrypt the electronic document. For example, the additional

keys might include one or more of a classifier key, a user or group key, or a security clearance level key.

The security information is typically provided in a header (or header portion) of the electronic document. The header is thus typically attached to the encrypted data portion. The header together with the encrypted data portion represents a secured electronic document. Typically, the security information would include access rules, a policy reference (classifier), a private state key and at least one key (e.g., file key). The at least one key can be encrypted by a public state key that corresponds to the state, as well as possibly one or more other keys. The at least one key is often secured by encrypting either the at least one key itself, or the security information more generally, through use of one or more various other keys (e.g., group key, content type key, and/or clearance key).

FIG. **5**A is a diagram of a document unsecuring system **500** according to one embodiment of the invention. The document unsecuring system **500** represents a counterpart to the document securing system **400** illustrated in FIG. **4**A.

The document unsecuring system **500** cooperates to receive a secured electronic document **502**. The secured electronic document typically includes an encrypted data portion **504** and a header **506**. Often, but not necessarily, the header **506** is encrypted. The header **506** includes a policy reference and at least one key, e.g., a file key, that is needed to decrypt the encrypted data portion **504**. The secured electronic document **502** is supplied to an unsecuring engine **508**. The unsecuring engine **508** examines the header **506** of the secured electronic document **502** to determine the policy reference. The policy reference identifies a process-based security policy **510**, or a state thereof, that governs the security of the secured document **502**. The unsecuring engine **508** also receives at least that portion of the process-based security policy that pertains to the state of the secured electronic document **502**. In other words, the unsecuring engine **508** needs the access restrictions for the current state of the process-driven security policy **510** to unsecure the secured electronic document **502**, and thus gain access to its contents. The unsecuring engine **508** then evaluates whether the secured electronic document **502** is permitted to be accessed by the requestor, based on the access restrictions so retrieved. When the unsecuring engine **508** determines that the requester is authorized to access the secured electronic document **502**, then the unsecuring engine **508** can decrypt the encrypted data portion **504** of the secured electronic document **502** (and also eliminate at least significant portions of the header **506**) to yield an electronic document **512** that is unsecured. In other words, the electronic document **512** is primarily (or exclusively) composed of the data portion of the encrypted data portion **504** after such has been decrypted. The decryption can involve the use of a number of keys (e.g., private keys) and decryption algorithms, one of such keys is the file key of the secured electronic document, and another of such keys is the private state key for the state of the secured electronic document.

FIGS. **5**B and **5**C are flow diagrams of a document access process **520** according to one embodiment of the invention. The document access process **520** operates to determine whether access to a particular document is permitted to a particular user (or group of users). The document access process **520** begins with a decision **522** that determines whether a request to access a secured electronic document has been received. When the decision **522** determines that such a request has not yet been received, the document access process **520** awaits such a request. Once the decision **522** determines that a request to access a secured electronic document has been received, the document access process **520** continues. In other words, the document access process **520** can be considered to be invoked once a request to access a secured electronic document has been received.

In any case, once a request to access a secured electronic document has been received, a policy reference for the secured electronic document to be accessed is determined **524**. In one embodiment, the security information portion of a secured electronic document contains the policy reference. Next, a process-driven security policy associated with the policy reference is determined **526**. Then, the current state of the process-driven security policy for the secured electronic document is determined **528**. In one embodiment, the policy reference (or other indicator) can indicate the current state of the state-based security policy. Next, access restriction are obtained **530** for the current state. Each of the different states of the process-driven security policy often has a different access restriction. Here, the state policy restrictions are those restrictions associated with the current state of a process-driven security policy.

Thereafter, a decision **542** determines whether the state policy restrictions are satisfied. In other words, the secured electronic document to be accessed is presently in the current state of the process-driven security policy. This current state has the access restriction associated therewith, that must be satisfied in order to gain access to the secured electronic document. Hence, the decision **542** determines whether the access restriction is satisfied by the requester (e.g., user or group of users) seeking access to the secured electronic document. When the decision **542** determines that the access restriction is not satisfied, access to the secured electronic document is denied **544**.

On the other hand, when the decision **542** determines that the access restriction has been satisfied, then a data portion of the secured electronic document is decrypted **546**. Then, the data portion of the electronic document is returned **548** to the requestor. Following the block **548**, as well as following the block **544**, the document access process **520** ends.

FIG. **5**D is a flow diagram of a decryption process **560** according to one embodiment of the invention. The decryption process **560** can, for example, pertain to detailed operations performed by the block **546** illustrated in FIG. **5**C. In any event, the decryption process **560** initially obtains **562** an encrypted file key from the security information portion of the secured electronic document. In addition, a private state key associated with the current state of the process-driven security policy for the secured electronic document is obtained **564**. Normally, only authorized users would be able to gain access to the private state key. The private state key is the private key of the same public and private cryptography key pair that provided the public state key that was used to encrypt the file key. Then, the encrypted file key is decrypted **566** using the private state key. Thereafter, the data portion of the secured electronic document is decrypted **568** using at least the file key. Consequently, the data portion of the secured electronic document is decrypted and is in the "clear" and thus usable by the requester. Following the block **568**, the decryption process **560** is complete and ends.

FIG. **6** is a flow diagram of a transition process **600** according to one embodiment of the invention. The transition process **600** pertains to processing that can be utilized to transition between states of a process-driven security policy. More particularly, the transition process **600** is, for example, suitable for use as the processing performed by the block **208** illustrated in FIG. **2**.

The transition process **600** initially obtains **602** an encrypted file key from the electronic document. Typically, the encrypted file key would be retrieved from the security

information portion of the electronic document. Then, a private state key is obtained **604**. Here, the private state key is associated with a previous state of a process-driven security policy that is imposed on the electronic document. After the private state key has been obtained **604**, the encrypted file key is decrypted **606** using the private state key. At this point, the file key has been decrypted and could be used to decrypt the data portion of the electronic document. However, the file key is instead re-encrypted in accordance with a next (current) state. More specifically, a public state key is then obtained **608**. The public state key is associated with the next state of the state-based security policy that is to be imposed on the electronic document. Then, using the public state key, the file key can be encrypted **610**. Thereafter, the electronic document is re-saved **612**. By re-saving **612** the electronic document, the security information portion of the electronic document is updated to include the new encrypted file key in accordance with the next state (or current state). Note that the data portion of the electronic document (which is secured by the file key) advantageously need not be decrypted in the transition process **600**; instead, the encryption of the file key is changed whenever a state transition occurs. Following the block **612**, the transition process **600** is complete.

In one embodiment, to effect a state transition, the user only needs permission to effect the state transition. Additionally, users authorized to effect state changes with respect to a document, might be quite different from users authorized to access the document.

FIG. **7** shows a basic security system **700** in which the invention may be practiced in accordance with one embodiment thereof. The security system **700** may be employed in an enterprise or inter-enterprise environment. It includes a first server **706** (also referred to as a central server) providing centralized access management for the enterprise. The first server **706** can control restrictive access to files secured by the security system **700**. To provide dependability, reliability and scalability of the system, one or more second servers **704** (also referred to as local servers, of which one is shown) may be employed to provide backup or distributed access management for users or client machines serviced locally. The server **704** is coupled to a network **708** and a network **710**. For illustration purposes, there are two client machines **701** and **702** being serviced by the local server **704**. Alternatively, one of the client machines **701** and **702** may be considered as a networked storage device.

Secured files may be stored in any one of the devices **701**, **702**, **704** and **706**. When a user of the client machine **701** attempts to exchange a secured file with a remote destination **712** being used by an external user, one or more of the processing **300**, **400**, **500** and **600** discussed above are activated to ensure that the requested secure file is delivered without compromising the security imposed on the secured file.

According to one embodiment, a created document is caused to go through an encryption process that is preferably transparent to a user. In other words, the created document is encrypted or decrypted under the authoring application so that the user is not aware of the process. One or more keys, such as a state key, a user key and/or a content type key, can be used to retrieve a file key to decrypt an encrypted document. Typically, the user key is associated with an access privilege for the user or a group of users, and the content type key is associated with the type of content of the created document. For a given secured document, only a user with proper access privileges can access the secured document.

In one setting, a secured document may be uploaded via the network **710** from the client computer **701** to a computing or storage device **702** that may serve as a central repository. Although not necessary, the network **710** can provide a private link between the computer **701** and the computing or storage device **702**. Such link may be provided by an internal network in an enterprise or a secured communication protocol (e.g., VPN and HTTPS) over a public network (e.g., the Internet). Alternatively, such link may simply be provided by a TCP/IP link. As such, secured documents on the computer **702** may be remotely accessed.

In another setting, the computer **701** and the computing or storage device **702** are inseparable, in which case the computing or storage device **702** may be a local store to retain secured documents or receive secured network resources (e.g., dynamic Web contents, results of a database query, or a live multimedia feed). Regardless of where the secured documents or secured resources are actually located, a user, with proper access privileges, can access the secured documents or resources from the client computer **701** or the computing or storage device **702** using an application (e.g., Microsoft Internet Explorer, Microsoft Word or Adobe Acrobat Reader).

Accordingly, respective local modules in local servers, in coordination with the central server, form a distributed mechanism to provide distributed access control enforcement. Such distributed access control enforcement ensures the dependability, reliability and scalability of centralized access control management undertaken by the central server for an entire enterprise or a business location.

FIG. **8** shows an exemplary data structure **820** of a secured file that may be used in one embodiment of the invention. The data structure **820** includes two portions: a header (or header portion) **822** and encrypted data (or an encrypted data portion) **824**. The header **822** can be generated in accordance with a security template associated with a data store and thus provides restrictive access to the data portion **824** which is an encrypted version of a plain file. Optionally, the data structure **820** may also include an error-checking portion **825** that stores one or more error-checking codes, for example, a separate error-checking code for each block of encrypted data **824**. These error-checking codes may also be associated with a Cyclical Redundancy Check (CRC) for the header **822** and/or the encrypted data **824**. The header **822** includes a flag bit or signature **827** and security information **826** that is in accordance with the security template for the store. According to one embodiment, the security information **826** is encrypted and can be decrypted with a user key associated with an authenticated user (or requestor).

The security information **826** can vary depending upon implementation. However, as shown in FIG. **8**, the security information **826** includes a user identifier (ID) **828**, access policy (access rules) **829**, a file key **830**, a classifier **831** and other information **832**. Although multiple user identifiers may be used, a user identifier **828** is used to identify a user or a group that is permitted to access the secured file. The access rules **829** provide restrictive access to the encrypted data portion **824**. The file key **830** is a cipher key that, once obtained, can be used to decrypt the encrypted data portion **824** and thus, in general, is protected. In one implementation of the data structure **820**, the file key **830** is encrypted in conjunction with the access rules **829**. In another implementation of the data structure **820**, the file key **830** is encrypted with a private state key and further protected by the access rules **829**. The other information **832** is an additional space for other information to be stored within the security information **826**. For example, the other information **832** may be used to include other information facilitating secure access to the secured file, such as version number or author identifier.

The invention is preferably implemented by software or a combination of hardware and software, but can also be imple-

mented in hardware. The invention can also be embodied as computer readable code on a computer readable medium. The computer readable medium is any data storage device that can store data which can thereafter be read by a computer system. Examples of the computer readable medium include read-only memory, random-access memory, CD-ROMs, DVDs, magnetic tape, optical data storage devices, and carrier waves. A tangible computer readable medium embodiment of the invention can be any data storage device that can store data which can thereafter be read by a computer system, but specifically excludes implementation only through the use of carrier waves, which are considered not to be tangible. The computer readable medium can also be distributed over network-coupled computer systems so that the computer readable code is stored and executed in a distributed fashion.

The various embodiments, implementations and features of the invention noted above can be combined in various ways or used separately. Those skilled in the art will understand from the description that the invention can be equally applied to or used in various other settings with respect to different combinations, embodiments, implementations or features as provided in the description herein.

The invention may be practiced in two broad approaches: one, where document moving asynchronously through a persistent workflow (here, the state changes are typically triggered by the users); and two, where documents move synchronously through a single-use workflow, a plurality of which however can be initiated from a workflow template (here, the state changes are typically due to administrator central command). The two approaches may be combined for use in a single enterprise. State changes due to external events may occur with both approaches.

The advantages of the invention are numerous. Different embodiments or implementations may yield one or more of the following advantages. One advantage of the invention is that file security systems are able to automatically enforce process-driven security policies on files (e.g., documents). The automatic nature of the enforcement of the process-driven security policies alleviates otherwise excessive burdens on an administrator. Another advantage of the invention is that changing of the security policies for files (e.g., documents) in accordance with a process allows greater flexibility in utilizing security policies. Still another advantage of the invention is that the process-driven security policies can be enforced centrally or locally. Still another advantage is that a workflow ordered through a centralized document management system (DMS) may be extended to a plurality of documents stored in a distributed fashion, thereby allowing a system administrator to use the well-known DMS interface.

The foregoing description of embodiments is illustrative of various aspects/embodiments of the present invention. Various modifications to the invention can be made to the preferred embodiments by those skilled in the art without departing from the true spirit and scope of the invention as defined by the appended claims. Accordingly, the scope of the present invention is defined by the appended claims rather than the foregoing description of embodiments.

What is claimed is:

1. A method for limiting access to an electronic document, comprising:

associating, by one or more computing devices, a classifier with a first state of a process-driven security policy having a plurality of states, with different states having different sets of access restrictions;

associating, by the one or more computing devices, an identifier representing a user or a group of users with the first state of the process-driven security policy;

associating, by the one or more computing devices, the electronic document with at least the first state of the process-driven security policy having a set of access restrictions on the electronic document;

limiting access to the electronic document by encrypting, by the one or more computing devices, at least a portion of the electronic document using a group key corresponding to the identifier and a state key corresponding to the classifier and requiring at least both the group key and the state key to decrypt at least the portion of the electronic document; and

changing, by the one or more computing devices, the state of the process-driven security policy for the electronic document automatically without user or administrator interaction from the first state to a second state in response to an internal or external system event, wherein the changed state is based on a transition rule associated with the event.

2. The method as recited in claim **1**, wherein the identifier is a user ID or a group ID.

3. The method as recited in claim **1**, wherein the process-driven security policy is provided as part of a document security system.

4. The method as recited in claim **1**, wherein said method further comprises:

creating the electronic document; and

assigning the identifier to the created electronic document.

5. The method as recited in claim **1**, wherein the process-driven security policy is provided as part of a document security system, and

wherein said method further comprises:

creating a plurality of electronic documents; and

assigning the identifier and the classifier to the created electronic documents associated with the first state.

6. The method of claim **1**, wherein the event occurs at or is received at a client machine.

7. The method as recited in claim **1**, wherein the event is a user-triggered event.

8. A method for imposing access restrictions on an electronic document, comprising:

associating, by one or more computing devices, an electronic document with at least a first state of a plurality of states of a process-driven security policy, the first state associated with a classifier and with an identifier representing a user or a group of users, with different states having different set of access restrictions;

imposing the set of access restrictions associated with the first state on the electronic document by encrypting, by the one or more computing devices, at least a portion of the electronic document using a group key corresponding to the identifier and a state key corresponding to the classifier and requiring at least both the group key and the state key to decrypt at least the portion of the electronic document; and

changing, by the one or more computing devices, the state of the process-driven security policy for the electronic document automatically without user or administrator interaction from the first state to a second state in response to an internal or external system event, wherein the changed state is based on a transition rule associated with the event.

9. The method as recited in claim **8**, wherein the event is a user-triggered event.

10. The method as recited in claim **8**, wherein the event occurs at or is received at the client machine.

11. The method as recited in claim **8**, wherein the electronic document includes security information, and the security

information includes at least an indication of the state of the process-driven security policy for the electronic document.

12. The method as recited in claim **8**, wherein said method is performed on a plurality of documents on a document-by-document basis.

13. The method as recited in claim **8**, wherein at the client machine, a plurality of electronic documents is in one of the states of the process-driven security policy.

14. A tangible computer-readable medium having stored thereon computer-executable instructions that, if executed by a computing device, cause the computing device to perform a method for imposing access restrictions on an electronic document, the method comprising:

associating an electronic document with at least a first state of a plurality of states of a process-driven security policy, the first state associated with a classifier and with an identifier representing a user or a group of users, with different states having different set of access restrictions;

imposing the set of access restrictions associated with the first state on the electronic document by encrypting at least a portion of the electronic document using a group key corresponding to the identifier and a state key corresponding to the classifier and requiring at least both the group key and the state key to decrypt at least the portion of the electronic document; and

changing the state of the process-driven security policy for the electronic document automatically without user or administrator interaction from the first state to a second state in response to an internal or external system event, wherein the changed state is based on a transition rule associated with the event.

15. The tangible computer-readable medium as recited in claim **14**, wherein the event is a user-triggered event.

16. The tangible computer-readable medium as recited in claim **14**, wherein the event occurs at or is received at the client machine.

17. The tangible computer-readable medium as recited in claim **14**, wherein the electronic document includes security information, and the security information includes at least an indication of the state of the process-driven security policy for the electronic document.

18. The tangible computer-readable medium as recited in claim **14**, wherein the process-driven security policy is imposed on a plurality of documents on a document-by-document basis.

19. The tangible computer-readable medium as recited in claim **14**, wherein at the client machine, a plurality of electronic documents is in one of the states of the process-driven security policy.

* * * * *

UNITED STATES PATENT AND TRADEMARK OFFICE

CERTIFICATE OF CORRECTION

PATENT NO. : 7,703,140 B2
APPLICATION NO. : 10/677049
DATED : April 20, 2010
INVENTOR(S) : Nath et al.

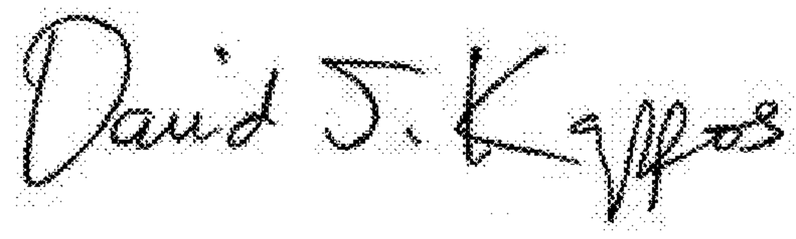
It is certified that error appears in the above-identified patent and that said Letters Patent is hereby corrected as shown below:

Title Page, Item (75), Inventor Section, please change "Los Altos, CA (US)" to --Mountain View, CA (US)--.

Column 16, line 46, claim 8, please change "different set" to --different sets--.

Column 17, line 18, claim 14, please change "different set" to --different sets--.

Signed and Sealed this
Twenty-second Day of March, 2011



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Director of the United States Patent and Trademark Office