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(54) **METASTABILITY INJECTOR FOR A  
CIRCUIT DESCRIPTION**

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**G06F 9/45** (2006.01)

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(58) **Field of Classification Search** ..... 716/5,  
716/6

See application file for complete search history.

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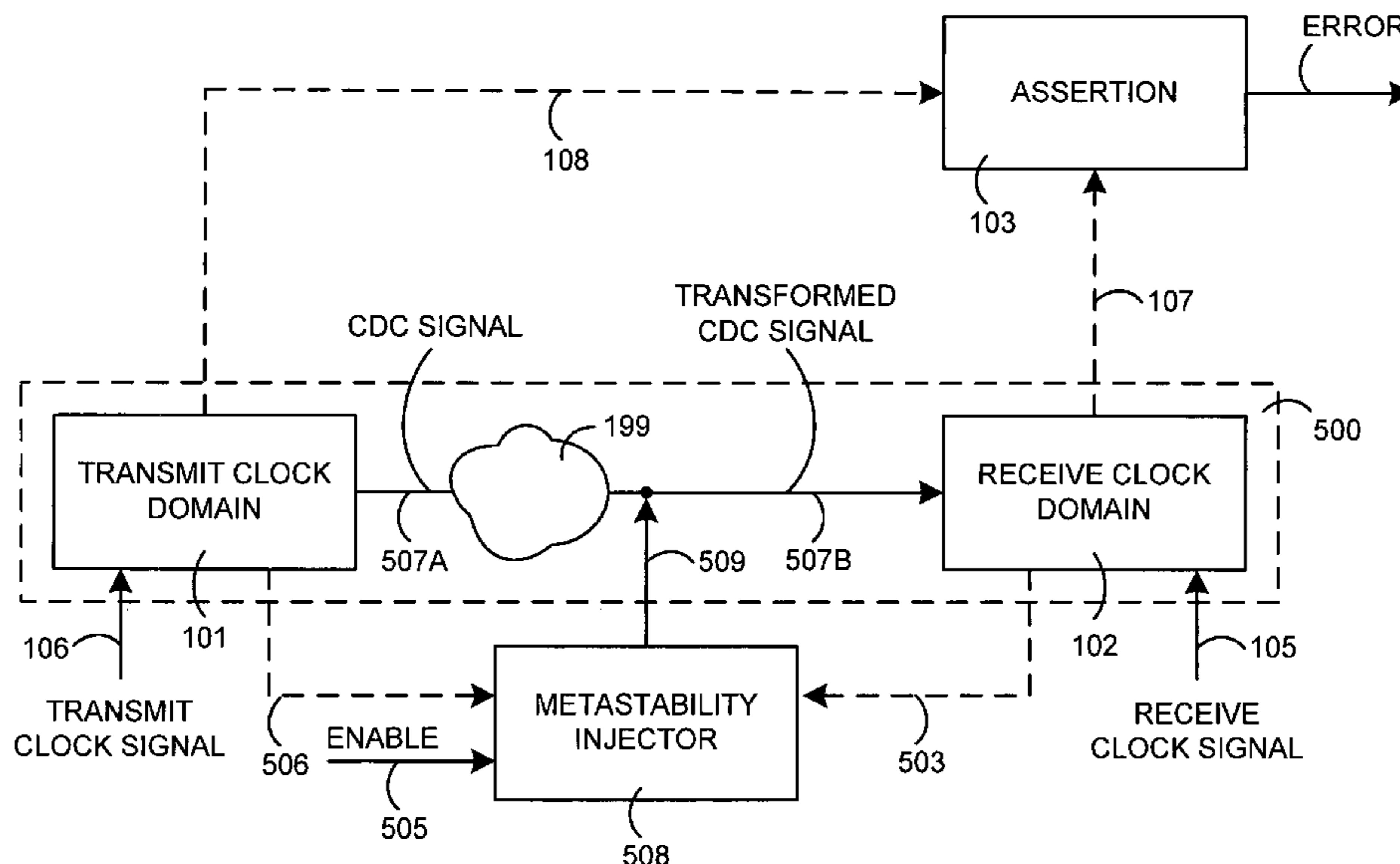
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(57) **ABSTRACT**

During verification of a description of a circuit containing a pre-determined assertion, in order to detect incorrect behavior of the circuit that may be caused by metastability occurring in signals that cross clock domains ("CDC" signals) in the circuit, the description of the circuit is automatically transformed by addition of circuitry to inject the effects of metastability into the CDC signals. The transformed description containing the circuitry to inject metastability is verified in the normal manner. Certain embodiments analyze the transformed description using a model checking method to determine a stimulus sequence that will cause the pre-determined assertion to be violated. The transformed circuit is then simulated in some embodiments, using the stimulus sequence from model checking, and an incorrect behavior of the circuit due to metastability is displayed, for diagnosis by the circuit designer. The circuit designer may revise the circuit description and iterate as noted above.

**24 Claims, 24 Drawing Sheets**



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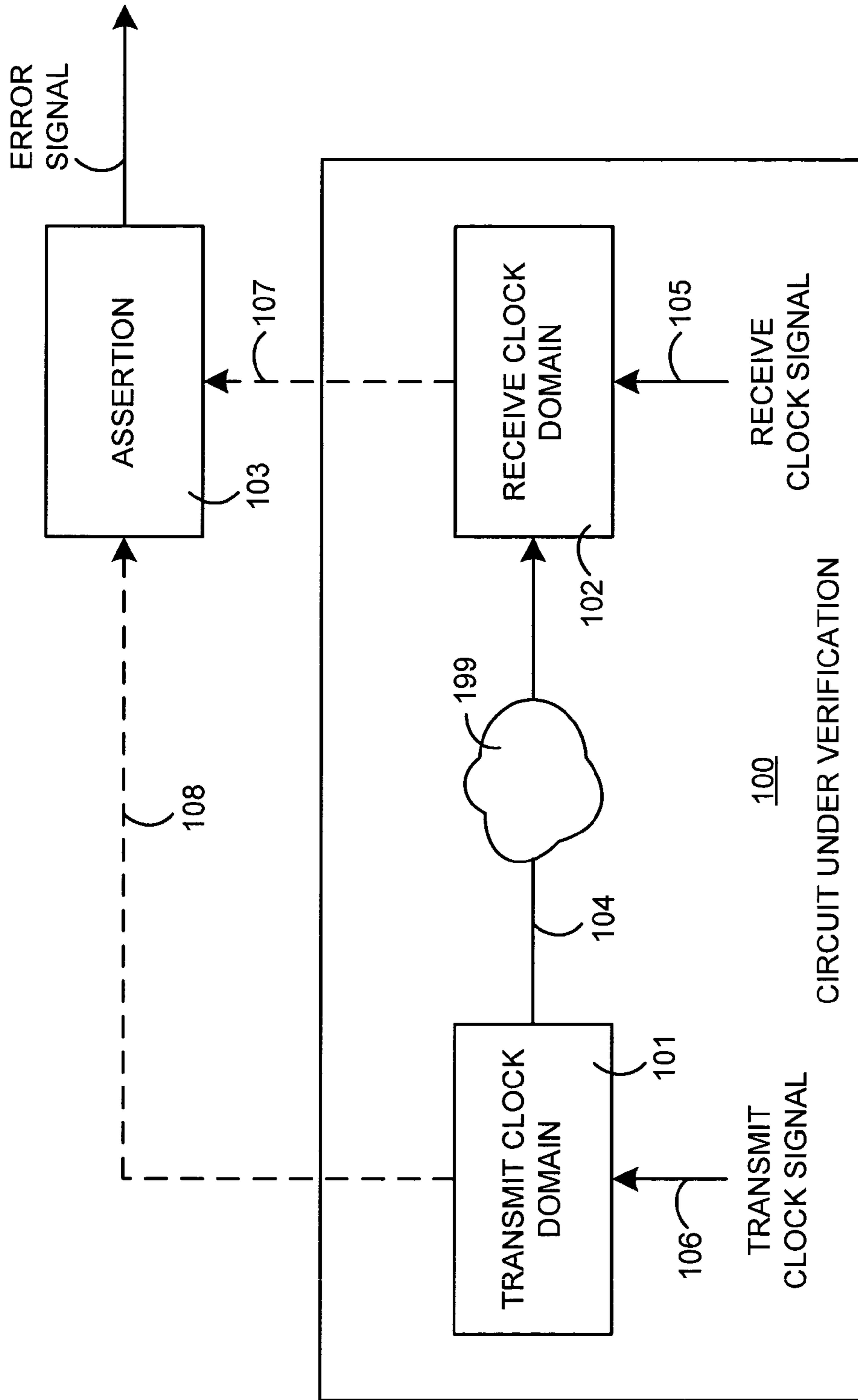


Fig. 1A  
(prior art)

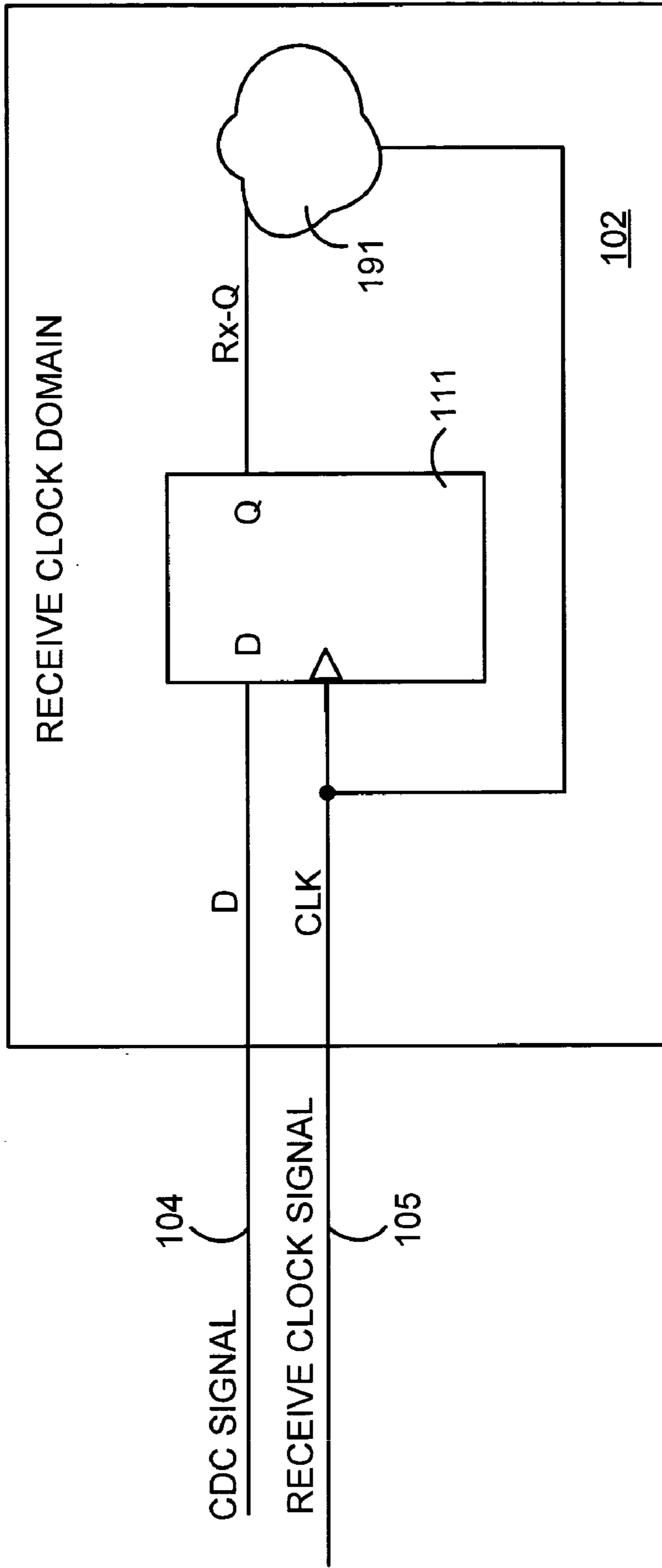


Fig. 1B  
(prior art)

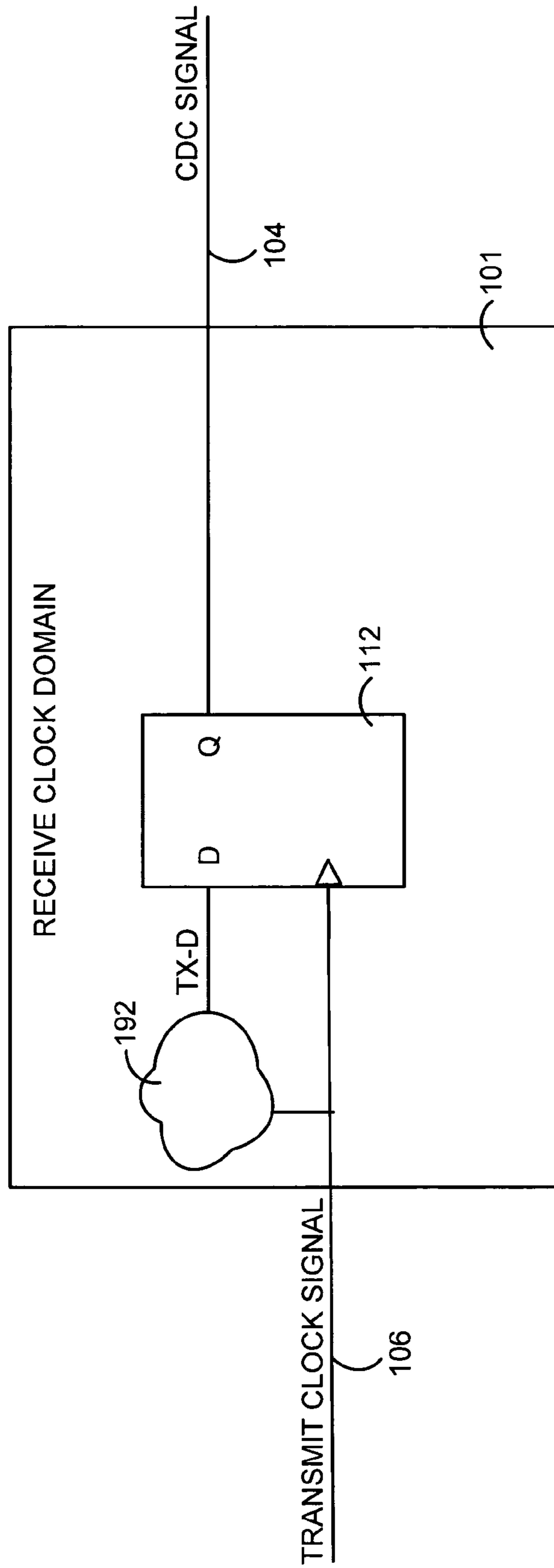


Fig. 1C  
(prior art)

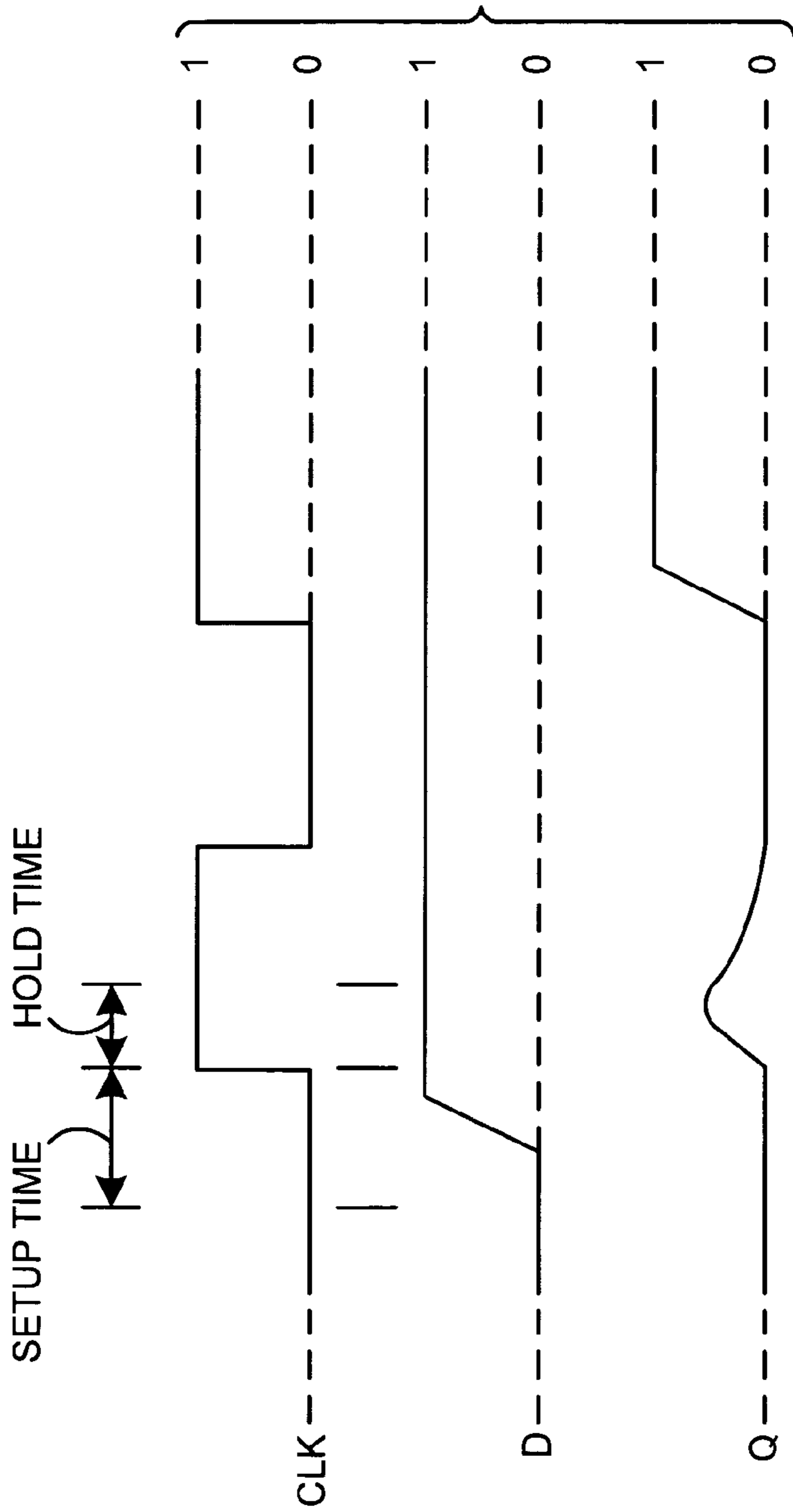


Fig. 2A  
(prior art)

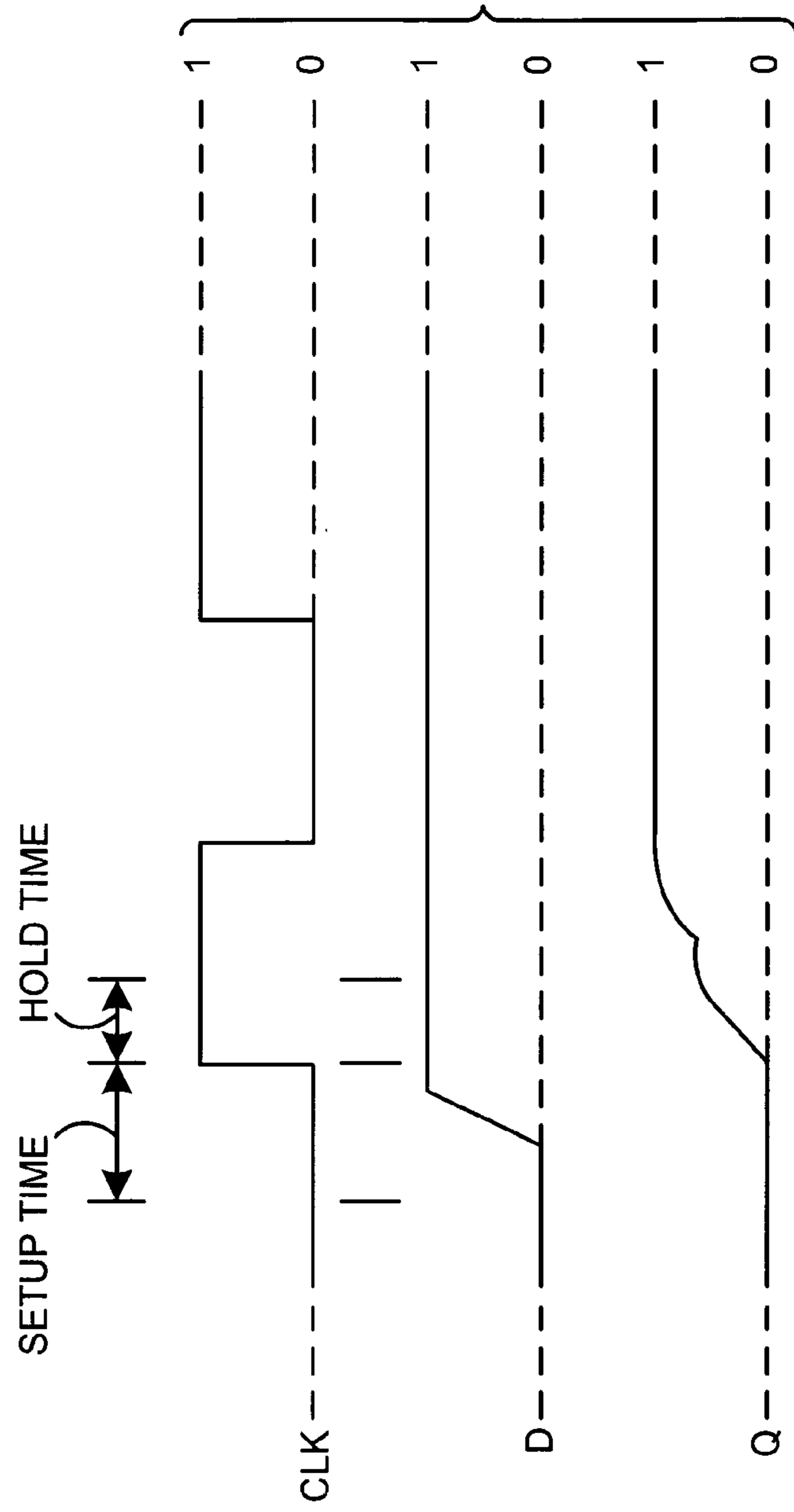


Fig. 2B  
(prior art)

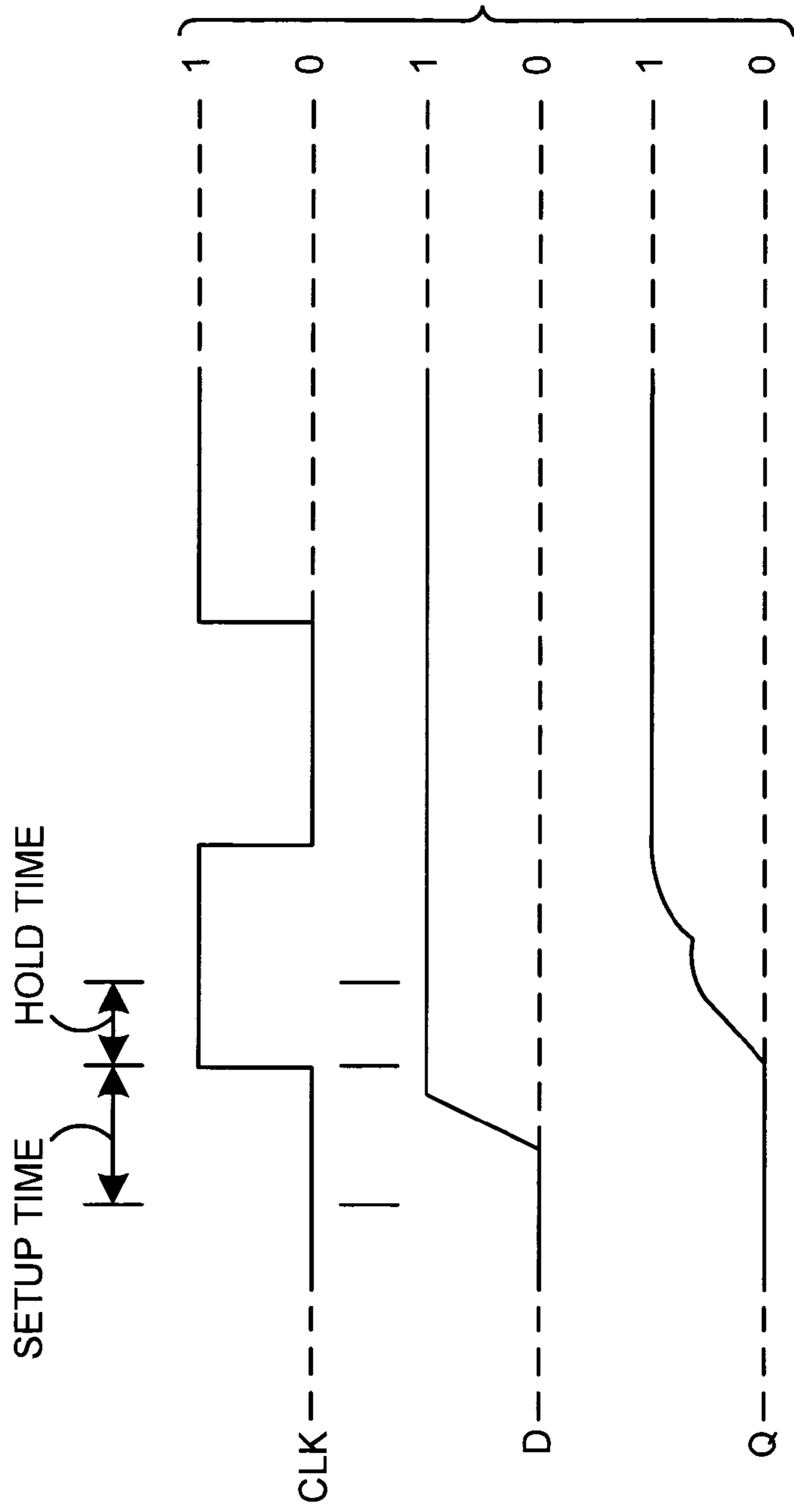


Fig. 2C  
(prior art)



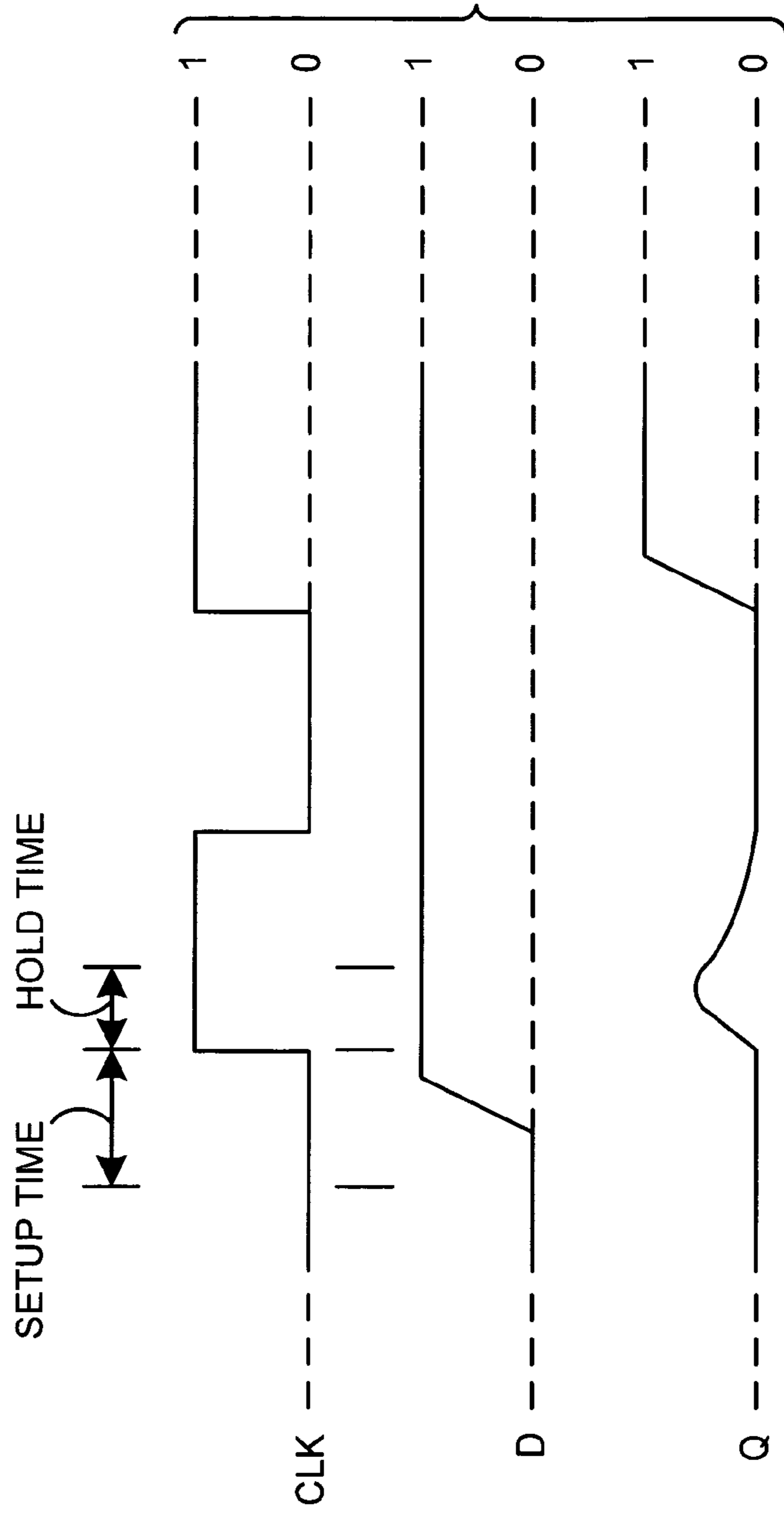


Fig. 2D  
(prior art)

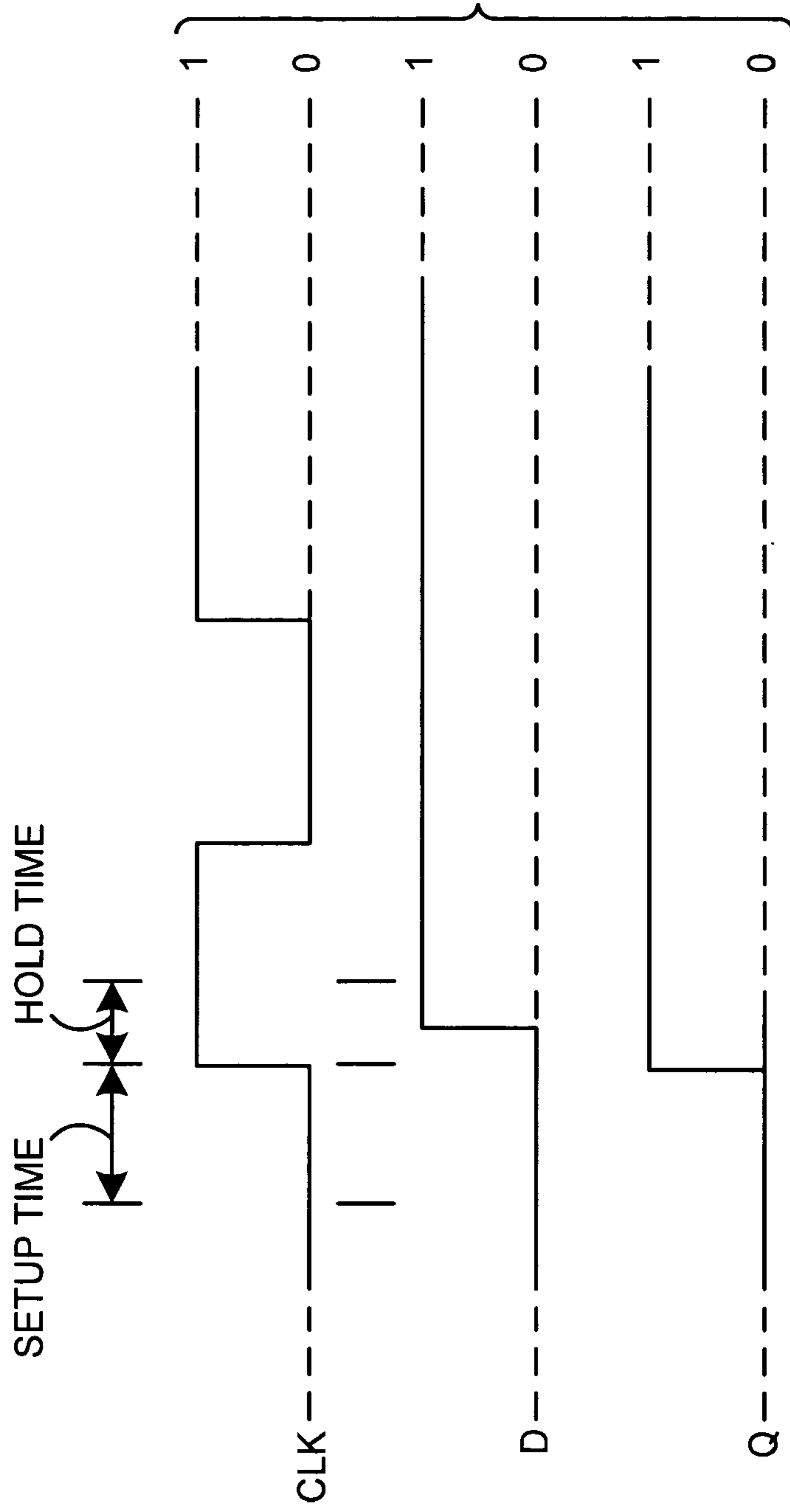


Fig. 3A  
(prior art)

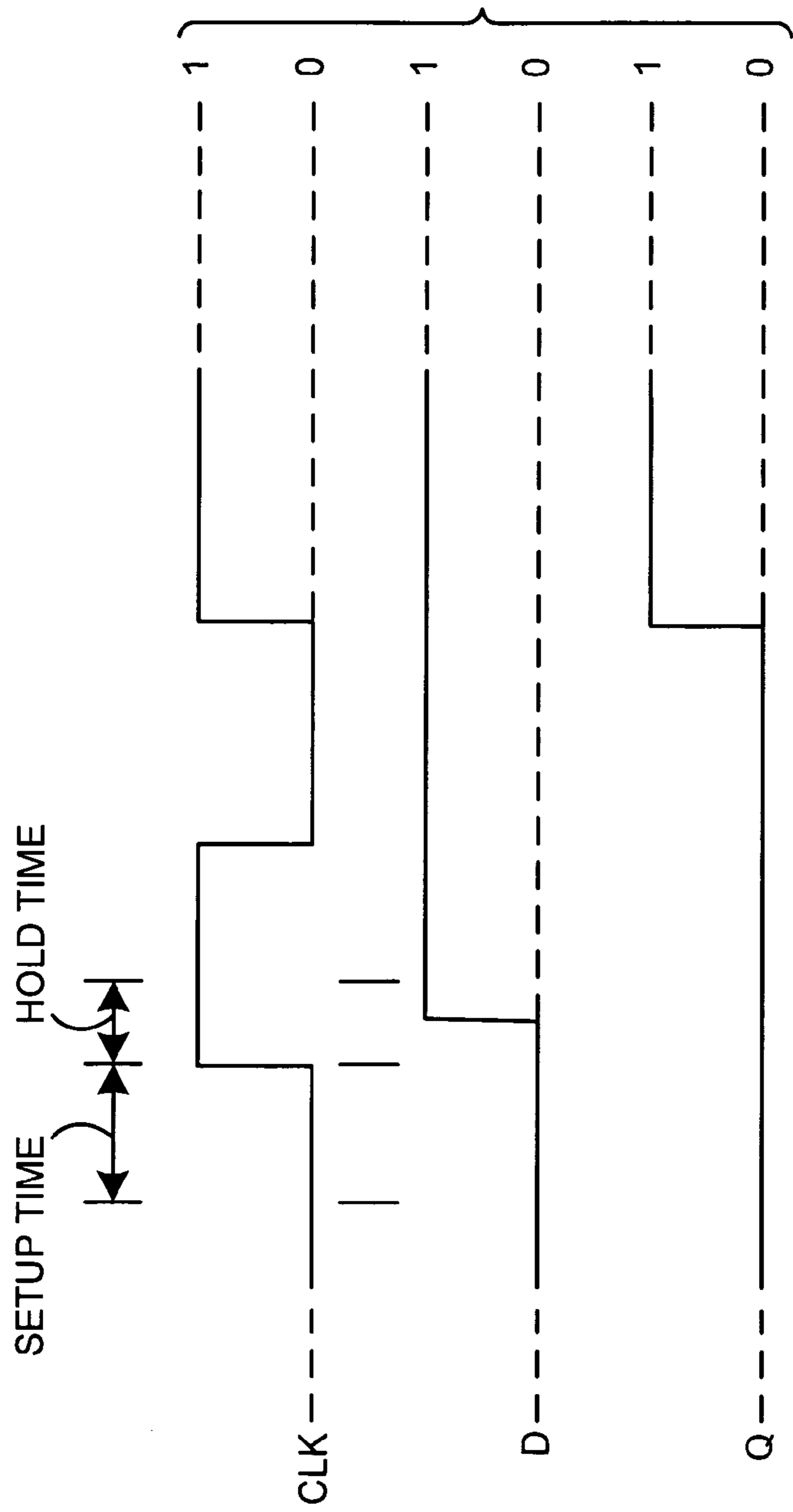


Fig. 3B  
(prior art)

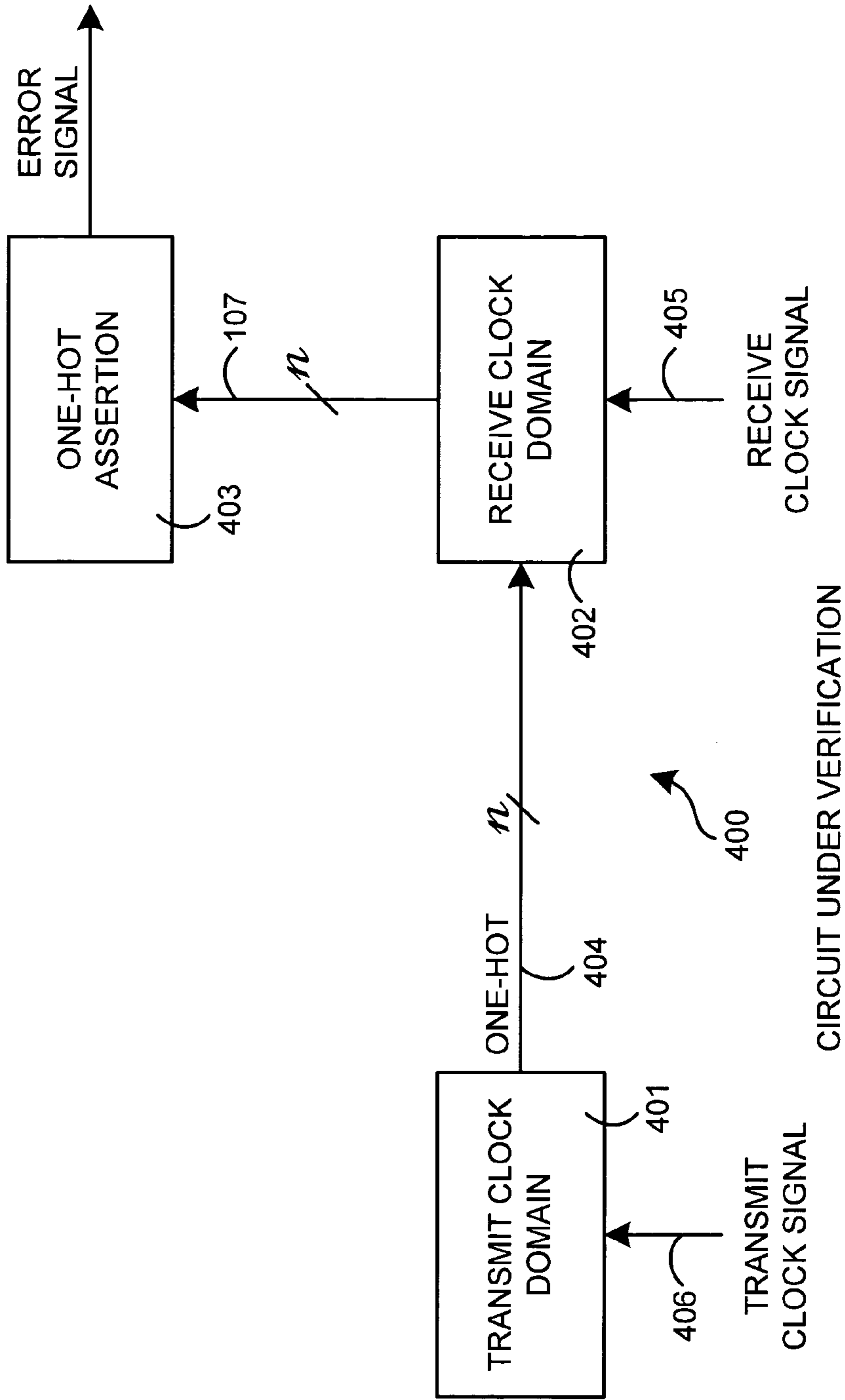


Fig. 4A  
(prior art)

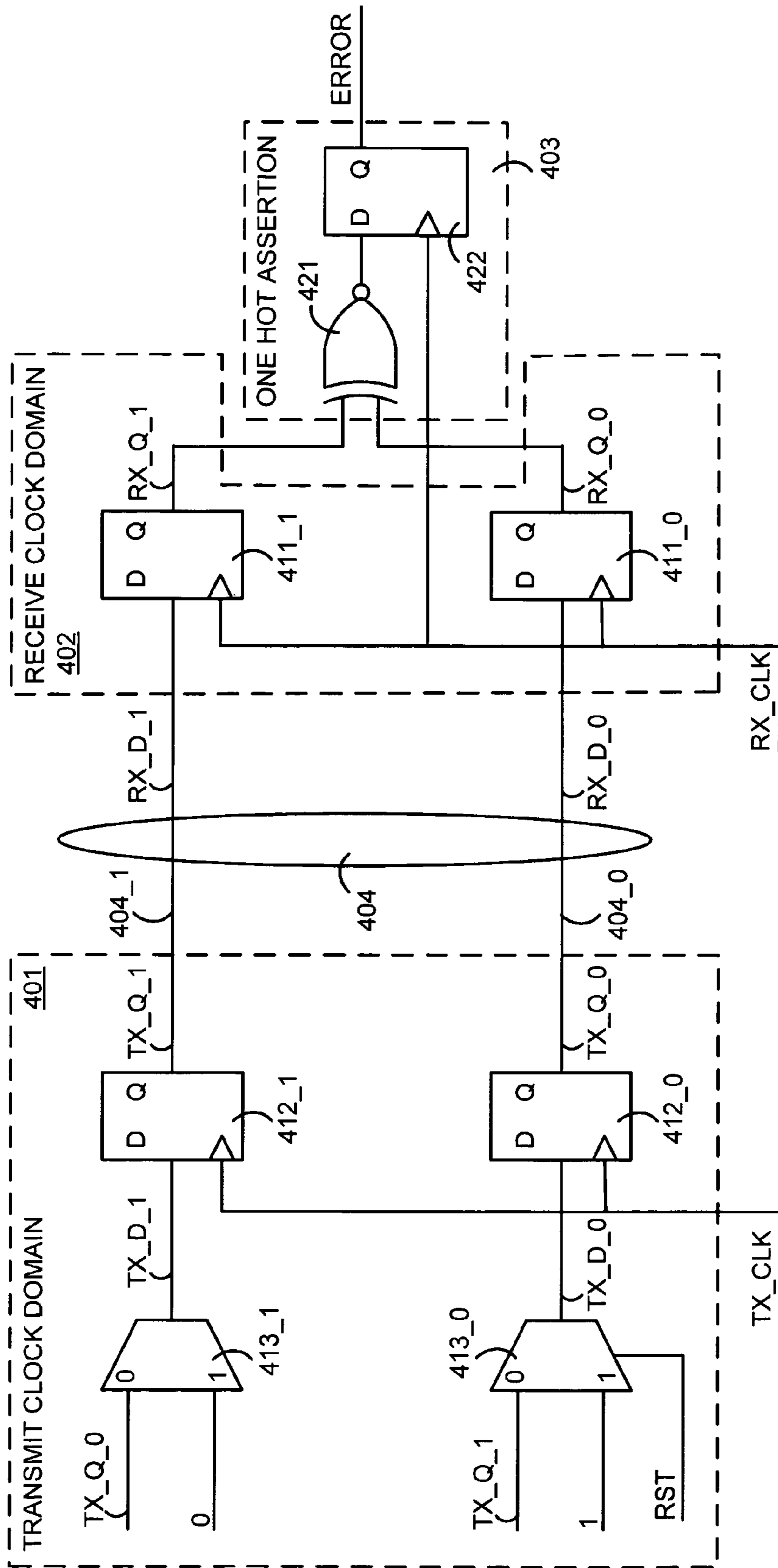


Fig. 4B  
(prior art)

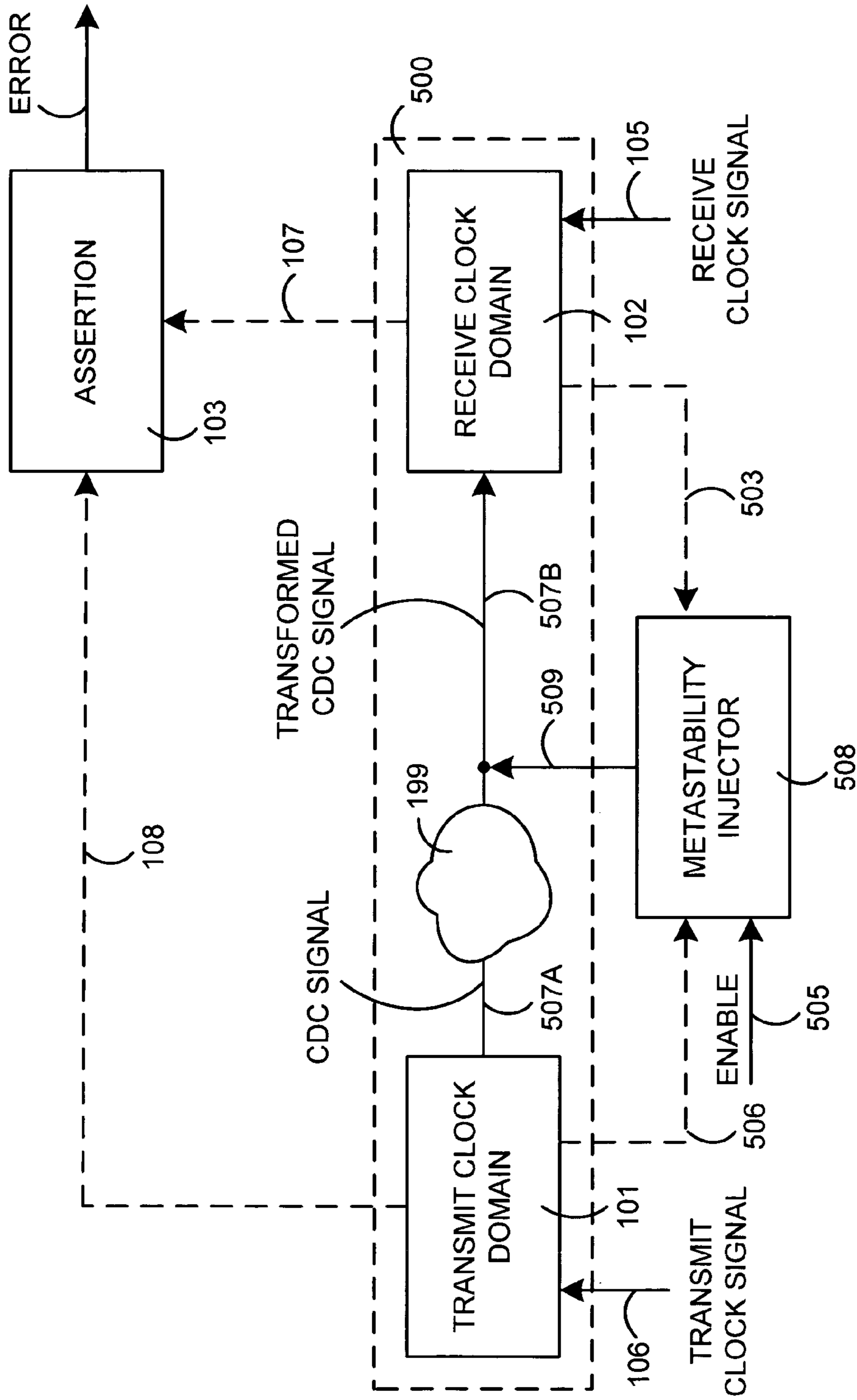


Fig. 5A

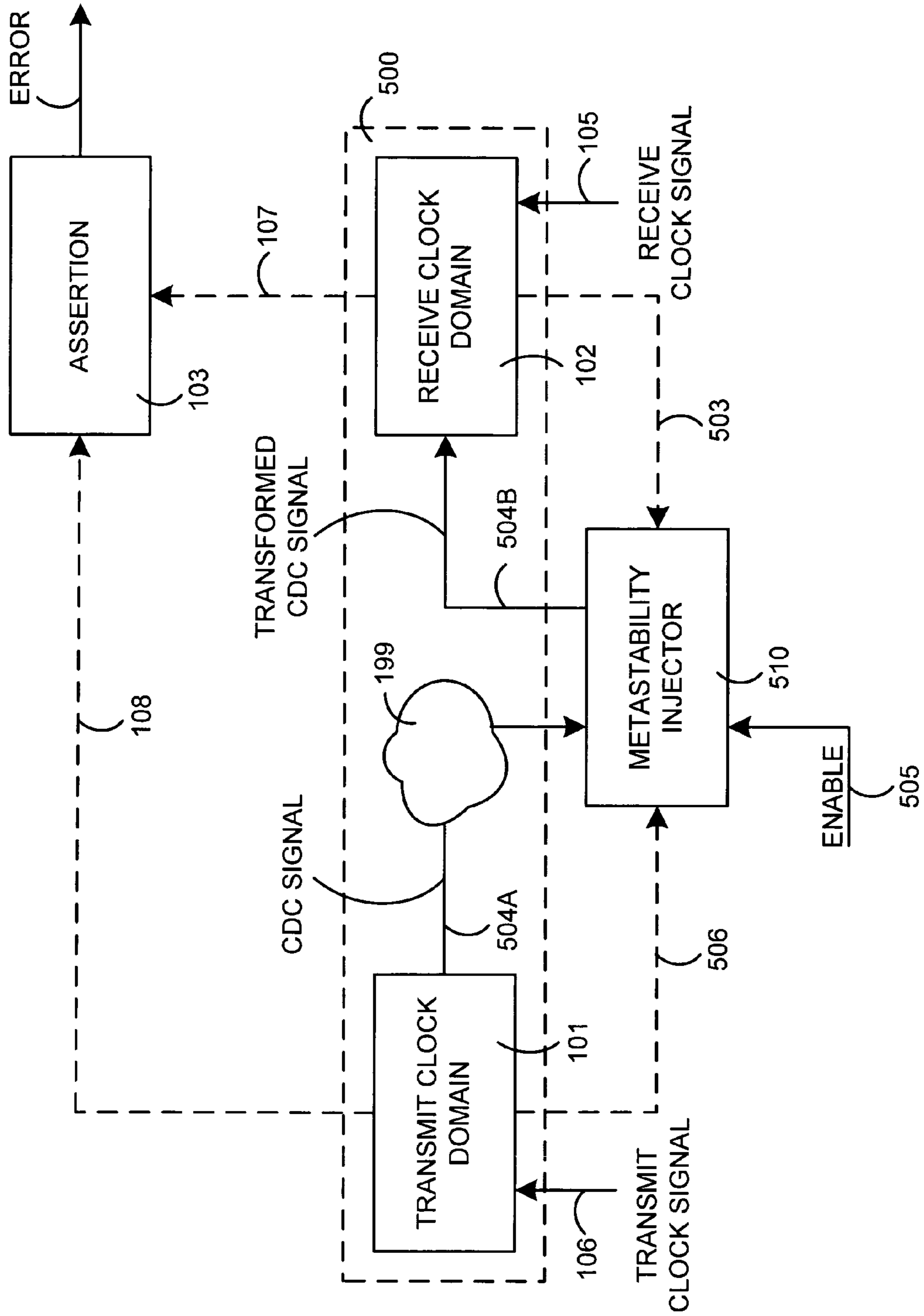


Fig. 5B

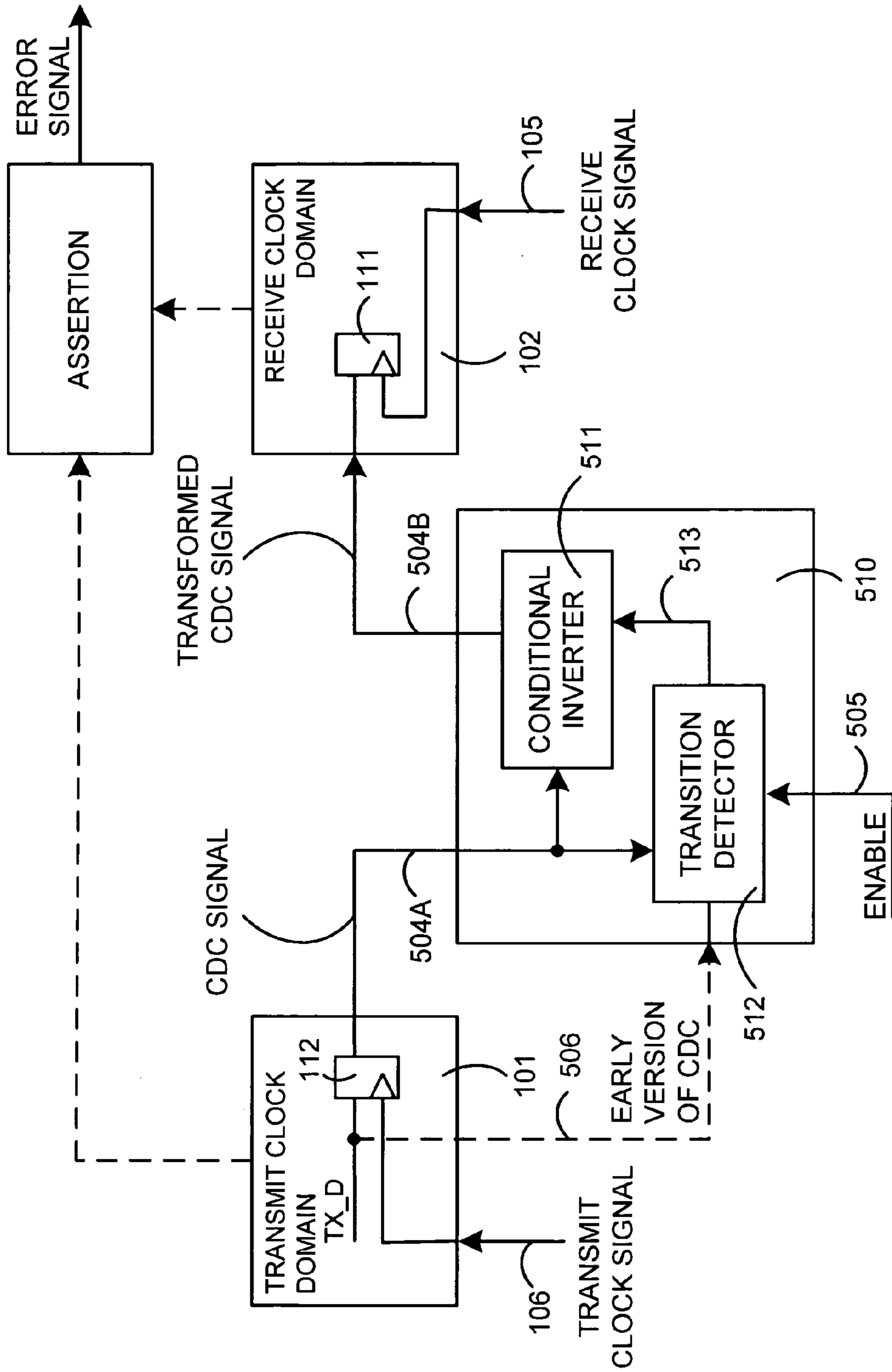


Fig. 5C



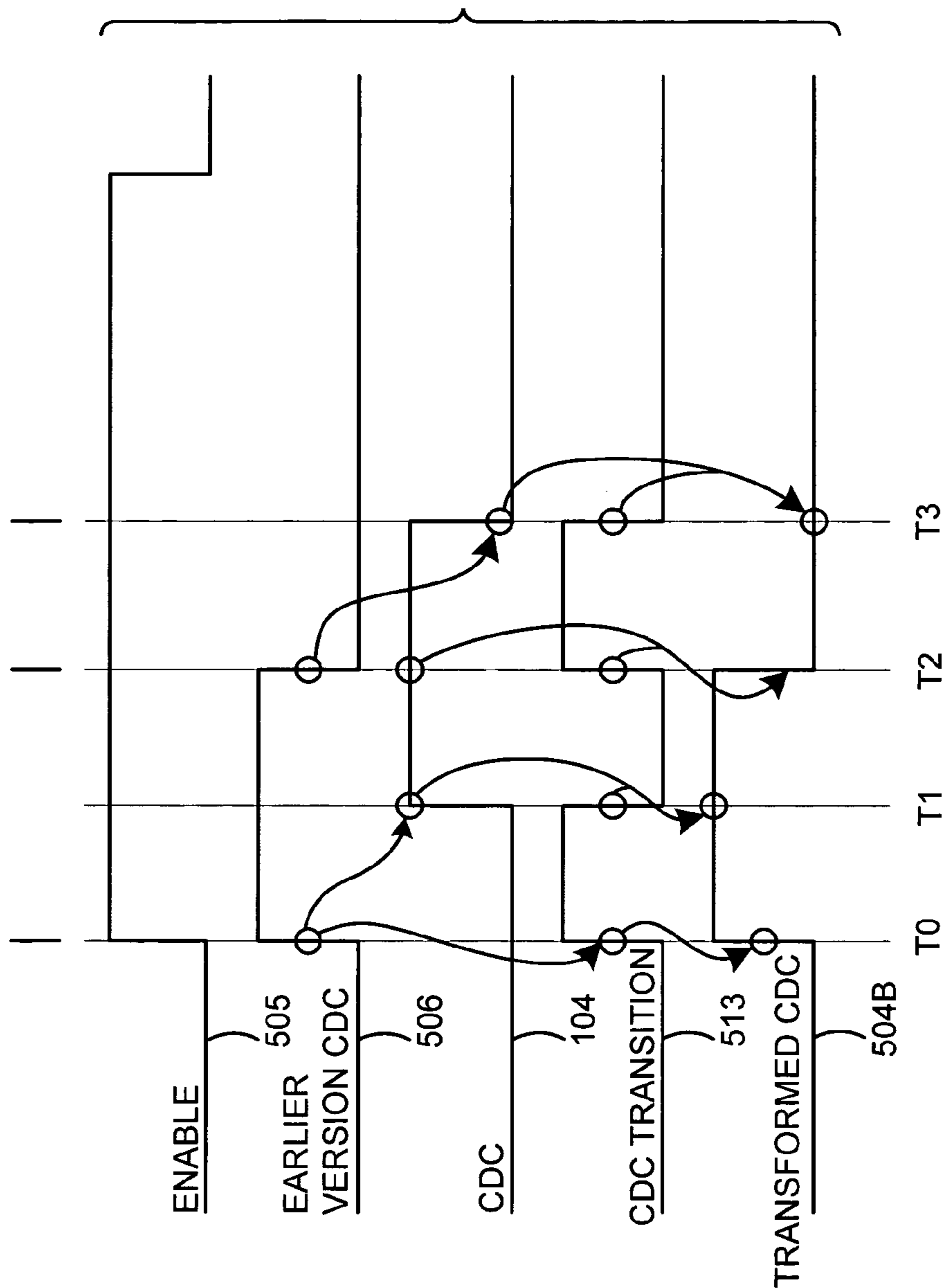


Fig. 5D

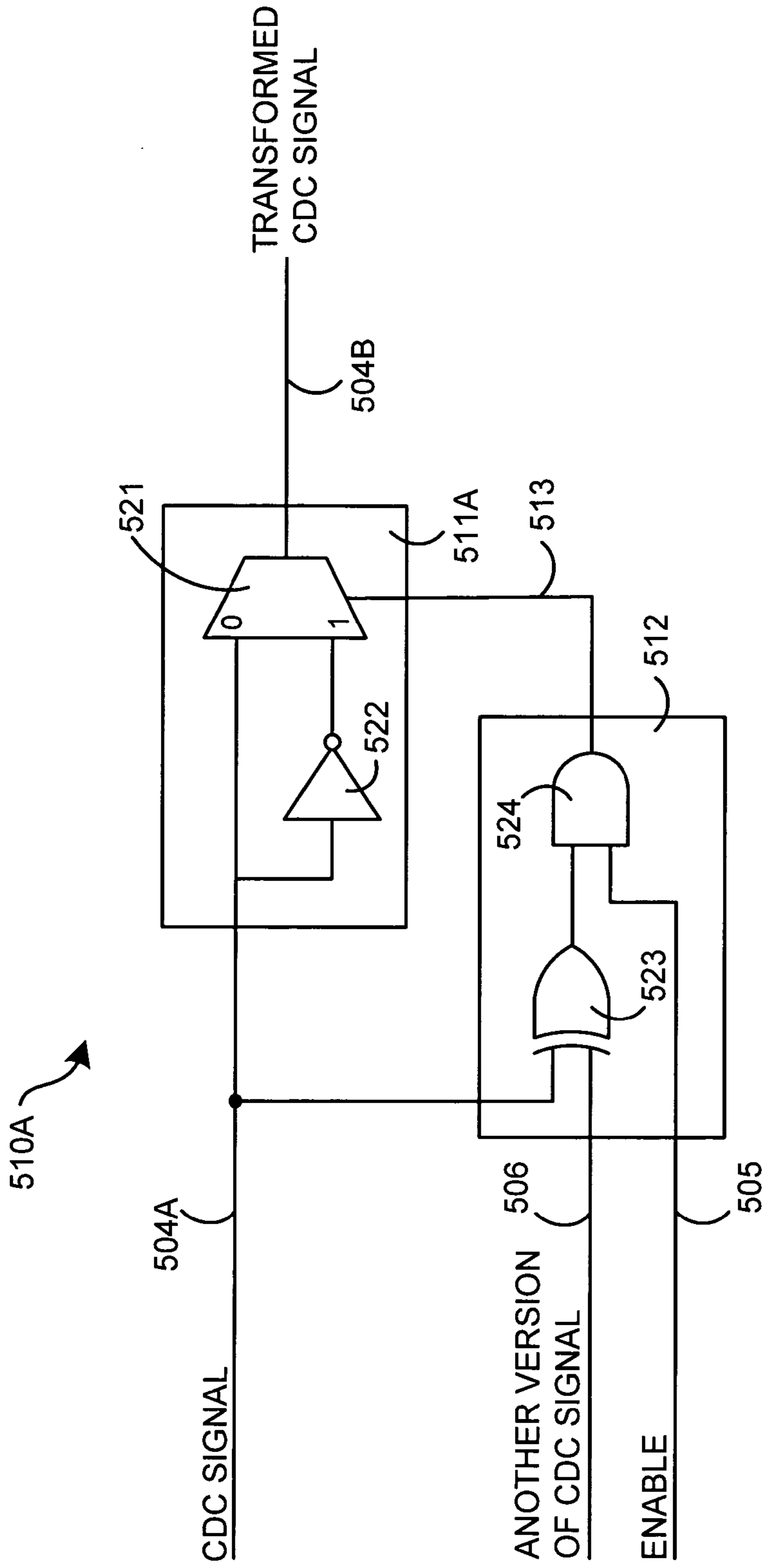


Fig. 5E

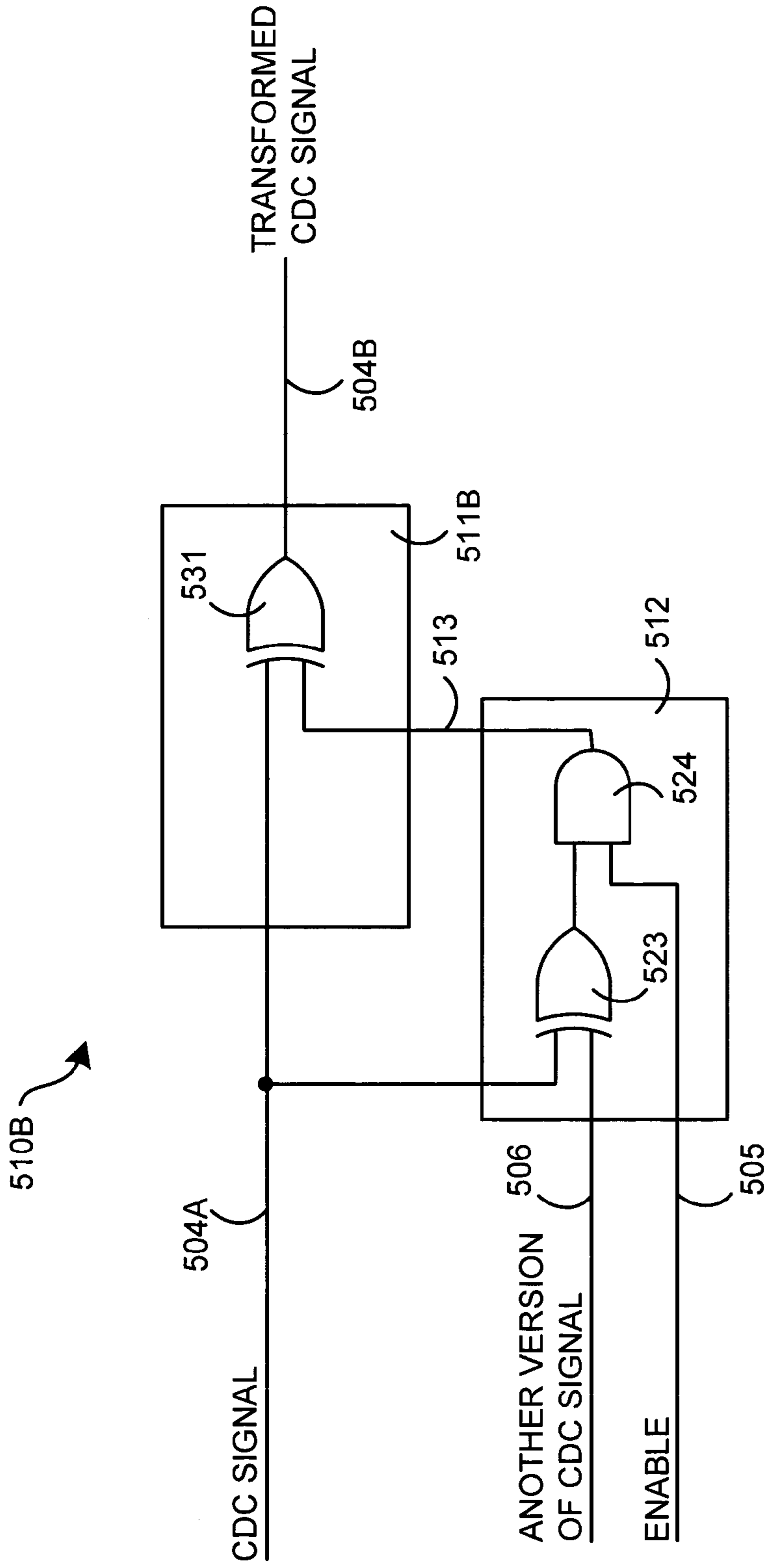


Fig. 5F

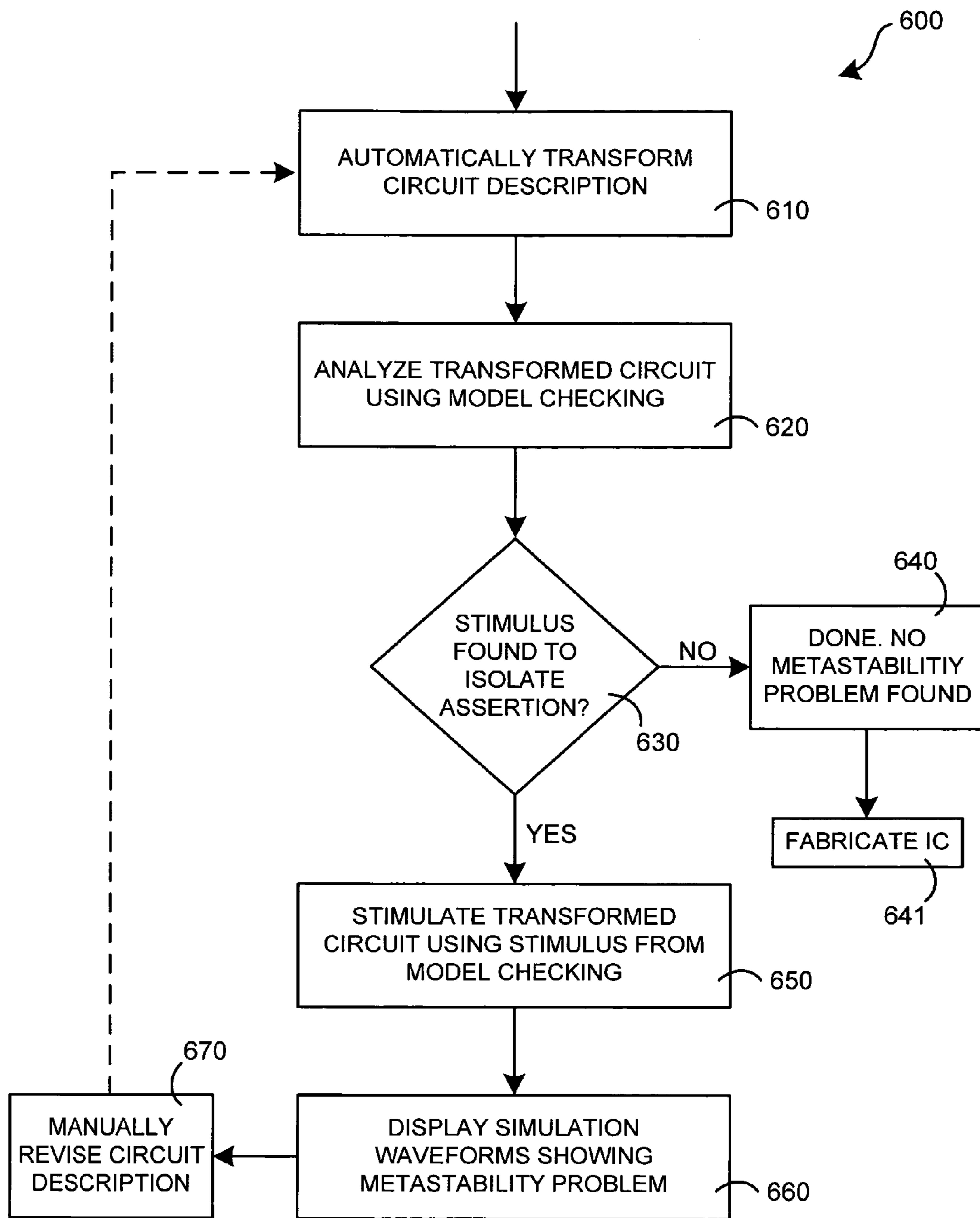


Fig. 6A

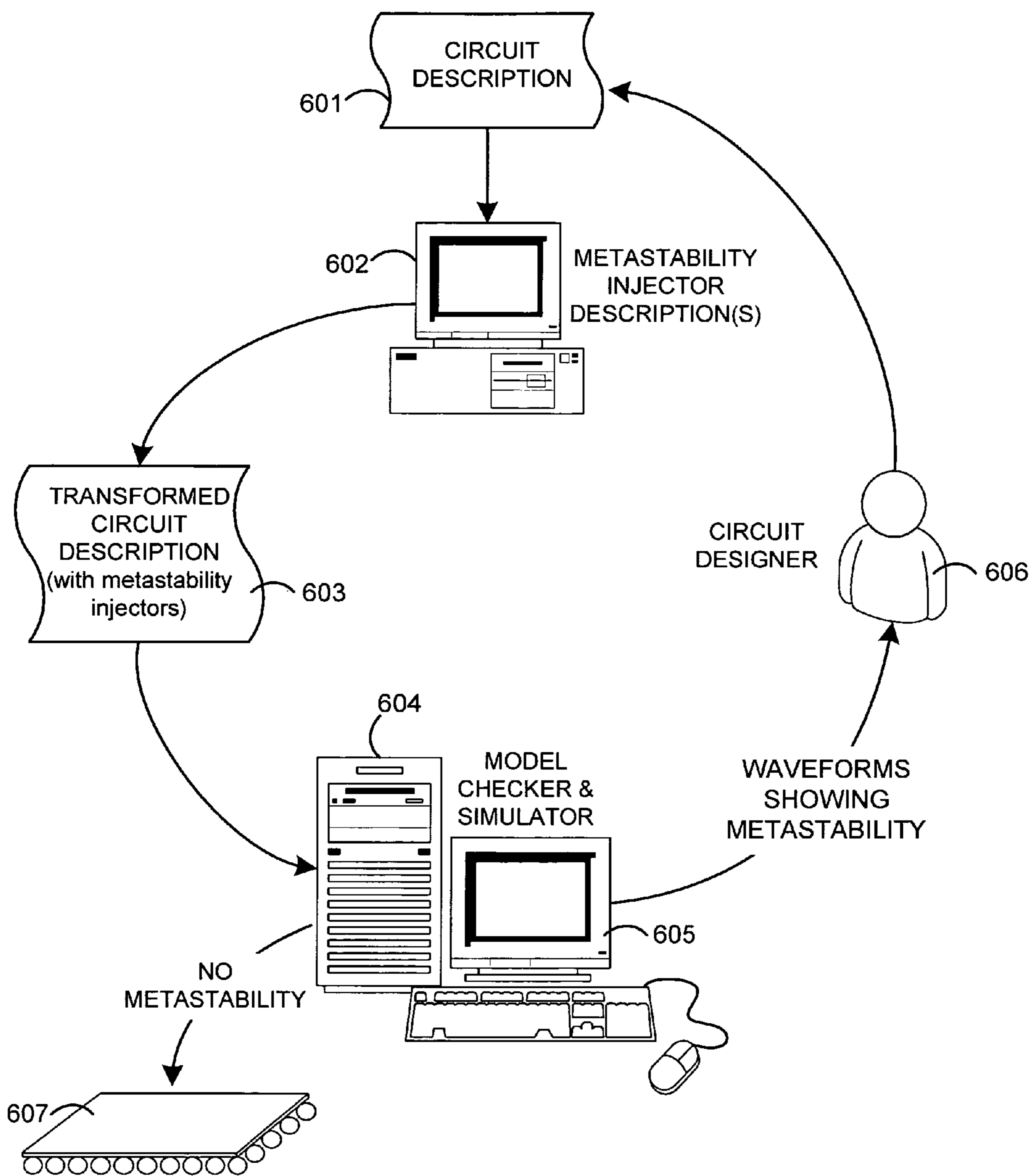


Fig. 6B

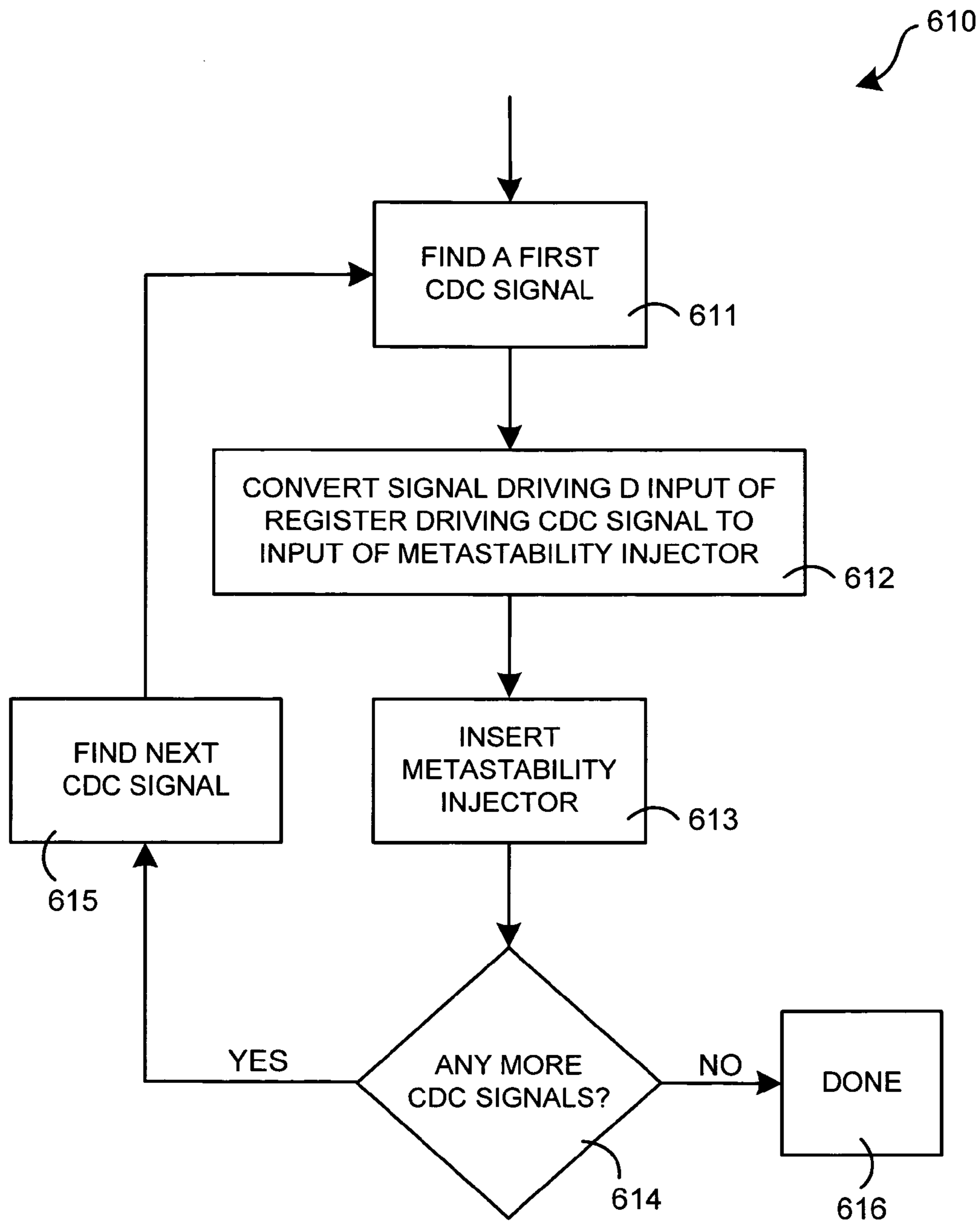


Fig. 6C

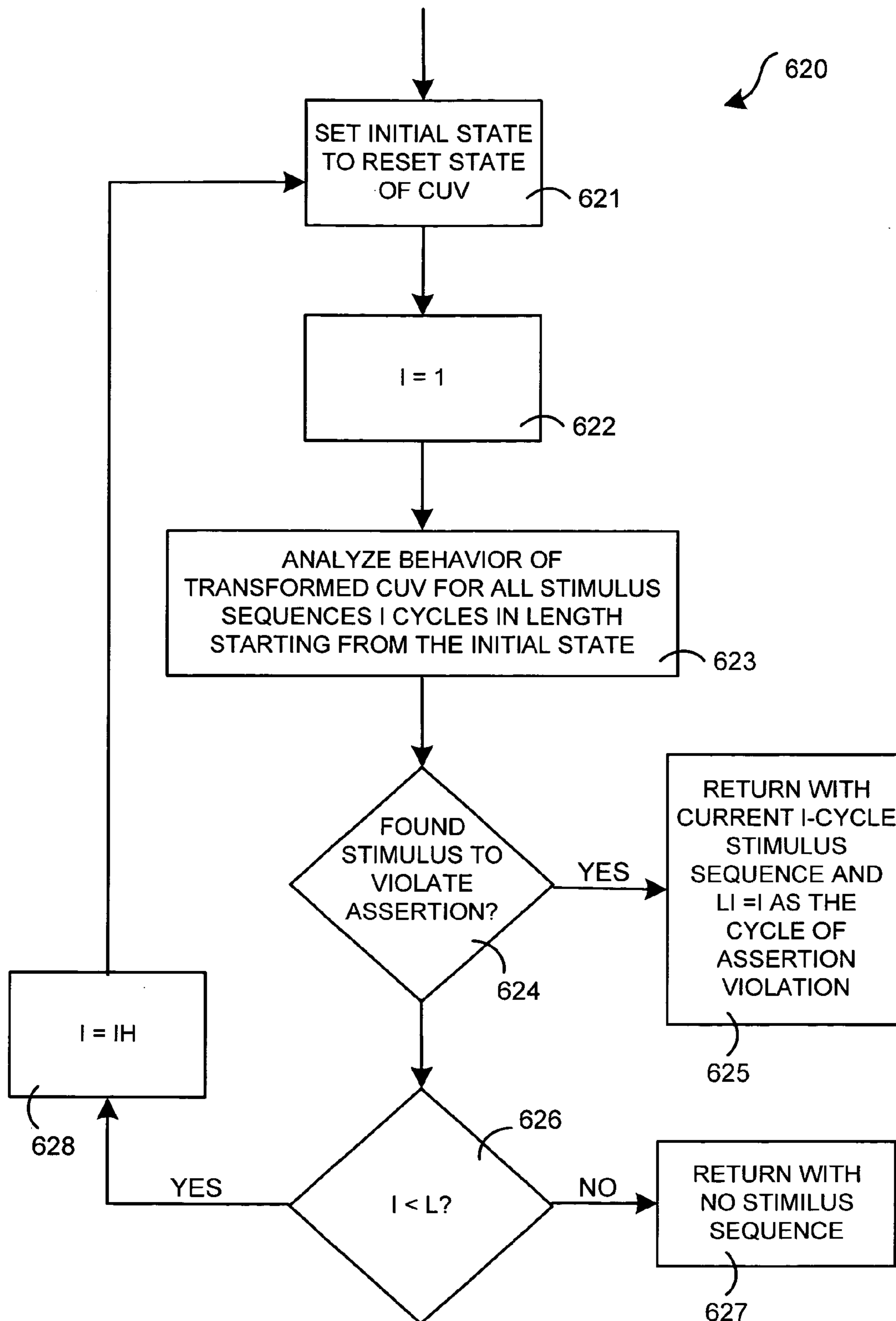


Fig. 6D

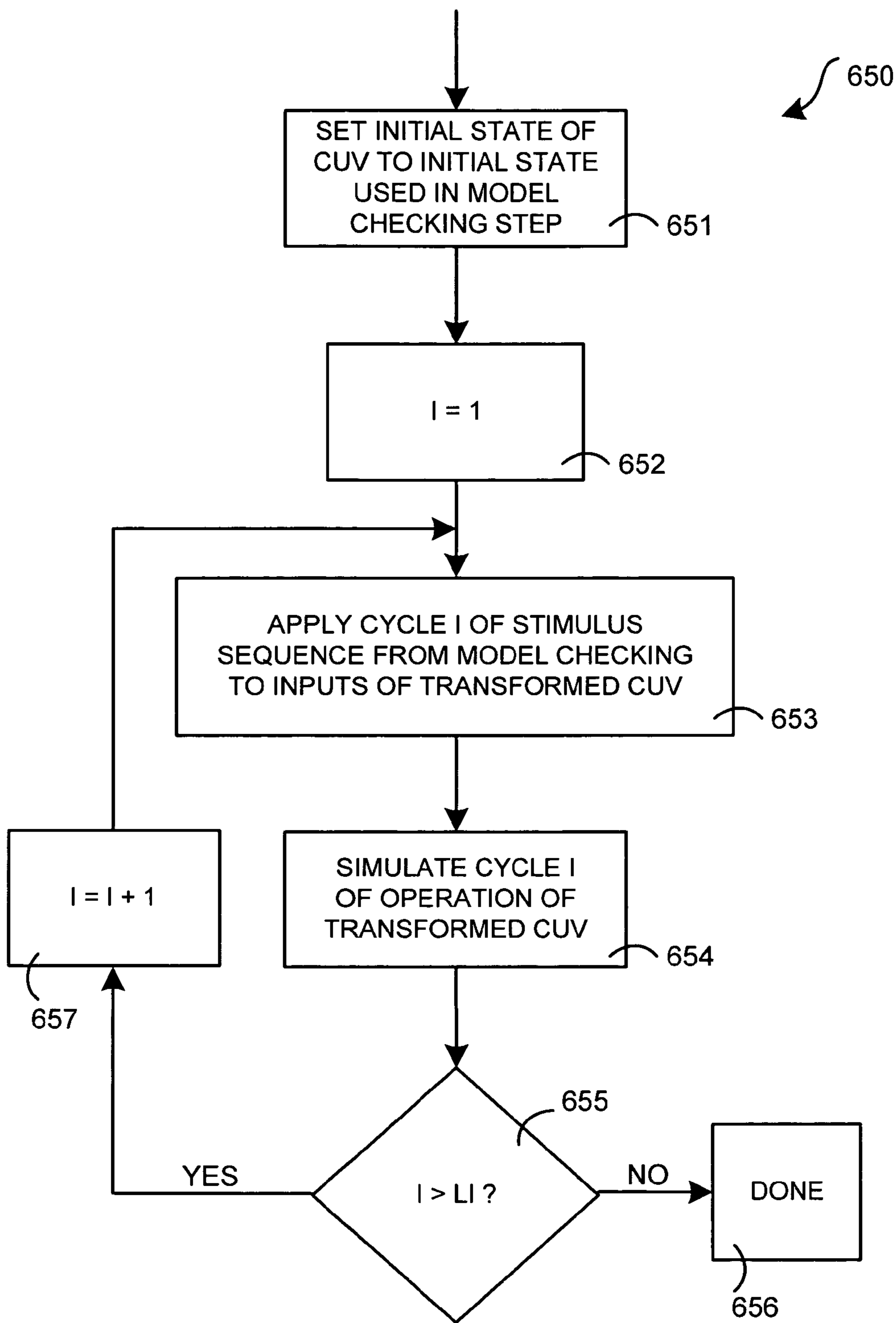


Fig. 6E



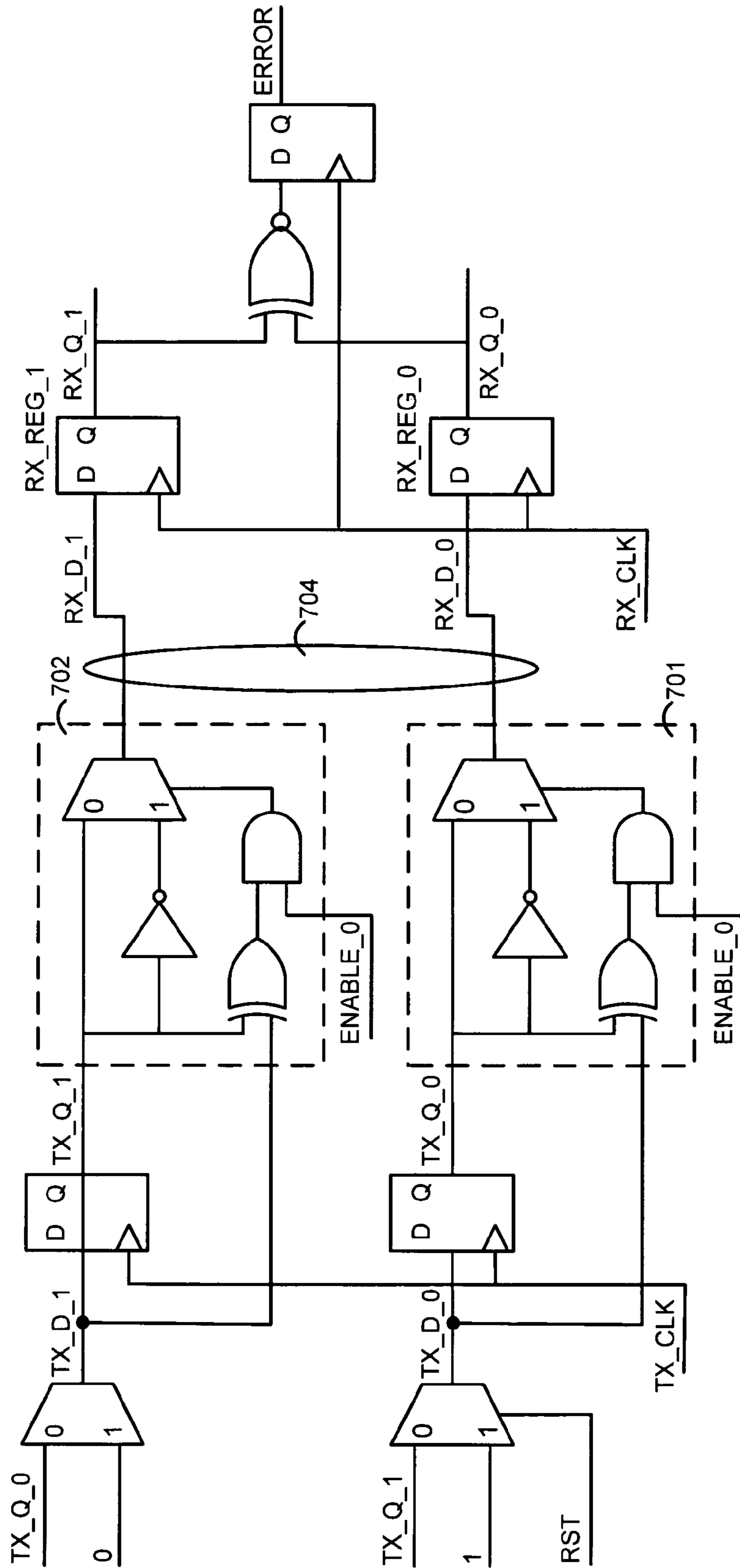


Fig. 7

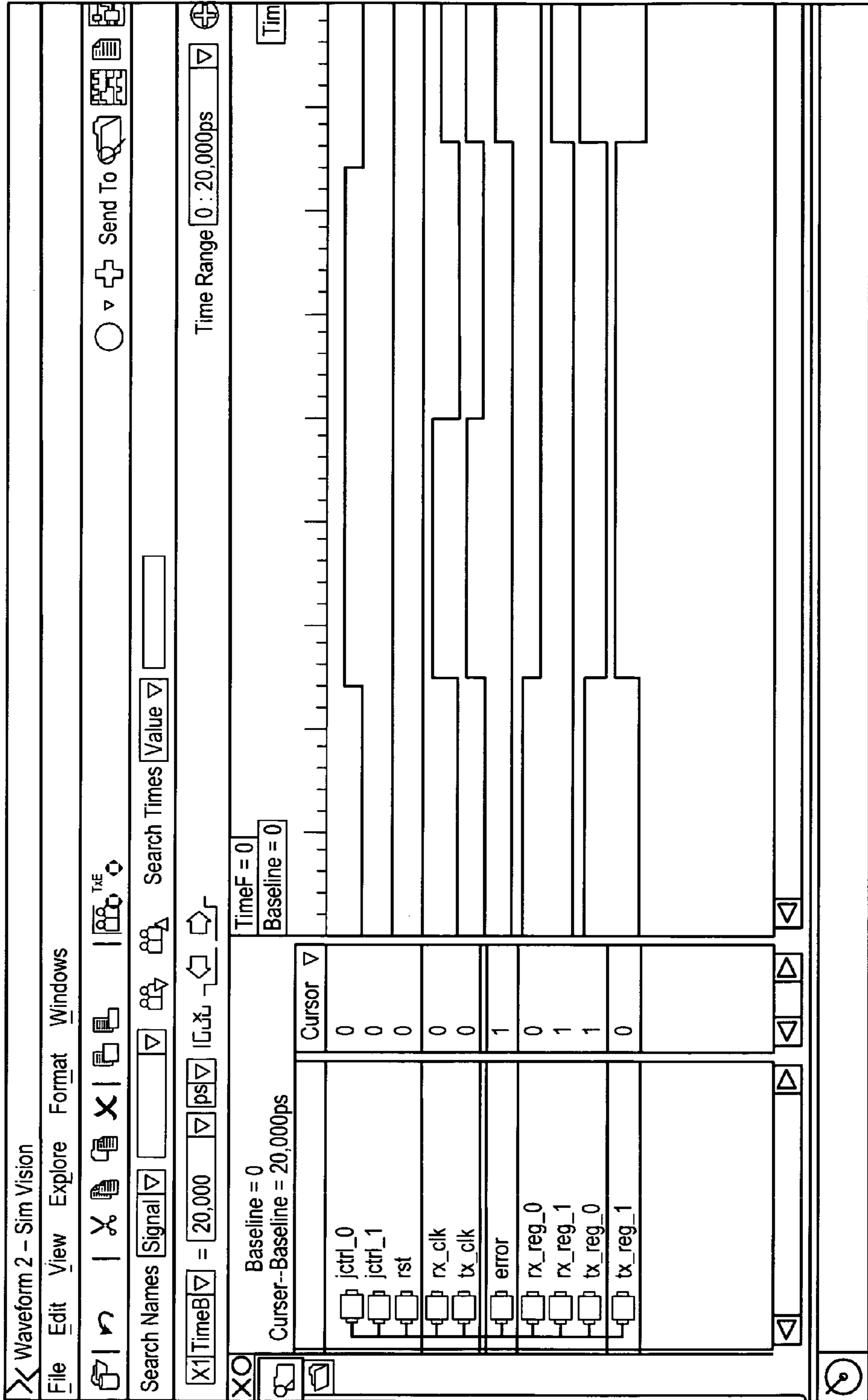


Fig. 8

## METASTABILITY INJECTOR FOR A CIRCUIT DESCRIPTION

### BACKGROUND

FIG. 1A shows a prior art circuit **100** containing two portions **101** and **102** and a path **104** that carries a signal from portion **101** to portion **102**. Path **104** may pass through any amount of combinational logic **199** (formed of logic elements but no storage elements). Registers in one portion **102** are clocked by a clock signal on a path **105** whereas registers in the other portion **101** are clocked by another clock signal on a different path **106**. Note that the two clock signals on the two paths **105** and **106** are different from one another, which makes the two portions **101** and **102** into two different clock domains, hereinafter referred to as transmit clock domain **101** and receive clock domain **102**. The difference in clock signals on paths **105** and **106** can be a difference in only frequency or only phase or both. For example, the clock signals on path **105** and **106** may have the respective frequencies 50 MHz and 37 MHz. A signal on path **104** crosses from clock domain **101** to clock domain **102**, and hence this signal (on path **104**) is hereinafter called a clock-domain-crossing (“CDC”) signal.

Although circuit **100** is illustrative of one clock-domain-crossing signal it is well known to the skilled artisan that today’s integrated circuits have 100s of 1000s of such clock-domain-crossing signals and have 100s of clock domains. Moreover, the clock-domain-crossing signal on path **104** may pass through any amount of combinational logic **109** when traveling from transmit domain **101** to receive domain **102**. Combinational logic **109** typically consists of any number of logic elements that are not clocked (i.e. there are no storage elements therein).

Each of portions **101** and **102** of circuit under verification **100** may contain any number of and any kind of circuit elements, e.g. storage elements that need to be clocked such as flip flops, as well as logic elements such as XOR gates and AND gates. For example, FIG. 1B shows a register **111** in the receive clock domain **102** of the circuit **100** of FIG. 1A. The D input of the register **111** of FIG. 1B is connected to the path **104** of FIG. 1A and therefore receives the clock-domain-crossing signal from domain **101**. Moreover, the Q output of the register **111** of FIG. 1B generates a signal RX\_Q that may be provided to any additional circuitry **191** in receive clock domain **102**.

FIG. 1C shows another register **112** that is located in the transmit clock domain **101** of the circuit **100** of FIG. 1A. Register **112** has a Q output which drives the clock-domain-crossing signal on path **104**. The D input of the register **112** of FIG. 1C receives a signal TX\_D from any additional circuitry **192** in the transmit clock domain **101**. The above-described additional circuitry **191** and **192** are each normally clocked by their respective clock signals on paths **105** and **106** respectively (and for this reason they belong to their respective clock domains).

It is well known in the art to verify the functional behavior of circuit **100** (which is also referred to as “circuit-under-verification”), based on a circuit description, by use of conventional register-transfer-level (hereinafter, RTL) simulators such as VCS (from Synopsys, Inc.) and Verilog NC (from Cadence Design Systems, Inc.). The circuit description for circuit **100** is normally articulated by a circuit designer in a Hardware Description Language (HDL), such as Verilog. Note that instead of a Verilog representation, circuit **100** may be described in any other HDL, such as VHDL, or in an internal representation (such as a graph

structure or a net list structure) in a programmed computer as will be apparent to the skilled artisan.

A designer of circuit **100** may additionally articulate a description of one or more assertions that monitor various signals in circuit **100** that normally occur during simulation. The assertions (also called “checkers”) are articulated to generate error signals when a certain combination of signals in circuit **100** cause a condition specified in the assertion to be violated during simulation. Assertions can receive signals from either or both portions **101** and **102** of circuit **100**, depending on the assertion.

FIG. 1A illustrates an assertion **103** that receives input signals on paths **108** and **107** respectively from each of the two clock domains **101** and **102**. Note that paths **107** and **108** are shown dashed in FIG. 1A to indicate that the paths are not necessarily present in a circuit description, e.g. assertion **103** may receive signals only on path **108** or only on path **107** or on both paths **107** and **108** depending on the circuit design and/or the assertion. For more information on assertions, see U.S. Pat. Nos. 6,175,946 and 6,609,229 granted to Ly et al that are incorporated by reference herein in their entirety.

During simulation of circuit **100** (FIG. 1A) with conventional RTL simulators, assertion **103** does not receive certain signals that result from the effects of metastability, because metastability is not modeled explicitly in prior art systems. In contrast, metastability effects are known to arise in physical implementations of circuit **100**, due to the difference in the two clock signals on paths **105** and **106**. Specifically, a physical register implemented in silicon, for example, the register in FIG. 1B that receives the clock-domain-crossing signal of FIG. 1A, is characterized by parameters called “setup time” and “hold time”. If a signal at the data input of the physical register changes logic values within the setup time before the active edge of the register’s clock signal, or within the hold time after the active edge of the register’s clock signal, then the output of the register becomes unpredictable, and may settle to either logic value 1 or logic value 0. For more information, see “Digital Systems Engineering,” Dally, W. J., and Poulton, J. W., Cambridge University Press, 1998, pp. 462–513.

A clock-domain-crossing signal on path **104** changes its logic value during the setup time or during the hold time of register **111** in the receive clock domain **102** due to the relative difference in times at which the two clock domains **101** and **102** are clocked by their respective clock signals on paths **106** and **105**. FIGS. 2A and 2B show representative electrical waveforms for the output of a physical register in the physical world that has been implemented in silicon (as an integrated circuit die), in situations where the clock-domain-crossing signal violates the setup time of this register **111**.

In FIG. 2A, a signal at the output of register **111** initially goes only part way to logic level 1 and then settles to logic level 0 whereas in FIG. 2B the same signal initially goes only part way to logic level 1 and then settles to logic level 1. Similarly, FIGS. 2C and 2D show the corresponding electrical waveforms when the hold time of the physical register **111** is violated and the output signal settles to logic level 1 and logic level 0 respectively. The logic level to which a signal settles in the physical world i at the output of a physical register **111** depends on a number of factors (such as thermal effects and/or voltages) that are not normally modeled in conventional RTL simulation.

FIGS. 3A and 3B show representative simulation waveforms produced by conventional RTL simulation of the circuit **100** of FIG. 1A in cases where a signal at the data

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input of a register **111** in the receive clock domain **102** violates the setup time and hold time parameters. As can be seen by comparing FIG. **3A** with FIGS. **2A** and **2B** and by comparing FIG. **3B** with FIGS. **2C** and **2D**, the electrical waveforms of the physical register may differ from the simulation waveforms produced by conventional RTL simulation when the setup or hold time parameter of the register is violated. Note that only one outcome is produced by the RTL simulator when the setup time is violated as shown in FIG. **3A**. Similarly only one outcome is produced when the hold time is violated as shown in FIG. **3B**. The outcome produced by the RTL simulator is also called the “correct” logic value, and the inversion of the outcome produced by the RTL simulator is also called the “incorrect” logic value.

In contrast, when a signal at the data input of a physical register in the physical world changes logic values within the setup time before the active edge of the register’s clock signal, then the signal at the output of the physical register in the physical world may settle to either a “correct” logic value (i.e., a value matching the value produced by conventional RTL simulation of the register), or an “incorrect” logic value (i.e., the inversion of the value produced by conventional RTL simulation of the register), as shown in FIGS. **2A** and **2B**. Similarly, two outcomes are possible when the signal changes within the hold time after the active edge of the register’s clock signal, as shown in FIGS. **2C** and **2D**.

An example circuit **400** shown in FIG. **4A** is similar or identical to the corresponding circuit **100** described above, except for the following differences. The reference numerals in FIG. **4A** are obtained from the corresponding reference numerals in FIG. **1A** by adding **300**. Circuit **400** includes multiple paths (e.g.  $n$  paths) in a bus **404** between the two clock domains **401** and **402**. In this example, the  $n$ -bit signal on bus **404** that crosses clock domains **401** and **402** happens to have been designed by the circuit designer to be one-hot, which satisfies the property that exactly one bit of the  $n$ -bit signal is asserted at all times during normal operation.

Note that in circuit **400** of FIG. **4A**, assertion **403** is coupled to only the receive clock domain **402** to receive therefrom a version of the  $n$ -bit signal after it has been clocked in by receive clock domain **402** (which receives this signal on path **404** from transmit clock domain **401**). Assertion **403** may be articulated by the designer of circuit **400** to be a one-hot assertion which checks that the signal on path **407** is in fact one hot (i.e. that exactly one bit of the  $n$ -bit signal is asserted at all times). Assertion **403** contains an XNOR gate **421** that receives signals  $RX\_Q\_1$  and  $RX\_Q\_0$  that are output by registers **411\_1** and **411\_0**. XNOR gate **421** supplies an error signal when its inputs are the same and this error signal is latched in a register **422** also included in assertion **403**. Note that assertion **403** is not connected to transmit clock domain **401** in this example although in other examples such an assertion may be connected to only transmit clock domain **401**, or to both clock domains.

An example of circuit **400**, for  $n=2$ , is described next, in reference to FIG. **4B**. Transmit clock domain **401** contains two registers that form a one-hot counter **412** (see registers **412\_1** and **412\_0**, together called “tx\_reg”). Counter **412** is clocked by the rising edge of the transmit clock signal  $TX\_CLK$ . When the reset signal  $RST$ , is asserted, register **412\_1** is set to 0 (deasserted) and register **412\_0** is set to 1 (asserted). At each rising edge of the transmit clock signal  $TX\_CLK$  after the reset signal  $RST$  is deasserted, the values stored in registers **412\_1** and **412\_0** are swapped. Therefore, the counter **412** (called “tx\_reg” which is a short form for “transmit register”) remains one-hot at all times after reset.

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In the example circuit of FIGS. **4B** and **4C**, the one hot signal from tx\_reg counter **412** (i.e. from registers **412\_1** and **412\_0**) is clocked into a counter **411** (formed by registers **411\_1** and **411\_0** that are together called “rx\_reg” which is a short form for “receive register”), on each rising edge of receive clock signal  $RX\_CLK$ . As described above, since input signal  $TX\_Q\_0$  is clocked into receiving register **411\_0** by a first clock signal ( $RX\_CLK$ ), transmitting register **412\_0** is in the combinational fanin of signal  $TX\_Q\_0$ , and register **412\_0** is clocked by a second clock signal ( $TX\_CLK$ ), it follows that signal  $TX\_Q\_0$  is a clock-domain-crossing (“CDC”) signal. Note that signal  $TX\_Q\_0$  transmitted by the transmit clock domain **401** on path **404\_0** is same as signal  $RX\_D\_0$  that is received by the receive clock domain **402** at the D input of register **411\_0**. In a similar manner, note that signal  $TX\_Q\_1$  is a CDC signal also, and is same as signal  $RX\_D\_1$  received at the D input of register **411\_1**.

A Verilog representation of circuit **400** of FIG. **4B** is shown in Appendix A which is located just before the claims in this patent application. Appendix A is an integral portion of this background section of this patent application and is incorporated by reference herein in its entirety. For a description of the Verilog language, see “The Verilog Hardware Description Language, Second Edition” Thomas, D. E., and Moorby, P. R., Kluwer Academic Publishers, 1995. In the Verilog of FIG. **4B**, registers **411\_1**, **411\_0**, **412\_1** and **412\_0** are represented as rx\_reg\_1, rx\_reg\_0, tx\_reg\_1, and tx\_reg\_0 respectively. Moreover, signal names shown in upper case letters in FIG. **4B** are replaced by corresponding names in lower case letters in Appendix A. Note that the initial state represented in the initial block of the Verilog shown in Appendix A corresponds to the reset state of the circuit under verification, i.e., tx\_reg\_1=0, tx\_reg\_0=1, rx\_reg\_1=0, and rx\_reg\_0=1.

As noted above, circuit **400** of FIG. **4B** contains assertion **403** to check that the value stored in the rx\_reg counter **411** is in fact one-hot (see lines **42–44** in Appendix A). The output of the one-hot assertion **403** becomes asserted when the assertion is “violated”, if and only if the value stored in rx\_reg counter **411** is not one-hot at the rising edge of the receive clock signal  $RX\_CLK$ . During conventional RTL simulation, a violation flagged by the one-hot assertion **403** indicates that the value of the rx\_reg counter **411** is not one-hot.

The just-described error in the rx\_reg counter **411** is treated by a circuit designer as an indication that an error occurred in the generation of the one-hot signal but not that the one-hot signal was corrupted during transmission across clock domains. This is because conventional RTL simulators such as VCS and NC Verilog do not accurately model metastability affecting the CDC signals. Therefore, during conventional RTL simulation of the example circuit **400** of FIG. **4B**, the value of the tx\_reg counter is modeled as being correctly transmitted to the rx\_reg counter, regardless of the violation of set up times (of registers **411\_0** and **411\_1**). For this reason, when the one-hot signal is correctly generated and stored in the tx\_reg counter **412** the one-hot assertion **403** that monitors the signal on path **107** is not violated during conventional RTL simulation.

As noted above, RTL simulation in the conventional manner produces only one outcome (i.e. one logic level) in the event of a setup time violation although two outcomes are possible. Moreover, RTL simulation also produces only one outcome (i.e. one logic level) in the event of a hold time violation, although two outcomes are possible. The inventors believe there is a need to take into account the outcomes

that are not conventionally produced by RTL simulation. Specifically, the inventors believe that explicit modeling of all outcomes could lead to detection of errors that are not otherwise detected by RTL simulation.

Incorporated by reference herein in its entirety as background is an article entitled “Using Assertion-Based Verification to Verify Clock Domain Crossing Signals” by Chris Ka-Kei Kwok, Vijay Vardhan Gupta and Tai Ly presented at Design and Verification Conference (DVCon 2003), February, 2003.

#### SUMMARY OF THE INVENTION

Prior to verification of a description of a circuit containing a pre-determined assertion, the circuit description is automatically transformed in accordance with the invention by addition of description(s) of one or more circuits (also called “metastability injectors”) to deliberately create one or more effects of metastability in the circuit. The transformed description (containing metastability injectors) is verified in the normal manner. Therefore, in some embodiments, use of metastability injectors during verification results in detection of incorrect behavior of the circuit (if present) that is caused by metastability in signals that cross clock domains in the circuit. Note that a circuit may be described in accordance with the invention (and the circuit description can be stored and used) in a programmed computer either in the form of HDL (such as Verilog or VHDL) or as an internal representation (such as a graph or a netlist).

During verification, certain embodiments analyze the transformed description using a formal verification method (such as model checking or bounded model checking) to identify one or more specific stimulus sequence(s) that will cause the pre-determined assertion to be violated in simulation. One specific stimulus sequence identified by formal verification is used in simulation of the transformed circuit, to display to the circuit designer one or more simulation waveforms (on a computer screen) that indicate an incorrect behavior of the circuit in the presence of metastability. The circuit designer may analyze such simulation waveforms, to determine one or more sources of error, and if necessary change the circuit description to eliminate the incorrect behavior in a future iteration of verification (in the above-described manner). The designer may change the circuit description in any manner, including but not limited to implementing protocols to correctly transmit information between clock domains in the presence of metastability.

#### BRIEF DESCRIPTION OF THE DRAWINGS

FIGS. 1A–1C illustrate, in block diagrams a prior art circuit with two clock domains, one CDC signal and one assertion.

FIG. 2A shows representative electrical waveforms for physical register 111 of FIG. 1B showing violation of the setup-time parameter followed by the Q output entering the metastable state followed by the Q output settling to the incorrect logic value.

FIG. 2B shows representative electrical waveforms for physical register 111 of FIG. 1B showing violation of the setup-time parameter followed by the Q output entering the metastable state followed by the Q output settling to the correct logic value.

FIG. 2C shows representative electrical waveforms for physical register 111 of FIG. 1B showing violation of the hold-time parameter followed by the Q output entering the metastable state followed by the Q output settling to the incorrect logic value.

FIG. 2D shows representative electrical waveforms for physical register 111 of FIG. 1B showing violation of the hold-time parameter followed by the Q output entering the metastable state followed by the Q output settling to the correct logic value.

FIG. 3A shows simulation waveforms from conventional RTL simulation of a model of register 111 of FIG. 1B showing violation of the setup-time parameter and the resulting Q output produced.

FIG. 3B shows simulation waveforms from conventional RTL simulation of a model of register 111 of FIG. 1B showing violation of the hold-time parameter and the resulting Q output produced.

FIG. 4A shows, in a block diagram, an example circuit having N CDC signals and a one-hot assertion.

FIG. 4B shows, in a detailed circuit diagram, the circuit of FIG. 4A having two CDC signals.

FIGS. 5A and 5B illustrate, in alternative embodiments, a metastability injector added in accordance with the invention, to the circuit-under-verification of FIG. 1A, to obtain a transformed circuit-under-verification.

FIG. 5C illustrates, in a lower level block diagram, a transition detector and a conditional inverter that are included in a metastability injector of some embodiments of the invention.

FIG. 5D illustrates, in a graph, waveforms of various signals in a metastability injector of one embodiment.

FIGS. 5E and 5F illustrate two exemplary implementations of a metastability injector of some embodiments.

FIG. 6A illustrates, in a flowchart, one embodiment of the method of the invention.

FIG. 6B illustrates, in a high-level block diagram, a flow of information when performing the method of FIG. 6A.

FIG. 6C illustrates, in a flowchart, one exemplary implementation of an act of transformation 610 in FIG. 6A.

FIG. 6D illustrates, in a flowchart, one exemplary implementation of an act of analysis 620 in FIG. 6A.

FIG. 6E illustrates, in a flowchart, one exemplary implementation of an act of simulation 650 in FIG. 6A.

FIG. 7 illustrates a transformed circuit obtained by performing the method of FIG. 6C on the circuit of FIG. 4B.

FIG. 8 illustrates a display on a computer screen of waveforms from simulation of a description (in Verilog) of the circuit of FIG. 7 while applying a stimulus sequence in accordance with the invention.

#### DETAILED DESCRIPTION

In accordance with the invention, a description of a circuit-under-verification (“CUV”) is automatically transformed so that it explicitly models the effects of metastability, resulting in a transformed description (hereinafter, also called the “transformed CUV”). The transformed description may be verified in any manner. Specifically, an original description of the CUV (which may be prepared by a circuit designer in the normal manner) is transformed, by insertion of an extra circuit to inject metastability effects into the path of a clock-domain-crossing signal. The extra circuit (also called “metastability injector”) has an enable input that is used to conditionally inject metastability effects into the transformed CUV.

FIG. 5A shows the result of transforming the original circuit 100 of FIG. 1A by adding a metastability injector 508. Note that the transmit clock domain 101 and receive clock domain 102 shown in FIG. 5A are similar or identical to the respective clock domains shown in FIG. 1A. Moreover, assertion 103 in FIG. 5A may also be same as or

similar to an assertion **103** that is already pre-existing in an original circuit **100** (see FIG. 1A). Note that in some embodiments, a circuit designer may pre-determine (and optionally add) one or more assertions **103** for use with circuit **100**, with such newly added assertions being intended to specifically detect errors resulting from metastability effects in the signal received on an input path **507B** of the receive clock domain **102**.

Note that in many embodiments, an assertion **103** that is violated as described herein is not deliberately selected but has only an indirect relationship to metastability (e.g. if the assertion is connected to the output of a register that is several sequential stages removed from the entry point of the CDC signal in the receive clock domain **102**). Furthermore, in several embodiments, assertion **103** may be any assertion that monitors a portion of circuit **100** located in the transitive sequential fanout of the signal received on an input path **507B** of the receive clock domain **102**. Note that the transitive sequential fanout of signal S is a set of registers, R, constructed as follows: (a) set the set R to contain all registers with inputs in the combinational fanout of S; (b) repeat the following until the set R does not grow any larger: for each register X in the design, if X is not already in R and an input of X is in the combinational fanout of some register in R, then add X to R.

In the embodiment of FIG. 5A, path **507A** originates in the transmit clock domain **101** in the manner described above in reference to path **104** of FIG. 1A. Note that the clock-domain-crossing (CDC) signal which is output by transmit clock domain **101** does not travel on path **507A** always unaltered to the receive clock domain **102** in FIG. 5A. Instead, the CDC signal on the path **507A** of FIG. 5A is modified by metastability effects that are output by metastability injector **508** on a path **509** that is connected to path **507A**. Input path **507B** of the receive clock domain **102** is connected to paths **507A** and **509**. Hence, in the circuit **500** of FIG. 5A, path **507B** carries the modified CDC signal (also called ‘metastable CDC signal’) to the receive clock domain **102**. Note that the above-described combinational logic **199** may be present between paths **507A** and **507B** at any location relative to path **509** (i.e. although in FIG. 5A path **507A** is shown passing through logic **199** and path **509** is directly connected to path **507B**, other embodiments may have path **507B** passing through combinational logic **199** with path **509** being directly connected to path **507A**, while still other embodiments may have both paths **507A** and **507B** passing through different portions of combinational logic **199**). Therefore, the specific connections among paths **507A**, **507B**, **509** and combinational logic **199** are different depending on the embodiment.

Note that metastability injector **508** may add any kind of metastability effect to the CDC signal generated by transmit clock domain **101**, depending on the embodiment. In some embodiments, metastability injector **508** simply inverts the CDC signal from path **507A**, whenever there is a transition in the CDC signal. An inversion forced by metastability injector **508** may be disabled (so the result is same as in RTL simulation) or the forced inversion may be timed to happen at various times relative to the set up and hold times of the receiving register (not shown in FIG. 5A; see register **111** in FIG. 1A).

Specifically, metastability injector **508** may be disabled (by de-asserting an enable signal on path **505**) all the time, in which case the CDC signal is left unaltered. Alternatively, metastability injector **508** may be disabled only until it becomes time for the transmit clock signal on path **106** to align with the receive clock signal on path **105** at which time

metastability injector **508** is enabled. In some embodiments, the enable signal on path **505** is output by an AND gate (not shown) that receives as input a signal that is asserted during alignment of the two clock signals, and as another input a signal indicating that injector **508** is activated. Note that AND gate **524** is a 3-input gate that additionally receives the signal that is asserted during alignment (in addition to the signal on path **505** and the output of gate **523**). The times at which two clock signals are considered to be aligned, depends on the particular embodiment. For example in some embodiments, the clock signals are considered to be aligned if the time between the rising edge of the transmit clock and the rising edge of the receive clock is less than the setup time of the receiving register, or if the time between the rising edge of the receive clock and the rising edge of the transmit clock is less than the hold time of the receiving register.

For example, in some embodiments, whenever there is a transition in the CDC signal (assuming it happens when the clocks are considered to be aligned), the modified CDC signal that is presented at the input of the receive clock domain **102** is obtained by the metastability injector inverting the CDC signal (i.e. the logic value is driven from 1 to 0 and from 0 to 1). Thus, when the next active edge of the receive clock occurs in the receive clock domain **102** (FIG. 5A), the value that is stored in the receiving register in clock domain **101** models the situation in which the physical receiving register (e.g. register **111** in FIG. 5C) enters the metastable state in the physical world and settles to the inverse of the logic value that would be produced by conventional RTL simulation of the non-transformed circuit.

In many embodiments of the type described herein, metastability injector **508** may be disabled by de-asserting a signal on a path **505** (FIG. 5A). Specifically, when the signal (also called ‘metastability enable’) on path **505** is asserted, the metastability injector of such embodiments forces the receive clock domain **102** to clock a modified CDC signal into the receiving register **111**. The modified CDC signal is either the unaltered version of the CDC signal when there is no transition in the CDC signal, or the inverted version of the CDC signal if there is a transition in the CDC signal.

When the metastability enable signal on path **505** is de-asserted, the metastability injector of such embodiments is disabled and hence it unconditionally allows the receive clock domain **102** to receive the unaltered version of the CDC signal (i.e. regardless of whether or not a transition is happening in the CDC signal). Thus, when an active edge of the receive clock occurs with the enable signal on path **505** deasserted, the value stored in the receiving register of receive clock domain **102** models the situation in which the physical receiving register enters the metastable state and settles to the same logic value as would be produced by conventional RTL simulation of the non-transformed circuit.

Note that a metastability enable signal of the type described above in reference to path **505** does not exist in the original description of circuit **100** (FIG. 1A), but this signal is used as a primary input during verification (e.g. in formal verification) to turn on and off metastability effects. Use of such a metastability enable signal introduces one primary input into the formal analysis for each clock-domain-crossing signal (i.e. over and beyond any primary inputs present in the original description of circuit **100**). Therefore, if transformed circuit **500** has an n-bit bus that connects transmit clock domain **101** to receive clock domain **102**, then ‘n’ enable inputs are now present, to conditionally turn on and off the metastability effects in each of the ‘n’ CDC signals. Transformed circuit **500** (FIG. 5A) models each of the two possible outcomes for receipt of each CDC signal in

a receiving register in clock domain **102** to model metastability effects, depending on whether the enable signal on path **505** is asserted or deasserted.

Note that, if assertion **103** is found to be not violated regardless of whether one or more metastability injector(s) **508** are enabled or disabled, then the design of circuit **500** is deemed to be verified to withstand metastability effects. On the other hand, if assertion **103** is violated when one of the metastability injector(s) **508** is enabled, the circuit designer may re-design circuit **500** to withstand metastability effects. Note that due to a change in the path between the two clock domains **101** and **102**, the clock domains **101** and **102** in FIG. **5A** are together referred to as circuit **500** (instead of circuit **100** as per FIG. **1A**).

Certain embodiments of metastability injector **508** that are responsive to a transition in the clock-domain-crossing (CDC) signal, may detect the transition, inter alia, by use of one or more signals on path **506** from the transmit clock domain **101** or one or more signals on path **503** from the receive clock domain **102** or signals on both paths **503** and **506**. Some embodiments of metastability injector **508** that are responsive to the transition in the CDC signal do not use any additional signals from clock domains **101** and **102**, and instead directly monitor the CDC signal alone, to detect the transition. In the just-described embodiments, the metastability injector **508** does not have paths **503** and **506**. The specific circuitry to be used in such a metastability injector **508** will be apparent to the skilled artisan in view of this detailed description

In another embodiment which is illustrated in FIG. **5B**, a path **504A** that carries the clock-domain-crossing (CDC) signal from the transmit clock domain **101** is not directly connected to a path **504B** that passes the modified CDC signal to the receive clock domain **102**. Instead, path **504A** terminates in metastability injector **510**, and another path **504B** originates in the metastability injector **510**. The modified CDC signal is supplied on path **504B** by injector **510**, and this signal includes metastability effects therein

In other embodiments, which are not shown, path **504B** of metastability injector **510** may be coupled to the transmit clock domain **101** (e.g. to insert metastability effects into an earlier version of the CDC signal between the additional circuitry **192** and the last register **112** in the transmit clock domain **101** of FIG. **1C**). In still other embodiments, which are also not shown, path **504B** may be coupled between the additional circuitry **191** and the first register **111** in the receive clock domain **102** of FIG. **1B** (e.g. to insert metastability effects into a later version of the CDC signal).

Metastability injector **510** of FIG. **5B** can be implemented in any manner depending on the embodiment. In some embodiments, metastability injector **510** enabled by the signal on path **505**, uses two versions of the clock-domain-crossing (CDC) signal to detect if a transition is happening in the CDC signal on path **504A**. The two versions of the CDC signal may be, for example, the CDC signal on path **504A** and an early version of the CDC signal obtained from the input of last register **112** in the transmit clock domain **101** (e.g. from the data input of the transmitting register in clock domain **101**, i.e., the TX\_D signal).

Such an early version of the CDC signal may be obtained via the above-described path **506** (FIG. **5C**) from transmit clock domain **101**. In such embodiments, the early version of the CDC signal is compared with the CDC signal itself, in circuitry **512** (called “transition detector”) that is located within metastability injector **510**. Any difference between the two versions of the CDC signal is indicated by transition detector **512** asserting a signal at its output on a path **513**

(FIG. **5C**). Note that in the embodiments illustrated in FIG. **5C**, the enable signal on path **505** is supplied directly to transition detector **512** to enable or disable its operation (when disabled, the signal on path **513** is deasserted regardless of transitions on path **504A**). Note also that although in some embodiments transition detector **512** has been described and illustrated as comparing two versions of the CDC signal (i.e. the logic values on paths **506** and **504A**), in other embodiments the transition detector may use only the CDC signal itself as noted above.

In some embodiments, transition detector **512** compares a current version of the CDC signal with an early version of the CDC signal, to detect whether a transition is going to happen in the CDC signal at the next clock cycle. The signal generated by transition detector **512** on path **513** (also called CDC transition) is illustrated in FIG. **5D** in the case of a CDC signal on path **504A** (FIG. **5C**) that is low at time **T0**, goes high at time **T1** and stays high for two clock cycles and goes low at time **T3**. As shown in FIG. **5D**, the early version of the CDC signal on path **506** exhibits the same behavior as the CDC signal on path **504A** but it is shifted earlier by one clock cycle. The CDC transition signal on path **513** goes high for one clock cycle before the CDC signal goes high i.e. between times **T0** and **T1**. The CDC transition signal on path **513** also goes high for one clock cycle before the CDC signal goes low i.e. between times **T2** and **T3**.

Referring to FIG. **5C**, the CDC transition signal on path **513** controls a circuit **511** that is also included in metastability detector **510** of these embodiments. Circuit **511** (also called “conditional inverter”) conditionally supplies on path **504B** (to the receive clock domain **102**), either an inverted version or an unaltered version of the CDC signal received on path **504A** (from the transmit clock domain **101**), depending on whether or not the signal on path **513** is asserted. As noted above, the signal on path **513** is asserted by the transition detector **512** whenever there is a transition in the CDC signal on path **504A**.

In the above-described example, the CDC transition is high between times **T0** and **T1** (as shown in FIG. **5D**) which causes the modified CDC signal on path **504B** to go high between times **T0** and **T1** (inverse of the low value in the CDC signal between times **T0** and **T1**). Note that at time **T1**, the CDC transition signal becomes low and hence the modified CDC signal goes high (due to pass-through of the high value in the CDC signal between times **T1** and **T2**). The remaining transitions at times **T2** and **T3** in the modified CDC signal are shown in FIG. **5D** and their behavior will be apparent to the skilled artisan in view of this detailed description.

Depending on the embodiment, an early version of the CDC signal for use in a transition detector **512** as described above may be obtained from an input of any storage element in the transmit clock domain **101**, in the transitive sequential fanin of CDC signal. Transitive sequential fanin of the CDC signal is consistent with use of this term in art, i.e. a set of registers, **R**, constructed as follows: (1) set the set **R** to contain all registers with outputs in the combinational fanin of **S**; (2) repeat the following until the set **R** does not grow any larger: for each register **X** in the design, if **X** is not already in **R** and the output of **X** is in the combinational fanin of some register in **R**, then add **X** to **R**. As noted above, some embodiments use as the early CDC signal a signal that is received from additional circuitry **192** (FIG. **1C**) at the input of the very last storage element **112** in the transmit clock domain **101** that supplies the CDC signal.

Also note that in other embodiments, instead of an early version of the CDC signal, a later version of the CDC signal

may be used in a transition detector in a manner identical to that described above (although the transition detection will occur later). As noted above, depending on the embodiment, a transition detector **512** in metastability injector **510** may use only the CDC signal itself as input (instead of two versions of the CDC signal).

FIGS. **5E** and **5F** illustrate two alternative embodiments of a metastability injector **510** of the type illustrated in FIG. **5C**, although numerous such embodiments will be apparent to the skilled artisan in view of this detailed description. Accordingly several features illustrated in FIGS. **5E** and **5F** are merely educational and are not intended to limit the scope of the invention. In both embodiments illustrated in FIGS. **5E** and **5F**, a transition detector **512** is implemented by an exclusive OR gate **523** that receives the two versions of the clock-domain-crossing (CDC) signal on paths **504A** and **506**, and an output signal from gate **523** is supplied to an AND gate **524** that also receives as input the above-described metastability enable signal on path **505**. The signal output by AND gate **524** is supplied as the CDC transition signal on path **513**.

Note, however, that conditional inverter **511** of metastability injector **510** is implemented differently in the two embodiments illustrated in FIGS. **5E** and **5F**, as noted next. Specifically, one conditional inverter **511A** which is illustrated in FIG. **5E** has a multiplexer **521** that is controlled by the CDC transition signal on path **513** to supply one of two inputs to path **504B**, namely either the CDC signal directly from path **504A** or an inverted form of the CDC signal generated by an inverter **522** (that is also coupled to path **504A**). Another conditional inverter **511B** which is illustrated in FIG. **5F** has an exclusive OR gate **531** that receives as its two inputs, the CDC transition signal on path **513** and the CDC signal directly from path **504A**. The output of the exclusive OR gate **531** is directly supplied as the modified CDC signal on path **504B**.

Many alternative embodiments of the metastability injector will be apparent to a person skilled in the art, including embodiments that use signals from the transmit clock domain other than the CDC signal and the TX\_D signal at the “D” input of the transmitting register **112** (see FIG. **1C**).

A verification method **600** used in some embodiments of the invention is illustrated in FIG. **6A**. Specifically, a description **601** (e.g. expressed in Verilog, or VHDL) of the circuit under verification (see FIG. **6B**) is automatically transformed by programmed computer **602** as per act **610** (FIG. **6A**), by adding description of circuitry to inject effects of metastability into the circuit under verification, resulting in a transformed description **603**. Specifically, one or more metastability injectors of the type described above in reference to FIGS. **5A** and **5B** may be added by computer **602** in act **610**, to obtain the transformed description **603**. Note that the metastability injectors are added for each CDC signal in circuit description **601**.

Each CDC signal in description **601** is found automatically as follows by computer **602** that is appropriately programmed as follows. Computer **602** looks at each register in description **601** and checks if the register’s combinational fanin contains another register and if so, whether these two registers have different clock signals. If they do have different clock signals, then the signal between the two registers is deemed to be a CDC signal. Next, a metastability injector is inserted in the manner described herein, for the just-found CDC signal.

Note that some embodiments build a netlist from the description **601**, and traverse the netlist for each register, to find all registers that drive the data input of the current

registers and if any of these registers are clocked by a different clock then the path between the two registers with different clocks is a CDC signal. Note that only combinational logic (in terms of logic elements) separates these two registers with different clocks.

Note that circuit description **601** may or may not contain one or more pre-determined assertion(s) of the type described above, depending on the embodiment. For example, in some embodiments, circuit description **601** does contain pre-determined assertions and these assertions remain unchanged in the transformed description **603**. In other embodiments, circuit description **601** does not contain pre-determined assertions and instead these assertions are held in a separate file, and they are added to the circuit description from the separate file after addition of metastability injectors as described above in reference to act **610**. Note that regardless of when added, transformed description **603** contains one or more pre-determined assertions and one or more metastability injectors.

The transformed description **603** is analyzed by a computer **605**, as per act **620**, using any method well known in the art. In many embodiments, act **620** involves performance of a formal verification method (such as bounded model checking in some particular embodiments). Note that although computer **605** is used in some embodiments to perform a formal verification method on description **603**, act **620** may be performed in other embodiments by computer **602** (that performed act **601**), or act **620** may even be performed manually in still other embodiments.

Note that when the same computer **602** performs both acts **610** and **620**, in some embodiments a transformed circuit description **603** is an internal representation (e.g. in the form of a graph) of the circuit **100** and is directly transformed by addition of metastability injectors as per act **610** and the resulting transformed internal representation is used directly during analysis in act **620**.

If the analysis in act **620** (FIG. **6A**) finds a stimulus sequence that will cause the assertion in the transformed circuit description **603** (FIG. **6B**) to be violated, during conventional RTL simulation of the transformed description **603** (as per act **630** in FIG. **6A**), then a metastability problem has been found. Model checking in act **620** may use a Boolean formula that is TRUE if and only if the underlying state transition system can realize a sequence of state transitions that reaches certain states of interest within a fixed number of transitions. If such a sequence cannot be found at a given length, *k*, the search is continued for larger *k*. If a limit “L” is reached without finding a sequence, the circuit designer may conclude that there are no errors sourced from metastability effects and therefore use the circuit description (without metastability injectors) to fabricate (as per act **641**) an integrated circuit die **607**, which may be tested and used in the normal manner.

In act **620** if a stimulus sequence to violate the assertion is found, the “yes” branch in act **630** (FIG. **6A**) is taken, and act **650** is performed by computer **605** in some embodiments or by computer **602** in other embodiments. In act **650**, the transformed circuit description **603** is simulated using the stimulus sequence determined in act **620** (FIG. **6A**).

Next, in act **660**, the simulation waveforms are displayed on a computer screen (e.g. screen **605** in FIG. **6B**) so that a circuit designer **606** may diagnose the metastability problem. The circuit designer **606** may then revise their original circuit description as per act **670**, and the revised description is optionally again subjected to acts **610–660** (e.g. act **610** is performed again by computer **602** this time on the revised



circuit description, to generate another transformed description followed by analysis as per act 620 and so on).

If the model checking performed in act 620 does not find any stimulus sequence that will cause the assertion to be violated in conventional RTL simulation of the transformed circuit description 603, then act 640 is performed, and the circuit description 601 (or the revised circuit description) is deemed to not have a metastability problem (and a message to this effect is displayed on the computer screen).

FIG. 6C illustrates one embodiment of an implementation of the transformation act 610 of FIG. 6A that automatically transforms the description of the circuit under verification by adding metastability injectors. Specifically, in act 611, a path of a clock domain crossing signal is found in the circuit description 601 that is to be transformed, and this path is set as the current path.

Next in act 612, the current path is replaced by (a) an input path to a metastability injector, (b) the metastability injector itself, and (c) an output path from the metastability injector. Additional connections that may be required, depending on the internal design of the metastability injector are also made in act 612, as appropriate. For example, a path carrying the TX\_D signal which is connected to the D input of register 112 (FIG. 5C) that drives the clock-domain-crossing (CDC) signal is determined and this path is connected (as per act 612 in FIG. 6C) to path 506 of metastability injector 510 (FIG. 5C).

Next, in act 613, the metastability injector 510 itself is inserted into the path of the CDC signal. Specifically, an input path 504A of the metastability injector 510 is connected to the Q output of register 112 (FIG. 5C). Finally, an output path 504B of metastability injector 510 is connected to the D input of the register of the receive clock domain 102 (see register 111 in FIG. 5C). Next, in act 614, a check is made as to whether any more CDC signals exist and if so, the path of the next CDC signal is set as the current path and control returns to act 612 (described above). If there no more CDC signals, then the transformation is completed (as per act 616 in FIG. 6C).

FIG. 6D illustrates an implementation of the analysis act 620 of FIG. 6A in some embodiments. In these embodiments, a model checking method is used to analyze the transformed circuit description 603 resulting from act 620 (or the transformed circuit model as noted above) to find a stimulus sequence that will cause the assertion to be violated during conventional RTL simulation of the transformed circuit description. In such embodiments, an initial state for the model checking analysis may be set (as per act 621) to a reset state (which may be any one of several reset states) that the transformed circuit under verification (also called “transformed CUV”) may have.

Depending on the circuit design, one of the reset states may be for all registers to be set to logic level zero, whereas other reset states may be for one or more of the registers to be set to logic value 1 while all other registers are set to logic value 0, or for some registers to be set to a state representing “don’t care” (i.e., the register can be assigned either logic value 0 or logic value 1 during formal analysis). Note that an initial state for use in model checking in act 620 may be manually selected by a user to be any state. Alternatively, an initial state may be obtained from test-benches used in simulation (e.g. in a commercially available simulator such as VCS from Synopsys, Mountain View, Calif.).

Next, a cycle identifier I is set to 1 in act 622 and control is transferred to act 623. In act 623, the behavior of the transformed circuit is analyzed for all stimulus sequences I cycles in length, starting from the initial state. As noted

above, in act 622 the cycle identifier was set to 1 and therefore the analysis in this first iteration is for only 1 cycle in length, although in later iterations that reach act 623 from act 628 the analysis becomes deeper (if no assertion is violated).

Then, in act 624, a check is made to see if a stimulus sequence is found that will cause the assertion to be violated, and if so the yes branch is taken and act 625 is performed. Specifically, the model checking method is concluded and the stimulus sequence is returned along with the current cycle (e.g. variable “LI” is set to I). If in act 624, the stimulus sequence is not found, then control is transferred to act 626. In act 626, a check is made as to whether a predetermined limit L on the cycle identifier I has been reached and if not then I is increased by one, and the process is iterated (returning to act 623). If the predetermined limit L was reached, then the model checking is concluded in act 627, and returns with no stimulus found.

Many alternative methods of selecting an initial state for the model checking method (as per act 621 in FIG. 6D) will be apparent to a person skilled in the art, including: selecting an initial state that is a non-reset state of the circuit under verification (“CUV”); selecting an initial state from simulation of the; selecting an initial state by analyzing waveforms from simulation of the CUV; selecting an initial state from simulation such that at least two simulated clock edges are aligned to allow metastability to occur according to the setup and hold time parameters of a register in the; and using a programmed computer to automatically determine an initial state. Similarly, it will be apparent to a person skilled in the art that an initial state for the model checking method can represent all reachable states of the CUV.

Many alternative embodiments of the model checking method 620 will be apparent to a person skilled in the art in view of this detailed description. Several such embodiments use one of the model checking methods described in “Model Checking”, E. Clarke, O. Grumberg, and D. Peled, MIT Press, 1999, and in “Bounded model checking using satisfiability solving,” E. Clarke, A. Biere, R. Raimi, and Y. Zhu, Formal Methods in Systems Design, 19(1):7–34, 2001 in place of the model checking method 620 shown in FIG. 6D. The just-described two papers are incorporated by reference herein in their entirety.

FIG. 6E shows one embodiment of a simulation act 650 of FIG. 6A. Specifically, in act 651, the initial state of the transformed CUV is set to the initial state used by the model checking step, and in act 651 the value of I is set to one. Note that the same cycle count I that was used in analysis act 620 is now used for simulation in act 650. Then, in act 653, stimulus for cycle number I from the model checking step of FIG. 6A is applied to the inputs of the transformed CUV and in act 654 cycle number I of operation of the transformed CUV is simulated. If I is less than a limit set by the variable LI returned from model checking, then I is increased by one and the process is iterated, otherwise the simulate step of FIG. 5A is complete. Note that LI is the cycle in which the model checking method finds the assertion to have been violated.

Performance of method 600 (FIG. 6A) on the example circuit of FIG. 4 is now described for some embodiments of the invention. Although in FIG. 7 the circuitry being inserted is shown by way of drawings, in several embodiments a Verilog description of the circuit under verification (“CUV”) shown in Appendix A is transformed by act 610 into the corresponding Verilog description of the transformed CUV in Appendix B (which is also located just before the claims). Note that Appendix B is an integral portion of this detailed

description of an embodiment of the invention, and is incorporated by reference herein in its entirety. Note also that in some embodiments, the Verilog description of Appendix A is transformed by act 610 into an internal representation equivalent to the Verilog description of Appendix B.

In act 610, a first metastability injector 701 (FIG. 7) is inserted in the path of CDC signal RX\_D\_0 of FIG. 4B, and a second metastability injector 702 (FIG. 7) is inserted in the path of the CDC signal RX\_D\_1 of FIG. 4B. Specifically, the circuitry in FIG. 4 is changed as follows, starting with circuit description 601. A module for the metastability injector in Verilog is copied and placed into description 601 twice, once for injector 701 and again for injector 702. When adding injectors 701 and 702 into the description 601, signals are appropriately renamed and/or the signal names are used in appropriate places for the injectors to become inserted into the paths of the CDC signals.

For example, whatever the CDC signal names are (such as signal names TX\_Q\_0 and TX\_Q\_1) these same names are used as the names of the CDC signal input to the respective injectors 701 and 702 (e.g. in injector 701 name TX\_Q\_0 may be used at each of (a) multiplexer input, (b) inverter input, and (c) XOR gate input). Moreover, whatever signal names are present at the data input of the transmitting registers (e.g. signal names TX\_D\_0 and TX\_D\_1) these names are used as the names of the early CDC signals at the respective injectors 701 and 702 (e.g. signal name TX\_D\_0 is used as a second input of the XOR gate). Finally, the names of signals that are output by injectors 701 and 702 are used as the signals input to receiving registers RX\_REG\_0 and RX\_REG\_1 in receive clock domain 102 (instead of the names of the CDC signals that were originally present in circuit description 601). In Appendix B, there are two new inputs in the transformed circuit description 603 that were not previously present in Appendix A, namely jitter\_control\_0 and jitter\_control\_1 which respectively represent two enable signals for the two metastability injectors 701 and 702. Note that one additional input for alignment between the receive and transmit clocks is not used in this embodiment (whose output is shown in Appendix B), although such an additional input is used in other embodiments.

Transformations of the type described in the previous paragraph, to add metastability injectors to a circuit description 601 can be done either directly in the Verilog language, or alternatively the transformations can be done on a schematic which is then translated into Verilog language. Moreover, such transformations can be done automatically in a computer 602 or alternatively the transformations can be done manually.

Note that when there are multiple metastability injectors in a transformed CUV, the enable signal of each metastability injector may be turned on or off independent of the other metastability injectors. Furthermore, even in the case of an “n” bit bus 704 whose signals are all stored in a single “n” bit register in a single device, note that each path for each bit in bus 704 has its own metastability injector, and each metastability injector may be independently enabled (so that each bit in the “n” bit register is made metastable independent of any other bit in the “n” bit register).

In some embodiment of the invention, the Verilog description in Appendix B is analyzed as per act 620 (FIG. 6A) by a computer programmed with a model checking program called “VIS” that is available via the Internet at “www-cad” dot “eecs” dot “berkeley” dot “edu” slash “Respep” slash “Research” slash “vis”, wherein the word “dot” should be replaced by “.” and the word “slash” is to be replaced by “/” to form the “http” address of a web page at the University

of California, Berkeley. Note that the VIS system is also described in an article entitled “VIS: A System for Verification and Synthesis”, The VIS Group, In the Proceedings of the 8th International Conference on Computer Aided Verification, p428–432, Springer Lecture Notes in Computer Science, #1102, Edited by R. Alur and T. Henzinger, New Brunswick, N.J., July 1996, which is incorporated by reference herein in its entirety.

Specifically, the VIS system is used to analyze the transformed CUV in Appendix B to determine stimulus to apply to the inputs of the transformed CUV during simulation of the transformed CUV using a conventional RTL simulator such as VCS or NC Verilog in order to violate the assertion. As described above, violation of the assertion during RTL simulation of the transformed CUV indicates that metastability in the physical CUV may cause incorrect behavior of the physical CUV.

In order to use the model checking method of the VIS system to determine the stimulus sequence to apply to the inputs of the transformed CUV in order to violate the assertion as per act 620 in FIG. 6A, the Verilog representation of the transformed CUV shown in Appendix B is placed in a file named “translate.v” and provided as input to a sequence of VIS system commands shown below:

```

25 vi2mv -c -F translate.v
   read_blif_mv translate.mv
   flatten_hierarchy
   static_order
   build_partition_mdds
30 check_invariant -f -d 1 -i -v 2 invar

```

The sequence of VIS system commands “vi2mv”, “read\_blif\_mv”, “flatten\_hierarchy”, “static\_order” and “build\_partition\_mdds” shown above create an internal representation of the transformed CUV in preparation for model checking. The VIS system command “check\_invariant” shown above performs model checking on the internal representation of the transformed CUV. The file “invar” in the VIS system command “check\_invariant -f -d 1 -i -v 2 invar” in the above set of commands contains a line “error=0”, directing the model checking program of the VIS system to find a counterexample for the invariant “error=0”, i.e., to find a counterexample for the one-hot checker in the transformed CUV.

In response to the sequence of commands shown in the previous paragraph, the model checking program of the VIS system produces an output file shown in Appendix C (which is located below, just before the claims). Appendix C forms an integral portion of this detailed description of some embodiments of the invention, and is incorporated by reference herein in its entirety. The output file shown in Appendix C represents the stimulus sequence to apply to the inputs of the transformed CUV during simulation using a conventional RTL simulator such as VCS or NC Verilog, starting from the reset state of the CUV, to violate the invariant “error=0”, i.e., to violate the one-hot assertion in the transformed CUV.

Thereafter, as per act 650, the VCS simulator is used to simulate the transformed CUV along with the stimulus sequence shown in Appendix C as input, starting from the reset state of the CUV. In addition, as per act 660, waveforms from the simulation are displayed on a computer screen (shown in FIG. 8), using a waveform viewer such as SimVision from Cadence Design Systems, Inc.

In the simulation waveforms shown in FIG. 8, at time 4800 ps, jitter control input jitter\_control\_0 of the transformed CUV becomes asserted, modeling the case in which

the physical register rx\_reg\_0 enters the metastable state at the rising edge of rx\_clk at time 5000 ps and then settles to the inverse of the logic value that would be produced by conventional RTL simulation of the non-transformed circuit, thus violating the one-hot assertion and causing the error signal to become asserted at time 15000 ps. As described above, violation of the one-hot assertion indicates that metastability in the physical CUV may cause incorrect behavior of the physical CUV.

Although, for illustrative purposes, the example circuit shown in FIG. 7 is small, containing only two clock signals and five registers, the method described above for the example circuit is also applied in the same manner to large circuits, for example, circuits containing hundreds of millions of registers and hundreds of thousands of clock signals. Numerous modifications and adaptations of the embodiments described herein will be apparent to a person of skill in the art of electronic design automation (EDA) in view of this disclosure.

Other embodiments of a method in accordance with the invention include one or more of the following steps: (1) automatically transforming a description of a CUV containing a predetermined assertion that is automatically inferred; (2) automatically transforming a description of a CUV containing a pre-determined assertion that is user-specified; (3) automatically transforming a description of a CUV containing a pre-determined assertion to detect incorrect behavior of the CUV due to metastability of a clock-domain-crossing (CDC) signal; (4) selecting an initial state for use by the model checking method that represents all reachable states of the CUV; (5) using a Verilog representation of the CUV as input to the model checking step; (6) using a VHDL representation of the CUV as input to the model checking step; (7) using a representation of the CUV stored in computer memory as input to the model checking step; (8) using a representation of the CUV stored on disk as input to the model checking step.

Although the present invention is illustrated in connection with specific embodiments for instructional purposes, the present invention is not limited thereto. Various adaptations and modifications may be made without departing from the scope of the invention. For example, although model checking is used in some embodiments, other embodiments use other formal verification methods to find a stimulus sequence that violates an assertion as noted above.

Tools for formal verification that may be used in act 620 are available in the prior art (either commercially or from public sources such as universities and laboratories), and may be based on any of a number of techniques, such as (1) symbolic model checking, (2) symbolic simulation, (3) explicit state enumeration, and (4) satisfiability (SAT). For background on each of the just-described techniques, see, for example, the following references, each of which is incorporated by reference herein in its entirety:

(1) an article by J. R. Burch, E. M. Clarke, K. L. McMillan, D. L. Dill, and J. Hwang, entitled "Symbolic model checking: 1020 states and beyond", published in *Information and Computation*, Vol. 98, no. 2, June 1992; another article entitled "Coverage Estimation for Symbolic Model Checking" by Yatin Hoskote, Timothy Kam, Pei-Hsin Ho, and Xudong Zhao, published in *Proceedings of DAC 1999 (Best Paper Award)*, pp. 300–305, and a PhD thesis by K. L. McMillan entitled "Symbolic model checking—an approach to the state explosion problem", Carnegie Mellon University, 1992;

(2) article entitled "Automatic Verification of Pipelined Microprocessor Control," by Jerry R. Burch and David L.

Dill, published in the proceedings of International Conference on Computer-Aided Verification, LNCS 818, Springer-Verlag, June 1994;

(3) article by E. M. Clarke, E. A. Emerson and A. P. Sistla entitled "Automatic verification of finite-state concurrent systems using temporal logic specifications" published in *ACM Transactions on Programming Languages and Systems*, 8(2):244–263, 1986; and article entitled "Protocol Verification as a Hardware Design Aid" by David Dill, Andreas Drexler, Alan Hu and C. Han Yang published in *Proceedings of the International Conference on Computer Design*, October 1992.

(4) article entitled "Bounded Model Checking Using Satisfiability Solving" by Edmund Clarke, Armin Biere, Richard Raimi, and Yunshan Zhu, published in *Formal Methods in System Design*, volume 19 issue 1, July 2001, by Kluwer Academic Publishers.

In addition, see U.S. Pat. No. 5,465,216 granted to Rotem, et al. on Nov. 7, 1995, and entitled "Automatic Design Verification" (that is incorporated by reference herein in its entirety) for an additional example of formal verification tool. See also U.S. Pat. No. 6,192,505 granted to Beer, et al. on Feb. 20, 2001, and entitled "Method and system for reducing state space variables prior to symbolic model checking" that is incorporated by reference herein in its entirety.

Formal verification tools available in the prior art for property checking include, for example, Symbolic Model Verification (SMV) software package available from Carnegie-Mellon University, and the coordinated specification analysis (COSPAN) software package available from Bell Laboratories (e.g. at ftp "dot" research "dot" att "dot" corn wherein the word "dot" is to be replaced by "." to form the ftp address).

For additional information on formal verification tools, see C. Kern and M. R. Greenstreet, "Formal Verification in Hardware Design: A Survey," in *ACM Trans. on Design Automation of Electronic Systems*, vol. 4, pp. 123–193, April 1999, that is incorporated by reference herein in its entirety.

Note also that some embodiments of the invention may be implemented as described in an article entitled "Formally Verifying Clock Domain Crossing Jitter Using Assertion-Based Verification", Design And Verification Conference, Tai Ly, Neil Hand and Chris Ka-kei Kwok, February 2004 that is incorporated by reference herein in its entirety.

Also, although formal verification is used in some embodiments of act 620, other embodiments may use other methods. For example, one alternative embodiment performs act 620 by simulation (either manually or using a simulator) of each and every possible stimulus (wherein the stimulus sequence is a sequence of vectors, with one vector of inputs for each cycle), for the number of cycles "L" and check if the assertion is violated during the simulation. So, in the example illustrated in FIG. 7, there are four possible inputs for the first cycle (as there are two enable signals for the two metastability injectors and each enable signal can be either asserted or deasserted). There are also four possible inputs for the second cycle. Therefore, if L is 2 in this example, then there are 16 possible sequences of inputs all of which are simulated.

Moreover, according to the method of the invention, an initial state represented in the Verilog can correspond to any state reachable by the circuit under verification during normal operation.

Furthermore, although transmission of a one-hot signal across clock domains, and checking by the pre-determined

assertion that the signal in the receive clock domain is in fact one hot has been described above in some embodiments, other embodiments may transmit signals with other properties across clock domains, and check their respective properties conform to the circuit designer's expectations. For example, some embodiments transmit a Gray coded signal for a count across clock domains, and the pre-determined assertion checks to confirm that the signal received in the receive clock domain is in fact Gray coded (e.g. that no more than one bit changes in each successive cycle).

Note that software (including instructions and data structures) for performing acts of the type illustrated in FIGS. 6A-6E may be embedded in computer readable storage media such as disk drives, magnetic tape, CDs (compact discs) and DVDs (digital versatile discs or digital video discs), and/or encoded in transmission media (with or without a carrier wave upon which the signals are modulated) including a communications network, such as the Internet.

Therefore, the spirit and scope of the appended claims should not be limited to the foregoing description.

## APPENDIX A

(PRIOR ART)

```

1
2 module design (tx_clk, rx_clk, rst, rx_q_0, rx_q_1, error);
3 input tx_clk, rx_clk, rst;
4 output rx_q_0, rx_q_1, error;
5
6 wire tx_d_0, tx_d_1, tx_q_0, tx_q_1, rx_d_0, rx_d_1;
7 reg tx_reg_0, tx_reg_1, rx_reg_0, rx_reg_1;
8 reg error;
9
10 assign tx_q_1 = tx_reg_1;
11 assign tx_q_0 = tx_reg_0;
12 assign tx_d_1 = tx_q_0;
13 assign tx_d_0 = tx_q_1;
14
15 initial begin
16   error = 1'b0;
17   tx_reg_1 = 1'b0;
18   tx_reg_0 = 1'b1;
19   rx_reg_1 = 1'b0;
20   rx_reg_0 = 1'b1;
21 end
22
23 always @(posedge tx_clk) begin
24   if (rst) begin
25     tx_reg_1 = 1'b0;
26     tx_reg_0 = 1'b1;
27   end
28   else begin
29     tx_reg_1 = tx_d_1;
30     tx_reg_0 = tx_d_0;
31   end
32 end
33
34 assign rx_q_1 = rx_reg_1;
35 assign rx_q_0 = rx_reg_0;
36
37 always @(posedge rx_clk) begin
38   rx_reg_1 = rx_d_1;
39   rx_reg_0 = rx_d_0;
40 end
41
42 always @(posedge rx_clk) begin
43   if (rx_q_0 == rx_q_1) error = 1'b1;
44 end
45
46 endmodule

```

## APPENDIX B

(example in one embodiment of invention)

```

5 1 module design (tx_clk, rx_clk, rst, rx_q_0, rx_q_1, error,
2     jitter_control_0, jitter_control_1);
3 input tx_clk, rx_clk, rst;
4 output rx_q_0, rx_q_1, error;
5 input jitter_control_0, jitter_control_1;
6 wire tx_d_0, tx_d_1, tx_q_0, tx_q_1, rx_d_0, rx_d_1;
10 7 reg tx_reg_0, tx_reg_1, rx_reg_0, rx_reg_1;
8 reg error;
9
10 assign tx_q_1 = tx_reg_1;
11 assign tx_q_0 = tx_reg_0;
12 assign tx_d_1 = tx_q_0;
15 13 assign tx_d_0 = tx_q_1;
14
15 initial begin
16   error = 1'b0;
17   tx_reg_1 = 1'b0;
18   tx_reg_0 = 1'b1;
19   rx_reg_1 = 1'b0;
20   rx_reg_0 = 1'b1;
21 end
22
23 always @(posedge tx_clk) begin
24   if (rst) begin
25     tx_reg_1 = 1'b0;
26     tx_reg_0 = 1'b1;
27   end
28   else begin
29     tx_reg_1 = tx_d_1;
30     tx_reg_0 = tx_d_0;
31   end
32 end
33
34 assign rx_q_1 = rx_reg_1;
35 assign rx_q_0 = rx_reg_0;
36 assign rx_d_1 =
37   (jitter_control_1 && (tx_reg_1 !== tx_d_1)) ?
38     !tx_q_1:tx_q_1;
39   (jitter_control_0 && (tx_reg_0 !== tx_d_0)) ? !tx_q_0:
40     tx_q_0;
41 always @(posedge rx_clk) begin
42   rx_reg_1 = rx_d_1;
43   rx_reg_0 = rx_d_0;
44 end
45
46 always @(posedge rx_clk) begin
47   if (rx_q_0 == rx_q_1) error = 1'b1;
48 end
49
50 endmodule

```

## APPENDIX C

(example in one embodiment of invention)

```

1
2 # INV: formula 1 failed - - - error=0
3 # INV: calling debugger
55 4 # INV: a sequence of states starting at an initial state leading to a
   bad state
5
6 --State 0:
7 error:0
8 rx_reg_0:1
9 rx_reg_1:0
60 10 tx_reg_0:1
11 tx_reg_1:0
12
13 --Goes to state 1:
14 error:0
65 15 rx_reg_0:0
16 rx_reg_1:0

```

## APPENDIX C-continued

(example in one embodiment of invention)

---

```

17 tx_reg_0:0
18 tx_reg_1:1
19 --On input:
20 jitter_control_0:1
21 jitter_control_1:0
22 rst:0
23 tx_clk:1
24 rx_clk:1
25
26 - -Goes to state 2:
27 error:0
28 rx_reg_0:0
29 rx_reg_1:0
30 tx_reg_0:0
31 tx_reg_1:1
32 --On input:
33 jitter_control_0:1
34 jitter_control_1:0
35 rst:0
36 tx_clk:0
37 rx_clk:0
38
39 --Goes to state 3:
40 error:1
41 rx_reg_0:0
42 rx_reg_1:1
43 tx_reg_0:1
44 tx_reg_1:0
45 --On input:
46 jitter_control_0:0
47 jitter_control_1:0
48 rst:0
49 tx_clk:1
50 rx_clk:1
51
52 # INV: Summary of invariant pass/fail
53 # INV: formula failed --- error=0

```

---

What is claimed is:

1. A method of verifying a design of a circuit under verification, the method comprising:

transforming an original description of the circuit under verification by adding description of circuitry to inject effects of metastability into the circuit under verification, resulting in a transformed description; and

analyzing the transformed description to determine a stimulus sequence for the transformed description to violate a pre-determined assertion,

outputting the stimulus sequence, for application during a simulation to a plurality of inputs described in the transformed description,

wherein said circuit under verification comprises a group of signals crossing between two clock domains and said plurality of inputs comprises a corresponding group of inputs of said circuitry added during said transforming, to inject effects of metastability, and on being asserted each input in said group of inputs enables said circuitry for a corresponding signal in said group of signals.

2. The method of claim 1, wherein said circuit under verification comprises a plurality of signals crossing between two clock domains, and wherein said transforming describes said circuitry being added as injecting effects of metastability into the output of at least one storage element receiving as input a signal in said plurality of signals.

3. The method of claim 2, wherein said transforming describes said circuitry being added as being further coupled to an input of a storage element in the transitive sequential fanin of at least one signal in said plurality of signals.

4. The method of claim 1, wherein said circuit under verification comprises a plurality of signals crossing between two clock domains, and wherein said transforming describes said circuitry being added as injecting effects of metastability into an input of at least one storage element that is described in the original description as receiving a signal in said plurality of signals as input.

5. The method of claim 4, wherein said transforming describes said circuitry being added as being further coupled to an input of a storage element in the transitive sequential fanin of at least one signal in said plurality of signals.

6. The method of claim 1, wherein the pre-determined assertion is described in the transformed description as monitoring behavior of a portion of said circuit under verification in the transitive sequential fanout of at least one signal in said plurality of signals.

7. The method of claim 1, further comprising: performing simulation of the transformed description in response to the stimulus sequence.

8. The method of claim 7, further comprising: automatically displaying a plurality of waveforms resulting from said simulation.

9. The method of claim 1, wherein said analyzing comprises formal analysis.

10. The method of claim 9, wherein said formal analysis comprises bounded model checking.

11. The method of claim 1, wherein the pre-determined assertion is already existing in the original description prior to said transforming.

12. The method of claim 1, wherein the pre-determined assertion does not exist in the original description prior to said transforming, and the method further comprises adding the pre-determined assertion to the transformed description after said transforming.

13. The method of claim 1, wherein:

said transforming is performed automatically in a computer; and said analyzing is also performed automatically in said computer.

14. A method of verifying a design of a circuit under verification, the method comprising:

transforming an original description of the circuit under verification by adding description of circuitry to inject effects of metastability into the circuit under verification, resulting in a transformed description; analyzing the transformed description using a model checking method, to automatically determine one or more stimulus sequences that violate a pre-determined assertion, wherein said analyzing comprises using as an initial state, a state derived from simulation of the circuit under verification; and outputting the one or more stimulus sequences that violate the pre-determined assertion for display or simulation of the circuit under verification.

15. The method of claim 14, wherein said analyzing determines stimulus sequence to be applied during simulation, to a plurality of inputs described in the transformed description.

16. The method of claim 15, further comprising: performing simulation of the transformed description, in response to the stimulus sequence.

17. The method of claim 16, further comprising: displaying a plurality of waveforms resulting from said simulation.

18. The method of claim 15, wherein: at least one input in the plurality of inputs is an input of the circuitry to inject effects of metastability.

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19. The method of claim 14 wherein said analyzing comprises using as an initial state, a reset state of the circuit under verification.

20. The method of claim 14 wherein said analyzing comprises using as an initial state, a non-reset state of the circuit under verification.

21. An apparatus configured to model a circuit under verification from a description thereof, wherein the apparatus comprises:

circuitry to inject effects of metastability into a receive clock domain, the circuitry comprising a transition detector coupled to a transmit clock domain, to receive a first signal output by the transmit clock domain, the circuitry further comprising a conditional inverter coupled to the receive clock domain, to supply a second signal being input to the receive clock domain, wherein the conditional inverter supplies one of (an unaltered version and an inverted version) of the first signal as the second signal, depending on an output of the transition detector;

a pre-determined asserter coupled to the receive clock domain to receive therefrom information related to the second signal; and

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instructions configured to perform model checking of said model, to automatically determine stimulus sequence to violate the predetermined assertion.

22. The apparatus of claim 21 wherein said circuitry is configured to:

supply one of the unaltered version and the inverted version when the enable signal is asserted; and unconditionally supply only the unaltered version when the enable signal is deasserted.

23. The apparatus of claim 21 wherein said transition detector is configured to:

enable a conditional inverter on detecting transitions, when the enable signal is asserted; and disable the conditional inverter when the enable signal is deasserted.

24. The apparatus of claim 21 further comprising: instructions configured to perform simulation of said model, in response to the stimulus sequence; and instructions configured to display a plurality of waveforms resulting from said simulation.

\* \* \* \* \*

UNITED STATES PATENT AND TRADEMARK OFFICE  
**CERTIFICATE OF CORRECTION**

PATENT NO. : 7,243,322 B1  
APPLICATION NO. : 10/859055  
DATED : July 10, 2007  
INVENTOR(S) : Ly et al.

Page 1 of 1

It is certified that error appears in the above-identified patent and that said Letters Patent is hereby corrected as shown below:

Title Page (76) Inventors: last name of Fourth named inventor spelled wrong delete "Ross Andrew Ander" and add --Ross Andrew Andersen--

Signed and Sealed this

Eleventh Day of December, 2007

A handwritten signature in black ink on a light gray dotted background. The signature reads "Jon W. Dudas" in a cursive style.

JON W. DUDAS

*Director of the United States Patent and Trademark Office*