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(54) **ENFORCING LIMITS ON A SELF-SERVE MODEL FOR PROVISIONING DATA VOLUMES FOR CONTAINERS RUNNING IN VIRTUAL MACHINES**

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This patent is subject to a terminal disclaimer.

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G06F 3/06 (2006.01)
G06F 9/455 (2018.01)

(52) **U.S. Cl.**

CPC ***G06F 3/0665*** (2013.01); ***G06F 3/0605*** (2013.01); ***G06F 3/067*** (2013.01); ***G06F 3/0608*** (2013.01); ***G06F 3/0619*** (2013.01); ***G06F 3/0631*** (2013.01); ***G06F 3/0632*** (2013.01); ***G06F 3/0653*** (2013.01); ***G06F 3/0689*** (2013.01); ***G06F 9/45558*** (2013.01); *G06F 2009/45562* (2013.01); *G06F 2009/45579* (2013.01); *G06F 2009/45583* (2013.01)

(58) **Field of Classification Search**

None

See application file for complete search history.

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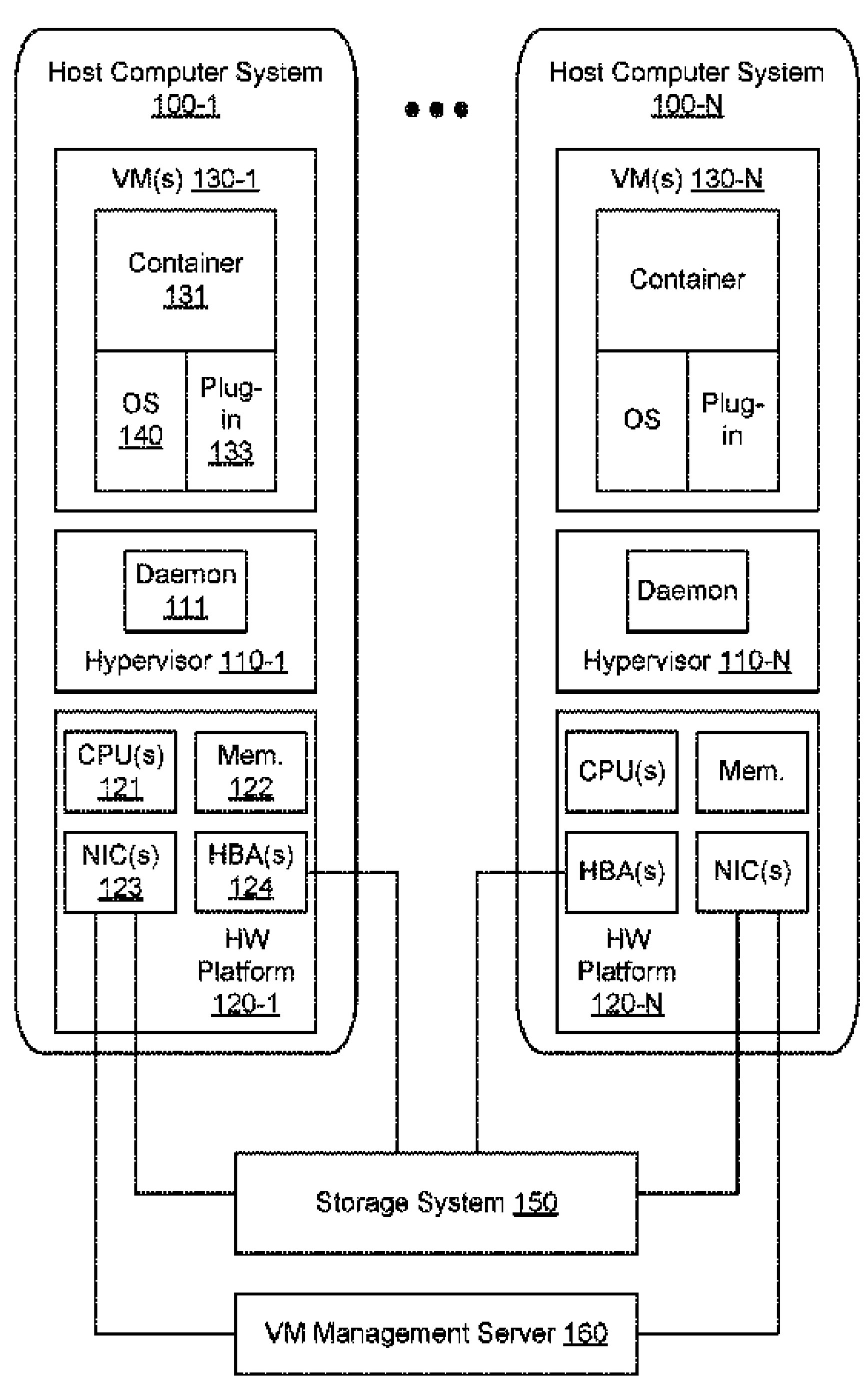
(57) **ABSTRACT**

A computer system has a virtualization software that supports execution of a virtual machine in which a container is run. A method of managing allocation of storage resources to the container includes the steps of monitoring a virtual socket, detecting, based on the monitoring, a request from a plug-in of the container to create a data volume, upon detecting the request, retrieving a storage resource limit that has been set for the first virtual machine and determining if creation of the data volume causes the storage resource limit to be exceeded, and communicating the request to the virtualization software to cause the virtualization software to create the data volume if the limit is determined to be not exceeded and returning an error if the limit is determined to be exceeded.

20 Claims, 4 Drawing Sheets



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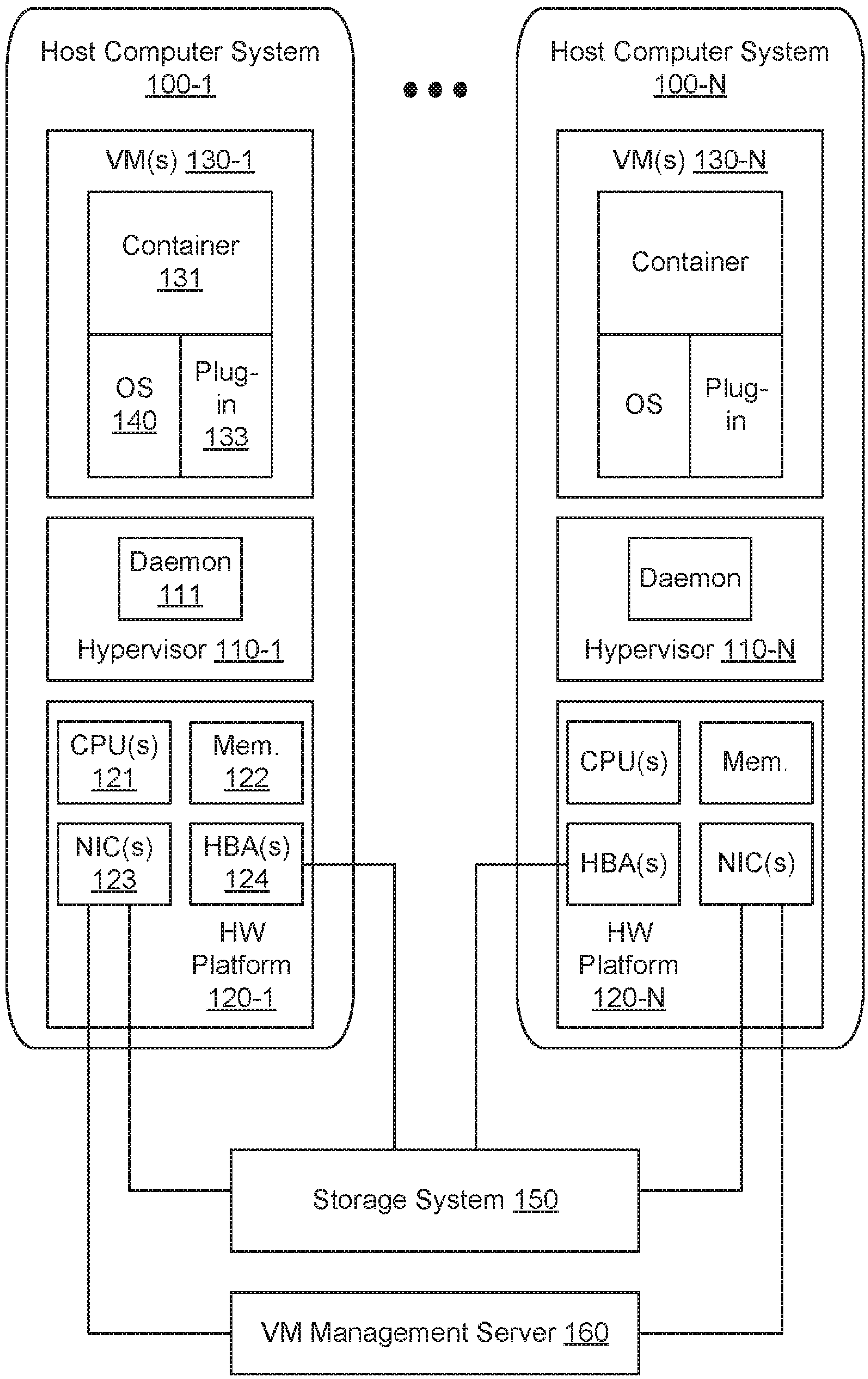
Host Computer System 100-1
VM(s) 130-1
Container 131
OS 140
Plug-in 133
Daemon 111
Hypervisor 110-1
CPU(s) 121
Mem. 122
NIC(s) 123
HBA(s) 124
HW Platform 120-1
Host Computer System 100-N
VM(s) 130-N
Container
OS
Plug-in
Daemon
Hypervisor 110-N
CPU(s)
Mem.
HBA(s)
NIC(s)
HW Platform 120-N
Storage System 150
VM Management Server 160

FIGURE 1

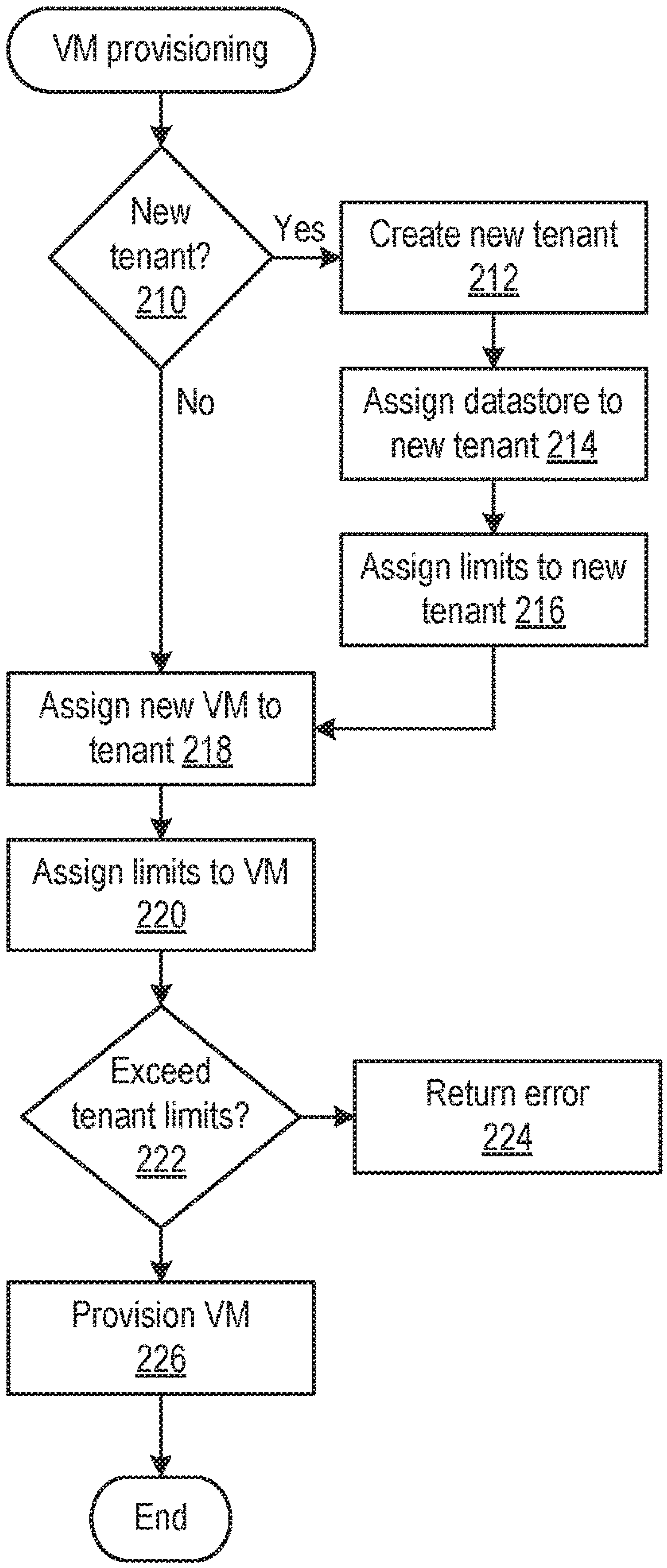

FIGURE 2A

Tenant	Datastore	Max Size	Max Disks	Max IOPS
ABC	Flash	100 GB	1000	50000
XYZ	Hybrid	1 TB	10000	20000
...	...	...	...	...

FIGURE 2B

VM	Tenant	Max Size	Max Disks	Max IOPS
0	ABC	10 GB	50	10000
1	XYZ	500 GB	100	8000
2	XYZ	100 GB	50	5000
3	ABC	50 GB	100	10000
...	...	...	...	...

FIGURE 2C

Volume	VM
radio2016	0
volume1	0
volume2	1
volume3	2
volume4	3
...	...

FIGURE 2D

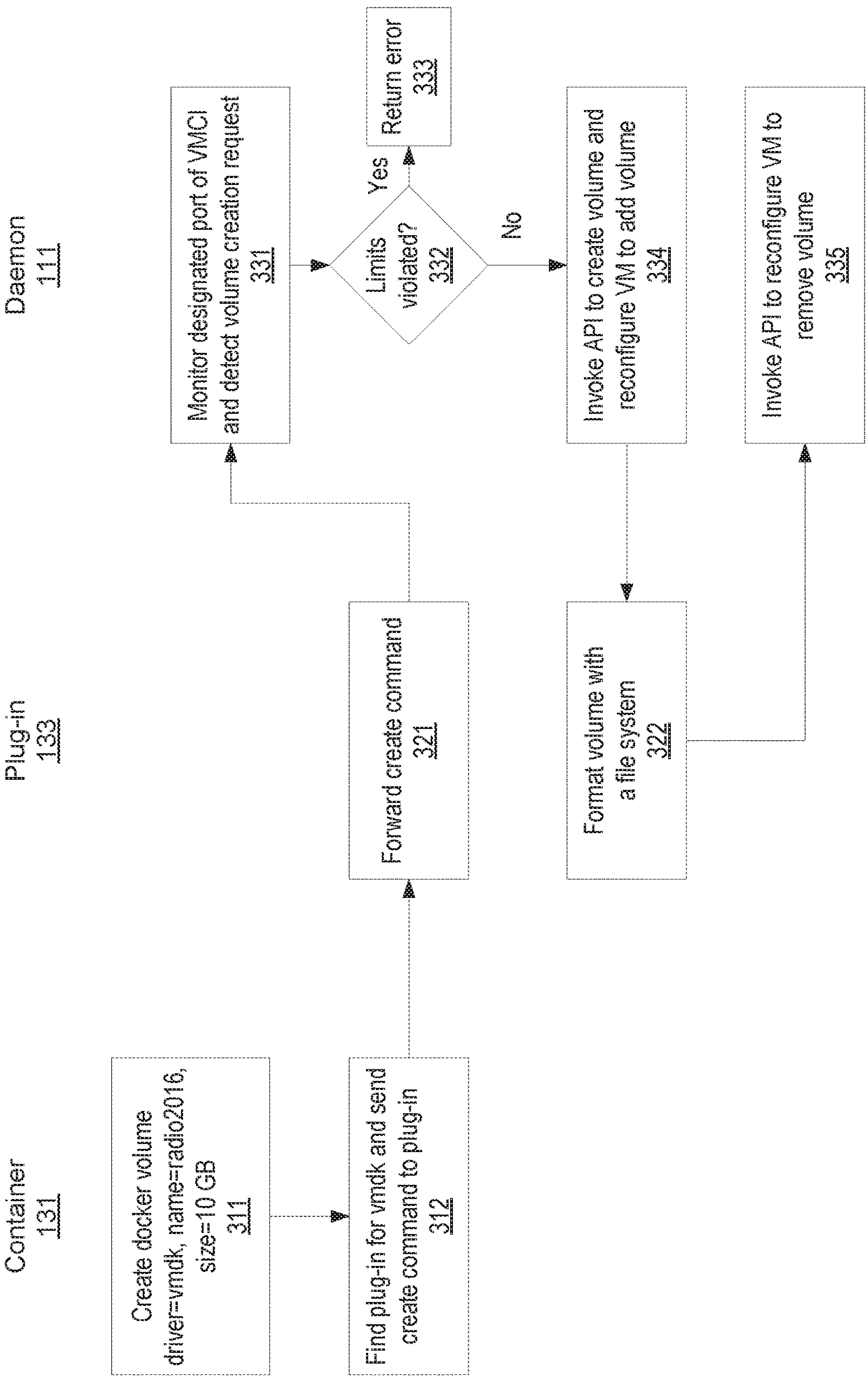
Container
131
Create docker volume driver=vmdk, name=radio2016, size=10 GB
311
Find plug-in for vmdk and send create command to plug-in
312
Plug-in
133
Forward create command
321
Format volume with a file system
322
Daemon
111
Monitor designated port of VMCI and detect volume creation request
331
Limits violated?
332
Yes
Return error
333
No
Invoke API to create volume and reconfigure VM to add volume
334
Invoke API to reconfigure VM to remove volume
335

FIGURE 3

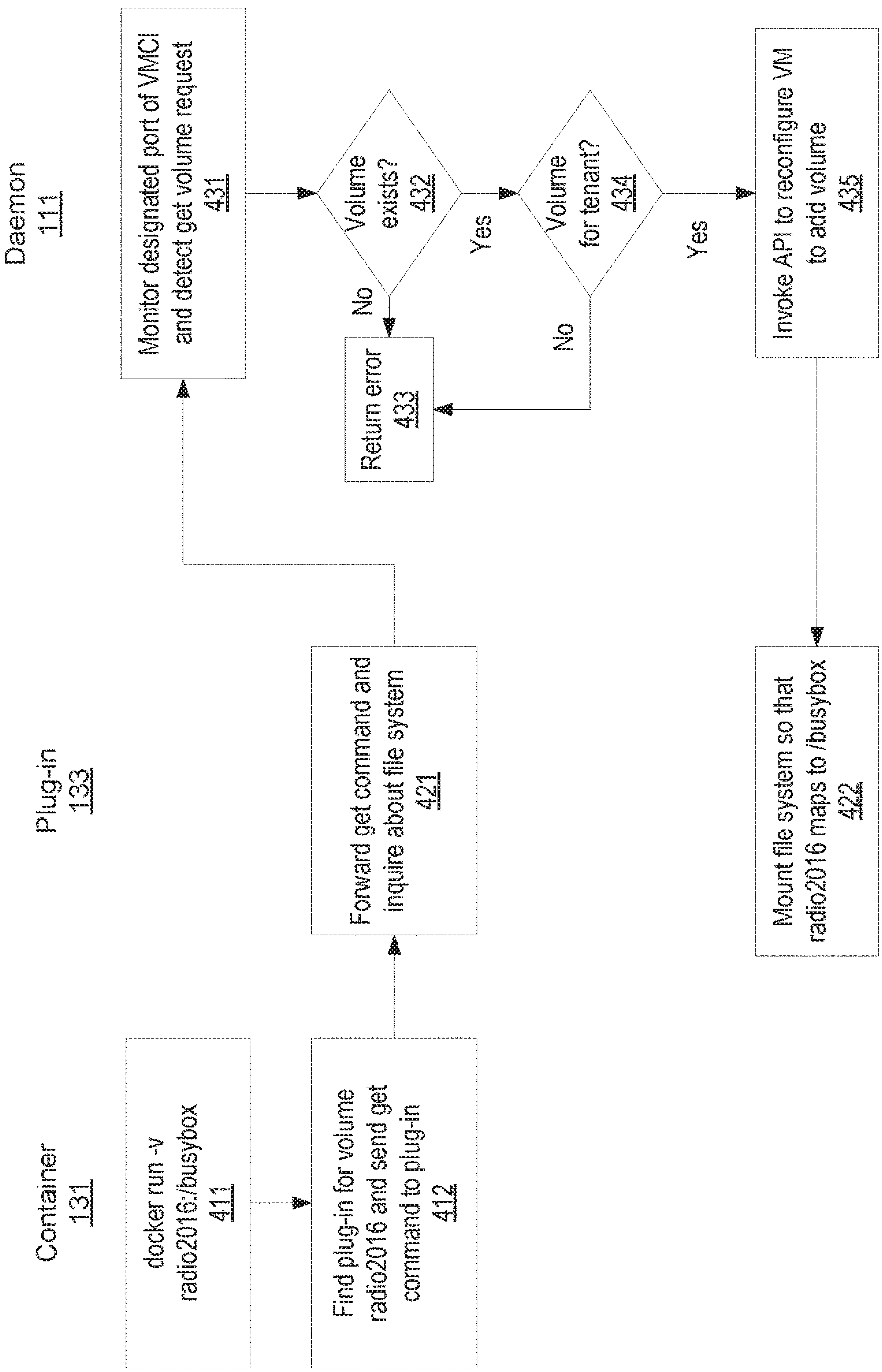
Container
131
docker run -v
radio2016:/busybox
411
Find plug-in for volume
radio2016 and send get
command to plug-in
412
Plug-in
133
Forward get command and
inquire about file system
421
Mount file system so that
radio2016 maps to /busybox
422
Daemon
111
Monitor designated port of VMCI
and detect get volume request
431
Volume
exists?
432
Yes
No
Return error
433
Volume
for tenant?
434
Yes
No
Invoke API to reconfigure VM
to add volume
435

FIGURE 4

ENFORCING LIMITS ON A SELF-SERVE MODEL FOR PROVISIONING DATA VOLUMES FOR CONTAINERS RUNNING IN VIRTUAL MACHINES

CROSS-REFERENCE TO RELATED APPLICATION(S)

This application claims the benefit of priority from U.S. Provisional Patent Application No. 62/343,780, filed May 31, 2016, which is incorporated by reference herein.

BACKGROUND

Increasingly, decisions to provision resources and manage resources are made by application logic, e.g., containers, running within virtual machines (VMs), and they typically require a self-serve-programmatic model for provisioning and management. Some frameworks can choose to create an instance of a container image and attach persistent storage (e.g., data volumes) to the container image, all within the VM.

However, there exist challenges when trying to meet the need for a self-serve-programmatic model. Some existing management stacks require manual steps, including opening up a user interface (UI) and directing the provisioning of data volumes through the UI. Other existing management stacks require invoking of a remote application programming interface (API) to a control plane for provisioning data volumes. This latter technique typically also requires per VM configuration.

SUMMARY

One or more embodiments provide a control plane for data volume management that can be invoked within a container that is spun up within a VM. One example of a data volume is a virtual disk. More generally, a "data volume" is a place where the container can store data persistently. The control plane is configured as a daemon or other service that is running in the user space of a hypervisor that is supporting the execution of the VM and listens in on a virtual socket provisioned within the VM.

Advantages of employing the control plane within the hypervisor, according to embodiments, are as follows. First, it does not require human intervention to carry out the data volume provisioning requested by the application administrator. Second, the control plane is local to the VM and does not require any additional configuration beyond the installation of the data volume plug-in software in the VM.

In one embodiment, to protect against untrusted plug-ins from sending control operations to a control plane within the hypervisor, the control plane requires control operations passed thereto to originate from software running in the root mode. As a result, only those plug-ins that are trusted software (e.g., signed with proper cryptographic keys) will be able to send control operations successfully to the control plane. For example, control operations sent to the control plane via third party plug-ins, which would be running in non-root mode, will be not be accepted by the control plane.

A method of method of managing allocation of storage resources to the container includes the steps of monitoring a virtual socket, detecting, based on the monitoring, a request from a plug-in of the container to create a data volume, upon detecting the request, retrieving a storage resource limit that has been set for the first virtual machine and determining if creation of the data volume causes the storage resource limit to be exceeded, and communicating the request to a virtualization software that supports the execution of a virtual machine in which the container is running, to cause the virtualization software to create the data volume if the limit is determined to be not exceeded and returning an error if the limit is determined to be exceeded.

Further embodiments include, without limitation, a non-transitory computer-readable medium that includes instructions that enable a processor to implement one or more aspects of the above method as well as a computer system having a processor, memory, and other components that are configured to implement one or more aspects of the above method.

BRIEF DESCRIPTION OF THE DRAWINGS

FIG. **1** is a block diagram of a virtualized computing environment in which embodiment may be practiced.

FIG. **2**A is a flow diagram of a method of setting storage allocation limits during provisioning of the virtual machine.

FIG. **2**B is a conceptual diagram of a data structure that is used to track storage allocation limits set for tenants.

FIG. **2**C is a conceptual diagram of a data structure that is used to track storage allocation limits set for virtual machines.

FIG. **2**D is a conceptual diagram of a data structure that is used to track data volumes that have been created for virtual machines.

FIG. **3** is a flow diagram of a method of creating a data volume according to embodiments.

FIG. **4** is a flow diagram of a method of mapping a data volume to a namespace according to embodiments.

DETAILED DESCRIPTION

FIG. **1** is a block diagram of a virtualized computing environment in which embodiments may be practiced. The virtualized computing environment of FIG. **1** includes a cluster of host computer systems **100-1** to **100**-N, where N is 2 or more. Alternatively, embodiments may be practiced in a virtualized computing environment that includes only a single host computer system. Host computer system **100-1** has a hardware platform **120-1** that includes one or more central processing units (CPUs) **121**, system memory **122** (typically volatile dynamic random access memory), one or more network interface controllers (NICs) **123**, and one or more host bus adapters (HBAs) **124**. Each of the other host computer systems **100**, including host computer system **100**-N which has a hardware platform **120**-N, includes the same (or similar) hardware components as hardware platform **120-1**. In addition, a hypervisor is installed in each of host computer systems **100** as system software. Hypervisor **110-1** supports the execution space of virtual machines (VMs) **130-1** and hypervisor **110**-N supports the execution space of VMs **130**-M. Hereinafter, VMs will be generally referred to as VM **130** or VMs **130** and the hypervisor supporting the VMs **130** will be generally referred to as hypervisor **110**.

As further illustrated in FIG. **1**, a container **131** runs inside VM **130-1** on top of an operating system (OS) **140** of VM **130-1**. One example of container **131** is a Docker® container that runs on top of a Linux® operating system. Typically, container **131** includes a management layer (known as a container engine) on top of OS **140** and one or more applications deployed therein to run on top of the management layer.

In FIG. **1**, a plug-in **133** is also illustrated. Plug-in **133**, which is implemented as part of the plug-in framework of the container (e.g., as part of Docker® plug-in framework for Docker® containers), is configured to communicate with hypervisor **110-1** over a virtual socket provisioned by hypervisor **110-1** as part of a virtual hardware platform for VM **130-1**. The virtual socket is also referred to as a back channel, and enables VM **130-1** to communicate with hypervisor **110-1**. In one embodiment, the virtual socket is implemented as shared memory, such as with virtual machine control interface (VMCI) employed in virtualization products available from VMware, Inc. of Palo Alto, Calif., and is accessed through VMCI ports. More specifically, daemon **111** runs in a user space of hypervisor **110-1** to listen in on this virtual socket and, in the embodiments, passes on control operations received through this virtual socket to hypervisor **110-1** for execution using standard APIs. Examples of these standard APIs include creating a data volume, deleting a data volume, attaching a data volume, and detaching a data volume. Accordingly, operations to create, delete, attach, or detach a data volume can be instigated within container **131** and such control operations are "plumbed" to plug-in **133** that forwards those control operations over the virtual socket to daemon **111**, which calls the standard APIs to perform control operations on the data volume.

A virtual machine management server (VMMS) **160** manages VMs across host computers systems **100**. The execution of the VMs is supported by the hypervisors of the respective host computer systems **100**. The standard APIs exposed by hypervisor **110** for creating, deleting, attaching, and detaching a data volume are made accessible through a user interface of VMMS **160** so that control operations for data volumes of VMs (e.g., virtual disks) can be instigated by a VM administrator.

The data volumes for the container or the VMs are stored in storage system **150**. In the embodiment illustrated in FIG. **1**, storage system **150** is a shared storage system, which is accessible from host computer systems **100** through their HBAs **124**. In another embodiment, storage system **150** may be network-attached storage (NAS) or virtual storage area network (SAN), which is accessible from host computer systems **100** over a network through their NICs **123**.

According to embodiments, the data volume control plane is implemented in hypervisor **110** through daemon **111** which is listening in on the virtual socket through which plug-in **133** forwards data volume control operations. As data volume control operations are passed down from container **131** to plug-in **133** and forwarded onto the virtual socket, daemon **111**, upon detection of the data volume control operation, invokes the standard APIs exposed by hypervisor **110** for provisioning data volumes. As a way to protect against untrusted applications or plug-ins from gaining access to the data volume control plane, any application or plug-in not running in root mode are blocked from gaining access to the data volume control plane. This is implemented by daemon **111** listening in on a privileged virtual socket, i.e., the virtual socket that is accessed through a privileged VMCI port. As such, any control operations forwarded onto a non-privileged virtual socket will be ignored by daemon **111**. Accordingly, in the embodiments, plug-in **133** is implemented as a secure module that runs in root mode. In order to preserve its image and to protect it against tampering, the executable code of this secure module is signed with cryptographic keys of a trusted entity.

In addition, the VM administrator who is managing the virtualized computing environment the infrastructure can set bounds on data volume provisioning. The application administrator is free to perform data volume control operations so long as they are within these bounds. The bounds include quotas (capacity), what kind of volumes, and how many volumes. Roles are also defined by the VM administrator. The roles specify which VMs may create or delete, which VMs may read or write. In addition, the VM administrator is given the ability to view and inspect the run time of the VMs (which data volumes were created by whom, who is consuming them, which volumes are unused, how much data was written, etc.)

FIG. **2**A is a flow diagram of a method of setting storage allocation limits during provisioning of the virtual machine. The steps of this method are carried out at VMMS **160** in response to an instruction to provision a new VM received through the UI of VMMS **160**.

At step **210**, VMMS **160** determines if a new tenant is requesting the provisioning of a new VM. If so, at step **212**, VMMS **160** creates the new tenant, leading to a creation of an entry for that tenant in the table of FIG. **2**B. According to the designation of a datastore made by the new tenant, at step **214**, VMMS **160** populates the "datastore" attribute of the tenant's entry with the designated datastore. At step **216**, storage limits of the designated datastore are populated into the corresponding attributes of the tenant's entry. In the embodiments illustrated herein, the storage limits defined for a tenant include maximum size (indicating maximum storage capacity), maximum number of disks (or data volumes), and maximum IOPS.

Then, VMMS **160** at step **218** assigns the new VM to be provisioned to the tenant, and at step **220** assigns storage limits for the new VM. In the embodiments illustrated herein, the storage limits for a VM include maximum size (indicating maximum storage capacity), maximum number of disks (or data volumes), and maximum IOPS. If the storage limits of the VM assigned at step **220** cause the storage limits for the tenant, which are determined from the table of FIG. **2**B, to be exceeded, VMMS **160** will not permit the new VM to be provisioned and returns an error at step **224**. On the other hand, if the storage limits of the VM assigned at step **220** do not cause the storage limits for the tenant to be exceeded, VMMS **160** provisions the new VM at **226** and the table of FIG. **2**C will be updated to include an entry for the newly provisioned VM. As illustrated, the entry for the newly provisioned VM includes the following attributes: VM ID, tenant ID, maximum size, maximum number of disks, and maximum TOPS.

FIG. **3** is a flow diagram of a method of creating a data volume according to embodiments. The method illustrated in FIG. **3** is carried out by container **131**, plug-in **133**, and daemon **111**. When the application administrator desires to create a data volume for container **131**, the application administrator enters command line instructions for creating the data volume at step **311**, e.g., "create docker volume, driver=vmdk, name=radio2016, size=10 GB." In response to the command line instruction entered at step **311**, container **131** searches for a plug-in of the driver indicated in the command, in this example, vmdk, and sends the create data volume command to the plug-in (step **312**).

At step **321**, the plug-in, e.g., plug-in **133**, upon receipt of the create data volume command from container **131**, forwards the create data volume command to daemon **111** through a virtual socket. In particular, plug-in **133** invokes a virtual socket API to forward the create data volume command to the virtual socket through a privileged VMCI port (e.g., a VMCI port that has been pre-designated as a privileged port).

Daemon **111** runs as a background process in the user space of hypervisor **110**, and listens in on (monitors) the privileged virtual socket for new requests at step **331**. Upon detecting a create data volume request, daemon **111** consults the table of FIG. **2**C to determine if the creation of the data volume of the size indicated violates the storage limits that have been set for the virtual machines that is hosting container **131**. If any of the storage limits is violated, e.g., exceeds maximum size or exceeds maximum number of data volumes, daemon **111** returns an error at step **333**. On the other hand, if none of the storage limits are violated, daemon **111** at step **334** invokes the standard APIs for (1) creating a data volume for the virtual machine that is hosting container **131**, and (2) reconfiguring the virtual machine to add the data volume (i.e., updating the virtual machine configuration file to include an identifier for the newly provisioned data volume). In response to the APIs invoked at step **332**, hypervisor **110** provisions a new data volume, and the newly provisioned data volume becomes attached to the virtual machine (i.e., the newly provisioned data volume is enumerated as one of the devices of the virtual machine). In addition, daemon **111** updates the table of FIG. **2**D, which is maintained in memory **122** and persisted in storage system **150**, to add the newly created data volume in association of with the virtual machine that is hosting container **131**.

At step **322**, plug-in **133** formats the data volume with a file system. A file system specified by the application administrator in the command line instructions may be used in formatting the data volume. If no such file system is specified, a default file system is used.

After the data volume has been formatted with the file system at step **322**, the control returns to daemon **111**, at which time daemon **111** at step **335** invokes the standard API for reconfiguring the virtual machine to detach the data volume (i.e., updating the virtual machine configuration file to remove the identifier for the newly provisioned data volume). In response to the API invoked at step **335**, the newly provisioned data volume becomes detached from the virtual machine (i.e., the newly provisioned data volume is no longer enumerated as one of the devices of the virtual machine).

FIG. **4** is a flow diagram of a method of mapping a data volume to a namespace according to embodiments. The method illustrated in FIG. **4** is carried out by container **131**, plug-in **133**, and daemon **111**, and in response to a container run command. When the application administrator desires to map a data volume to a namespace for container **131**, the application administrator enters command line instructions to run the container at step **411**, e.g., "docker run, radio2016:/busybox." When this particular command line instruction is executed within container **131**, container **131** is spun up using data volume, radio2016, mapped to the namespace/busybox. Also, in response to the command line instruction entered at step **411**, container **131** locates the plug-in corresponding to the data volume indicated in the command, in this example, radio2016, and sends a get data volume command to the plug-in (step **412**).

At step **421**, the plug-in, e.g., plug-in **133**, upon receipt of the get data volume command from container **131**, forwards the get data volume command to daemon **111** through a virtual socket. In particular, plug-in **133** invokes a virtual socket API to forward the get data volume command to the virtual socket through the privileged VMCI port.

Daemon **111** listens in on (monitors) the privileged virtual socket for new requests at step **431**. Upon detecting a get data volume request, daemon **111** at step **432** checks the table of FIG. **2**D to see if the data volume exists. If no such data volume exists, daemon **111** returns an error at step **433**. If the data volume exists, daemon **111** at step **434** checks the table of FIG. **2**C to see if the data volume belongs to the same tenant to whom the virtual machine hosting container **131** is assigned. If so, the flow proceeds to step **435**. If not, daemon **111** returns an error at step **433**.

At step **435**, daemon **111** invokes the standard APIs for reconfiguring the virtual machine to add the data volume (i.e., updating the virtual machine configuration file to include an identifier for the data volume). In response to the APIs invoked at step **435**, the data volume becomes attached to the virtual machine (i.e., the data volume is enumerated as one of the devices of the virtual machine).

In response to the virtual socket API invoked at step **421**, plug-in **133** at step **422** receives a device ID corresponding to the data volume from daemon **111**, maps the device ID to the data volume, and mounts the file system of the data volume into the namespace used by container **131** so that the data volume can be mapped to a folder accessible by container **131**, e.g., so that the volume, radio2016, can be mapped to the/busybox folder.

In the example given above, a container that instigated creation of a data volume may be the same or different from a container that is run using that data volume. In addition, a container that instigated creation of a data volume may be running in a first virtual machine and a container that is run using that data volume may be running in a second virtual machine so long as the two virtual machine are assigned to the same tenant. The first and second virtual machines may be executed in the same or different host computer systems so long as the host computer systems are accessing the same storage system in which the data volume is provisioned.

Certain embodiments as described above involve a hardware abstraction layer on top of a host computer. The hardware abstraction layer allows multiple contexts or emulated computing instances to share the hardware resource. In one embodiment, these emulated computing instances are isolated from each other, each having at least a user application running therein. The hardware abstraction layer thus provides benefits of resource isolation and allocation among the emulated computing instances. In the foregoing embodiments, emulated machines are used as an example for the emulated computing instances and hypervisors as an example for the hardware abstraction layer. As described above, each emulated machine includes a guest operating system in which at least one application runs.

The various embodiments described herein may employ various computer-implemented operations involving data stored in computer systems. For example, these operations may require physical manipulation of physical quantities usually, though not necessarily, these quantities may take the form of electrical or magnetic signals, where they or representations of them are capable of being stored, transferred, combined, compared, or otherwise manipulated. Further, such manipulations are often referred to in terms, such as producing, identifying, determining, or comparing. Any operations described herein that form part of one or more embodiments of the invention may be useful machine operations. In addition, one or more embodiments of the invention also relate to a device or an apparatus for performing these operations. The apparatus may be specially constructed for specific required purposes, or it may be a general purpose computer selectively activated or configured by a computer program stored in the computer. In particular, various general purpose machines may be used with computer programs written in accordance with the teachings herein, or it may be more convenient to construct a more specialized apparatus to perform the required operations.

The various embodiments described herein may be practiced with other computer system configurations including hand-held devices, microprocessor systems, microprocessor-based or programmable consumer electronics, minicomputers, mainframe computers, and the like.

One or more embodiments of the present invention may be implemented as one or more computer programs or as one or more computer program modules embodied in one or more computer readable media. The term computer readable medium refers to any data storage device that can store data which can thereafter be input to a computer system computer readable media may be based on any existing or subsequently developed technology for embodying computer programs in a manner that enables them to be read by a computer. Examples of a computer readable medium include a hard drive, network attached storage (NAS), read-only memory, random-access memory (e.g., a flash memory device), a CD (Compact Discs) CD-ROM, a CD-R, or a CD-RW, a DVD (Digital Versatile Disc), a magnetic tape, and other optical and non-optical data storage devices. The computer readable medium can also be distributed over a network coupled computer system so that the computer readable code is stored and executed in a distributed fashion.

Although one or more embodiments of the present invention have been described in some detail for clarity of understanding, it will be apparent that certain changes and modifications may be made within the scope of the claims. Accordingly, the described embodiments are to be considered as illustrative and not restrictive, and the scope of the claims is not to be limited to details given herein, but may be modified within the scope and equivalents of the claims. In the claims, elements and/or steps do not imply any particular order of operation, unless explicitly stated in the claims.

Plural instances may be provided for components, operations or structures described herein as a single instance. Finally, boundaries between various components, operations and data stores are somewhat arbitrary, and particular operations are illustrated in the context of specific illustrative configurations. Other allocations of functionality are envisioned and may fall within the scope of the invention(s). In general, structures and functionality presented as separate components in exemplary configurations may be implemented as a combined structure or component. Similarly, structures and functionality presented as a single component may be implemented as separate components. These and other variations, modifications, additions, and improvements may fall within the scope of the appended claims.

What is claimed is:

1. In a computer system having a hypervisor on which a first virtual machine is run, the hypervisor supporting execution of the first virtual machine in which a container is run, the first virtual machine comprising a guest operating system on which the container executes, the container comprising one or more applications deployed therein, a method of managing allocation of storage resources to the container, comprising:

monitoring, by the hypervisor, a virtual socket of the first virtual machine, wherein the virtual socket enables the first virtual machine to communicate with the hypervisor;

based on said monitoring, detecting, by a component running within the hypervisor, a request from a plug-in of the container to create a data volume, the data volume backed by a storage device, wherein the plug-in is configured to provide an interface between the container and the hypervisor, and further wherein the hypervisor is configured to provide an interface between the plug-in and the storage device;

upon detecting the request, retrieving a storage resource limit that has been set for the first virtual machine and determining if creation of the data volume causes the storage resource limit to be exceeded; and

communicating, by the component, the request to the hypervisor to cause the hypervisor to create the data volume if the storage resource limit is determined to be not exceeded and returning an error if the storage resource limit is determined to be exceeded.

2. The method of claim **1**, wherein the storage resource limit is a total size of all data volumes created for the first virtual machine.

3. The method of claim **1**, wherein the storage resource limit is a total number of data volumes created for the first virtual machine.

4. The method of claim **1**, further comprising:

based on said monitoring, detecting a request from a plug-in of a container running in a second virtual machine to perform a control operation on the data volume;

determining if the first virtual machine and the second virtual machine are assigned to a same tenant;

if the first virtual machine and the second virtual machine are assigned to the same tenant, performing the control operation on the data volume; and

if the first virtual machine and the second virtual machine are not assigned to the same tenant, returning an error.

5. The method of claim **4**, further comprising:

determining that the data volume has been created prior to performing the control operation on the data volume.

6. The method of claim **5**, wherein the control operation is one of attaching the data volume to the second virtual machine, and deleting the data volume.

7. The method of claim **1**, further comprising:

for each new virtual machine to be provisioned, setting the storage resource limit for the new virtual machine.

8. The method of claim **7**, wherein provisioning of a new virtual machine for a tenant is blocked if a total of storage resource limits of all virtual machines of the tenant, including the new virtual machine to be provisioned, exceeds an aggregate limit set for the tenant.

9. The method of claim **1**, wherein the container comprises an application logic run within the first virtual machine.

10. A non-transitory computer readable medium comprising instructions to be executed in a computer system having a hypervisor on which a first virtual machine is run, the hypervisor supporting execution of the first virtual machine in which a container is run, the first virtual machine comprising a guest operating system on which the container executes, the container comprising one or more applications deployed therein, wherein the instructions when executed cause the computer system to carry out a method of managing allocation of storage resources to the container, said method comprising:

monitoring, by the hypervisor, a virtual socket of the first virtual machine, wherein the virtual socket enables the first virtual machine to communicate with the hypervisor;

based on said monitoring, detecting, by a component running within the hypervisor, a request from a plug-in of the container to create a data volume, the data volume backed by a storage device, wherein the plug-in is configured to provide an interface between the container and the hypervisor, and further wherein the hypervisor is configured to provide an interface between the plug-in and the storage device;

upon detecting the request, retrieving a storage resource limit that has been set for the first virtual machine and determining if creation of the data volume causes the storage resource limit to be exceeded; and

communicating, by the component, the request to the hypervisor to cause the hypervisor to create the data volume if the storage resource limit is determined to be not exceeded and returning an error if the storage resource limit is determined to be exceeded.

11. The non-transitory computer readable medium of claim **10**, wherein the storage resource limit is a total size of all data volumes created for the first virtual machine.

12. The non-transitory computer readable medium of claim **10**, wherein the storage resource limit is a total number of data volumes created for the first virtual machine.

13. The non-transitory computer readable medium of claim **10**, wherein the storage resource limit is a total of TOPS allocated for data volumes created for the first virtual machine.

14. The non-transitory computer readable medium of claim **10**, wherein the method further comprises:

based on said monitoring, detecting a request from a plug-in of a container running in a second virtual machine to perform a control operation on the data volume;

determining if the first virtual machine and the second virtual machine are assigned to a same tenant;

if the first virtual machine and the second virtual machine are assigned to the same tenant, performing the control operation on the data volume; and

if the first virtual machine and the second virtual machine are not assigned to the same tenant, returning an error.

15. The non-transitory computer readable medium of claim **14**, wherein the method further comprises:

determining that the data volume has been created prior to performing the control operation on the data volume.

16. The non-transitory computer readable medium of claim **15**, wherein the control operation is one of attaching the data volume to the second virtual machine, and deleting the data volume.

17. The non-transitory computer readable medium of claim **10**, wherein the method further comprises:

for each new virtual machine to be provisioned, setting the storage resource limit for the new virtual machine.

18. The non-transitory computer readable medium of claim **17**, wherein provisioning of a new virtual machine for a tenant is blocked if a total of storage resource limits of all virtual machines of the tenant, including the new virtual machine to be provisioned, exceeds an aggregate limit set for the tenant.

19. A computer system having a first host computer system including a first virtualization software on which a first virtual machine is run, the first virtualization software supporting execution of the first virtual machine in which a first container is run, the first virtual machine comprising a guest operating system on which the first container executes, the first container comprising one or more applications deployed therein, and a second host computer system including a second virtualization software supporting execution of a second virtual machine in which a second container is run, wherein the first virtualization software has a background process running therein to perform the steps of:

monitoring, by the first virtualization software, a first virtual socket of the first virtual machine, wherein the first virtual socket enables the first virtual machine to communicate with the first virtualization software;

based on said monitoring, detecting a request from a plug-in of the first container to create a data volume, the data volume backed by a storage device, wherein the plug-in is configured to provide an interface between the container and the first virtualization software, and further wherein the first virtualization software is configured to provide an interface between the plug-in and the storage device;

upon detecting the request, retrieving a storage resource limit that has been set for the first virtual machine and determining if creation of the data volume causes the storage resource limit to be exceeded; and

communicating the request to the first virtualization software to cause the first virtualization software to create the data volume if the storage resource limit is determined to be not exceeded and returning an error if the storage resource limit is determined to be exceeded.

20. The computer system of claim **19**, wherein the second virtualization software has a background process running therein to perform the steps of:

monitoring a virtual socket;

based on said monitoring, detecting a request from a plug-in of the second container to perform a control operation on the data volume;

determining if the first virtual machine and the second virtual machine are assigned to a same tenant;

if the first virtual machine and the second virtual machine are assigned to the same tenant, performing the control operation on the data volume; and

if the first virtual machine and the second virtual machine are not assigned to the same tenant, returning an error.

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