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**Ou**

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(54) **CONFIGURABLE BUFFER MANAGER  
ARBITRATION REQUEST THROTTLING  
MECHANISM**

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**G06F 12/0802** (2016.01)

(52) **U.S. Cl.**  
CPC ..... **G06F 3/0634** (2013.01); **G06F 3/0605** (2013.01); **G06F 3/0656** (2013.01); **G06F 3/0679** (2013.01); **G06F 12/0802** (2013.01)

(58) **Field of Classification Search**  
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See application file for complete search history.

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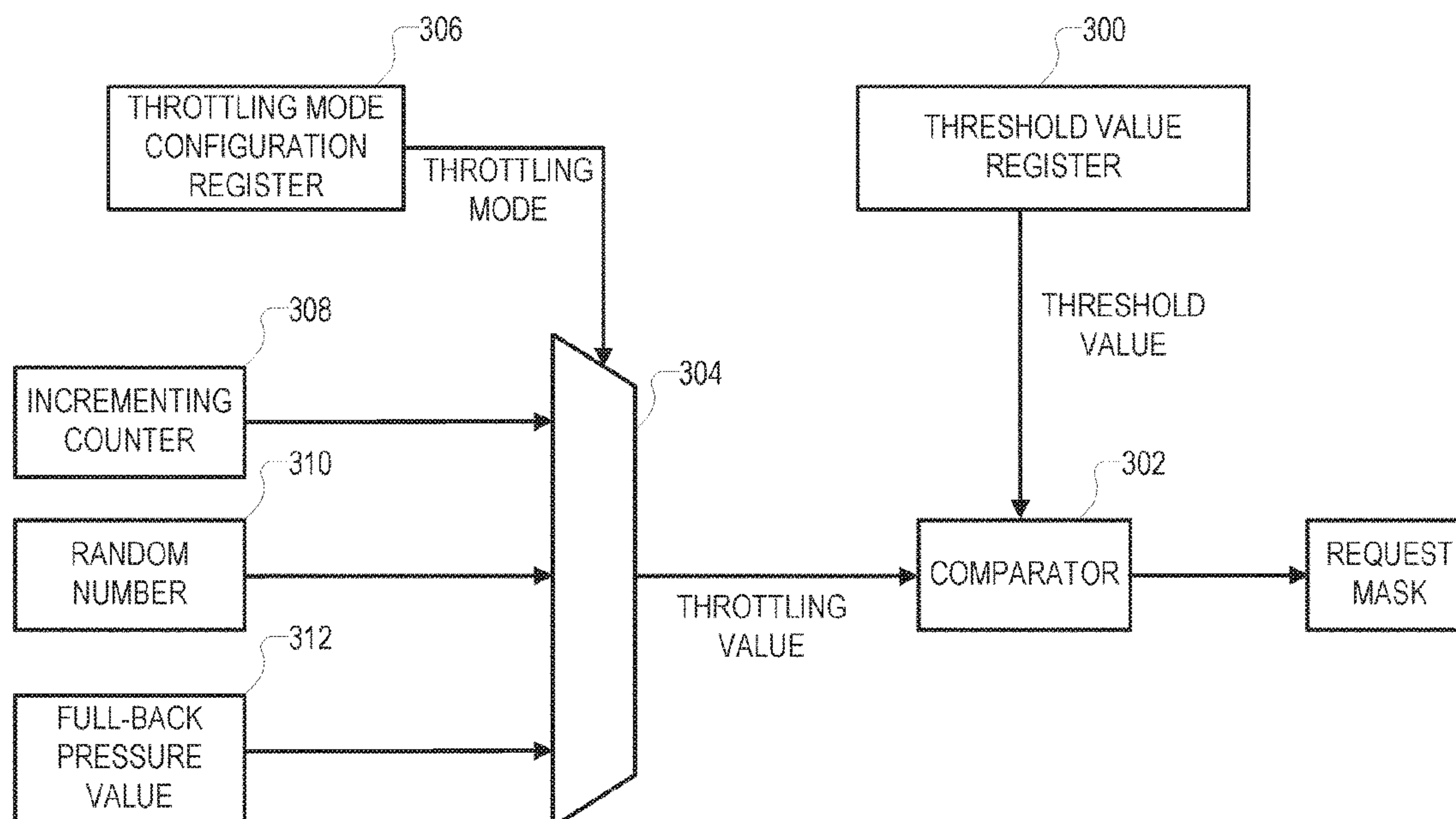
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(57) **ABSTRACT**

A throttling value is determined according to one of multiple throttling modes. A threshold value is determined. The throttling value is compared to the threshold value. A request mask is generated based on a result of the comparison of the throttling value to the threshold value. In response to detecting the request mask, an arbitration request is masked using the request mask prior to passing the arbitration request to an arbitration component that manages access of multiple client requestors to a buffer. The arbitration request is generated in response to receiving an access request for the buffer from a client requestor of the multiple client requestors.

**20 Claims, 7 Drawing Sheets**

210



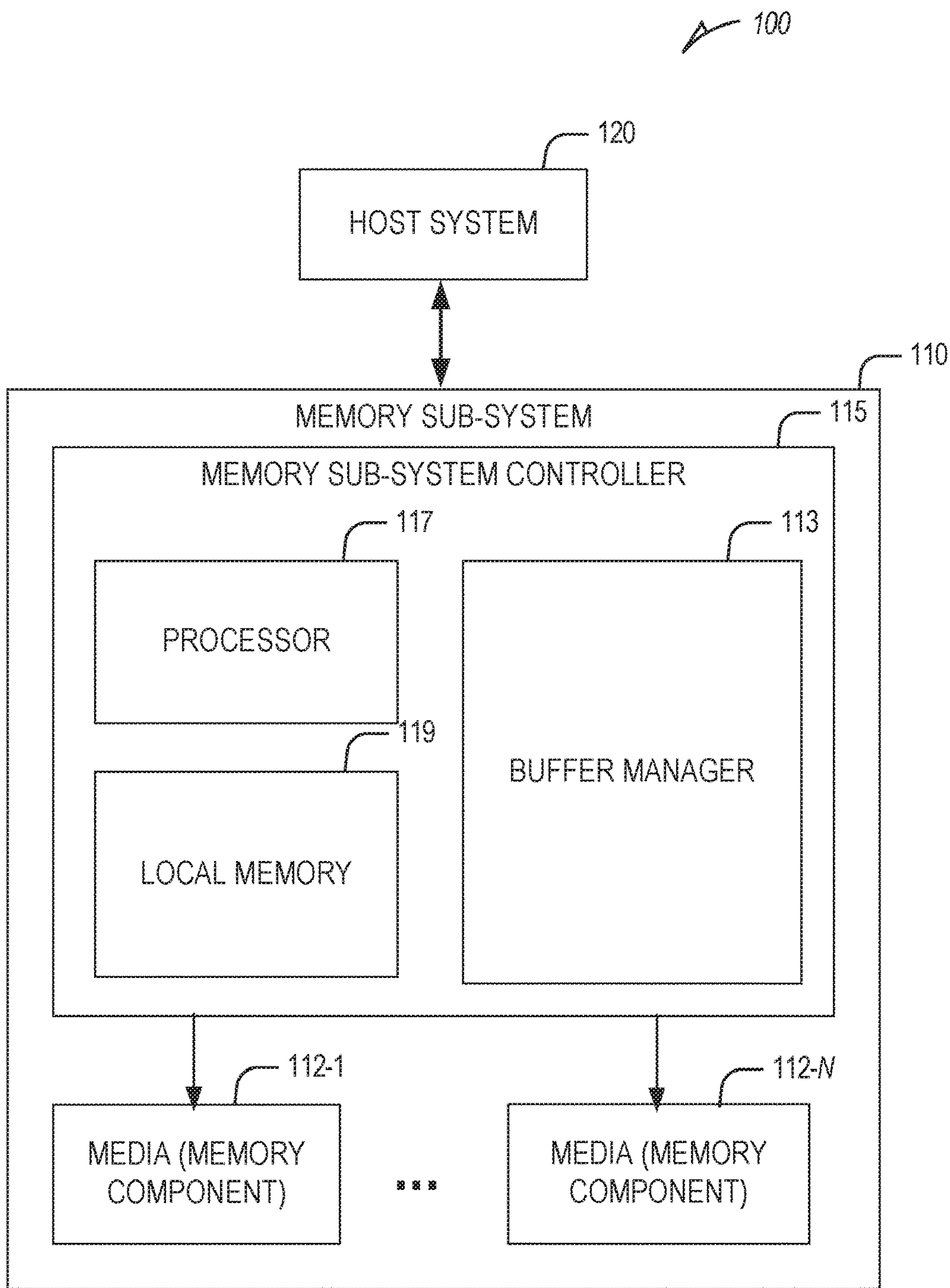


FIG. 1



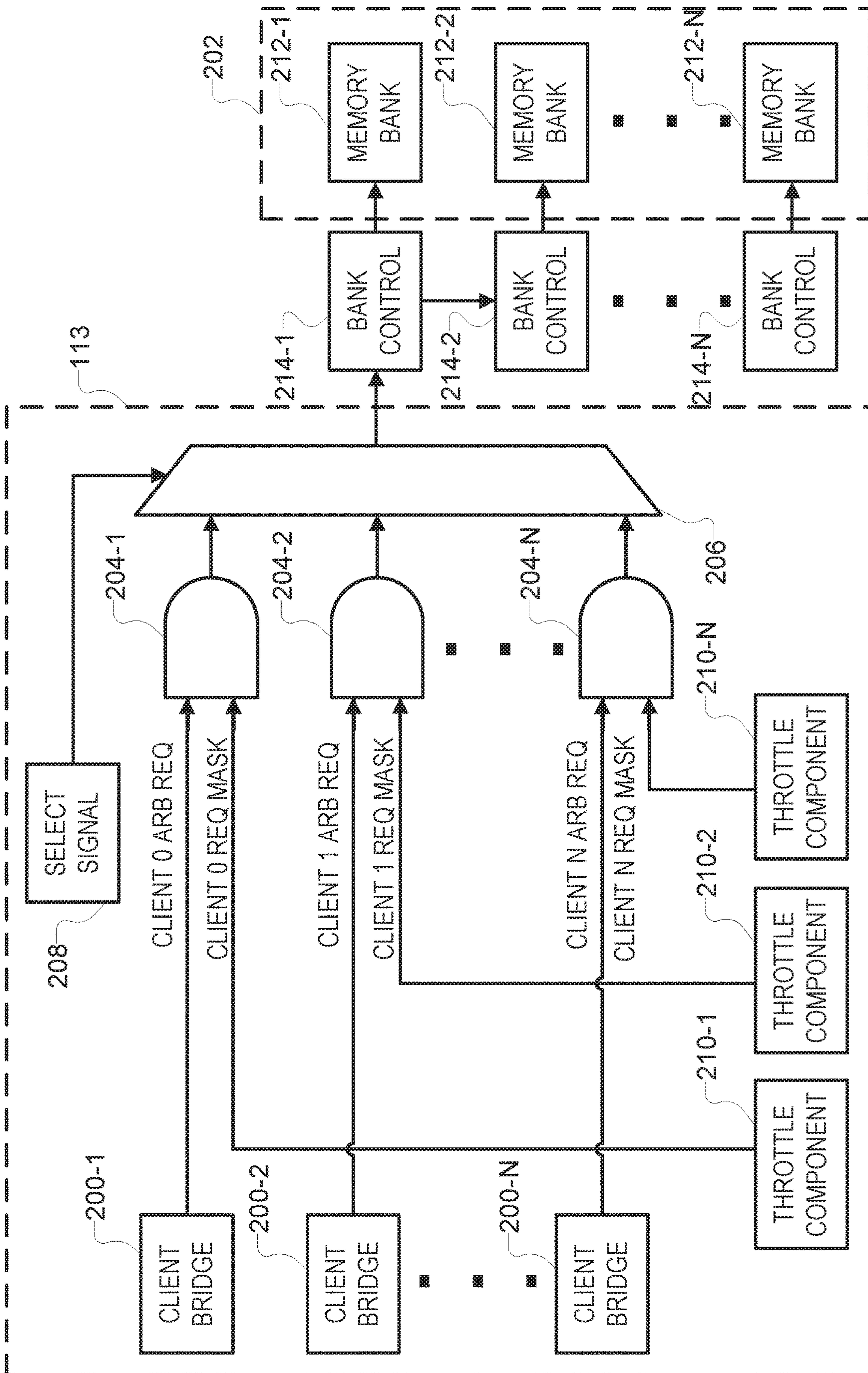


FIG. 2

210

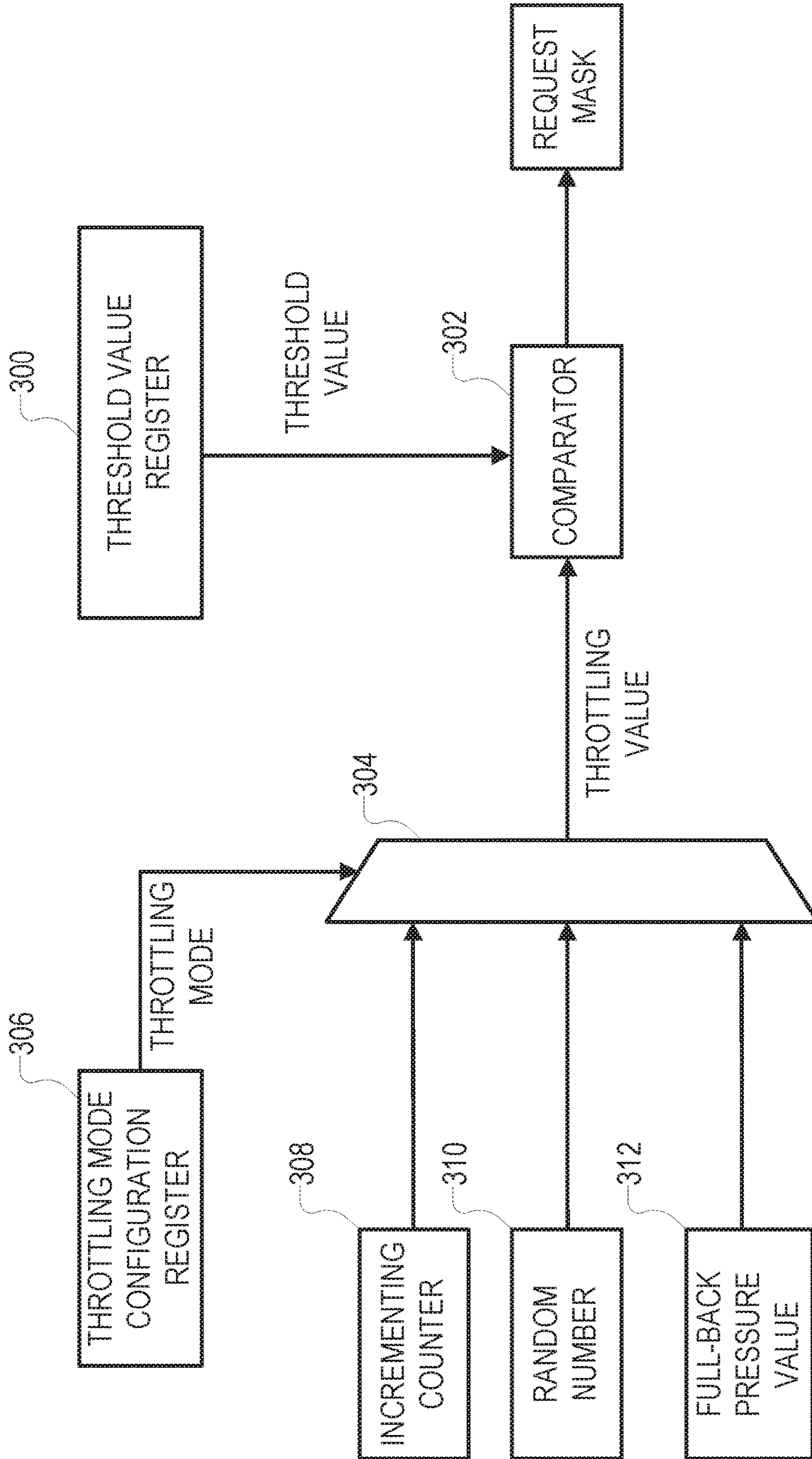


FIG. 3

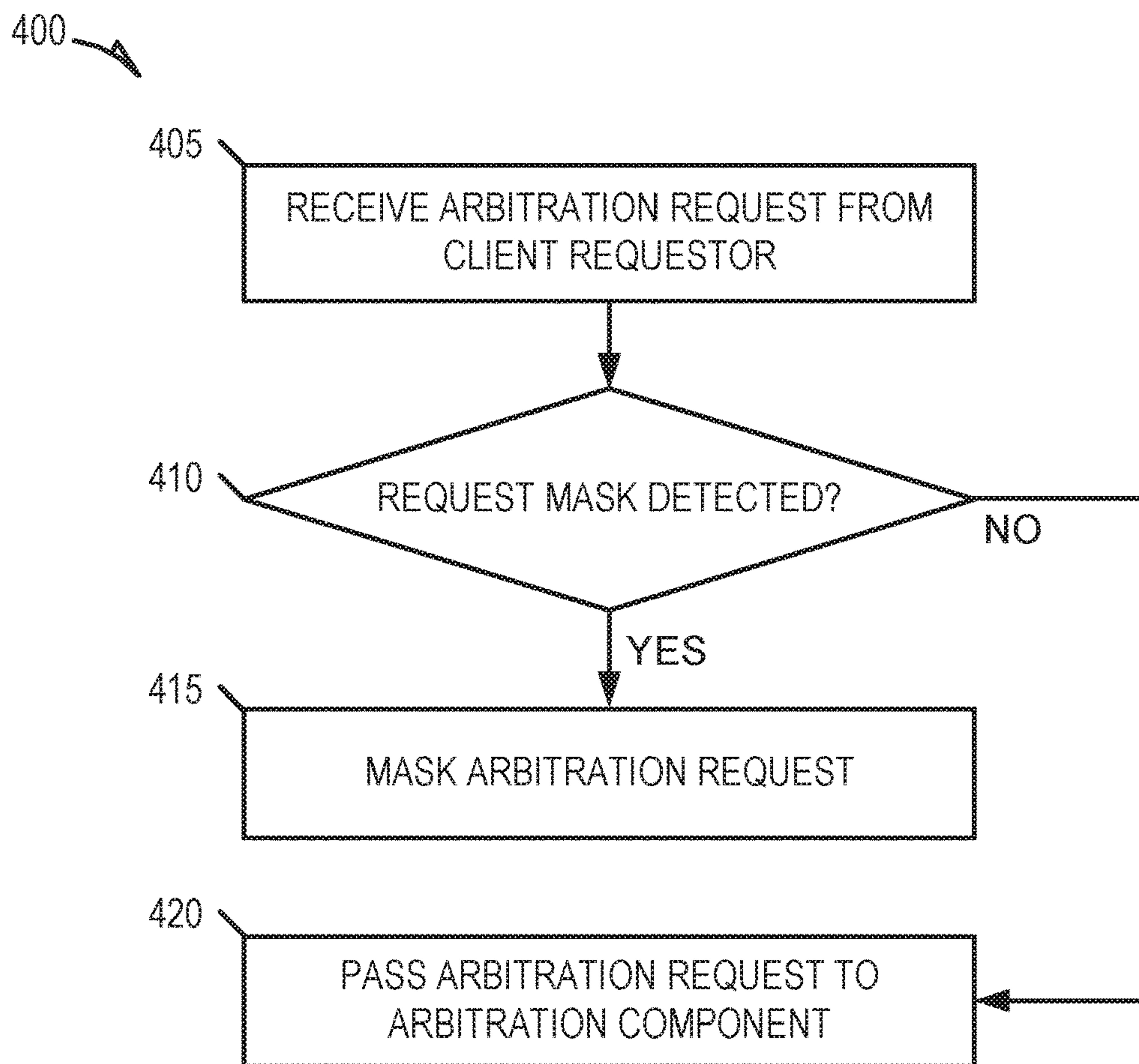


FIG. 4



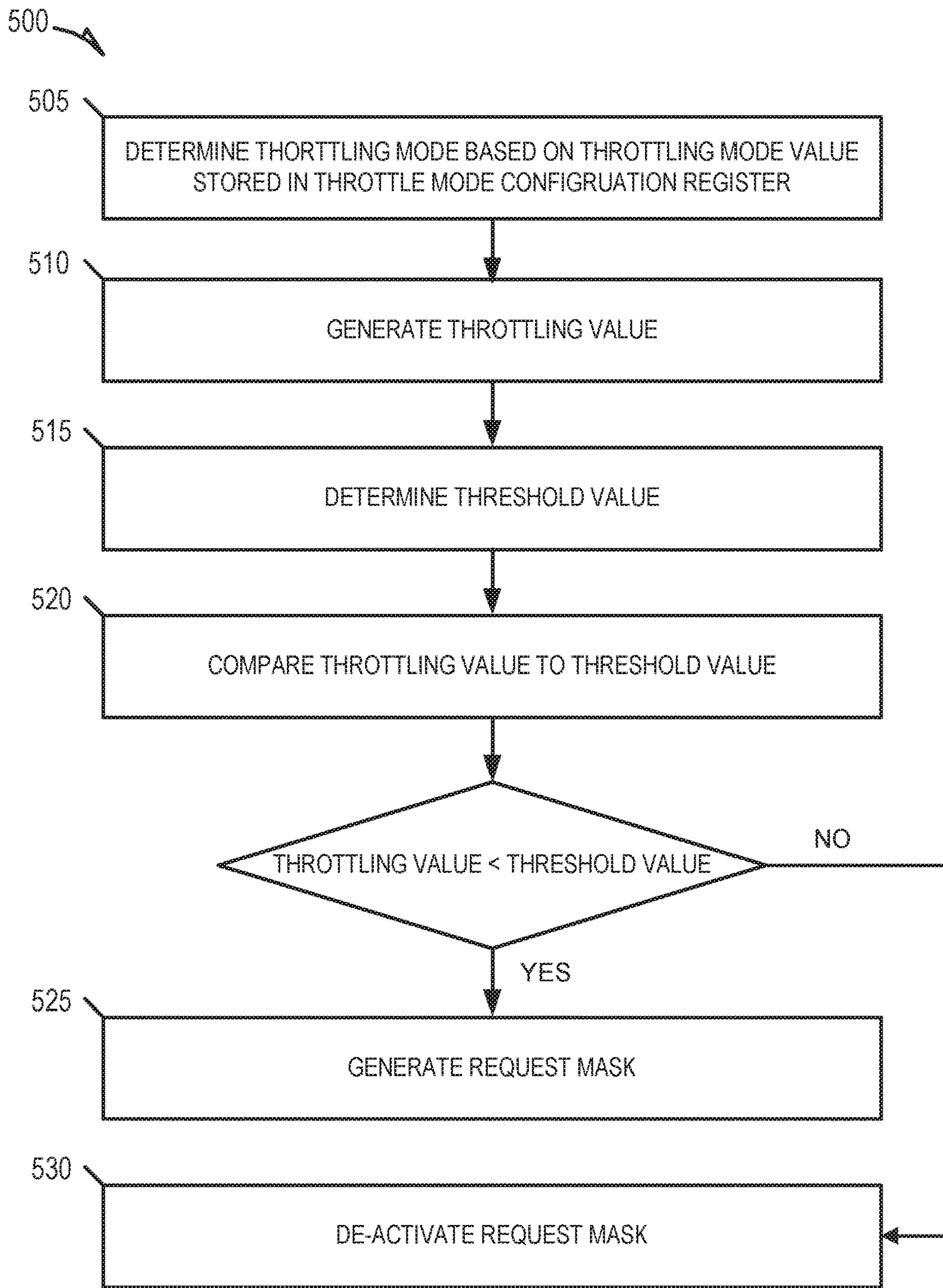


FIG. 5

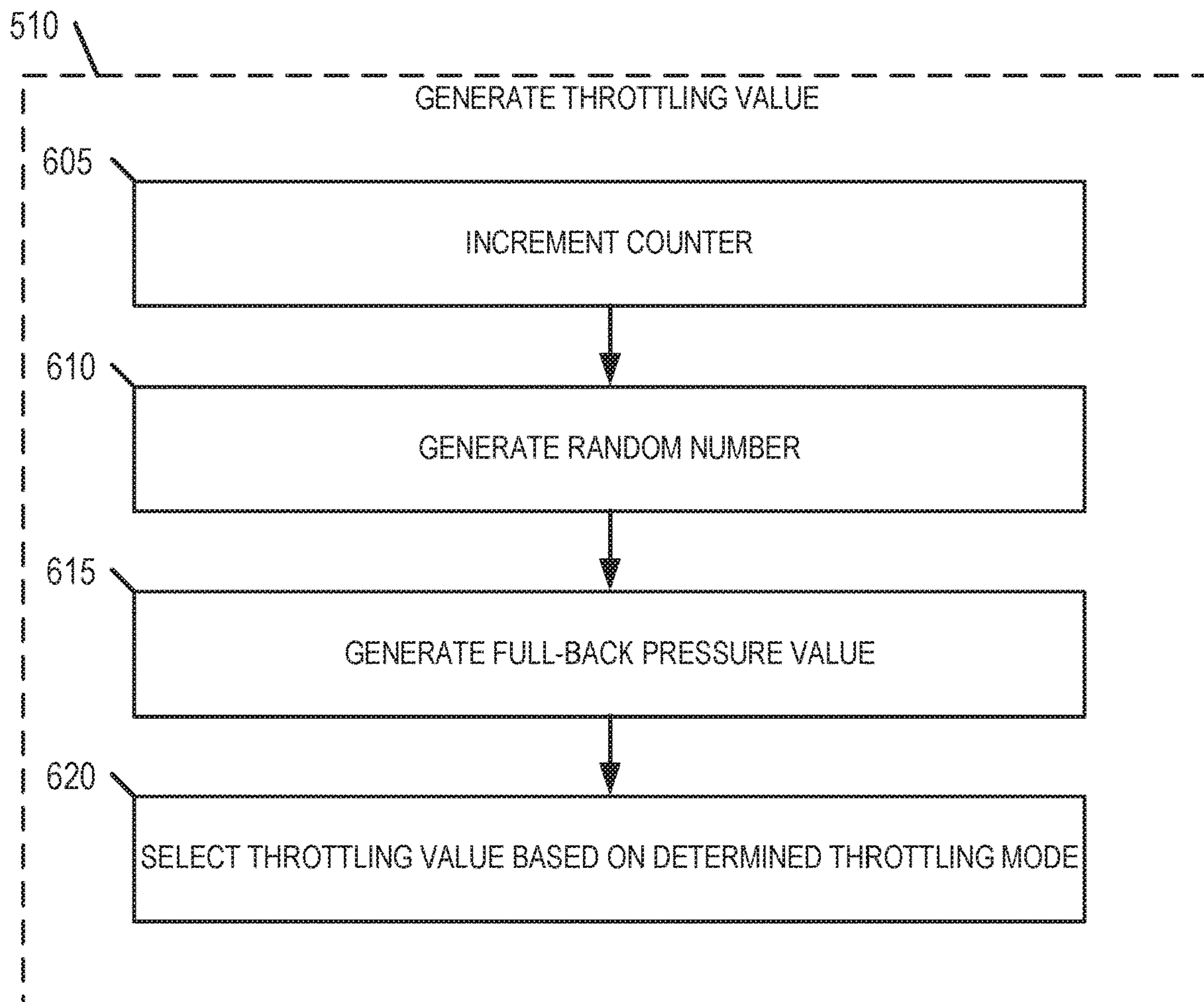


FIG. 6

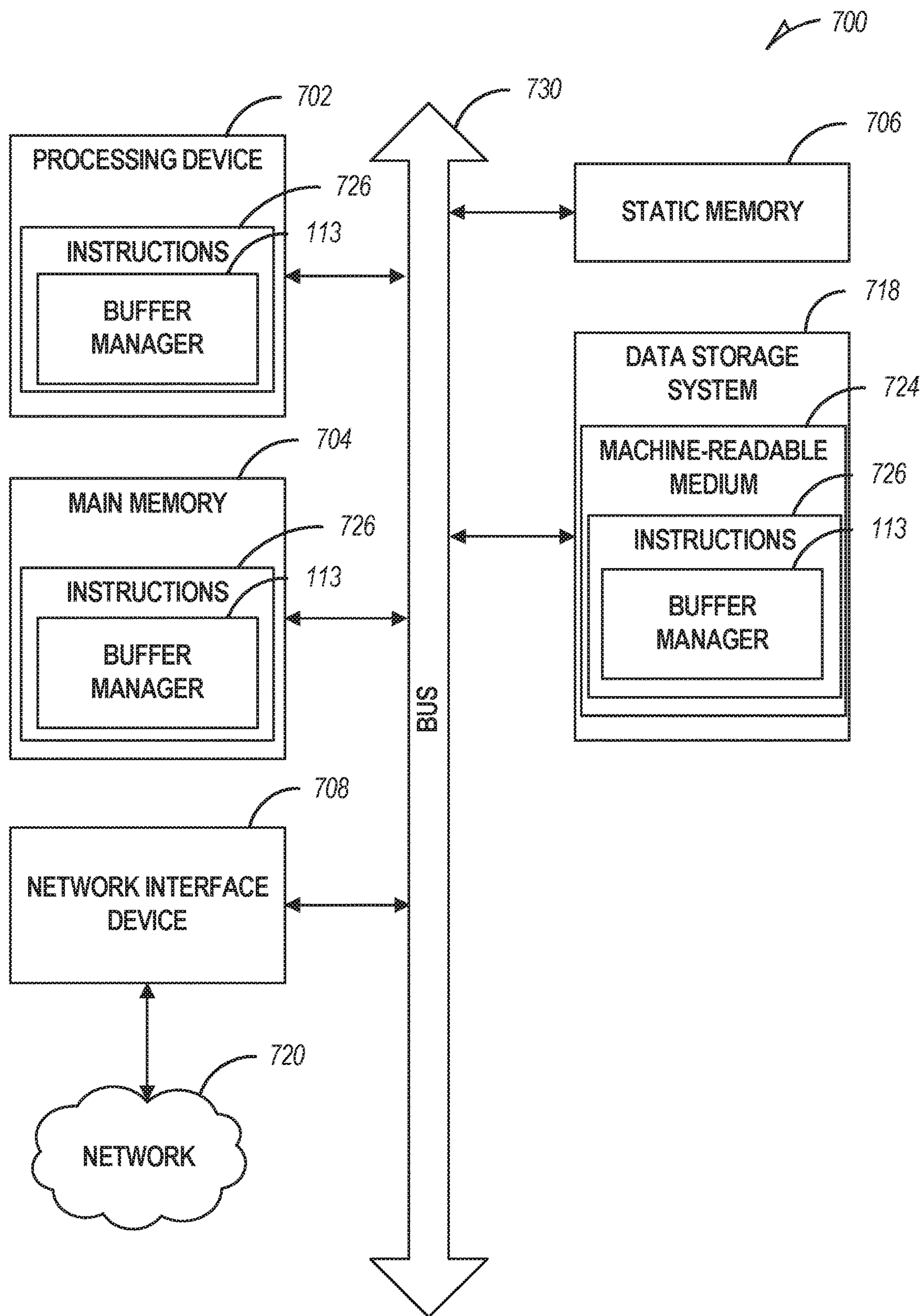


FIG. 7



## 1

**CONFIGURABLE BUFFER MANAGER  
ARBITRATION REQUEST THROTTLING  
MECHANISM**

TECHNICAL FIELD

Embodiments of the disclosure relate generally to memory sub-systems and, more specifically, to a configurable buffer manager arbitration request throttling mechanism for a memory sub-system.

BACKGROUND

A memory sub-system can be a storage system, such as a solid-state drive (SSD), and can include one or more memory components that store data. The memory components can be, for example, non-volatile memory components and volatile memory components. In general, a host system can utilize a memory sub-system to store data at the memory components and to retrieve data from the memory components.

BRIEF DESCRIPTION OF THE DRAWINGS

The present disclosure will be understood more fully from the detailed description given below and from the accompanying drawings of various embodiments of the disclosure.

FIG. 1 illustrates an example computing environment that includes a memory sub-system, in accordance with some embodiments of the present disclosure.

FIG. 2 is a block diagram illustrating components of a buffer manager, in accordance with some embodiments of the present disclosure.

FIG. 3 is a block diagram illustrating sub-components of a client throttling component, which is provided as part of the buffer manager, in accordance with some embodiments of the present disclosure.

FIG. 4 is a flow diagram of an example method to throttle arbitration requests, in accordance with some embodiments of the present disclosure.

FIG. 5 is a flow diagram of an example method to activate a request mask, in accordance with some embodiments of the present disclosure.

FIG. 6 is a flow diagram of an example method to generate a throttling value, in accordance with some embodiments of the present disclosure.

FIG. 7 is a block diagram of an example computer system in which embodiments of the present disclosure can operate.

DETAILED DESCRIPTION

Aspects of the present disclosure are directed to class-based dynamic memory slot allocation in a memory sub-system. A memory sub-system is also hereinafter referred to as a “memory device.” An example of a memory sub-system is a storage system, such as a SSD. In some embodiments, the memory sub-system is a hybrid memory/storage sub-system. In general, a host system can utilize a memory sub-system that includes one or more memory components. The host system can provide data to be stored at the memory sub-system and can request data to be retrieved from the memory sub-system.

In many instances, one or more of the memory components of the memory sub-system can be configured as a buffer memory to store user and system data. Such memory components can be shared by multiple clients. The memory sub-system can include a buffer manager to control the

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buffer memory components. The buffer manager is responsible for managing read and write transfer requests received from the multiple clients. The buffer manager can include an arbiter component to select which client has access to the buffer manager during each clock cycle.

In a configurable parameterizable buffer manager architecture, a configurable number of client requestors are connected to the buffer manager to make requests to a shared memory component. The requested access is for a fixed portion of the memory component. In buffer configuration with multiple memory banks, for example, the requested access can be for a memory “stripe.” The term “stripe” may be used to refer to a single slice of memory with the same address across all memory banks. For example, in an 8-bank memory configuration with each bank being 16-bytes wide, a stripe request would consist of an access from the same memory bank address across all banks, which would be 128 bytes in this configuration. In buffer configurations with multiple memory banks, the requests are pipelined along the memory bank structure so that each request reads from a single memory bank during one clock cycle, and multiple requests from different client requestors occur in parallel across the multiple memory banks.

In many scenarios, client requestors are both independent and greedy, and often request as much of the shared memory resource as possible without regard to other client requests. When the number of requests exceeds the number of available resources, the buffer manager will not be able to provide full throughput to every client. This means that all clients may experience reduced throughput. As example, consider a buffer manager with 8 individual memory banks, each which can process up to 16 bytes per clock cycles and which run at 500 Mhz. Each individual bank can be accessed simultaneously by one client, so up to 8 banks can be accessed by 8 different clients. This provides a maximum theoretical throughput of 8 GB (500 Mhz\*16 bytes) per second, which when multiplied across the 8 banks equates to 64 GB throughput per second. If there are 16 clients all actively requesting, since only 8 out of 16 clients will be able to access the 8 banks in a given clock cycle, each client will only achieve 50% bandwidth or 4 GB max.

Reduced throughput from all requesting clients is often an accepted condition in many memory sub-systems, but in certain situations, some clients may need a minimum guaranteed level of throughput while other clients may accept a lower level of throughput. The minimum guaranteed level of throughput for various clients may also dynamically change during system operation, so there is a need for system firmware to be able to consistently make adjustments on-the-fly depending on current request and performance profiles.

Additionally, from a test and verification perspective, it is desirable to provide some controllability to mimic different types of traffic patterns for buffer manager arbitration and data flow. Testing a full system running with all clients active is usually more difficult to control and manage, and simulation performance is significantly slower. Many problems can be found by running a smaller level of testing through a single client or a small subset of active clients in the test bench. However, in these smaller client configuration test benches, there needs to be a way to mimic the different burst and random traffic patterns as seen by the single or small subset of requesting clients.

Aspects of the present disclosure address the above and other deficiencies by providing a buffer manager having a control mechanism to throttle client requests. In particular, a throttling component conditionally masks out requests from a client to allow other client requests to have arbitration



granted over the masked out requests. The conditional masking is performed by the throttling component on clock cycle-by-clock cycle basis. Since masking is performed on a clock cycle-by-clock cycle basis, the throttling component improves upon conventional buffer managers by allowing fine grain control over how requests are issued to the buffer manager's arbiter, thus allowing the ability to create custom profiles on how requests are masked off or granted. As a further improvement to conventional buffer managers, the per-clock cycle request and arbitration grant described as part of the present disclosure also allows for on-the-fly control of how masking is enabled or disabled, since a configuration change will be automatically reflected in the following clock cycle.

The buffer manager can include a throttling control component for each client requestor. Each throttling control can generate a request mask during each clock cycle, which can mask off the client request when asserted. This simplifies the interaction between the client bridge and the arbiter compared to conventional buffer managers since the mask signal hides the request from the arbiter allowing it to select between any remaining valid requesting client bridges. Each throttling control component is capable of operating in multiple throttling modes, which dictate how and when the throttling component masks request. The throttling mode of any given throttling component can be configured on-the-fly to accommodate for changes in arbitration request frequency.

The multiple throttling modes can be particularly advantageous in testing and qualification in design verification contexts. For example, in system-wide testing, a complete memory system with all requesting clients and a buffer manager is simulated to model the real system interaction. However, the large scale of testing is often cumbersome and extremely slow since modeling the entire system and all its active components consumes an enormous amount of computer resources and memory. It is also hard to configure and program all system components to mimic how they will function in the final product, as the block interactions can be more difficult to control.

To speed up testing and simulation, the test environment is often split up into very small subsystems, often consisting of a single or a small handful of clients interacting with the buffer manager. Smaller scale testing runs much faster and is much more controllable. However, the problem with the smaller test environment is that these systems do not model real arbitration delays. More specifically, the number of requestors to the shared buffer manager is usually much less than the available resources, meaning the requestors may never see real arbitration delays that would be present in a larger system. For example, in a single requestor test environment, the requestor will always get arbitration and never have to wait for other requestors to get their requests processed.

Configuring a throttling control component with one of the multiple throttling modes acts as a much better model of the actual system delays that a single requestor can encounter. This can help to uncover real system problems in the block level test bench. Certain throttling modes can also allow for testing how a single requestor in its block environment handles the situation of other high priority requestors consuming all system bandwidth.

FIG. 1 illustrates an example computing environment 100 that includes a memory sub-system 110, in accordance with some embodiments of the present disclosure. The memory sub-system 110 can include media, such as memory components 112-1 to 112-N. The memory components 112-1 to

112-N can be volatile memory components, non-volatile memory components, or a combination of such. In some embodiments, the memory sub-system is a storage system. An example of a storage system is a SSD. In some embodiments, the memory sub-system 110 is a hybrid memory/storage sub-system. In general, the computing environment 100 can include a host system 120 that uses the memory sub-system 110. For example, the host system 120 can write data to the memory sub-system 110 and read data from the memory sub-system 110.

The host system 120 can be a computing device such as a desktop computer, laptop computer, network server, mobile device, or such computing device that includes a memory and a processing device. The host system 120 can include or be coupled to the memory sub-system 110 so that the host system 120 can read data from or write data to the memory sub-system 110. The host system 120 can be coupled to the memory sub-system 110 via a physical host interface. As used herein, "coupled to" generally refers to a connection between components, which can be an indirect communicative connection or direct communicative connection (e.g., without intervening components), whether wired or wireless, including connections such as electrical, optical, magnetic, etc. Examples of a physical host interface include, but are not limited to, a serial advanced technology attachment (SATA) interface, a peripheral component interconnect express (PCIe) interface, universal serial bus (USB) interface, Fibre Channel. Serial Attached SCSI (SAS), etc. The physical host interface can be used to transmit data between the host system 120 and the memory sub-system 110. The host system 120 can further utilize an NVM Express (NVMe) interface to access the memory components 112-1 to 112-N when the memory sub-system 110 is coupled with the host system 120 by the PCIe interface. The physical host interface can provide an interface for passing control, address, data, and other signals between the memory sub-system 110 and the host system 120.

The memory components 112-1 to 112-N can include any combination of the different types of non-volatile memory components and/or volatile memory components. An example of non-volatile memory components includes a negative-and (NAND) type flash memory. Each of the memory components 112-1 to 112-N can include one or more arrays of memory cells such as single level cells (SLCs) or multi-level cells (MLCs) (e.g., triple level cells (TLCs) or quad-level cells (QLCs)). In some embodiments, a particular memory component can include both an SLC portion and a MLC portion of memory cells. Each of the memory cells can store one or more bits of data (e.g., data blocks) used by the host system 120. Although non-volatile memory components such as NAND type flash memory are described, the memory components 112-1 to 112-N can be based on any other type of memory such as a volatile memory. In some embodiments, the memory components 112-1 to 112-N can be, but are not limited to, random access memory (RAM), read-only memory (ROM), dynamic random access memory (DRAM), synchronous dynamic random access memory (SDRAM), phase change memory (PCM), magneto random access memory (MRAM), negative-or (NOR) flash memory, electrically erasable programmable read-only memory (EEPROM), and a cross-point array of non-volatile memory cells. A cross-point array of non-volatile memory can perform bit storage based on a change of bulk resistance, in conjunction with a stackable cross-gridded data access array. Additionally, in contrast to many flash-based memories, cross-point non-volatile memory can perform a write in-place operation, where a



non-volatile memory cell can be programmed without the non-volatile memory cell being previously erased. Furthermore, as noted above, the memory cells of the memory components **112-1** to **112-N** can be grouped as data blocks that can refer to a unit of the memory component used to store data.

A memory sub-system controller **115** (hereinafter referred to as “controller”) can communicate with the memory components **112-1** to **112-N** to perform operations such as reading data, writing data, or erasing data at the memory components **112-1** to **112-N** and other such operations. The controller **115** can include hardware such as one or more integrated circuits and/or discrete components, a buffer memory, or a combination thereof. The controller **115** can be a microcontroller, special purpose logic circuitry (e.g., a field programmable gate array (FPGA), an application specific integrated circuit (ASIC), etc.), or other suitable processor. The controller **115** can include a processor (processing device) **117** configured to execute instructions stored in local memory **119**. In the illustrated example, the local memory **119** of the controller **115** includes an embedded memory configured to store instructions for performing various processes, operations, logic flows, and routines that control operation of the memory sub-system **110**, including handling communications between the memory sub-system **110** and the host system **120**. In some embodiments, the local memory **119** can include memory registers storing memory pointers, fetched data, etc. The local memory **119** can also include read-only memory (ROM) for storing micro-code. While the example memory sub-system **110** in FIG. **1** has been illustrated as including the controller **115**, in another embodiment of the present disclosure, a memory sub-system **110** may not include a controller **115**, and may instead rely upon external control (e.g., provided by an external host, or by a processor or controller separate from the memory sub-system).

In general, the controller **115** can receive commands or operations from the host system **120** and can convert the commands or operations into instructions or appropriate commands to achieve the desired access to the memory components **112-1** to **112-N**. The controller **115** can be responsible for other operations such as wear leveling operations, garbage collection operations, error detection and error-correcting code (ECC) operations, encryption operations, caching operations, and address translations between a logical block address and a physical block address that are associated with the memory components **112-1** to **112-N**. The controller **115** can further include host interface circuitry to communicate with the host system **120** via the physical host interface. The host interface circuitry can convert the commands received from the host system into command instructions to access the memory components **112-1** to **112-N** as well as convert responses associated with the memory components **112-1** to **112-N** into information for the host system **120**.

The memory sub-system **110** can also include additional circuitry or components that are not illustrated. In some embodiments, the memory sub-system **110** can include a cache or buffer (e.g., DRAM) and address circuitry (e.g., a row decoder and a column decoder) that can receive an address from the controller **115** and decode the address to access the memory components **112-1** to **112-N**.

The memory sub-system **110** includes a buffer manager **113** that manages access to a buffer memory and services multiple client requestors. The buffer memory can be included in the controller **115** or any one or more of the memory components **112-1** to **112-N**. In some embodiments,

the controller **115** includes at least a portion of the buffer manager **113**. For example, the controller **115** can include a processor **117** (processing device) configured to execute instructions stored in local memory **119** for performing the operations described herein. In some embodiments, the buffer manager **113** is part of the host system **110**, an application, or an operating system.

The buffer manager **113** manages read and write transfer requests for buffer memory components received from a configurable number of requestors (also referred to as “clients”). The buffer manager **113** includes an arbiter component that selects which client has access to the buffer manager each clock cycle. For each client, the arbiter component includes a throttling control component that is configured to conditionally mask requests from the client to allow requests from other clients to be serviced. The throttling control component determines whether to mask requests on a clock cycle by clock cycle basis. That is, at each clock cycle, the throttling control component can decide to generate a request mask to mask an incoming request. The throttling control component is capable of being programmed according to one of multiple throttling modes and the throttling control component decides whether to generate a request mask at each clock cycle in accordance with a throttling mode to which the throttling control component is programmed. Further details with regards to the operations of the buffer manager **113** are described below.

FIG. **2** is a block diagram illustrating components of the buffer manager **113**, in accordance with some embodiments of the present disclosure. As shown, the buffer manager **113** includes client bridges **200-1** to **200-N**. The client bridges **200-1** to **200-N** present a uniform interface to the external clients of the buffer manager **113**. A client bridge **200** exists for each external client, and provides an interface to the buffer manager **113** for the corresponding client. That is, the clients make stripe requests (e.g., read or write) to the buffer manager **113** through their connected client bridge **200**.

Each client bridge **200** packs incoming requests into a 64-bit command word and stores this in an input command/data first-in, first-out (FIFO) for processing. For write requests, the write data is also stored in the command/data FIFO. In response to receiving a client request, the receiving client bridge **200** outputs an arbitration request. Once arbitration is granted, the write data is immediately sent out of the command/data FIFO in the following clock cycle and stored in a buffer memory **202**. When some or all data has been processed for that request, any subsequent requests stored in the FIFO can be executed.

For read requests, once the arbitration request is granted by the buffer manager **113**, the data transfer from the buffer memory **202** immediately begins. Data being read from the memory components **112-1** to **112-N** will be stored in an incoming FIFO. This data will be then streamed out and sent to the requesting client on its clock domain.

Arbitration requests output by each of the client bridges **200-1** are received at a corresponding masking component. For example, as shown, the buffer manager **113** includes masking components **204-1** to **204-N**. Each of the masking components **204-1** to **204-N** can receive two inputs—an arbitration request and a request mask. If a request mask is provided at its input, the masking component **204** masks the arbitration request. If no request mask is provided at its input, the masking component **204** passes the arbitration request on to an arbiter **206**.

For each of the masking components **204-1** to **204-N**, a throttling component **210** is provided to conditionally generate request masks for arbitration requests sent by the



corresponding masking component **204**. For example, as shown, the buffer manager **113** includes throttling components **210-1** to **210-N**. Each of the throttling components **210-1** to **210-N** are programmable according to one of multiple throttling modes, which dictate the conditions under which the request mask is generated. Further details regarding the throttling component **210** are discussed below in reference to FIG. 3.

The arbiter **206** selects which client has access to the buffer memory **204** at each clock cycle. As shown, the arbiter **206** receives a select signal **208**. The select signal **208** provided to the arbiter **206** dictates which client has access to the buffer memory **202** at each clock cycle. The select signal **208** can, for example, employ a round robin scheme. If an unmasked arbitration request is provided by the selected client at an input of the arbiter **206**, the request is granted and the client is granted access to the buffer memory **202** for that clock cycle. Otherwise, the selected client is denied access to the buffer memory **202** for that clock cycle.

The buffer memory **202** comprises memory banks **212-1** to **212-N**. The memory banks **212-1** to **212-N** can be spread across one or more of the memory components **112-1** to **112-N**. Each stripe request received from a client is for the same memory address across all memory banks **212-1** to **212-N**. Memory bank controls **214-1** to **214-N** respectively process requests received from the buffer manager **113** for the memory banks **212-1** to **212-N**.

FIG. 3 is a block diagram illustrating sub-components of a client throttling component **210**, which is provided as part of the requestor arbitration control, in accordance with some embodiments of the present disclosure.

The client throttling component **210** includes a threshold value register **300** that stores a threshold value. During each clock cycle, a comparator **302** compares the threshold value to a throttling value output by a throttling value selector **304**. If the throttling value is less than the threshold value, the comparator **302** generates and outputs a request mask during that clock cycle. Otherwise, no request mask is generated for that clock cycle and, if an arbitration request is received from the corresponding client during that clock cycle, the arbitration is passed on the arbiter **206**, as discussed above in reference to FIG. 2.

The throttling value output by the throttling value selector **304** depends on a programmable throttling mode of the client throttling component **210**, which is represented by a throttling mode value stored in a throttling mode configuration register **306**. As noted above, each client throttling component **210** is capable of being programmed according to one of multiple throttling modes. The throttling modes can, for example, include a counter-based mode, a random number-based mode, and a full-back pressure mode.

In a counter-based throttling mode, the throttling value selector **304** outputs a counter **308** as the throttling value. The client control throttling component **210** can increment the counter **308** on a per clock cycle or per request basis. In other words, the client control throttling component **210** can increment the counter **308** at each clock cycle or each time an arbitration request is received from a corresponding client bridge **200**. The counter **308** is incremented until it reaches a maximum value (based on counter width) and then it rolls over back to zero.

In the counter-based throttling mode, if the counter **308** is less than the compare threshold value, the comparator **302** generates a request mask to gate off a request to the arbiter **206**. When the counter **308** is greater than the threshold value, the request mask is de-activated to allow a request to pass through to the arbiter **206**.

The counter-based throttling mode allows requests to be completely throttled across a large time window. For example, if the maximum value of the counter **308** is 64 and the compare value threshold is set to 32, then for the first 32 arbitration clock cycles, the client request that is being handled by the client throttling component **210** will get no arbitration. In the next 32 clock cycles, the client request would get its normal arbitration. This would model the situation where a requestor did not get arbitration granted across longer time windows.

The counter-based throttling mode also allows for a lower bandwidth priority client to have its requests scaled back when firmware of the memory sub-system **110** firmware deems it necessary to allow other higher bandwidth priority clients to obtain more bandwidth. For example, in an 8-bank, 16-byte wide memory bank system running at 500 Mhz, with 4 higher priority clients and 8 lower priority clients all making simultaneous requests, the normal achieved bandwidth spread for all 12 requestors spread across 8-banks would be 67% (i.e., 8/12) of maximum throughput for the buffer manager **113** (5.33 GB per client). However, the counter-based throttling mode could be used in this situation to set the 8 lower priority clients with a threshold of 50% or more of the maximum counter value, which would reduce their requests by 50%. This means the effective requests being presented to the arbiter **206** would be 8 (4 higher priority requests+(8 requestors\*0.50 requests due to throttling)), which would allow the higher priority clients to get a maximum bandwidth of 8 GB per client and would allow the 8 lower priority clients to get less bandwidth (4 GB per client).

In a random number based mode, the throttling value selector **304** outputs a random (or pseudo-random) number **310** as the throttling value. The random (or pseudo-random) number **310** is generated by a random number generator (not shown). In some embodiments, the random number generator is a linear feedback shift register (LFSR) based random number generator that includes one or more LFSRs. Consistent with these embodiments, the LFSR-based random number generator generates a pseudo-random number sequence, with the random number range based on the width of the LFSR registers used. During each clock cycle, the comparator **302** compares the threshold value against the generated random number **310**. When the random number **310** is lower than the threshold value, the comparator **302** generates a request mask for that clock cycle. If the random number **310** is higher than the threshold value, the request mask is de-activated, and a request is allowed to pass through to the arbiter **206**.

An advantage of the random number-based mode is that it can simulate a more random traffic profile, where occasionally the requestor may get its requests partially throttled at any given clock cycle, as opposed to within a larger time window as seen in the counter-based mode. The percentage of throttling is based on the programmable threshold value. For example, if the random number range is 0-63 and the compare threshold value is set to 16, then 25% (i.e., 16/64) of the requests are throttled in a random manner across the entire time window. If the request threshold is set to 48, then 75% (i.e., 48/64) of the requests are throttled in a random manner across the entire time window. This same method is also usable as an alternate throttling mechanism similar to the counter-based method, where higher bandwidth priority clients are able to get more guaranteed throughput over lower bandwidth priority clients by having the lower priority clients have their threshold values set to some higher percentage to throttle their requests.



In a full-back pressure mode, a full-back pressure value **312** that is less than the compare threshold value is used as the throttling value. In this manner, while in the full-back pressure mode, the comparator **302** will generate a request mask at each clock cycle so that all requests will be blocked. In this mode, the requestor gets no access to the arbiter **206** until the throttling mode is changed by the firmware of the memory sub-system **110**. This mode can be used to restrict a client from taking up any amount of bandwidth until the firmware decides to allow that client to proceed.

Since the threshold value and throttling value are compared each clock cycle and there is no dependency on previous clock cycles, the values can be changed at any time by the CPU to control throttling on-the-fly. In addition, the throttling mode can also be changed at any time on the fly. The threshold value and the throttling mode are configurable through a standard CPU peripheral programming interface. In some embodiments, the interface is an advanced micro-controller bus architecture (AMBA) Advanced Peripheral Bus (APB).

FIG. **4** is a flow diagram of an example method to throttle arbitration requests, in accordance with some embodiments of the present disclosure. The method **400** can be performed by processing logic that can include hardware (e.g., processing device, circuitry, dedicated logic, programmable logic, microcode, hardware of a device, integrated circuit, etc.), software (e.g., instructions run or executed on a processing device), or a combination thereof. In some embodiments, the method **400** is performed by the buffer manager **113** of FIG. **1**. Although shown in a particular sequence or order, unless otherwise specified, the order of the processes can be modified. Thus, the illustrated embodiments should be understood only as examples, and the illustrated processes can be performed in a different order, and some processes can be performed in parallel. Additionally, one or more processes can be omitted in various embodiments. Thus, not all processes are required in every embodiment. Other process flows are possible.

At operation **405**, the processing device receives an arbitration request from a client requestor. Prior to passing the arbitration request to an arbitration component (e.g., the arbiter **206**), the processing device, at operation **410**, determines whether a request mask is detected. If, at operation **410**, the processing device detects a request mask, the processing device, at operation **415**, masks the arbitration request using the request mask prior to passing the arbitration request to the arbitration component. Because the arbitration request is masked, the arbitration component will not see the arbitration request, and thus the client requestor will not be granted arbitration for the clock cycle. If, on the other hand, the processing device does not detect the request mask, the processing device, at operation **420**, passes the arbitration request to the arbitration component for arbitration.

It shall be understood that the method **400** is described above in reference to a single client requestor during a single clock cycle, though it shall be appreciated that the method **400** can be performed for each client requestor each clock cycle. Thus, while a certain client requestor may be denied arbitration during a particular clock cycle due to requesting masking, the client requestor may be granted arbitration in subsequent clock cycles. Further, while a certain client requestor's request may be masked during a particular clock cycle, one or more other client requestors' requests may be passed on to the arbitration component during the same clock cycle without masking.

FIG. **5** is a flow diagram of an example method to activate a request mask, in accordance with some embodiments of the present disclosure. The method **500** can be performed by processing logic that can include hardware (e.g., processing device, circuitry, dedicated logic, programmable logic, microcode, hardware of a device, integrated circuit, etc.), software (e.g., instructions run or executed on a processing device), or a combination thereof. In some embodiments, the method **500** is performed by the buffer manager **113** of FIG. **1**. Although shown in a particular sequence or order, unless otherwise specified, the order of the processes can be modified. Thus, the illustrated embodiments should be understood only as examples, and the illustrated processes can be performed in a different order, and some processes can be performed in parallel. Additionally, one or more processes can be omitted in various embodiments. Thus, not all processes are required in every embodiment. Other process flows are possible.

At operation **505**, the processing device determines a throttling mode that indicates a manner of throttling arbitration requests to an arbiter that manages access of multiple clients to a buffer memory. The processing device is responsible for throttling arbitration requests of a particular client. As discussed above, each arbitration request is generated in response to receiving an access request from the particular client requesting access to the buffer memory. The processing device can determine the throttling mode by accessing a throttling mode value stored in a throttling mode configuration register. The throttling mode value can be set based on external input or based on input provided by system firmware.

At operation **510**, the processing device generates a throttling value. The processing device can be configured to operate in multiple throttling modes, and the processing device generates the throttling value according to the throttling mode to which the processing device is configured. For example, in a counter-based throttling mode, the processing device uses an incrementing counter as the throttling value. As another example, in a random-number-based throttling mode, the processing device generates a random (or pseudo random) number as the throttling value. Further details regarding the generation of the throttling value are discussed below in reference to FIG. **6**.

At operation **515**, the processing device determines a threshold value. The threshold value is stored in a threshold value register (e.g., threshold value register **300**), and the processing device can determine the threshold value by accessing the value from the threshold value register. The threshold value can be determined by system firmware or based on external input (e.g., received via a CPU peripheral programming interface).

At operation **520**, the processing device compares the throttling value to the threshold value. If, based on the comparison, the processing device determines that the throttling value is less than the threshold value, the processing device generates a request mask at operation **525**. As discussed above, an incoming arbitration request is masked based on detecting the request mask.

If, based on the comparison, the processing device determines that the throttling value is equal to or greater than the threshold value, the processing device de-activates the request mask at operation **530**. As discussed above, if no request mask is detected, an incoming arbitration request mask is passed to an arbitration component (e.g., the arbiter **206**) by the processing device.

As shown in FIG. **6**, the method **500** can, in some embodiments, include operations **605**, **610**, **615**, and **620**.



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Consistent with these embodiments, the operations **605**, **610**, **615**, and **620** can be performed as part of operation **510**, where the processing device generates the throttling value.

At operation **605**, the processing device increments a counter. The counter can be incremented on a per-clock cycle or per-request basis. That is, the processing device can increment the counter each clock cycle or upon receipt of each arbitration request.

At operation **610**, the processing device generates a random (or pseudo random) number. As discussed above, the processing device can utilize a LFSR-based random number generator to generate a pseudo-random number sequence, where the random number range is based on the width of the LFSR registers used. Accordingly, the generating of the random number can include accessing an output of a LFSR-based random number generator.

At operation **615**, the processing device generates a full-back pressure value. As noted above, the full-back pressure value is, by design, less than the threshold value. Accordingly, in generating the full-back pressure value, the processing device accesses the threshold value from the threshold value register and decrements the threshold value by one or more units to generate the full-back pressure value.

At operation **620**, the processing device selects the throttling value based on the determined throttling mode. In a first example, the processing device selects the counter as the throttling value while in a counter-based throttling mode. In a second example, the processing device selects the random number as the throttling value while in a random number-based throttling mode. In a third example, the processing device selects the full-back pressure value as the throttling value while in the full-back pressure mode.

FIG. 7 illustrates an example machine of a computer system **700** within which a set of instructions, for causing the machine to perform any one or more of the methodologies discussed herein, can be executed. In some embodiments, the computer system **700** can correspond to a host system (e.g., the host system **120** of FIG. 1) that includes, is coupled to, or utilizes a memory sub-system (e.g., the memory sub-system **110** of FIG. 1) or can be used to perform the operations of a controller (e.g., to execute an operating system to perform operations corresponding to the buffer manager **113** of FIG. 1). In alternative embodiments, the machine can be connected (e.g., networked) to other machines in a local area network (LAN), an intranet, an extranet, and/or the Internet. The machine can operate in the capacity of a server or a client machine in client-server network environment, as a peer machine in a peer-to-peer (or distributed) network environment, or as a server or a client machine in a cloud computing infrastructure or environment.

The machine can be a personal computer (PC), a tablet PC, a set-top box (STB), a Personal Digital Assistant (PDA), a cellular telephone, a web appliance, a server, a network router, a switch or bridge, or any machine capable of executing a set of instructions (sequential or otherwise) that specify actions to be taken by that machine. Further, while a single machine is illustrated, the term “machine” shall also be taken to include any collection of machines that individually or jointly execute a set (or multiple sets) of instructions to perform any one or more of the methodologies discussed herein.

The example computer system **700** includes a processing device **702**, a main memory **704** (e.g., read-only memory (ROM), flash memory, dynamic random access memory (DRAM) such as synchronous DRAM (SDRAM) or Ram-

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bus DRAM (RDRAM), etc.), a static memory **706** (e.g., flash memory, static random access memory (SRAM), etc.), and a data storage system **718**, which communicate with each other via a bus **730**.

Processing device **702** represents one or more general-purpose processing devices such as a microprocessor, a central processing unit, or the like. More particularly, the processing device can be a complex instruction set computing (CISC) microprocessor, reduced instruction set computing (RISC) microprocessor, very long instruction word (VLIW) microprocessor, or a processor implementing other instruction sets, or processors implementing a combination of instruction sets. Processing device **702** can also be one or more special-purpose processing devices such as an ASIC, a FPGA, a digital signal processor (DSP), network processor, or the like. The processing device **702** is configured to execute instructions **726** for performing the operations and steps discussed herein. The computer system **700** can further include a network interface device **708** to communicate over the network **720**.

The data storage system **718** can include a machine-readable storage medium **724** (also known as a computer-readable medium) on which is stored one or more sets of instructions **726** or software embodying any one or more of the methodologies or functions described herein. The instructions **726** can also reside, completely or at least partially, within the main memory **704** and/or within the processing device **702** during execution thereof by the computer system **700**, the main memory **704** and the processing device **702** also constituting machine-readable storage media. The machine-readable storage medium **724**, data storage system **718**, and/or main memory **704** can correspond to the memory sub-system **110** of FIG. 1.

In one embodiment, the instructions **726** include instructions to implement functionality corresponding to a memory allocation system (e.g., the buffer manager **113** of FIG. 1). While the machine-readable storage medium **724** is shown in an example embodiment to be a single medium, the term “machine-readable storage medium” should be taken to include a single medium or multiple media that store the one or more sets of instructions. The term “machine-readable storage medium” shall also be taken to include any medium that is capable of storing or encoding a set of instructions for execution by the machine and that cause the machine to perform any one or more of the methodologies of the present disclosure. The term “machine-readable storage medium” shall accordingly be taken to include, but not be limited to, solid-state memories, optical media, and magnetic media.

Some portions of the preceding detailed descriptions have been presented in terms of algorithms and symbolic representations of operations on data bits within a computer memory. These algorithmic descriptions and representations are the ways used by those skilled in the data processing arts to most effectively convey the substance of their work to others skilled in the art. An algorithm is here, and generally, conceived to be a self-consistent sequence of operations leading to a desired result. The operations are those requiring physical manipulations of physical quantities. Usually, though not necessarily, these quantities take the form of electrical or magnetic signals capable of being stored, combined, compared, and otherwise manipulated. It has proven convenient at times, principally for reasons of common usage, to refer to these signals as bits, values, elements, symbols, characters, terms, numbers, or the like.

It should be borne in mind, however, that all of these and similar terms are to be associated with the appropriate physical quantities and are merely convenient labels applied



to these quantities. The present disclosure can refer to the action and processes of a computer system, or similar electronic computing device, that manipulates and transforms data represented as physical (electronic) quantities within the computer system's registers and memories into other data similarly represented as physical quantities within the computer system memories or registers or other such information storage systems.

The present disclosure also relates to an apparatus for performing the operations herein. This apparatus can be specially constructed for the intended purposes, or it can include a general purpose computer selectively activated or reconfigured by a computer program stored in the computer. Such a computer program can be stored in a computer readable storage medium, such as, but not limited to, any type of disk including floppy disks, optical disks, CD-ROMs, and magnetic-optical disks, read-only memories (ROMs), random access memories (RAMs), EPROMs, EEPROMs, magnetic or optical cards, or any type of media suitable for storing electronic instructions, each coupled to a computer system bus.

The algorithms and displays presented herein are not inherently related to any particular computer or other apparatus. Various general purpose systems can be used with programs in accordance with the teachings herein, or it can prove convenient to construct a more specialized apparatus to perform the method. The structure for a variety of these systems will appear as set forth in the description below. In addition, the present disclosure is not described with reference to any particular programming language. It will be appreciated that a variety of programming languages can be used to implement the teachings of the disclosure as described herein.

The present disclosure can be provided as a computer program product, or software, that can include a machine-readable medium having stored thereon instructions, which can be used to program a computer system (or other electronic devices) to perform a process according to the present disclosure. A machine-readable medium includes any mechanism for storing information in a form readable by a machine (e.g., a computer). In some embodiments, a machine-readable (e.g., computer-readable) medium includes a machine (e.g., a computer) readable storage medium such as a ROM, RAM, magnetic disk storage media, optical storage media, flash memory components, and the like.

In the foregoing specification, embodiments of the disclosure have been described with reference to specific example embodiments thereof. It will be evident that various modifications can be made thereto without departing from the broader scope of embodiments of the disclosure as set forth in the following claims. The specification and drawings are, accordingly, to be regarded in an illustrative sense rather than a restrictive sense.

#### EXAMPLES

Example 1 is a system comprising: a first memory component; and a processing device, operatively coupled to the first memory component, configured to perform operations comprising: determining a throttling mode that indicates a manner of throttling arbitration requests to an arbiter that manages access to the buffer memory by multiple clients, each arbitration request being generated in response to receiving an access request for the buffer memory from a particular client from among the multiple clients; generating a throttling value based on the throttling mode; comparing

the throttling value to a threshold value; generating a request mask based on a result of comparing the throttling value to the threshold value; and in response to detecting the request mask, masking, using the request mask, an arbitration request from the client prior to passing the arbitration request to the arbiter.

In Example 2, the subject matter of Example 1 optionally comprises a threshold comparison register to store the threshold value, and the processing device determines the threshold value by accessing the threshold comparison register.

In Example 3, the subject matter of any one or more of Examples 1 and 2 optionally comprises a throttling mode configuration register to store a throttling mode value, and the determining of the throttling mode is based on the throttling mode value stored in the throttling mode configuration register.

In Example 4, the generating of the throttling value of any one or more of Examples 1-3 optionally comprises selecting one of a counter or a random number as the throttling value based on the determined throttling mode.

In Example 5, the generating of the throttling value of any one or more of Examples 1-4 optionally comprises incrementing the counter, and generating the random number.

In Example 6, the generating of the random number in Example 5 optionally comprises accessing an output of a linear feedback shift register (LFSR) based random number generator.

In Example 7, the throttling mode of any one or more of Examples 1-6, optionally comprises one of: a counter-based throttling mode, a random number-based throttling mode, and a full-back pressure throttling mode.

In Example 8, the subject matter of any one of Examples 1-7 optionally comprises generating the throttling value based on a counter-based throttling mode by accessing a counter that is incremented each clock cycle

In Example 9, the subject matter of any one of Examples 1-7 optionally comprises generating the throttling value based on a random number-based throttling mode by accessing a random number.

In Example 10, the subject matter of any one of Examples 1-7 optionally comprises generating the throttling value based on a full-back pressure throttling mode by generating a full-back pressure value by decrementing the threshold value by at least one unit; and selecting the full-back pressure value as the throttling value.

Example 11 includes the subject matter of any one of the Examples 1-10 wherein the throttling value is optionally a first throttling value generated at a first clock cycle; the arbitration request is optionally a first arbitration request; and the operations optionally comprise: generating, at a second clock cycle, a second throttling value based on with the throttling mode; comparing the second throttling value to the threshold value; de-activating the request mask based on a result of comparing the second throttling value to the threshold value; and passing, during the second clock cycle, a second arbitration request to the arbitration component.

Example 12 is a method comprising accessing a throttling mode configuration register to determine a throttling mode that defines a manner of throttling arbitration requests to an arbiter that manages access to a buffer memory by multiple clients, each arbitration request being generated in response to receiving an access request for the buffer memory from a particular client from among the multiple clients; generating a throttling value in accordance with the throttling mode; determining a threshold value by accessing a threshold comparison register, comparing the throttling value to the



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threshold value; generating a request mask based on a result of comparing the throttling value to the threshold value; and in response to detecting the request mask, masking, using the request mask, an arbitration request from the client prior to passing the arbitration request to the arbiter.

In example 13, the generating of the throttling value of Example 12 optionally comprises selecting one of a counter or a random number as the throttling value based on the determined throttling mode.

In example 14, the generating of the random number in Example 13 optionally comprises accessing an output of a LFSR-based random number generator.

In example 15, the throttling mode of any one or more of Examples 12-14, optionally comprises one of: a counter-based throttling mode, a random number-based throttling mode, and a full-back pressure throttling mode.

In example 16, the subject matter of any one or more of Examples 12-15 optionally comprises generating the throttling value in accordance with a counter-based throttling mode by accessing a counter that is incremented each clock cycle.

In example 17, the subject matter of any one or more of Examples 12-16 optionally comprises generating the throttling value in accordance with a random number-based throttling mode by accessing a random number.

In example 18, the subject matter of any one or more of Examples 12-17 optionally comprises generating the throttling value in accordance with a full-back pressure throttling mode by: generating a full-back pressure value by decrementing the threshold value by at least one unit; and selecting the full-back pressure value as the throttling value.

Example 19 includes the subject matter of any one of Examples 12-18 wherein the throttling value is optionally a first throttling value generated at a first clock cycle; the arbitration request is optionally a first arbitration request; and the method optionally further comprises: generating, at a second clock cycle, a second throttling value in accordance with throttling mode; comparing the second throttling value to the threshold value; de-activating the request mask based on a result of comparing the second throttling value to the threshold value; and passing, during the second clock cycle, a second arbitration request to the arbitration component.

Example 20, A non-transitory computer-readable storage medium comprising instructions that, when executed by a processing device, cause the processing device to perform operations comprising: accessing a throttling mode configuration register to determine a throttling mode that defines a manner of throttling arbitration requests to an arbiter that manages access to a buffer memory by multiple clients, each arbitration request being generated in response to receiving an access request for the buffer memory from a particular client from among the multiple clients; generating a throttling value in accordance with the throttling mode; determining a threshold value by accessing a threshold comparison register, comparing the throttling value to the threshold value; de-activating a request mask based on a result of comparing the throttling value to the threshold value; and based on the request mask being de-activated, passing an arbitration request to the arbiter for arbitration.

What is claimed is:

1. A method comprising:

accessing a throttling mode value from a throttling mode configuration register to determine a throttling mode that defines a manner of throttling arbitration requests to an arbiter that manages access to a buffer memory by multiple clients, each arbitration request being gener-

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ated in response to receiving an access request for the buffer memory from a particular client from among the multiple clients;

selecting a throttling value from a group of values using the throttling mode value, the group of values comprising a counter incremented at each clock cycle, a random number, and a full-back pressure value; determining a threshold value by accessing a threshold comparison register;

comparing the throttling value to the threshold value; generating a request mask based on a result of comparing the throttling value to the threshold value; and in response to detecting the request mask, masking, using the request mask, an arbitration request from the client prior to passing the arbitration request to the arbiter.

2. The method of claim 1, comparing of the throttling value to the threshold value includes determining the throttling value is less than the threshold value.

3. The method of claim 1, wherein determining the throttling value comprises accessing a random number output by a linear feedback shift register (LFSR) based random number generator.

4. The method of claim 1, wherein the throttling mode is selected from multiple throttling modes, the multiple throttling modes including:

a counter-based throttling mode, a random number-based throttling mode, and a full-back pressure throttling mode.

5. The method of claim 4, wherein determining the throttling value in accordance with the counter-based throttling mode comprises accessing the counter that is incremented each clock cycle.

6. The method of claim 4, wherein the determining the throttling value in accordance with the random number-based throttling mode comprises accessing a random number.

7. The method of claim 4, wherein the determining the throttling value in accordance with the full-back pressure throttling mode comprises:

generating the full-back pressure value by decrementing the threshold value by at least one unit; and selecting the full-back pressure value as the throttling value.

8. The method of claim 1, wherein:

the throttling value is a first throttling value generated at a first clock cycle;

the arbitration request is a first arbitration request; and the method further comprises:

generating, at a second clock cycle, a second throttling value in accordance with throttling mode; comparing the second throttling value to the threshold value;

de-activating the request mask based on a result of comparing the second throttling value to the threshold value; and

passing, during the second clock cycle, a second arbitration request to the arbitration component.

9. A non-transitory computer-readable storage medium comprising instructions that, when executed by a processing device, cause the processing device to perform operations comprising:

accessing a throttling mode value from a throttling mode configuration register to determine a throttling mode that defines a manner of throttling arbitration requests to an arbiter that manages access to a buffer memory by multiple clients, each arbitration request being gener-



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ated in response to receiving an access request for the buffer memory from a particular client from among the multiple clients;

selecting a throttling value from a group of values using the throttle mode value, the group of values comprising: a counter incremented at each clock cycle, a random number, or a full-back pressure value;

determining a threshold value by accessing a threshold comparison register;

comparing the throttling value to the threshold value;

de-activating a request mask based on a result of comparing the throttling value to the threshold value; and based on the request mask being de-activated, passing an arbitration request to the arbiter for arbitration.

**10.** A system comprising:

a plurality of memory components, at least one of the memory components being configured as a buffer memory; and

a processing device, operatively coupled with the plurality of memory components, to perform operations comprising:

determining, based on a throttling mode value stored in a throttling mode configuration register, a throttling mode that indicates a manner of throttling arbitration requests to an arbiter that manages access to the buffer memory by multiple clients, each arbitration request being generated in response to receiving an access request for the buffer memory from a particular client from among the multiple clients;

selecting a throttling value from a group of values using the throttling mode value, the group of values comprising a counter incremented at each clock cycle, a random number, and a full-back pressure value;

comparing the throttling value to a threshold value;

generating a request mask based on a result of comparing the throttling value to the threshold value; and

in response to detecting the request mask, masking, using the request mask, an arbitration request from the client prior to passing the arbitration request to the arbiter.

**11.** The system of claim **10**, further comprising a threshold comparison register to store the threshold value, wherein the processing device is to determine the threshold value by accessing the threshold comparison register.

**12.** The system of claim **10**, further comprising a throttling mode configuration register to store a throttling mode

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value, wherein the determining of the throttling mode is based on the throttling mode value stored in the throttling mode configuration register.

**13.** The system of claim **10**, wherein the comparing of the throttling value to the threshold value includes determining the throttling value is less than the threshold value.

**14.** The system of claim **13**, wherein the determining of the throttling value further comprises:

incrementing the counter; and

generating the random number.

**15.** The system of claim **14**, wherein generating the random number comprises accessing an output of a linear feedback shift register (LFSR).

**16.** The system of claim **10**, wherein the throttling mode comprises one of: a counter-based throttling mode, a random number-based throttling mode, and a full-back pressure throttling mode.

**17.** The system of claim **16**, wherein determining the throttling value based on the counter-based throttling mode comprises selecting the counter that is incremented each clock cycle.

**18.** The system of claim **16**, wherein the determining the throttling value based on the random number-based throttling mode comprises selecting the random number.

**19.** The system of claim **16**, wherein the determining the throttling value based on the full-back pressure throttling mode comprises:

generating the full-back pressure value by decrementing the threshold value by at least one unit; and

selecting the full-back pressure value as the throttling value.

**20.** The system of claim **10**, wherein:

the throttling value is a first throttling value generated at a first clock cycle;

the arbitration request is a first arbitration request; and the operations further comprise:

generating, at a second clock cycle, a second throttling value based on the throttling mode;

comparing the second throttling value to the threshold value;

de-activating the request mask based on a result of comparing the second throttling value to the threshold value; and

passing, during the second clock cycle, a second arbitration request to the arbitration component.

\* \* \* \* \*