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Billa et al.

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(54) **NSP MANAGER**

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(*) Notice: Subject to any disclaimer, the term of this patent is extended or adjusted under 35 U.S.C. 154(b) by 716 days.

(21) Appl. No.: **14/207,933**

(22) Filed: **Mar. 13, 2014**

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(51) **Int. Cl.**
G06F 17/30 (2006.01)
G06F 7/00 (2006.01)
(Continued)

(52) **U.S. Cl.**
CPC ***G06F 17/30327*** (2013.01); ***G06F 3/067*** (2013.01); ***G06F 3/0619*** (2013.01);
(Continued)

(58) **Field of Classification Search**
CPC G06F 17/30961; G06F 17/30327; G06F 17/30864; G06F 17/30595;
(Continued)

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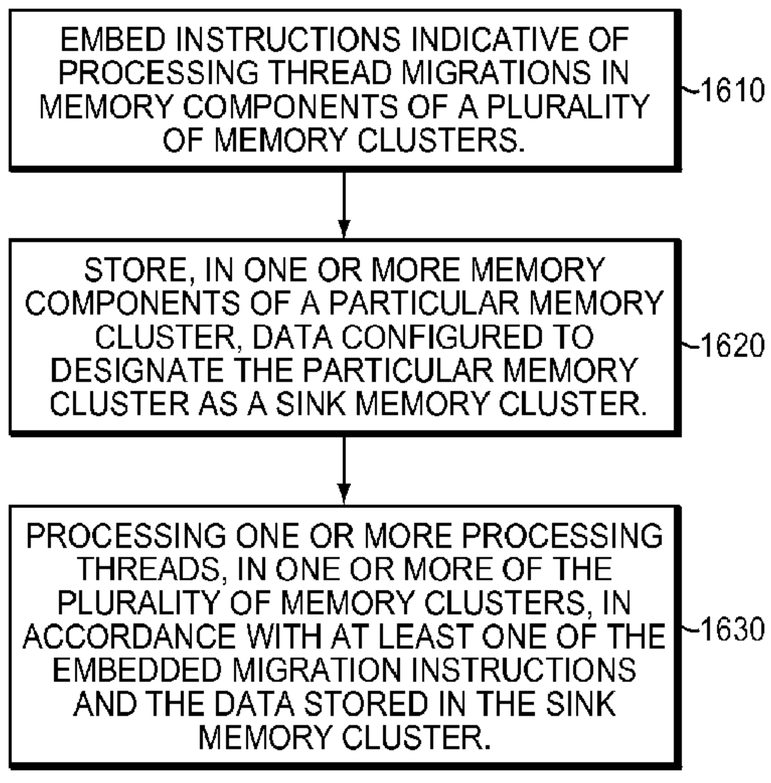
Assistant Examiner — Raquel Perez-Arroyo
(74) *Attorney, Agent, or Firm* — Hamilton, Brook, Smith & Reyonlds, P.C.

(57) **ABSTRACT**

In an embodiment, a method of updating a memory with a plurality of memory lines, the memory storing a tree, a plurality of buckets, and a plurality of rules, can include maintaining a copy of the memory with a plurality of memory lines. The method can further include writing a plurality of changes to at least one of the tree, the plurality of buckets, and the plurality of rules to the copy. The method can additionally include determining whether each of the plurality of changes is an independent write or a dependent write. The method can further include merging independent writes to the same line of the copy. The method further includes transferring updates from the plurality of lines of the copy to the plurality of lines of the memory.

21 Claims, 40 Drawing Sheets



(51) **Int. Cl.**
G06F 3/06 (2006.01)
G06F 11/16 (2006.01)

(52) **U.S. Cl.**
CPC ***G06F 3/0646*** (2013.01); ***G06F 11/167*** (2013.01); *G06F 11/1666* (2013.01)

(58) **Field of Classification Search**
CPC G06F 17/30625; G06F 17/30584; G06F 3/061; G06F 3/0611; G06F 3/0613; G06F 3/0635; G06F 3/0631; G06Q 10/06; H04L 45/745
USPC 707/797
See application file for complete search history.

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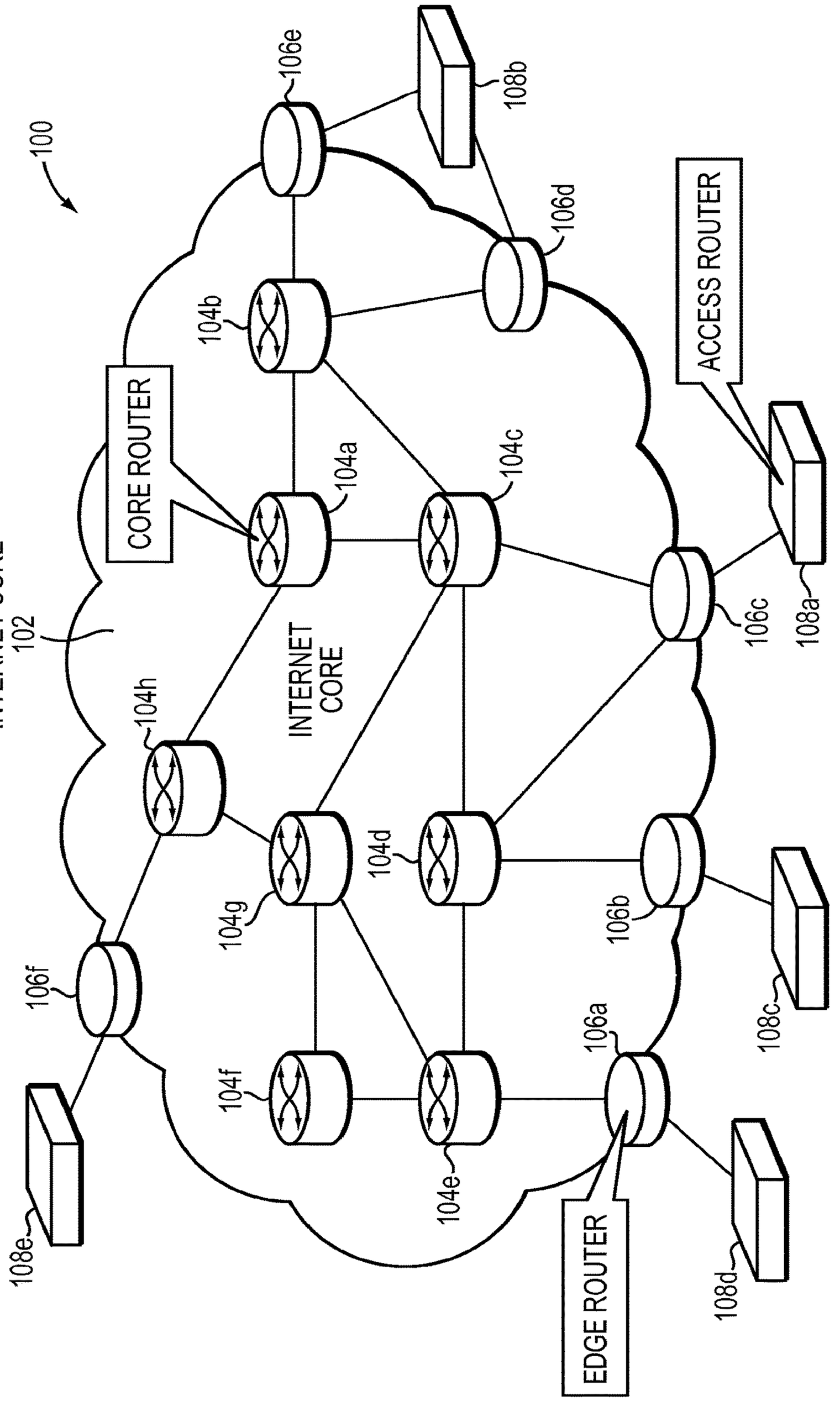
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INTERNET CORE
102
CORE ROUTER
ACCESS ROUTER
EDGE ROUTER
INTERNET CORE
104a
104b
104c
104d
104e
104f
104g
104h
106a
106b
106c
106d
106e
106f
108a
108b
108c
108d
108e

FIG. 1

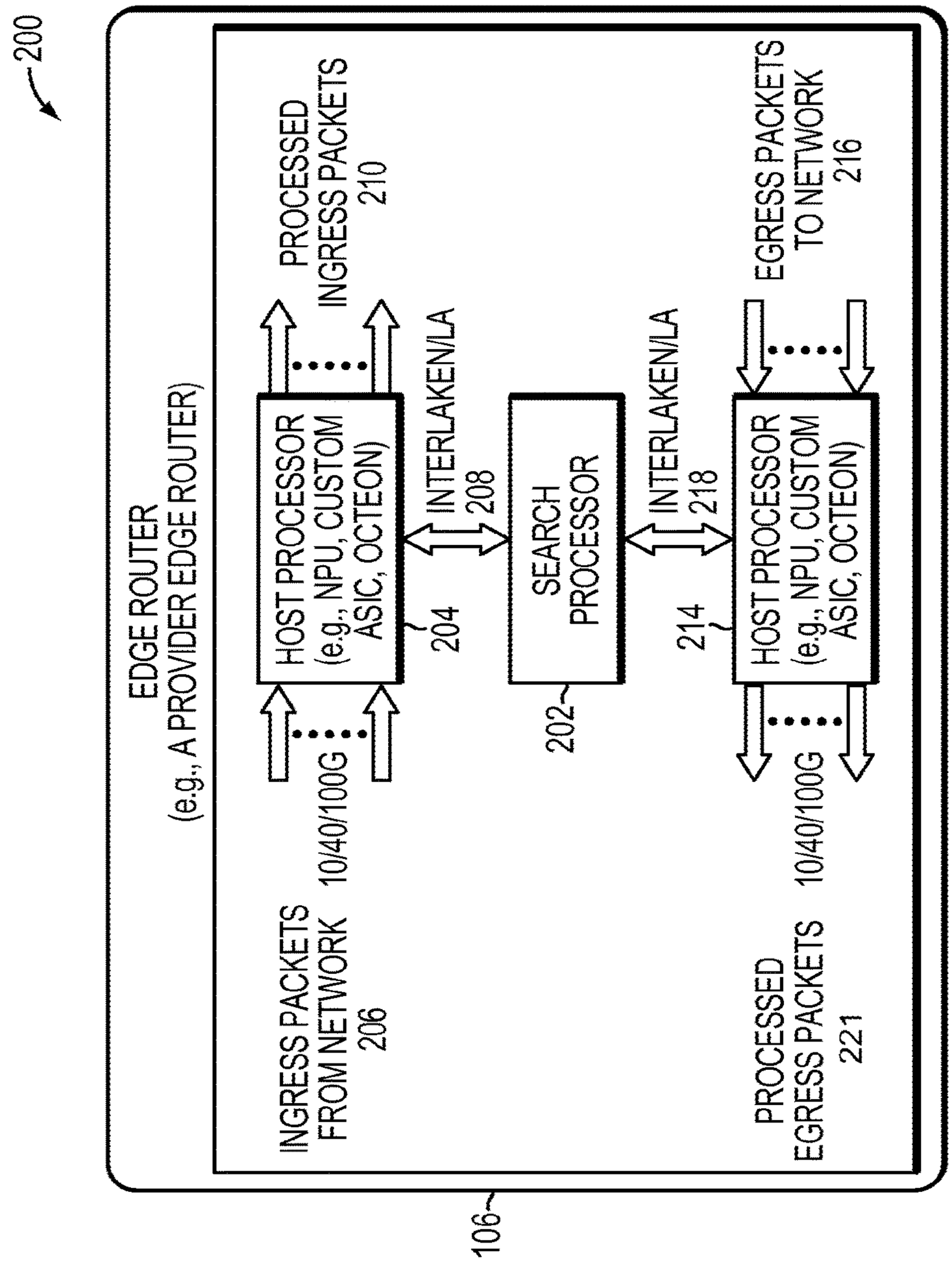

200
EDGE ROUTER
(e.g., A PROVIDER EDGE ROUTER)
INGRESS PACKETS FROM NETWORK 206
10/40/100G
HOST PROCESSOR (e.g., NPU, CUSTOM ASIC, OCTEON)
204
INTERLAKEN/LA 208
SEARCH PROCESSOR
202
INTERLAKEN/LA 218
214
HOST PROCESSOR (e.g., NPU, CUSTOM ASIC, OCTEON)
10/40/100G
PROCESSED EGRESS PACKETS 221
PROCESSED INGRESS PACKETS 210
EGRESS PACKETS TO NETWORK 216
106

FIG. 2A

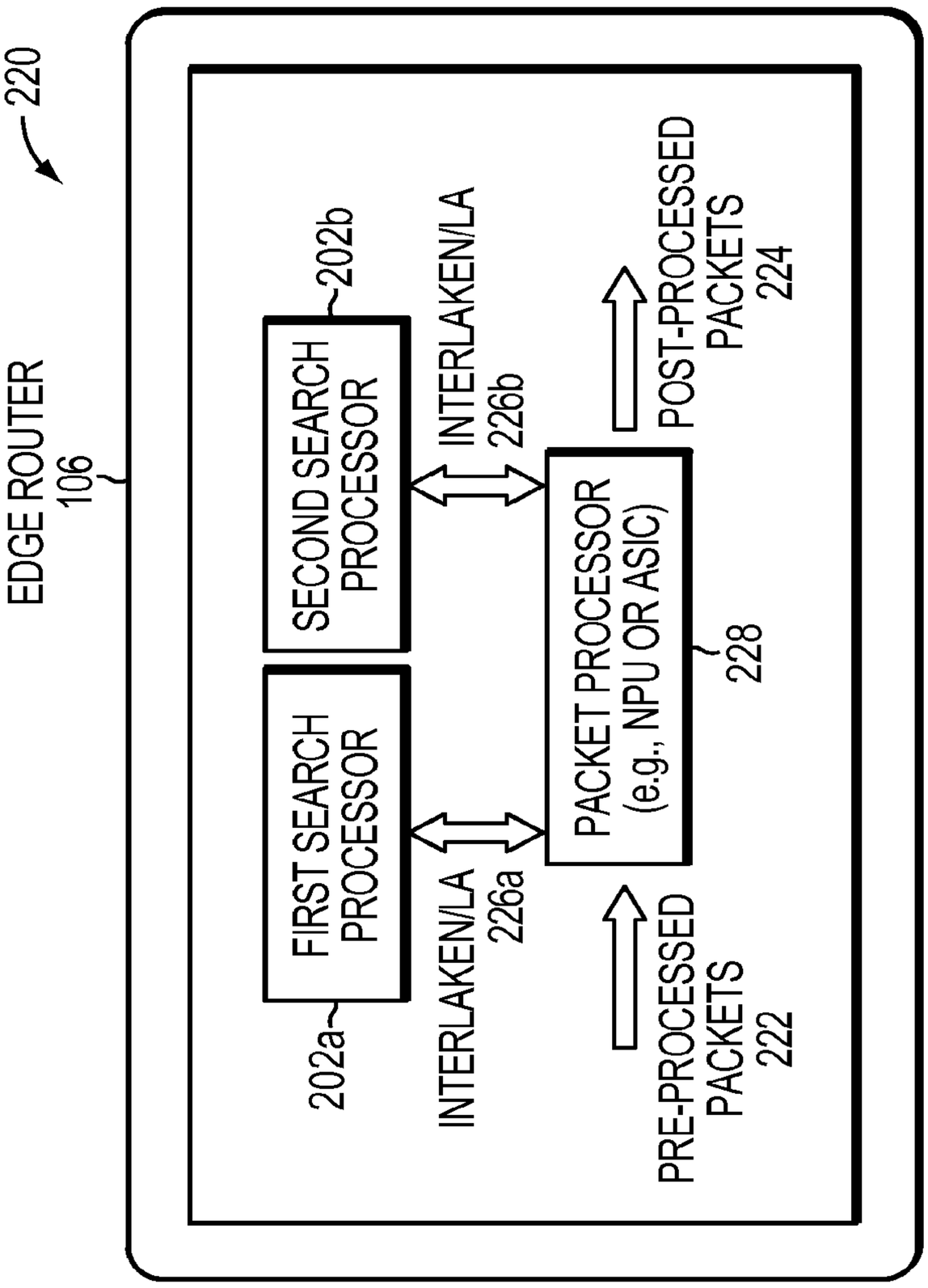

220
EDGE ROUTER
106
FIRST SEARCH PROCESSOR
202a
SECOND SEARCH PROCESSOR
202b
INTERLAKEN/LA
226a
INTERLAKEN/LA
226b
PACKET PROCESSOR (e.g., NPU OR ASIC)
228
PRE-PROCESSED PACKETS
222
POST-PROCESSED PACKETS
224

FIG. 2B

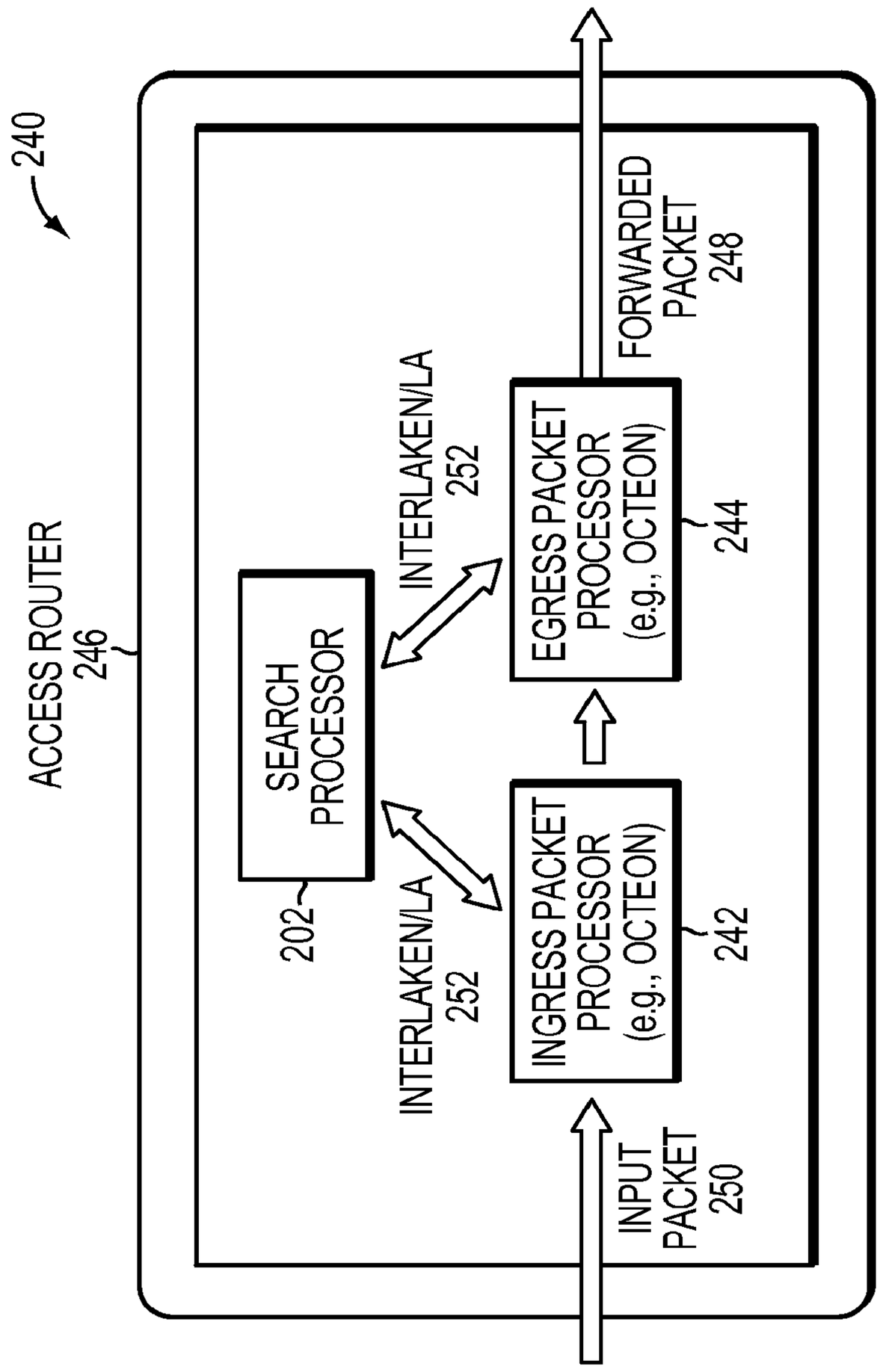

240
ACCESS ROUTER
246
SEARCH PROCESSOR
202
INTERLAKEN/LA
252
INTERLAKEN/LA
252
INGRESS PACKET PROCESSOR (e.g., OCTEON)
242
EGRESS PACKET PROCESSOR (e.g., OCTEON)
244
FORWARDED PACKET
248
INPUT PACKET
250

FIG. 2C

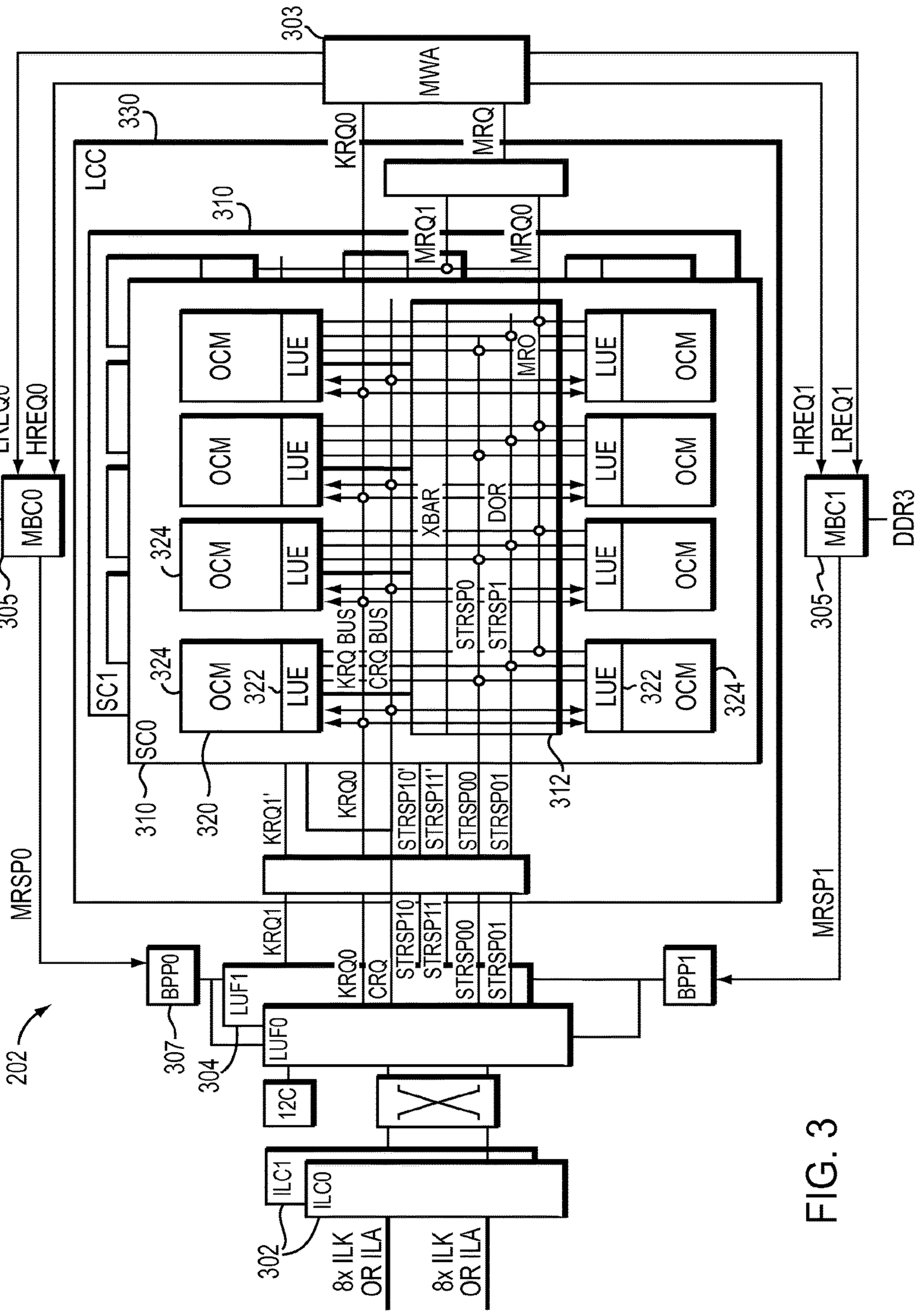

303
MWA
330
LCC
KRQ0
MRQ
310
MRQ1
MRQ0
OCM
LUE
XBAR
DOR
MRO
324
STRSP0
STRSP1
KRQ BUS
CRQ BUS
322
SC1
SC0
320
312
LREQ0
HREQ0
DDR3
MBC0
305
MRSP0
HREQ1
LREQ1
MBC1
MRSP1
KRQ1'
KRQ0
STRSP10'
STRSP11'
STRSP00
STRSP01
KRQ1
CRQ
STRSP10
STRSP11
BPP0
BPP1
LUF1
LUF0
307
304
202
12C
ILC1
ILC0
302
8x ILK OR ILA

FIG. 3

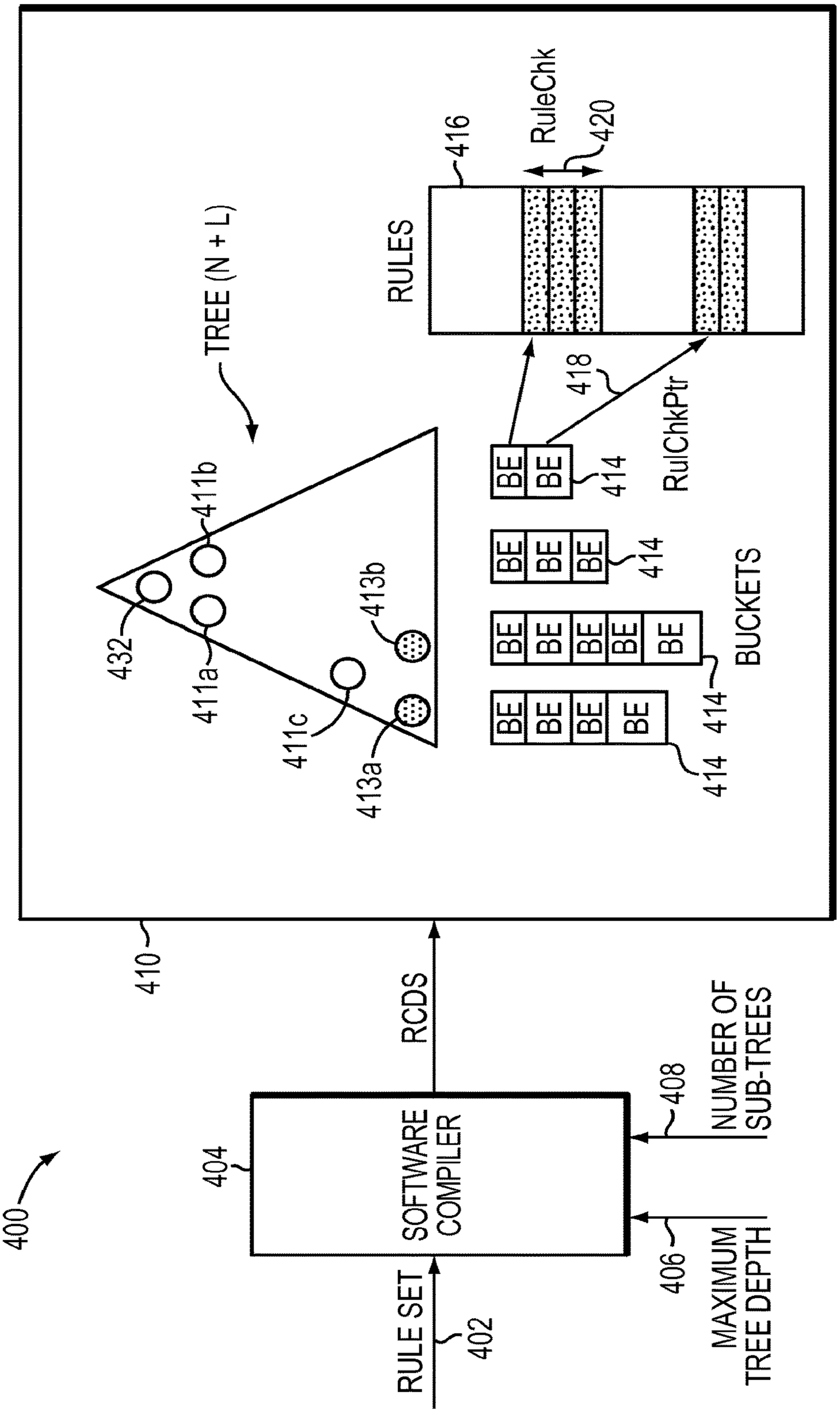

400
RULE SET
402
SOFTWARE COMPILER
404
MAXIMUM TREE DEPTH
406
NUMBER OF SUB-TREES
408
RCDS
410
TREE (N + L)
RULES
416
RuleChk
420
418
RulChkPtr
BE
414
BUCKETS
432
411a
411b
411c
413a
413b

FIG. 4

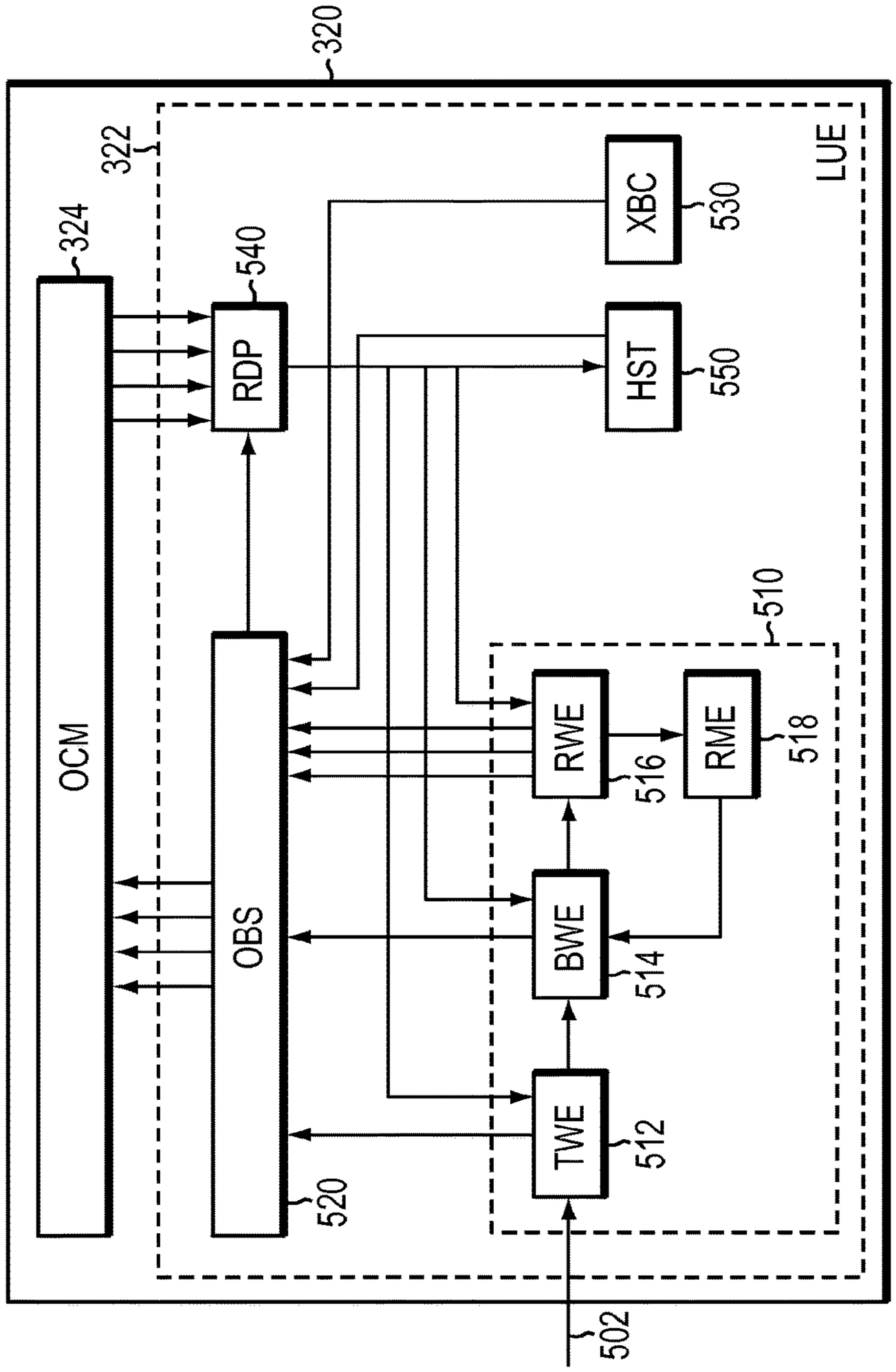

320
322
LUE
324
OCM
OBS
RDP
540
XBC
530
HST
550
510
RWE
516
RME
518
BWE
514
TWE
512
520
502

FIG. 5

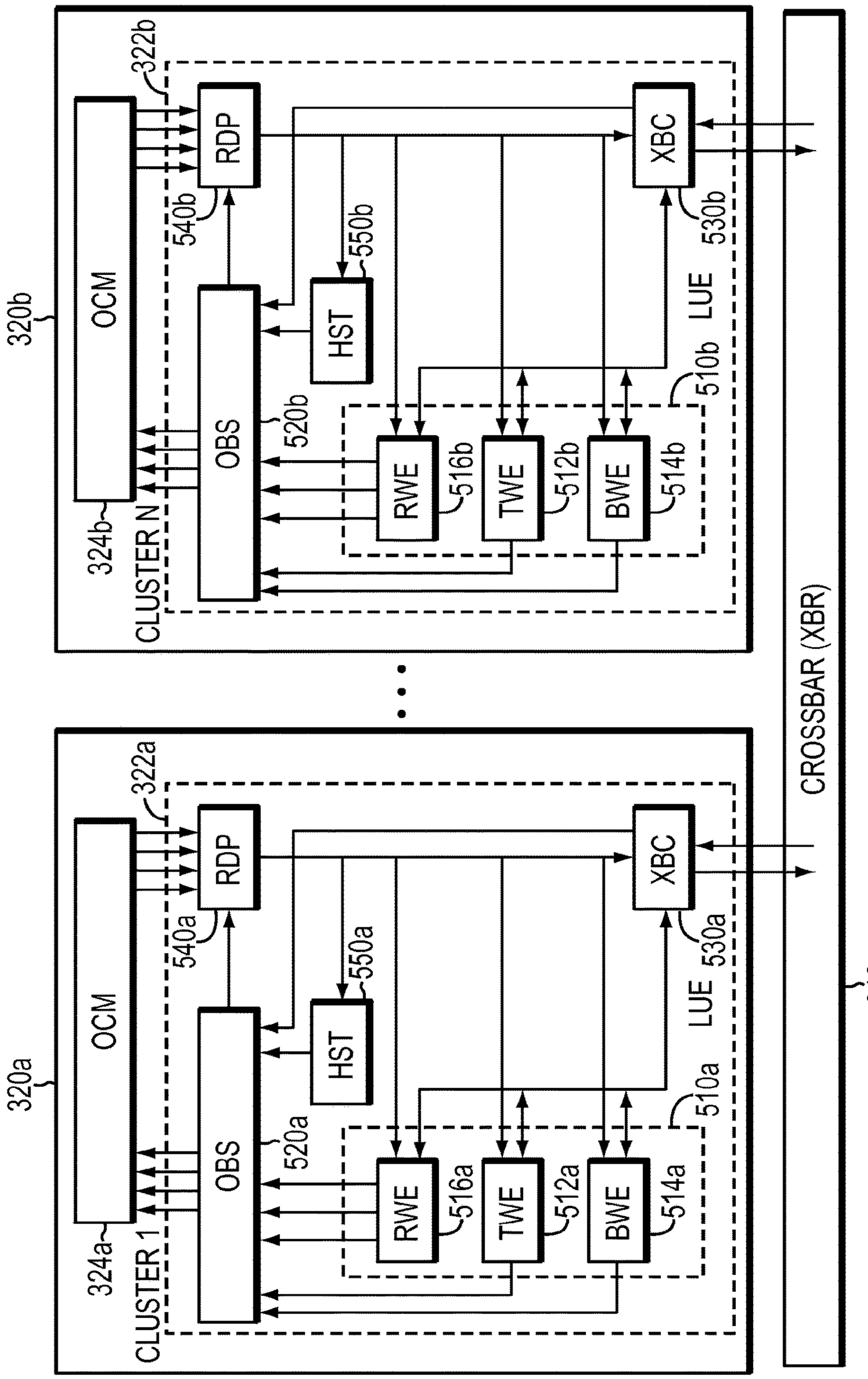

320b
322b
RDP
XBC
540b
550b
530b
OCM
HST
OBS
520b
510b
LUE
RWE
516b
TWE
512b
BWE
514b
324b
CLUSTER N
CROSSBAR (XBR)
320a
322a
540a
550a
530a
520a
510a
516a
512a
514a
324a
CLUSTER 1
312

FIG. 6A

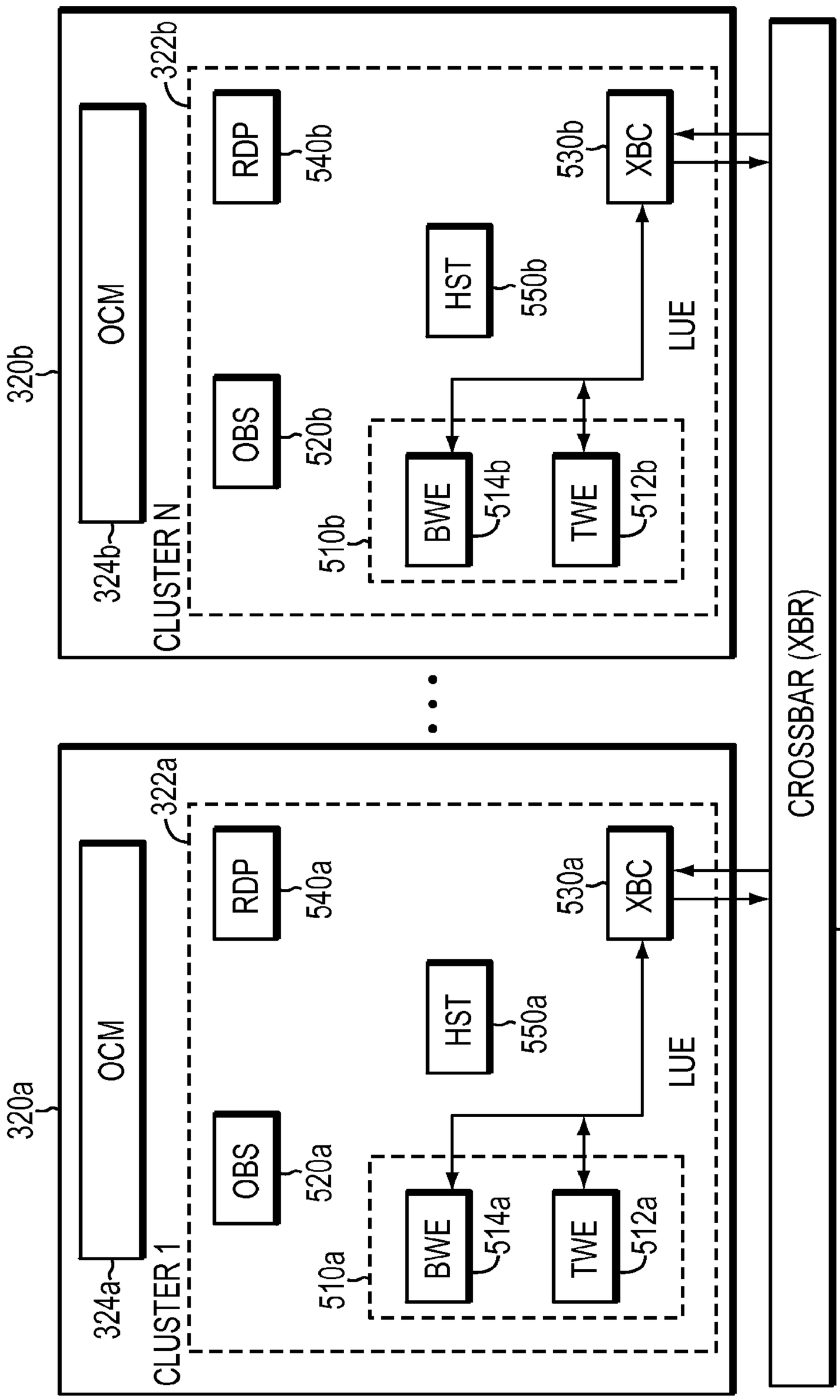

CLUSTER N
322b
RDP
540b
530b
XBC
HST
550b
LUE
OCM
320b
OBS
520b
510b
BWE
514b
TWE
512b
324b
CLUSTER 1
322a
RDP
540a
530a
XBC
HST
550a
LUE
OCM
320a
OBS
520a
510a
BWE
514a
TWE
512a
324a
CROSSBAR (XBR)
312

FIG. 6B

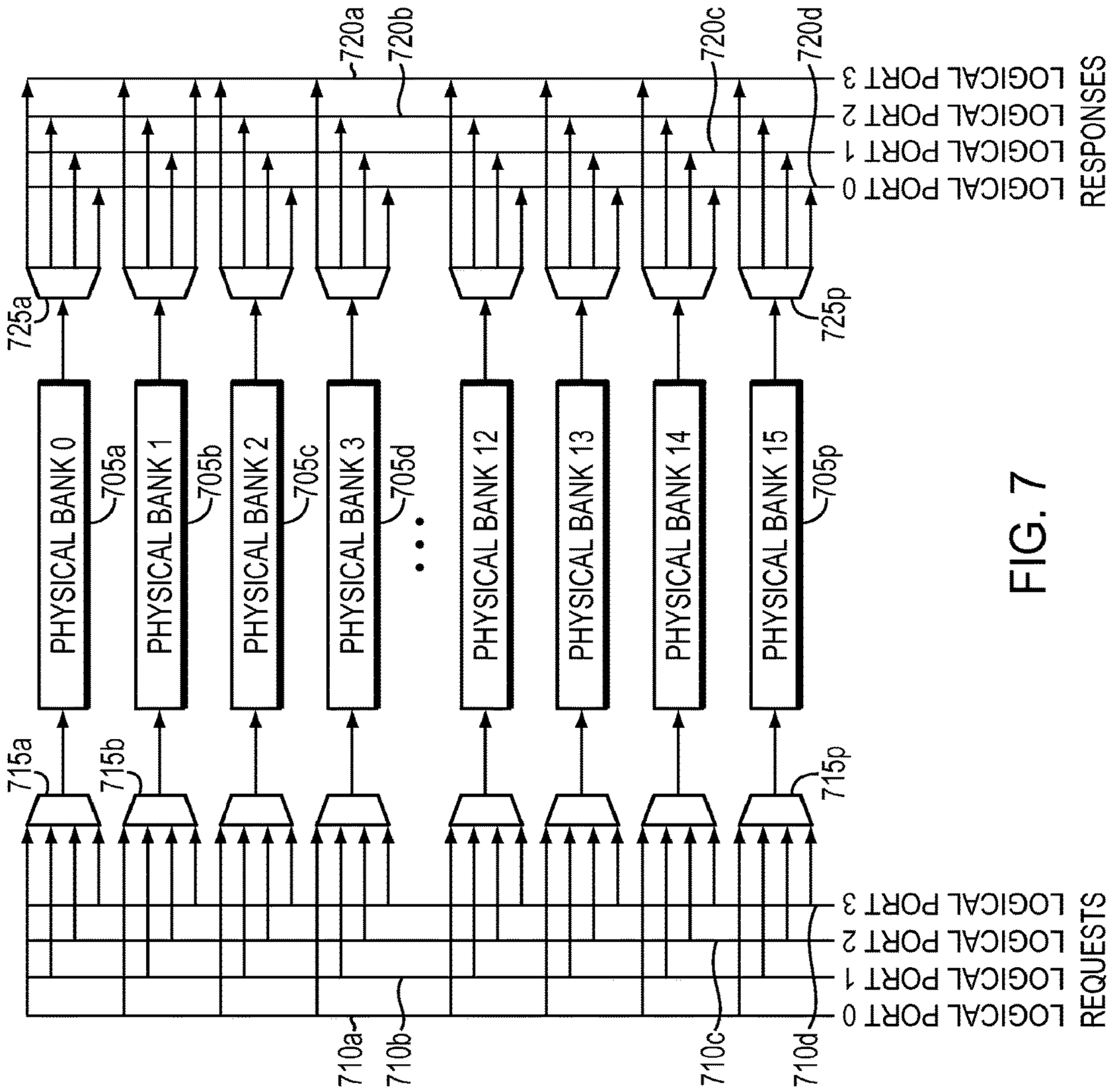

720a
720b
720c
720d
LOGICAL PORT 3
LOGICAL PORT 2
LOGICAL PORT 1
LOGICAL PORT 0
RESPONSES
725a
725p
PHYSICAL BANK 0
PHYSICAL BANK 1
PHYSICAL BANK 2
PHYSICAL BANK 3
PHYSICAL BANK 12
PHYSICAL BANK 13
PHYSICAL BANK 14
PHYSICAL BANK 15
705a
705b
705c
705d
705p
715a
715b
715p
LOGICAL PORT 3
LOGICAL PORT 2
LOGICAL PORT 1
LOGICAL PORT 0
REQUESTS
710a
710b
710c
710d

FIG. 7

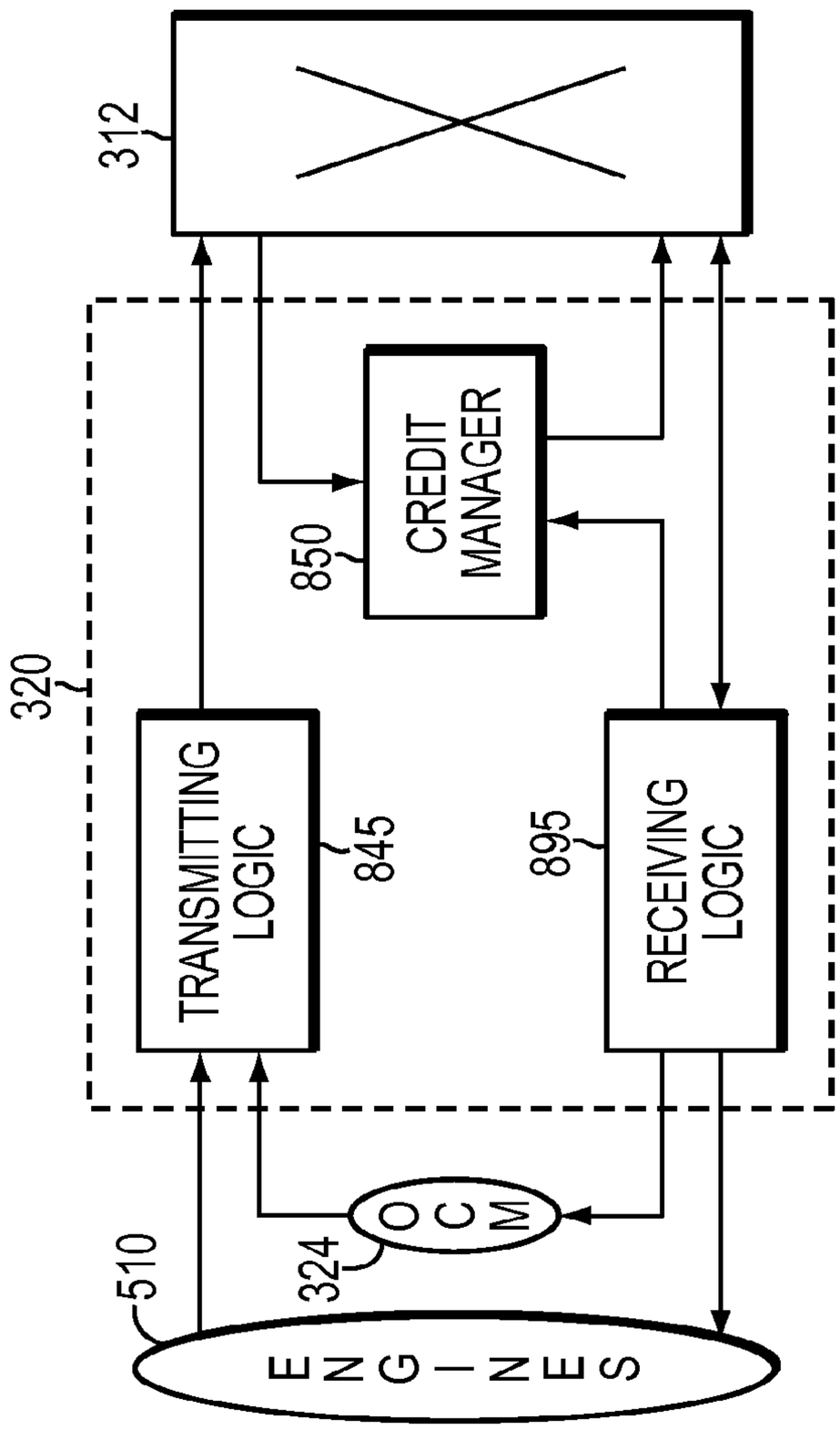

312
320
CREDIT MANAGER
850
TRANSMITTING LOGIC
845
RECEIVING LOGIC
895
OCM
324
510
ENG—NES

FIG. 8A

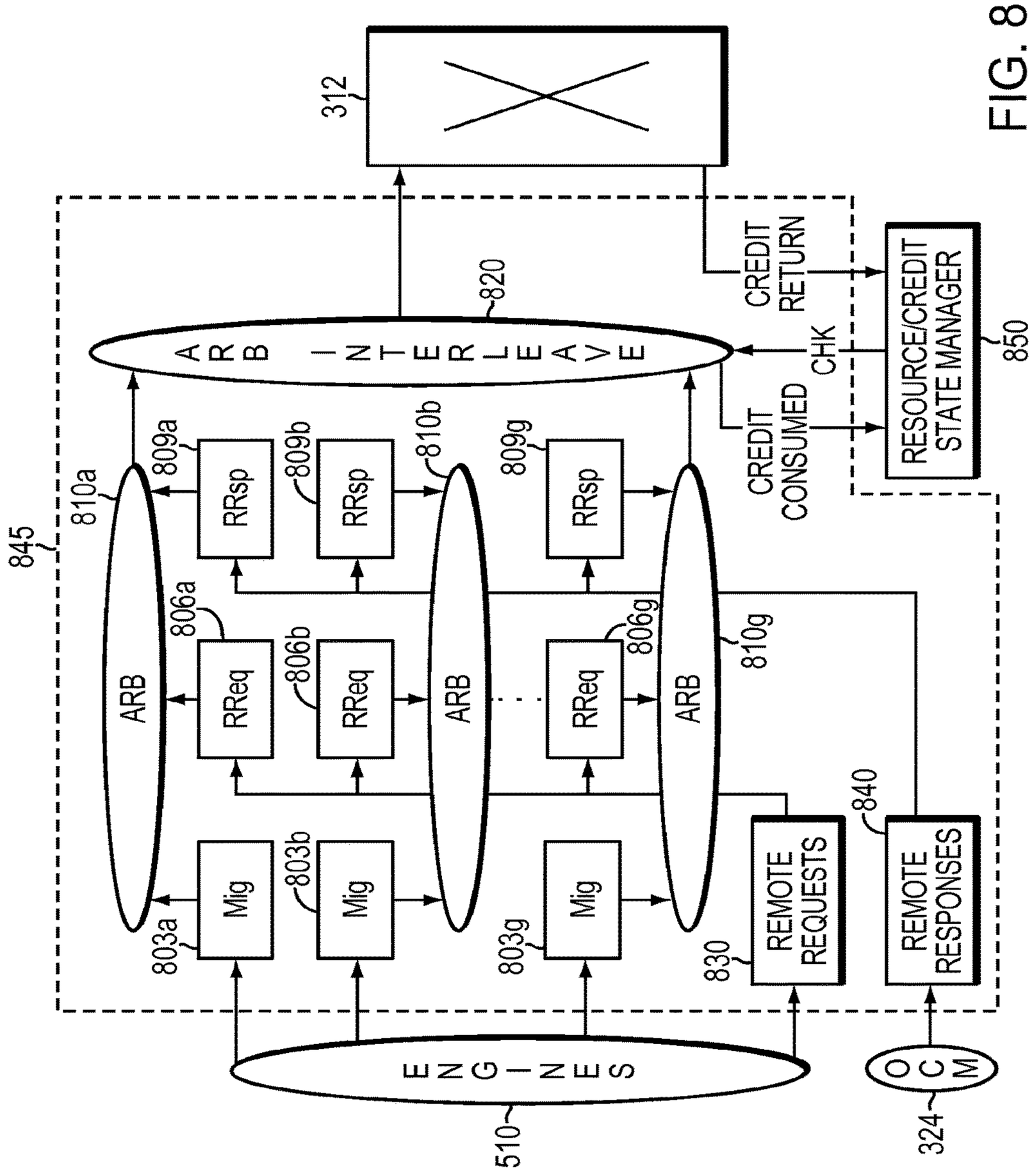

312
820
ARB INTERLEAVE
CREDIT RETURN
CHK
CREDIT CONSUMED
RESOURCE/CREDIT STATE MANAGER
850
845
810a
809a
RRsp
809b
810b
809g
806a
RReq
806b
806g
810g
ARB
Mig
803a
803b
803g
840
REMOTE REQUESTS
830
REMOTE RESPONSES
ENG-INES
510
OCM
324

FIG. 8B

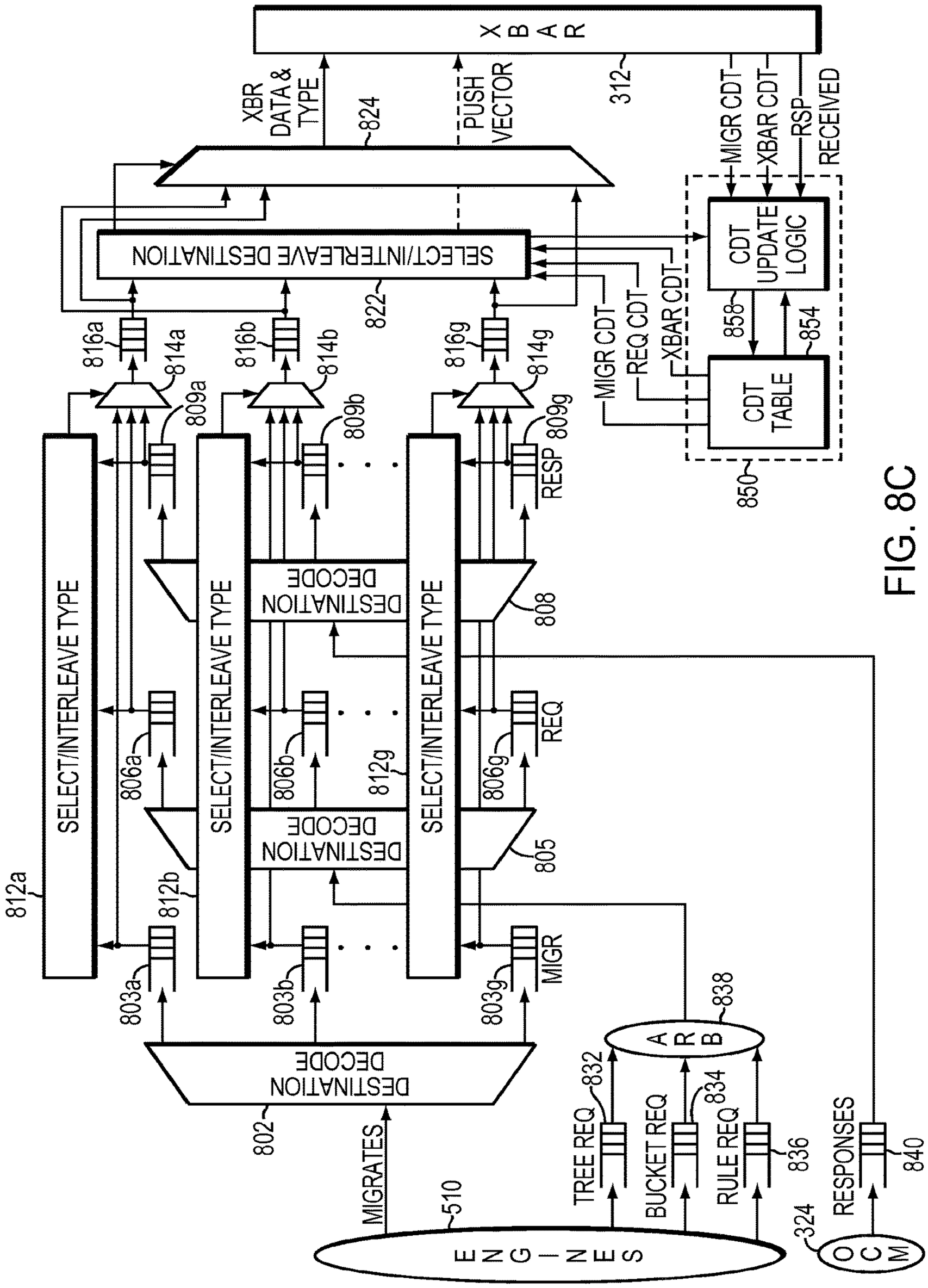

XBAR
XBR DATA & TYPE
824
PUSH VECTOR
312
MIGR CDT
XBAR CDT
RSP RECEIVED
SELECT/INTERLEAVE DESTINATION
CDT UPDATE LOGIC
816a
814a
816b
814b
816g
814g
809a
809b
809g
RESP
858
854
CDT TABLE
850
MIGR CDT
REQ CDT
XBAR CDT
SELECT/INTERLEAVE TYPE
DESTINATION DECODE
808
806a
806b
806g
REQ
812a
812b
812g
805
803a
803b
803g
MIGR
838
ARB
832
834
836
840
TREE REQ
BUCKET REQ
RULE REQ
RESPONSES
802
MIGRATES
510
324
ENGINES
COM

FIG. 8C

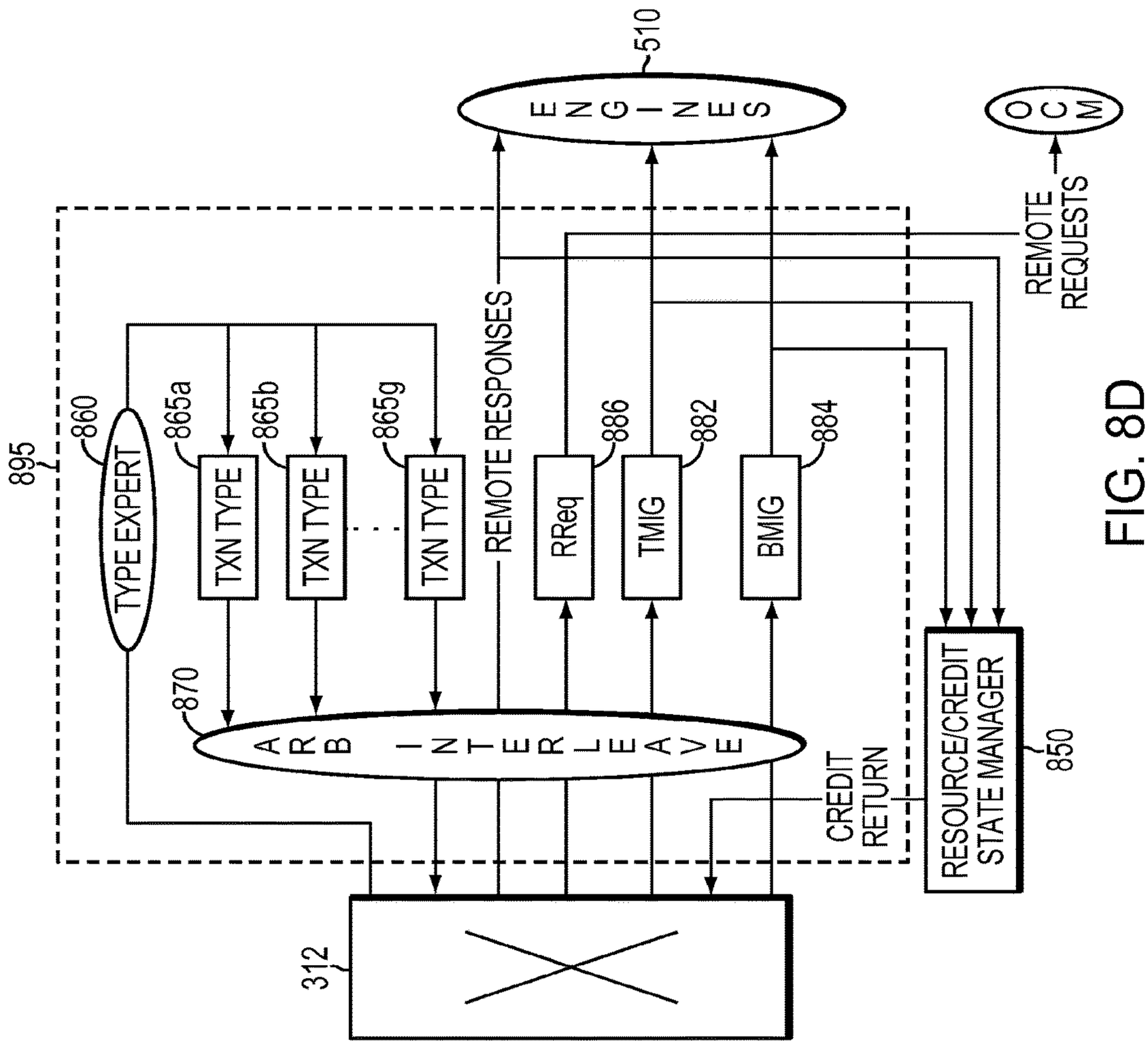
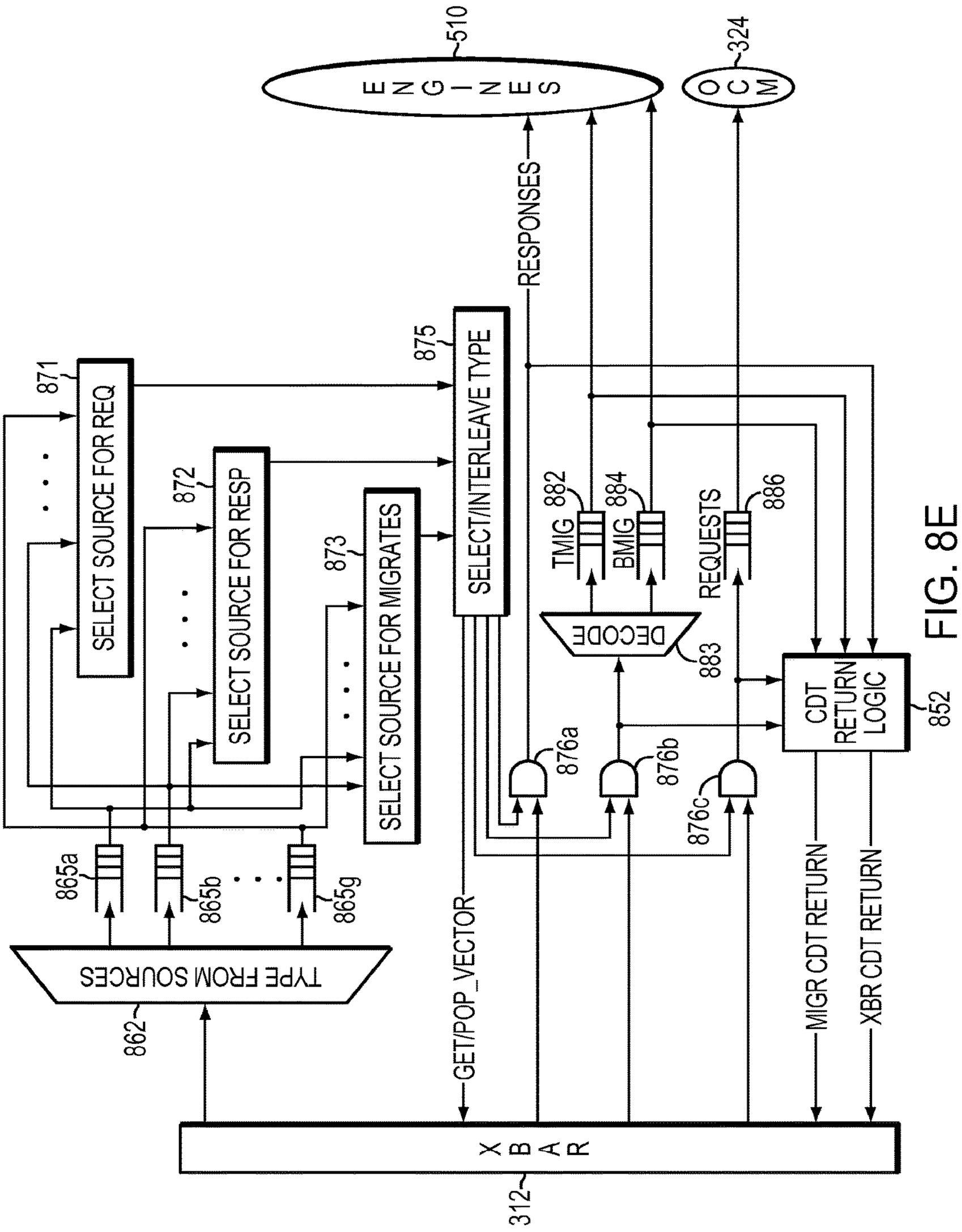

510
ENG—NES
OCM
REMOTE REQUESTS
895
860
TYPE EXPERT
865a
TXN TYPE
865b
TXN TYPE
865g
TXN TYPE
REMOTE RESPONSES
886
RReq
882
TMIG
884
BMIG
870
ARB—NTERLEAVE
CREDIT RETURN
RESOURCE/CREDIT STATE MANAGER
850
312

FIG. 8D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ENG-NES
324
OCM
RESPONSES
871
SELECT SOURCE FOR REQ
872
SELECT SOURCE FOR RESP
873
SELECT SOURCE FOR MIGRATES
875
SELECT/INTERLEAVE TYPE
TMIG
882
BMIG
884
REQUESTS
886
DECODE
883
CDT RETURN LOGIC
852
876a
876b
876c
865a
865b
865g
TYPE FROM SOURCES
862
GET/POP_VECTOR
MIGR CDT RETURN
XBR CDT RETURN
XBAR
312

FIG. 8E

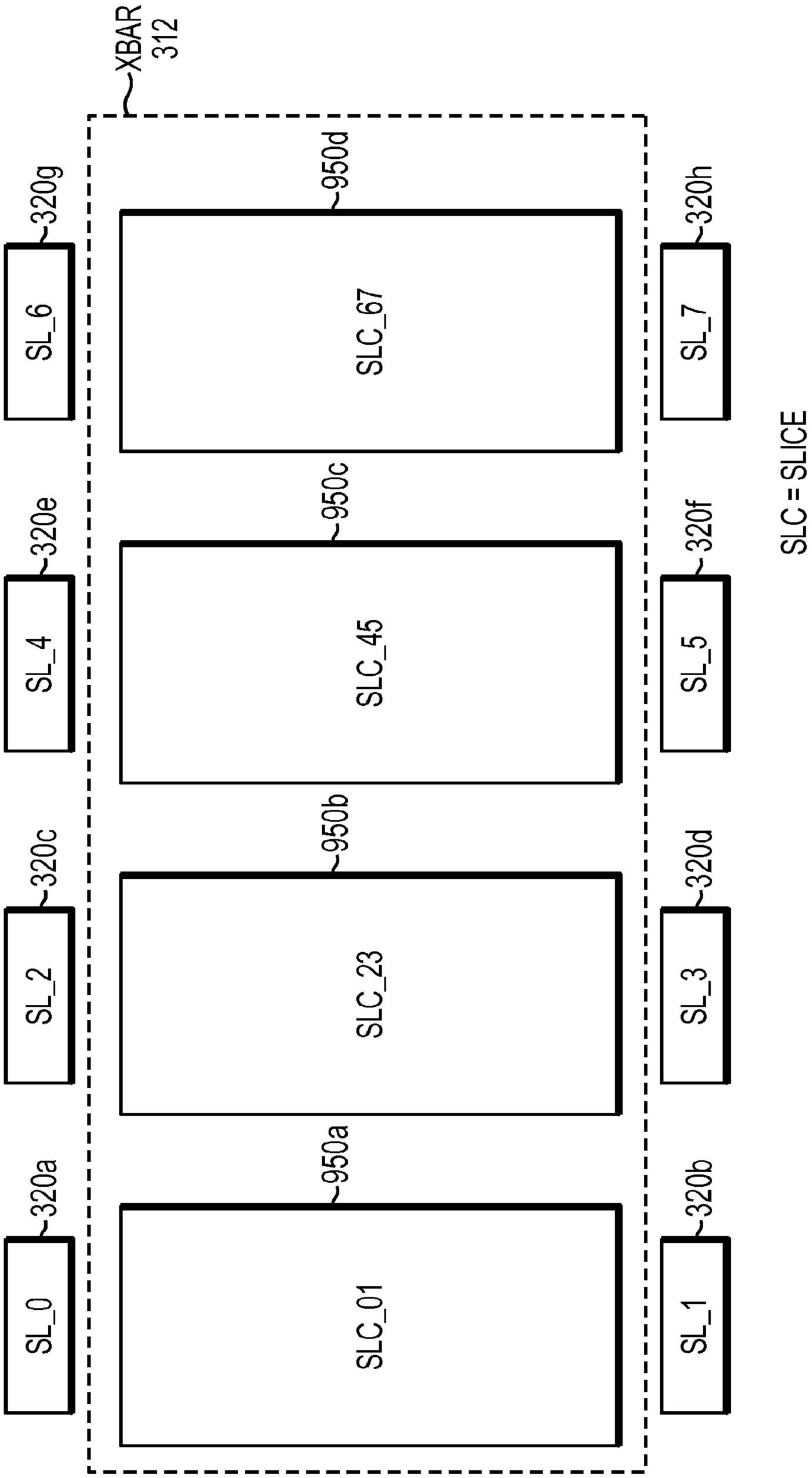
XBAR
312
950d
SLC_67
320g
SL_6
320h
SL_7
950c
SLC_45
320e
SL_4
320f
SL_5
950b
SLC_23
320c
SL_2
320d
SL_3
950a
SLC_01
320a
SL_0
320b
SL_1
SLC = SLICE

FIG. 9A

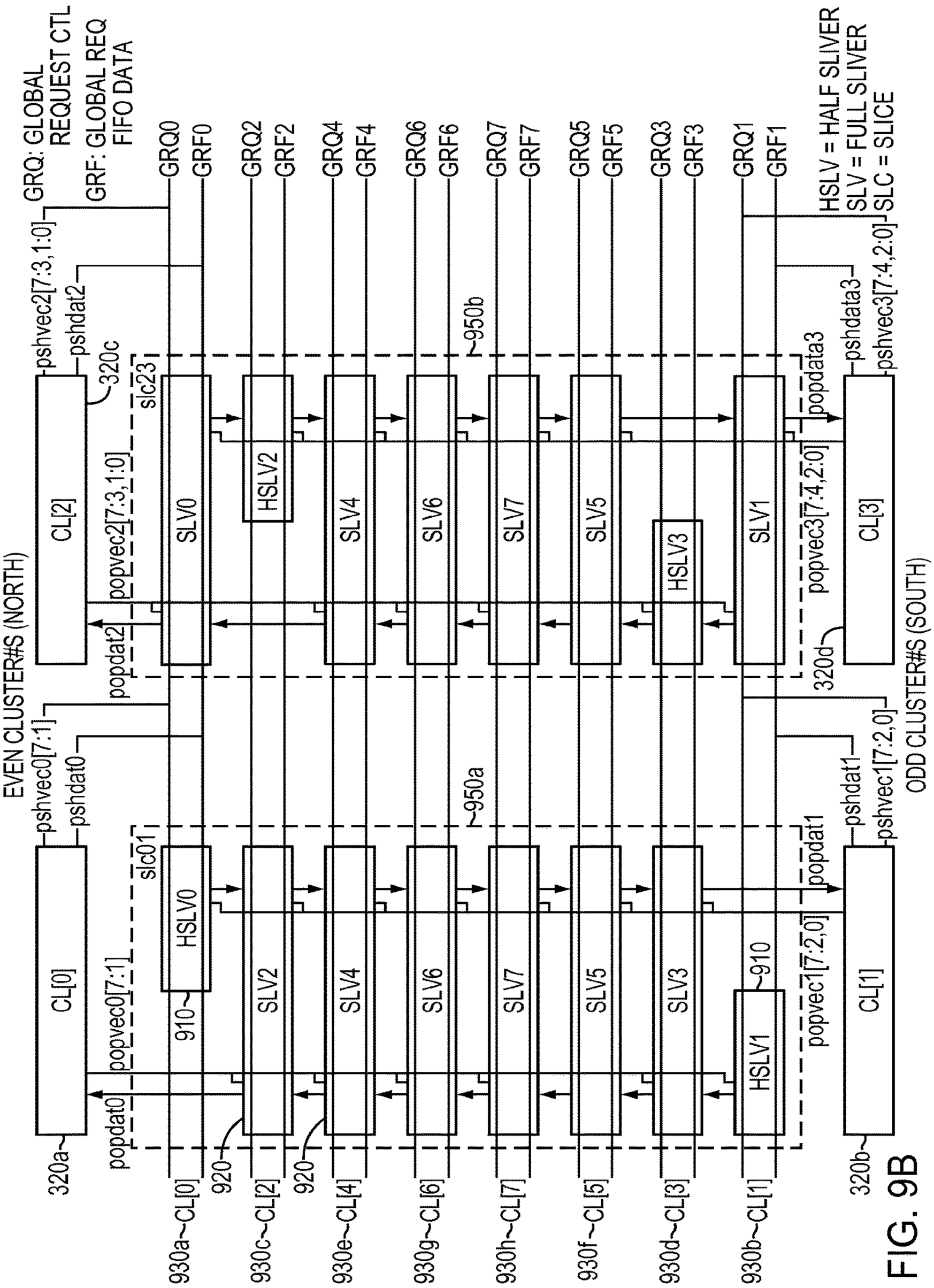
GRQ: GLOBAL REQUEST CTL
GRF: GLOBAL REQ FIFO DATA
HSLV = HALF SLIVER
SLV = FULL SLIVER
SLC = SLICE
EVEN CLUSTER#S (NORTH)
ODD CLUSTER#S (SOUTH)
pshvec2[7:3,1:0]
pshdat2
320c
slc23
popvec2[7:3,1:0]
popdat2
CL[2]
SLV0
HSLV2
SLV4
SLV6
SLV7
SLV5
HSLV3
SLV1
950b
950a
pshvec0[7:1]
pshdat0
slc01
popvec0[7:1]
popdat0
CL[0]
HSLV0
SLV2
SLV4
SLV6
SLV7
SLV5
SLV3
HSLV1
910
920
320a
930a~CL[0]
930c~CL[2]
930e~CL[4]
930g~CL[6]
930h~CL[7]
930f~CL[5]
930d~CL[3]
930b~CL[1]
popvec3[7:4,2:0]
popdata3
pshdata3
pshvec3[7:4,2:0]
CL[3]
320d
pshvec1[7:2,0]
pshdat1
popvec1[7:2,0]
popdat1
CL[1]
320b

FIG. 9B

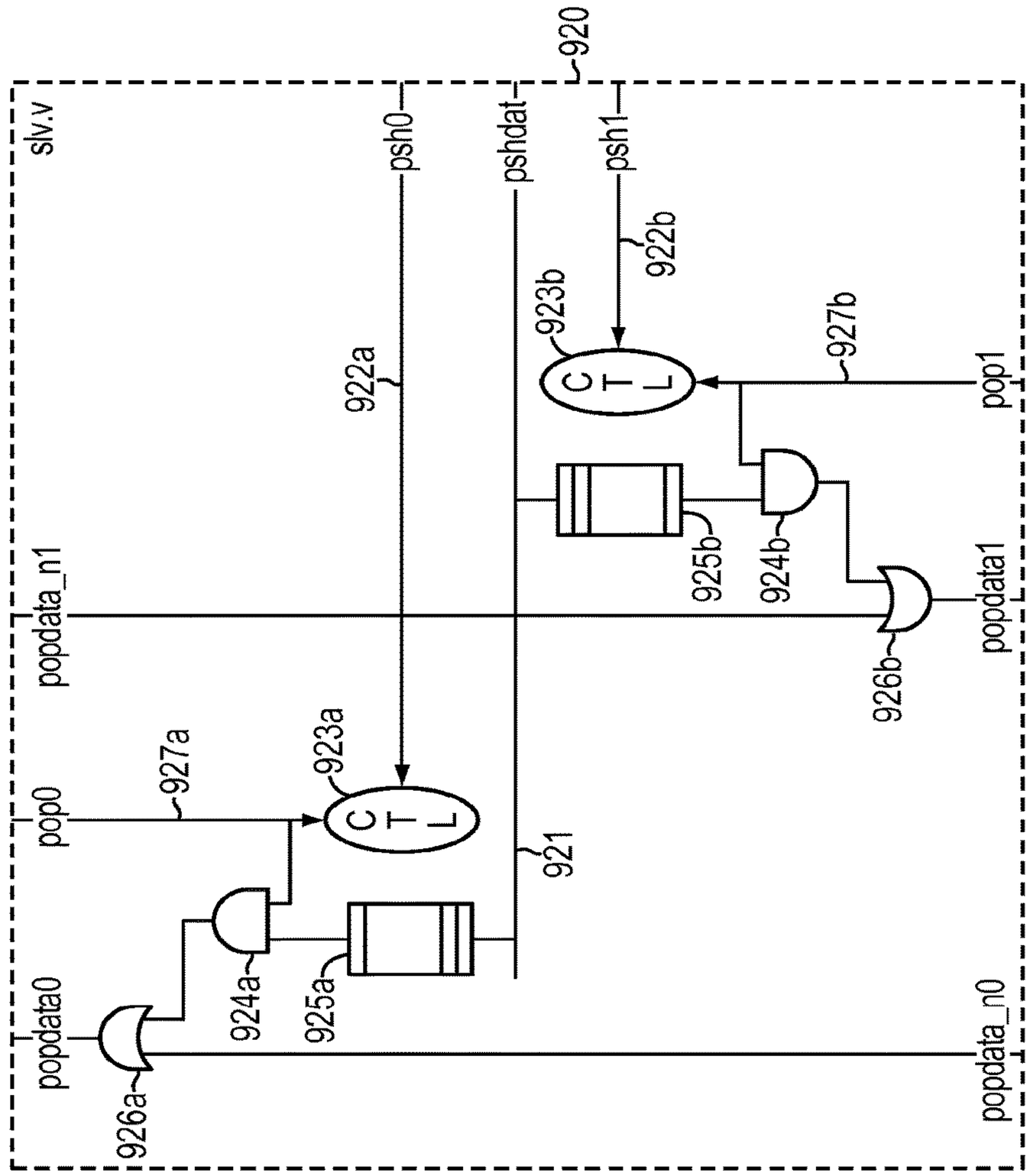
920
slv.v
psh0
pshdat
psh1
922b
922a
923b
927b
pop1
popdata_n1
925b
924b
popdata1
926b
pop0
927a
923a
921
popdata0
924a
925a
926a
popdata_n0

FIG. 9C

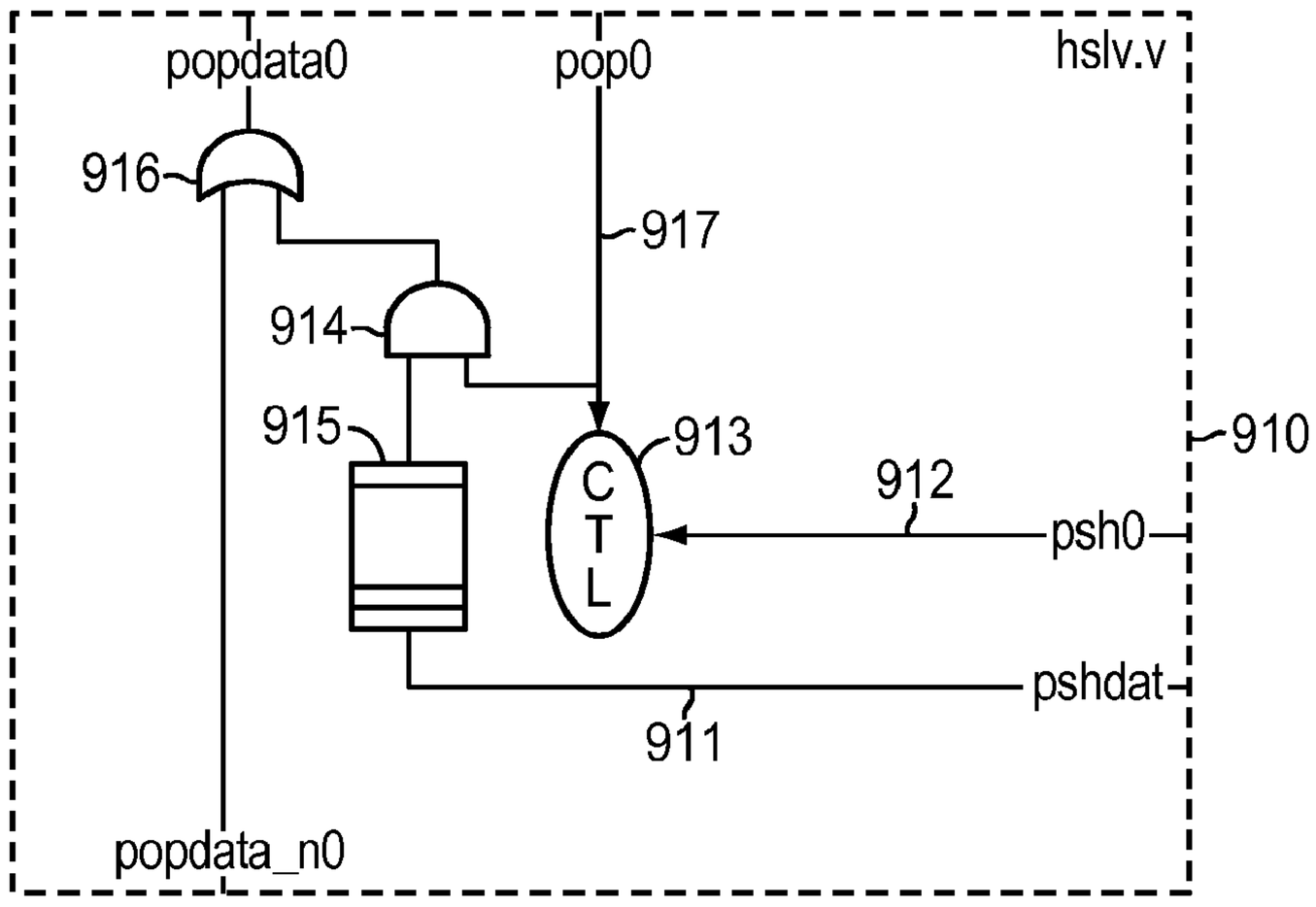
popdata0
pop0
hslv.v
916
917
914
915
913
910
912
C
T
L
psh0
pshdat
911
popdata_n0

FIG. 9D

CLUSTER INDEX	UNIFIED MIGRATION CREDITS	REMOTE REQUEST CREDIT	XBR CREDIT
0	–	–	–
1	5	8	16
2	5	8	16
3	5	8	16
4	5	8	16
5	5	8	16
6	5	8	16
7	5	8	16

FIG. 10A

CLUSTER INDEX	TREE MIGRATION CREDITS	BUCKET MIGRATION CREDITS	REMOTE REQUEST CREDIT	XBR CREDIT
0	–	–	–	–
1	5	5	8	16
2	5	5	8	16
3	5	5	8	16
4	5	5	8	16
5	5	5	8	16
6	5	5	8	16
7	5	5	8	16

FIG. 10B

FIG. 11A

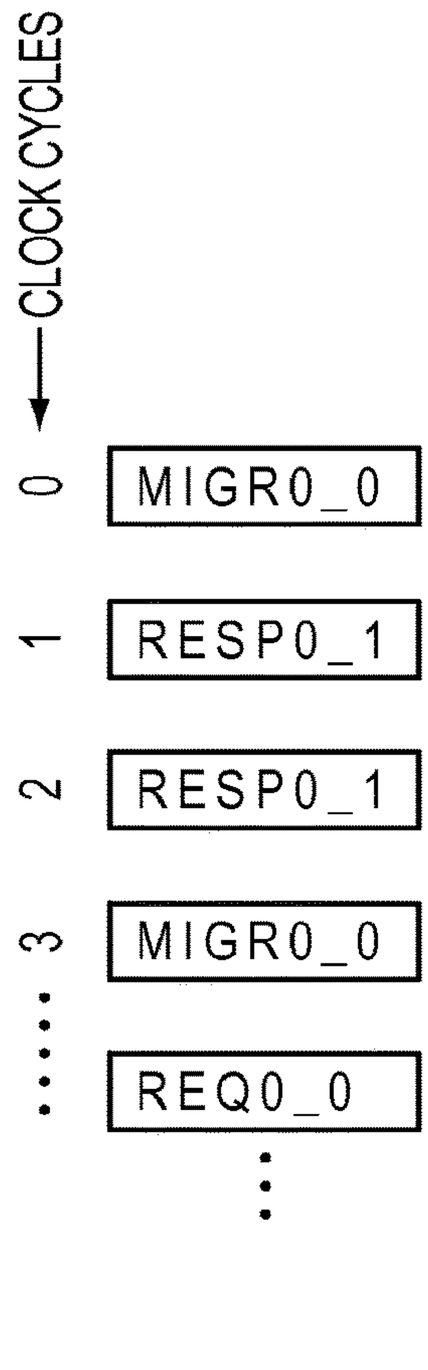

FIG. 11B

EXAMPLE_1: INTERLEAVING OF TYPES

MIGR0_0: MIGRATION DATA BELONGING TO MIGRATE 0, DESTINATION CLUSTER = 0

REQ0_0: REQUEST 0, DESTINATION CL = 0

REQ1_0: REQUEST 1, DESTINATION CL = 0

EXAMPLE_2: INTERLEAVING OF TYPES AND DESTINATION

MIGR0_0: MIGRATE 0, DEST CL = 0

RESP0_1: RESPONSE 0, DEST CL = 1

REQ0_0: REQUEST 0, DEST CL = 0

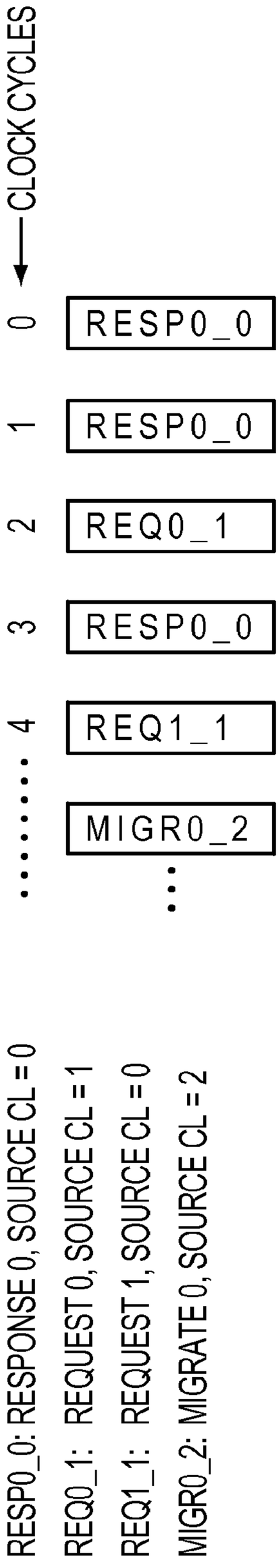
CLOCK CYCLES
0
RESP0_0
1
RESP0_0
2
REQ0_1
3
RESP0_0
4
REQ1_1
MIGR0_2
AND SOURCES
EXAMPLE_1: INTERLEAVING TYPES
RESP0_0: RESPONSE 0, SOURCE CL = 0
REQ0_1: REQUEST 0, SOURCE CL = 1
REQ1_1: REQUEST 1, SOURCE CL = 0
MIGR0_2: MIGRATE 0, SOURCE CL = 2

FIG. 11C

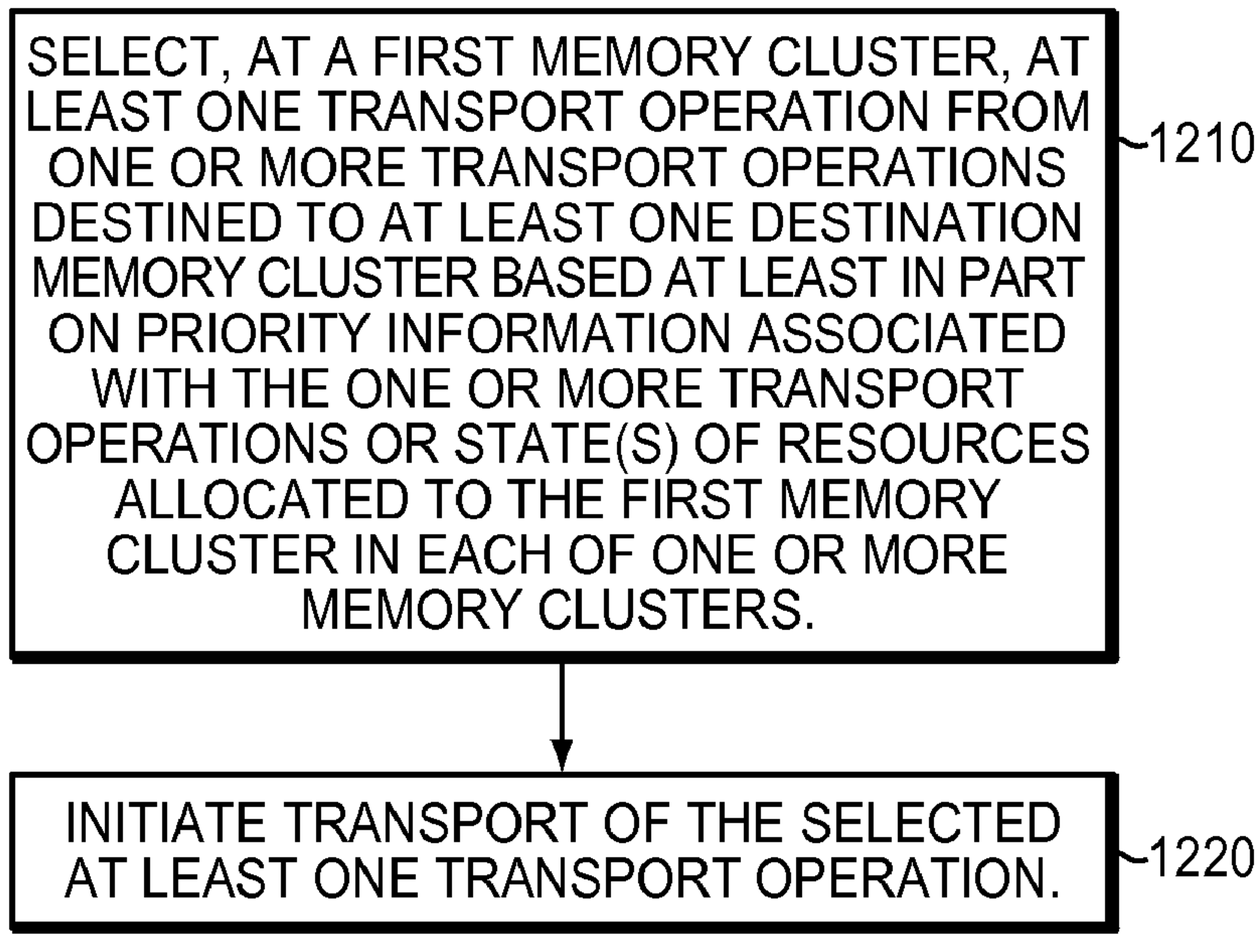
SELECT, AT A FIRST MEMORY CLUSTER, AT LEAST ONE TRANSPORT OPERATION FROM ONE OR MORE TRANSPORT OPERATIONS DESTINED TO AT LEAST ONE DESTINATION MEMORY CLUSTER BASED AT LEAST IN PART ON PRIORITY INFORMATION ASSOCIATED WITH THE ONE OR MORE TRANSPORT OPERATIONS OR STATE(S) OF RESOURCES ALLOCATED TO THE FIRST MEMORY CLUSTER IN EACH OF ONE OR MORE MEMORY CLUSTERS.
1210
INITIATE TRANSPORT OF THE SELECTED AT LEAST ONE TRANSPORT OPERATION.
1220

FIG. 12A

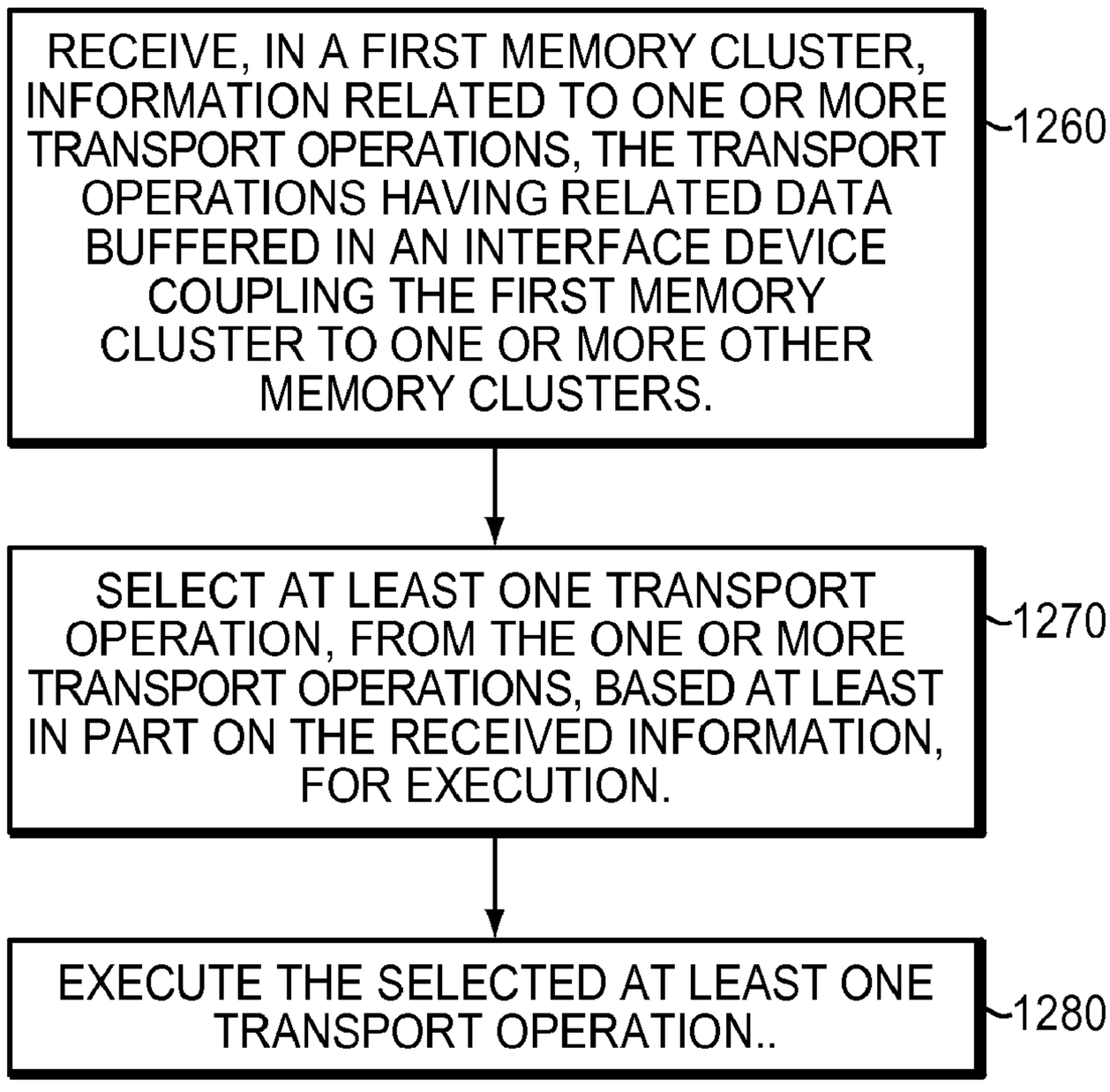
RECEIVE, IN A FIRST MEMORY CLUSTER, INFORMATION RELATED TO ONE OR MORE TRANSPORT OPERATIONS, THE TRANSPORT OPERATIONS HAVING RELATED DATA BUFFERED IN AN INTERFACE DEVICE COUPLING THE FIRST MEMORY CLUSTER TO ONE OR MORE OTHER MEMORY CLUSTERS.
1260
SELECT AT LEAST ONE TRANSPORT OPERATION, FROM THE ONE OR MORE TRANSPORT OPERATIONS, BASED AT LEAST IN PART ON THE RECEIVED INFORMATION, FOR EXECUTION.
1270
EXECUTE THE SELECTED AT LEAST ONE TRANSPORT OPERATION..
1280

FIG. 12B

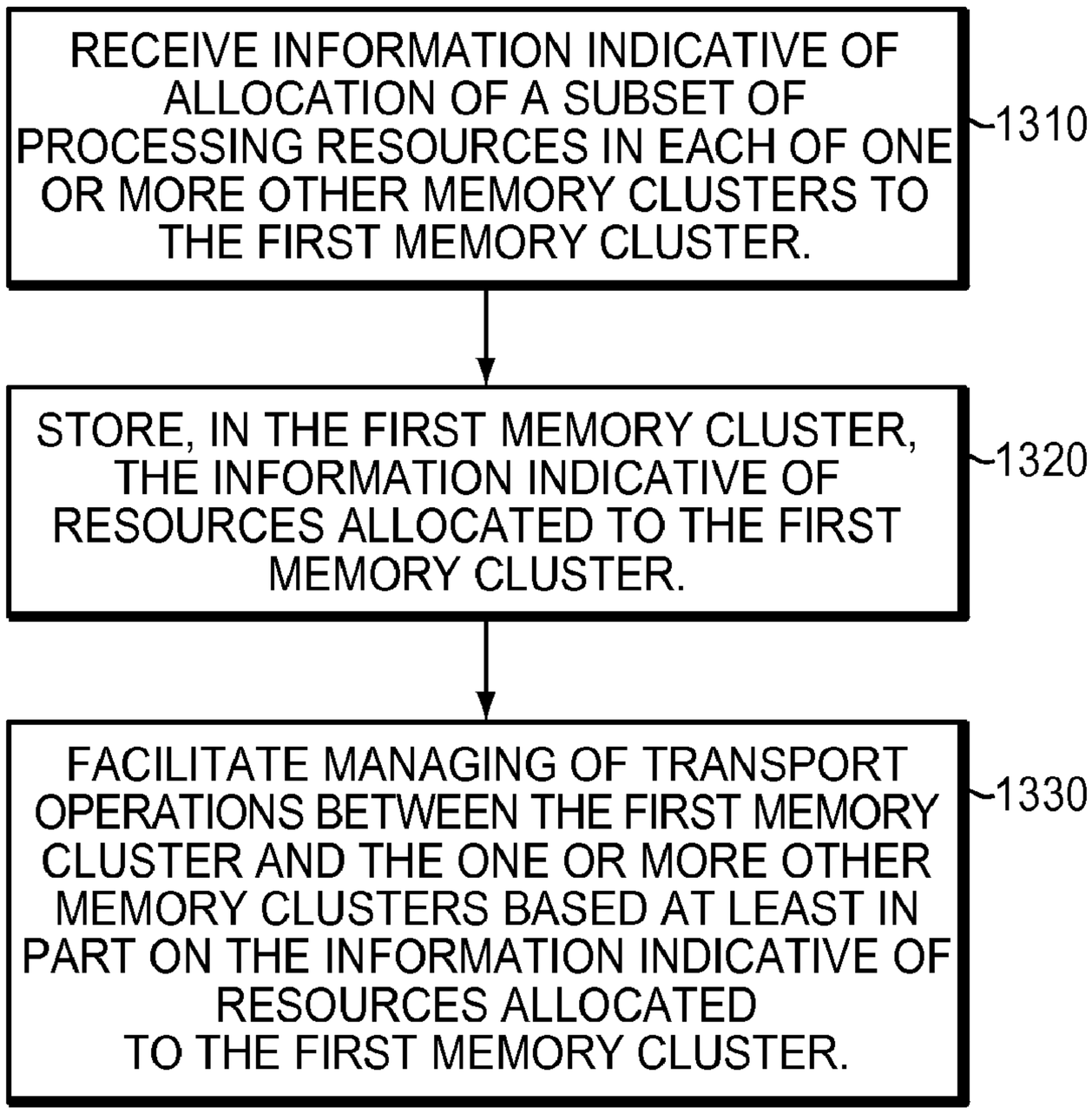
RECEIVE INFORMATION INDICATIVE OF ALLOCATION OF A SUBSET OF PROCESSING RESOURCES IN EACH OF ONE OR MORE OTHER MEMORY CLUSTERS TO THE FIRST MEMORY CLUSTER.
1310
STORE, IN THE FIRST MEMORY CLUSTER, THE INFORMATION INDICATIVE OF RESOURCES ALLOCATED TO THE FIRST MEMORY CLUSTER.
1320
FACILITATE MANAGING OF TRANSPORT OPERATIONS BETWEEN THE FIRST MEMORY CLUSTER AND THE ONE OR MORE OTHER MEMORY CLUSTERS BASED AT LEAST IN PART ON THE INFORMATION INDICATIVE OF RESOURCES ALLOCATED TO THE FIRST MEMORY CLUSTER.
1330

FIG. 13

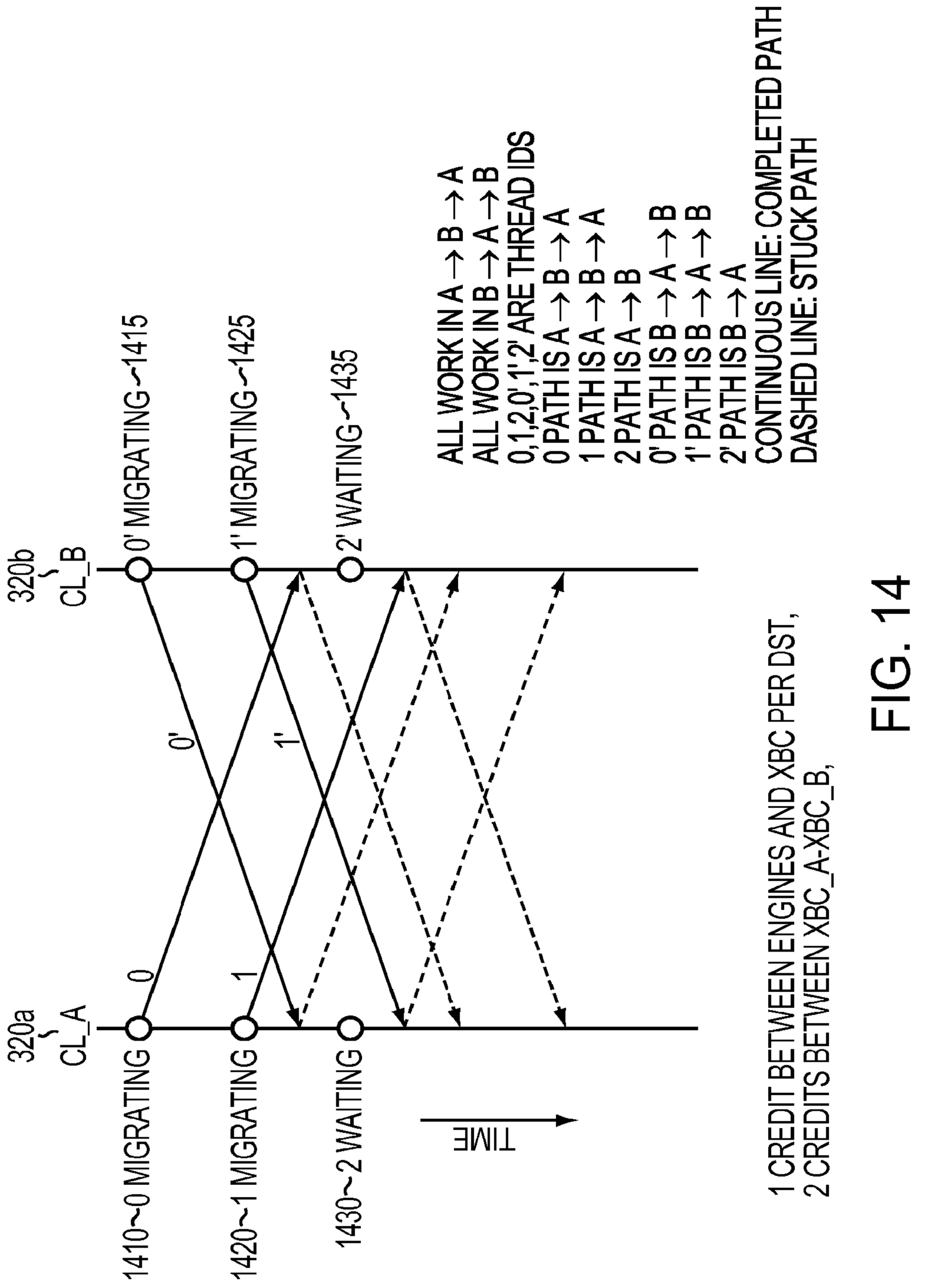
320b
CL_B
0' MIGRATING ~1415
1' MIGRATING ~1425
2' WAITING ~1435
0'
1'
320a
CL_A
0
1
1410~0 MIGRATING
1420~1 MIGRATING
1430~2 WAITING
TIME
ALL WORK IN A → B → A
ALL WORK IN B → A → B
0,1,2,0',1',2' ARE THREAD IDS
0 PATH IS A → B → A
1 PATH IS A → B → A
2 PATH IS A → B
0' PATH IS B → A → B
1' PATH IS B → A → B
2' PATH IS B → A
CONTINUOUS LINE: COMPLETED PATH
DASHED LINE: STUCK PATH
1 CREDIT BETWEEN ENGINES AND XBC PER DST,
2 CREDITS BETWEEN XBC_A-XBC_B,

FIG. 14

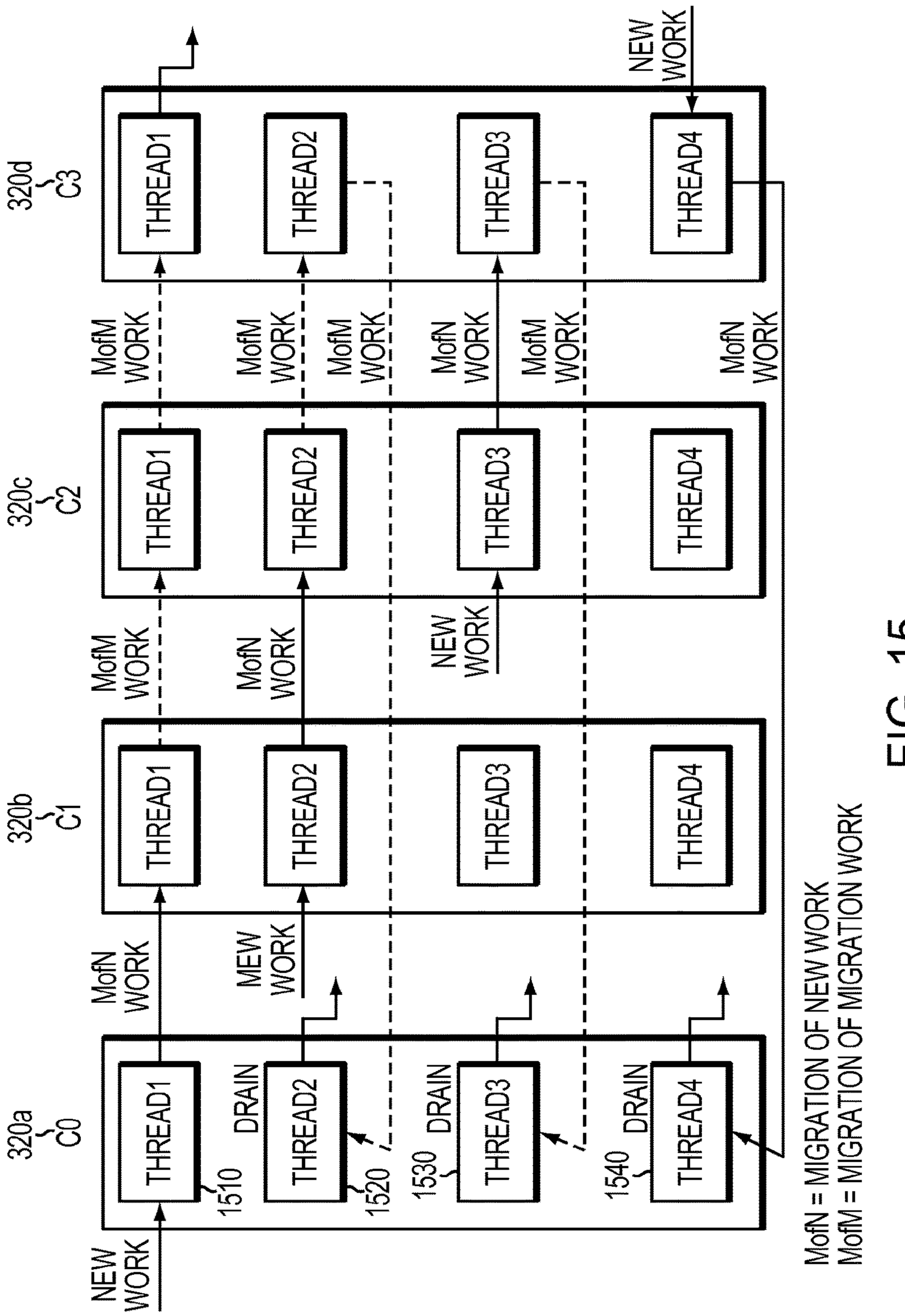
320d
C3
THREAD1
THREAD2
THREAD3
THREAD4
NEW WORK
MofM WORK
MofM WORK
MofM WORK
MofN WORK
MofM WORK
MofN WORK
320c
C2
THREAD1
THREAD2
THREAD3
THREAD4
MofM WORK
MofN WORK
NEW WORK
320b
C1
THREAD1
THREAD2
THREAD3
THREAD4
MofN WORK
MEW WORK
320a
C0
THREAD1
DRAIN
THREAD2
DRAIN
THREAD3
DRAIN
THREAD4
NEW WORK
1510
1520
1530
1540
MofN = MIGRATION OF NEW WORK
MofM = MIGRATION OF MIGRATION WORK

FIG. 15

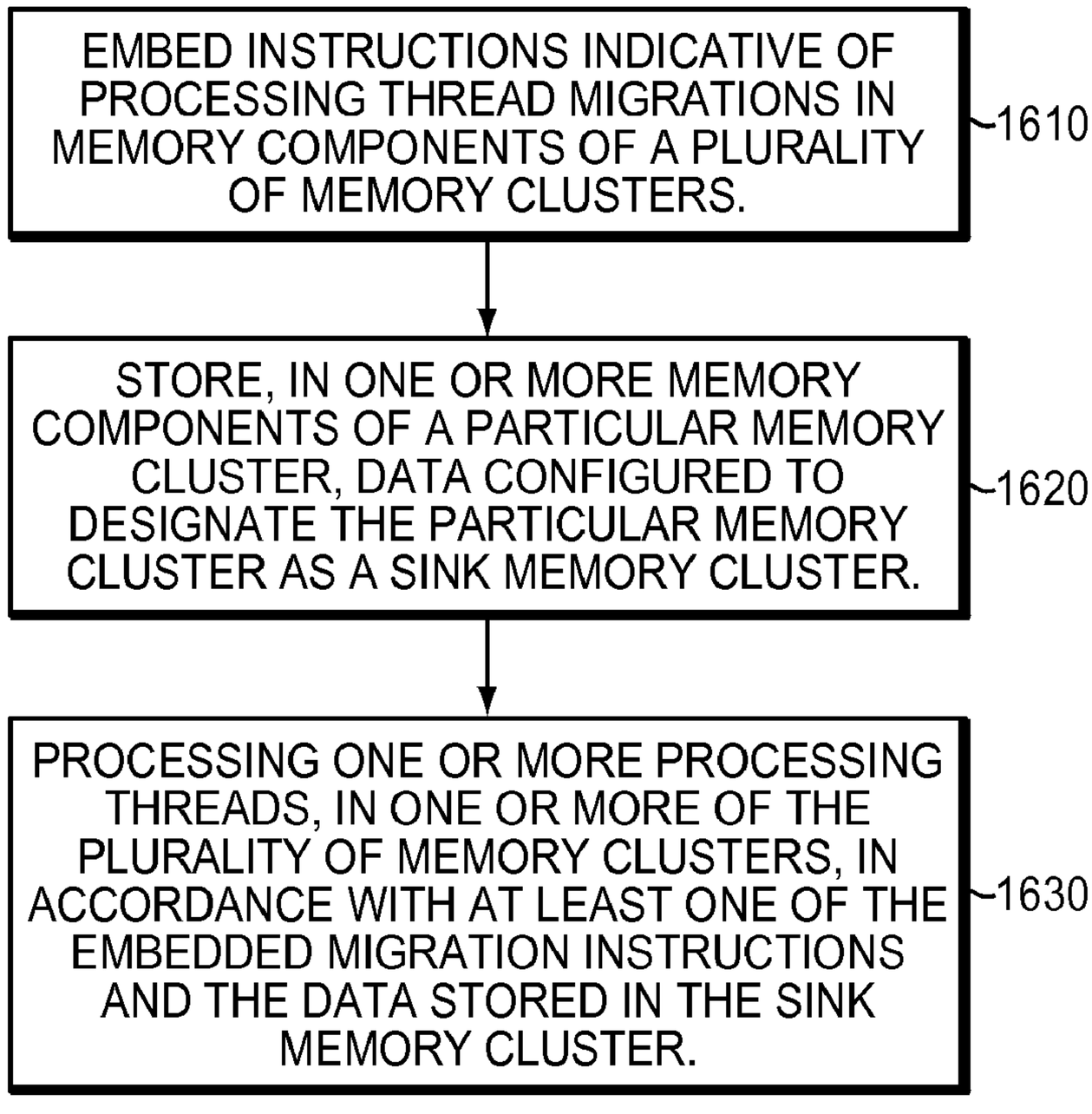
EMBED INSTRUCTIONS INDICATIVE OF PROCESSING THREAD MIGRATIONS IN MEMORY COMPONENTS OF A PLURALITY OF MEMORY CLUSTERS.
1610
STORE, IN ONE OR MORE MEMORY COMPONENTS OF A PARTICULAR MEMORY CLUSTER, DATA CONFIGURED TO DESIGNATE THE PARTICULAR MEMORY CLUSTER AS A SINK MEMORY CLUSTER.
1620
PROCESSING ONE OR MORE PROCESSING THREADS, IN ONE OR MORE OF THE PLURALITY OF MEMORY CLUSTERS, IN ACCORDANCE WITH AT LEAST ONE OF THE EMBEDDED MIGRATION INSTRUCTIONS AND THE DATA STORED IN THE SINK MEMORY CLUSTER.
1630

FIG. 16

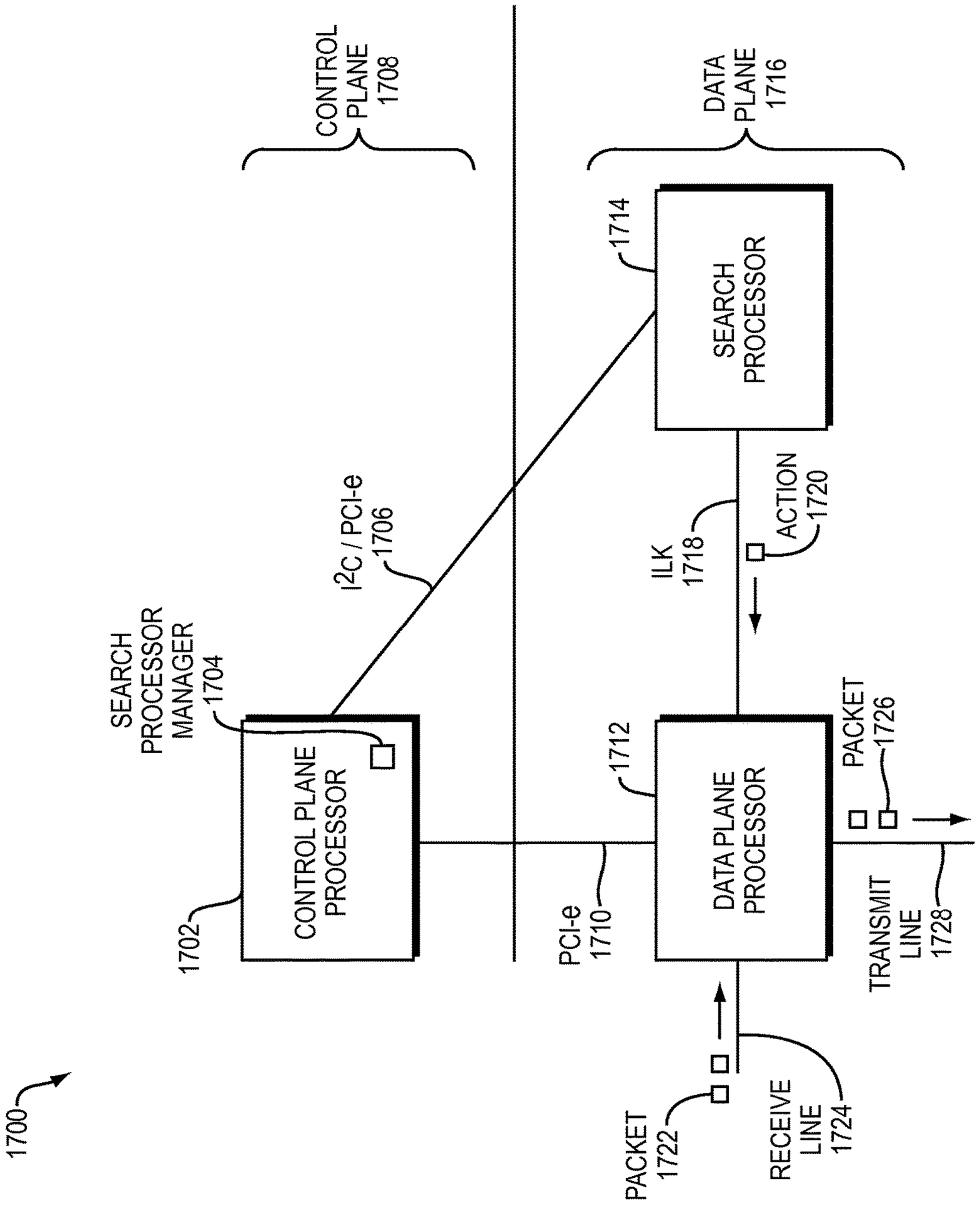
CONTROL PLANE 1708
DATA PLANE 1716
SEARCH PROCESSOR 1714
I2C / PCI-e 1706
SEARCH PROCESSOR MANAGER 1704
ILK 1718
ACTION 1720
CONTROL PLANE PROCESSOR
1702
PCI-e 1710
1712
DATA PLANE PROCESSOR
PACKET 1726
TRANSMIT LINE 1728
PACKET 1722
RECEIVE LINE 1724
1700

FIG. 17

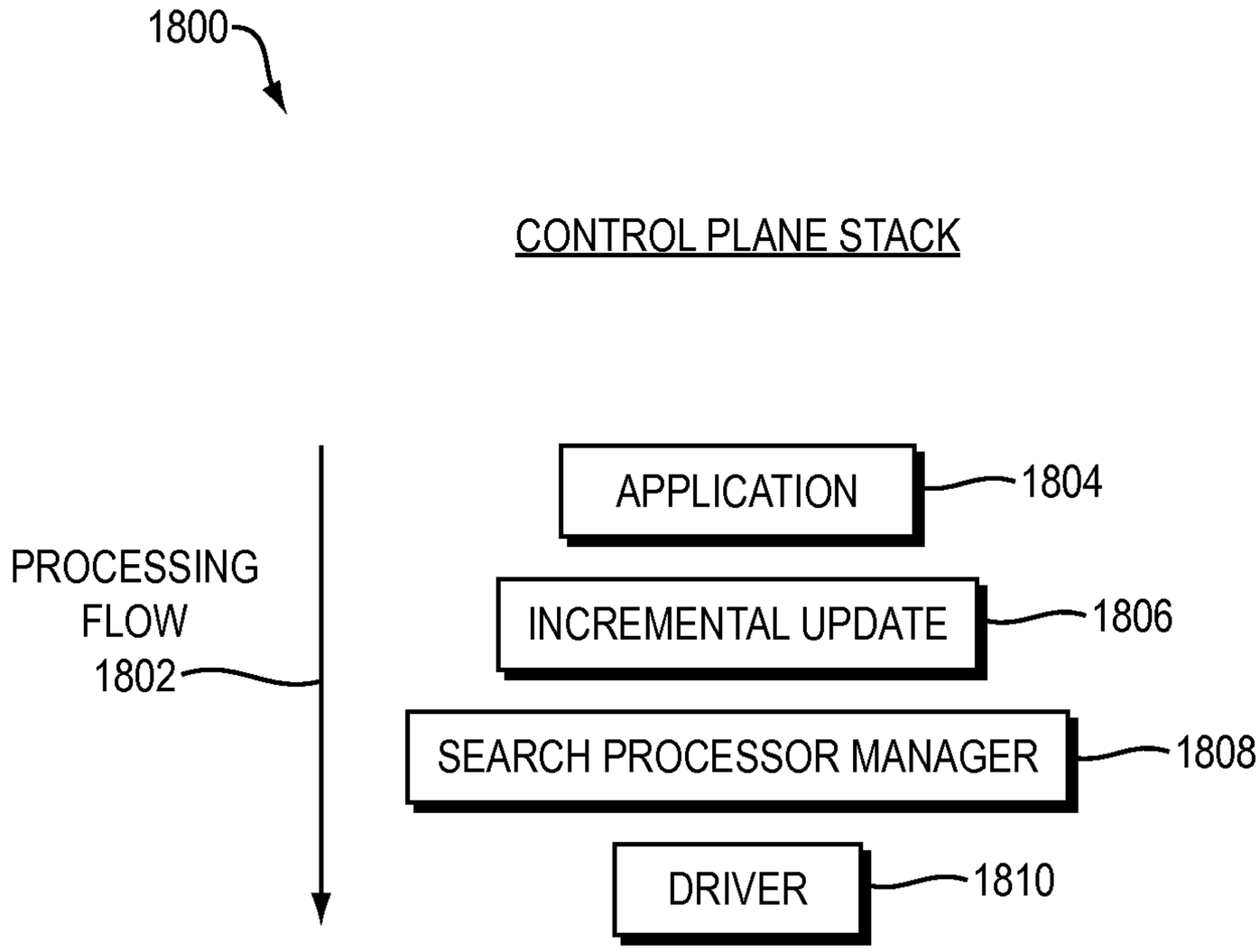
1800
CONTROL PLANE STACK
PROCESSING FLOW
1802
APPLICATION
1804
INCREMENTAL UPDATE
1806
SEARCH PROCESSOR MANAGER
1808
DRIVER
1810

FIG. 18

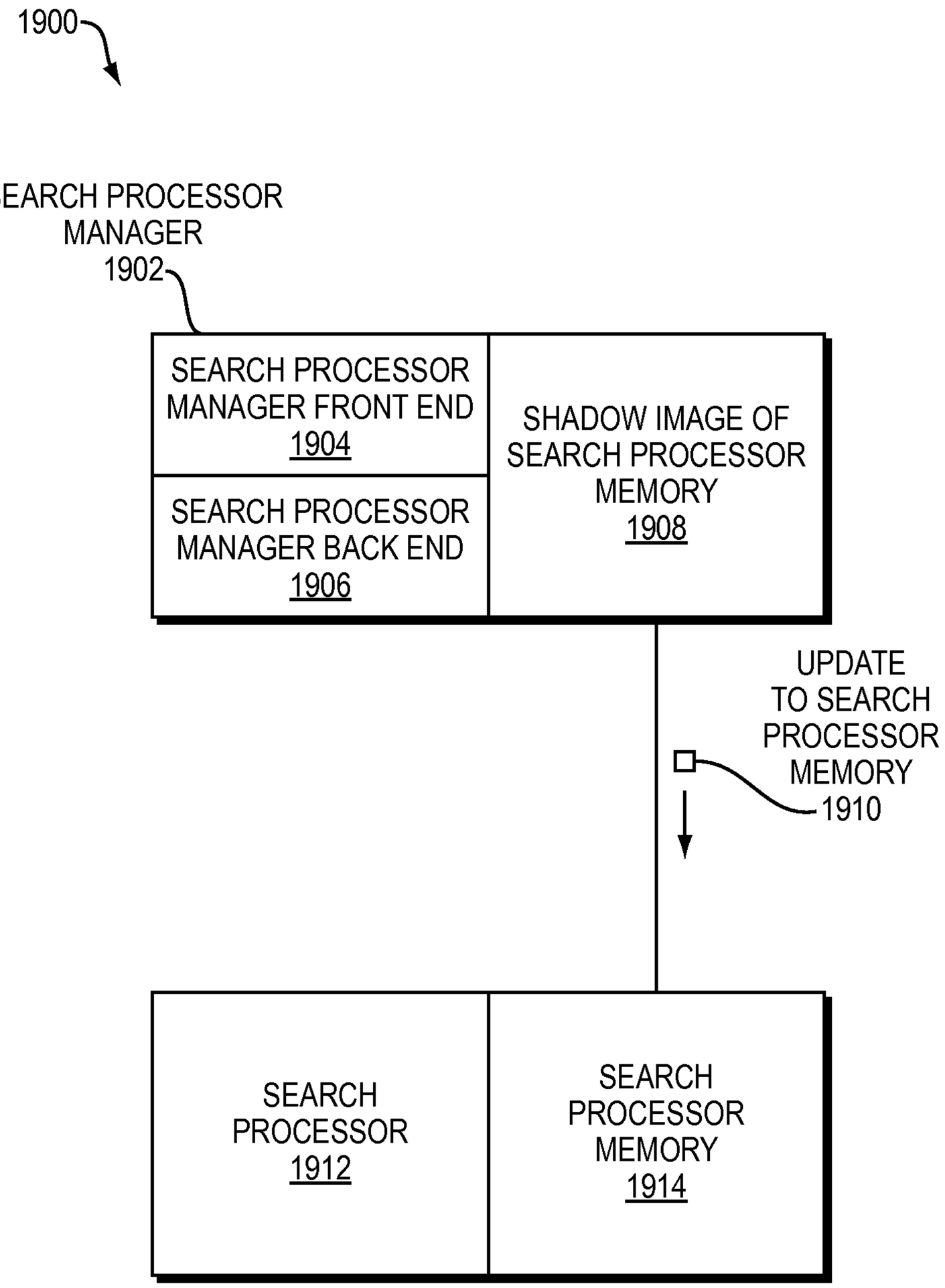
1900
SEARCH PROCESSOR MANAGER
1902
SEARCH PROCESSOR MANAGER FRONT END
1904
SEARCH PROCESSOR MANAGER BACK END
1906
SHADOW IMAGE OF SEARCH PROCESSOR MEMORY
1908
UPDATE TO SEARCH PROCESSOR MEMORY
1910
SEARCH PROCESSOR
1912
SEARCH PROCESSOR MEMORY
1914

FIG. 19

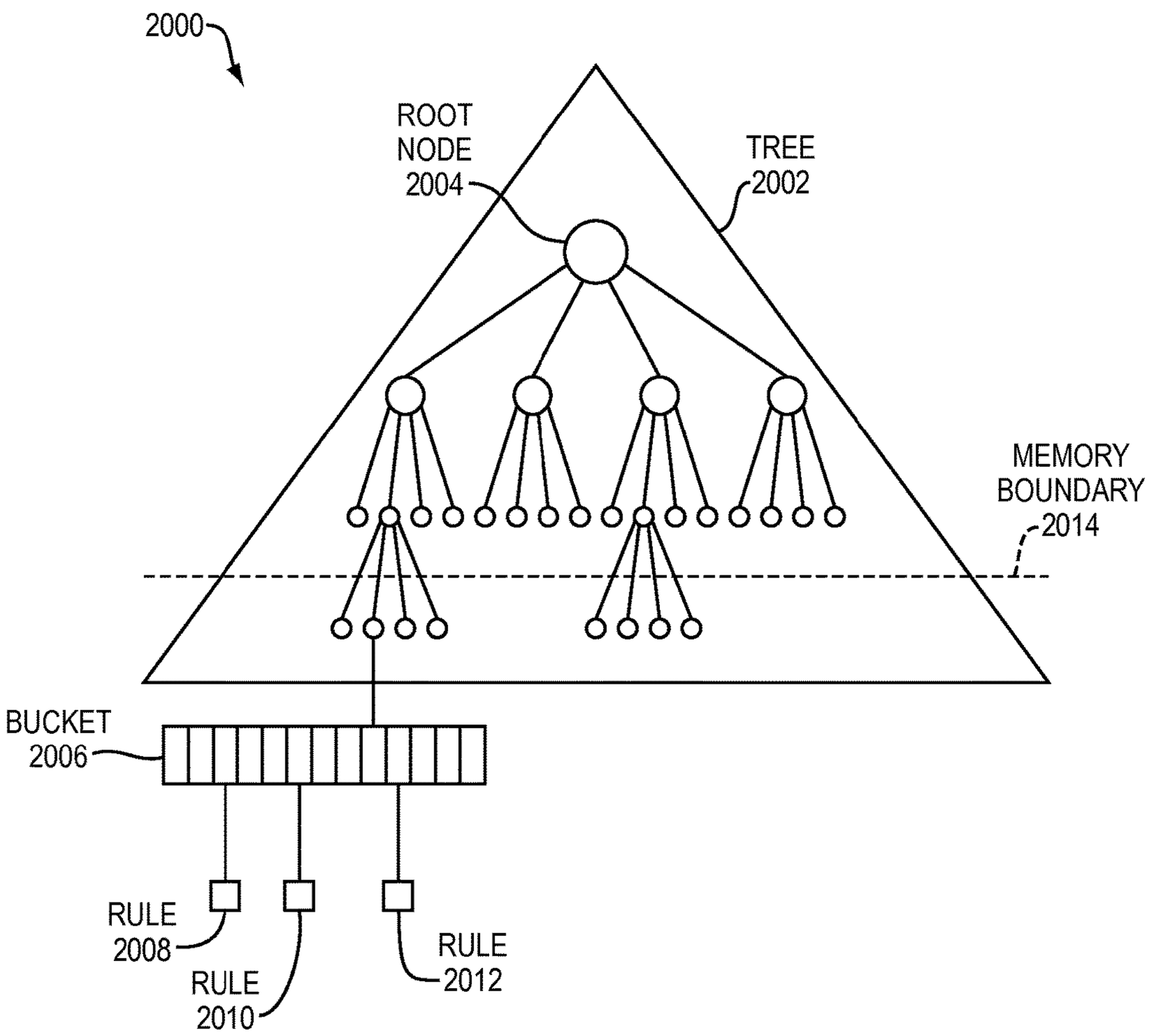
2000
ROOT NODE 2004
TREE 2002
MEMORY BOUNDARY 2014
BUCKET 2006
RULE 2008
RULE 2010
RULE 2012

FIG. 20A

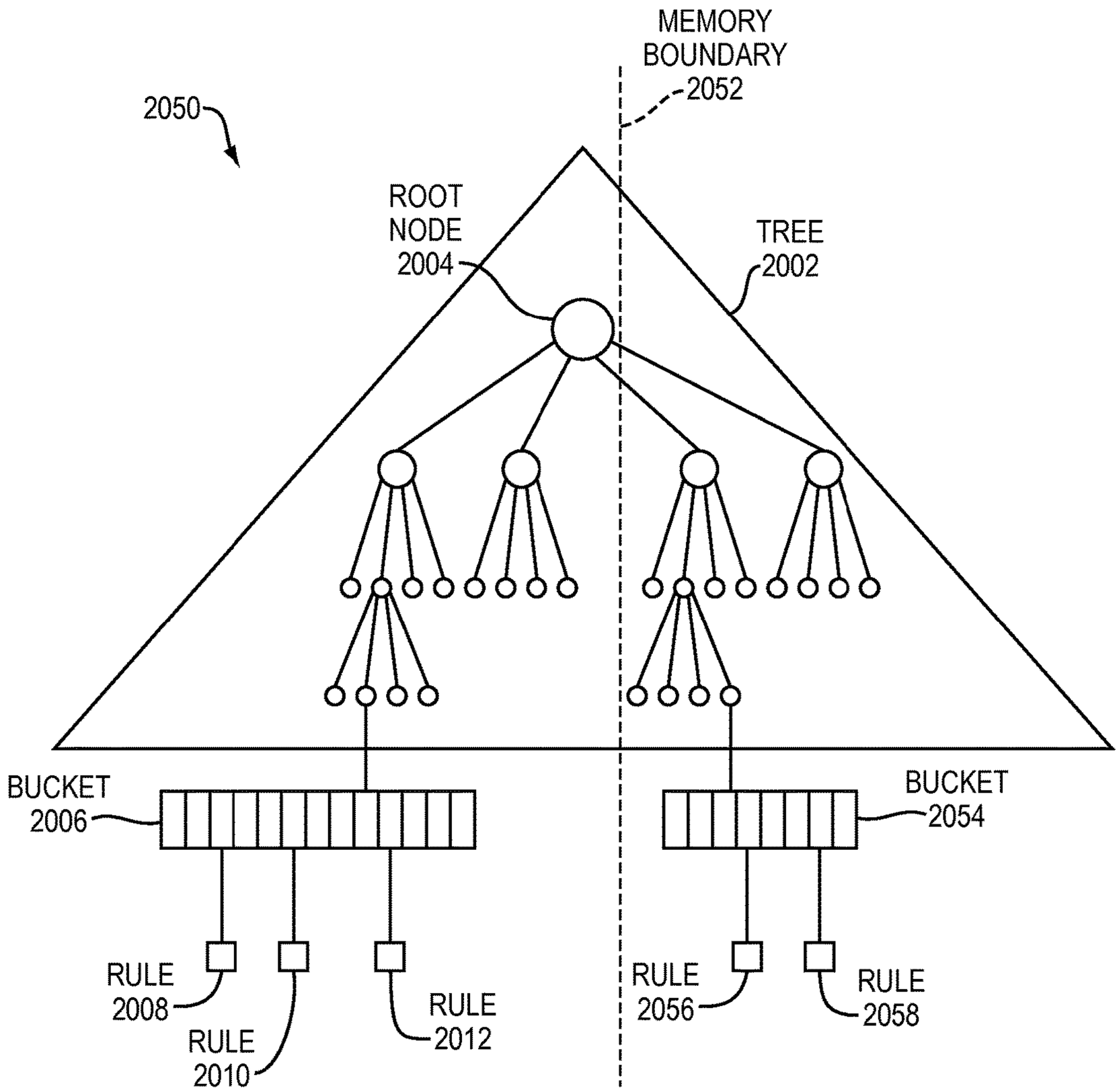
2050
MEMORY
BOUNDARY
2052
ROOT
NODE
2004
TREE
2002
BUCKET
2006
BUCKET
2054
RULE
2008
RULE
2010
RULE
2012
RULE
2056
RULE
2058

FIG. 20B

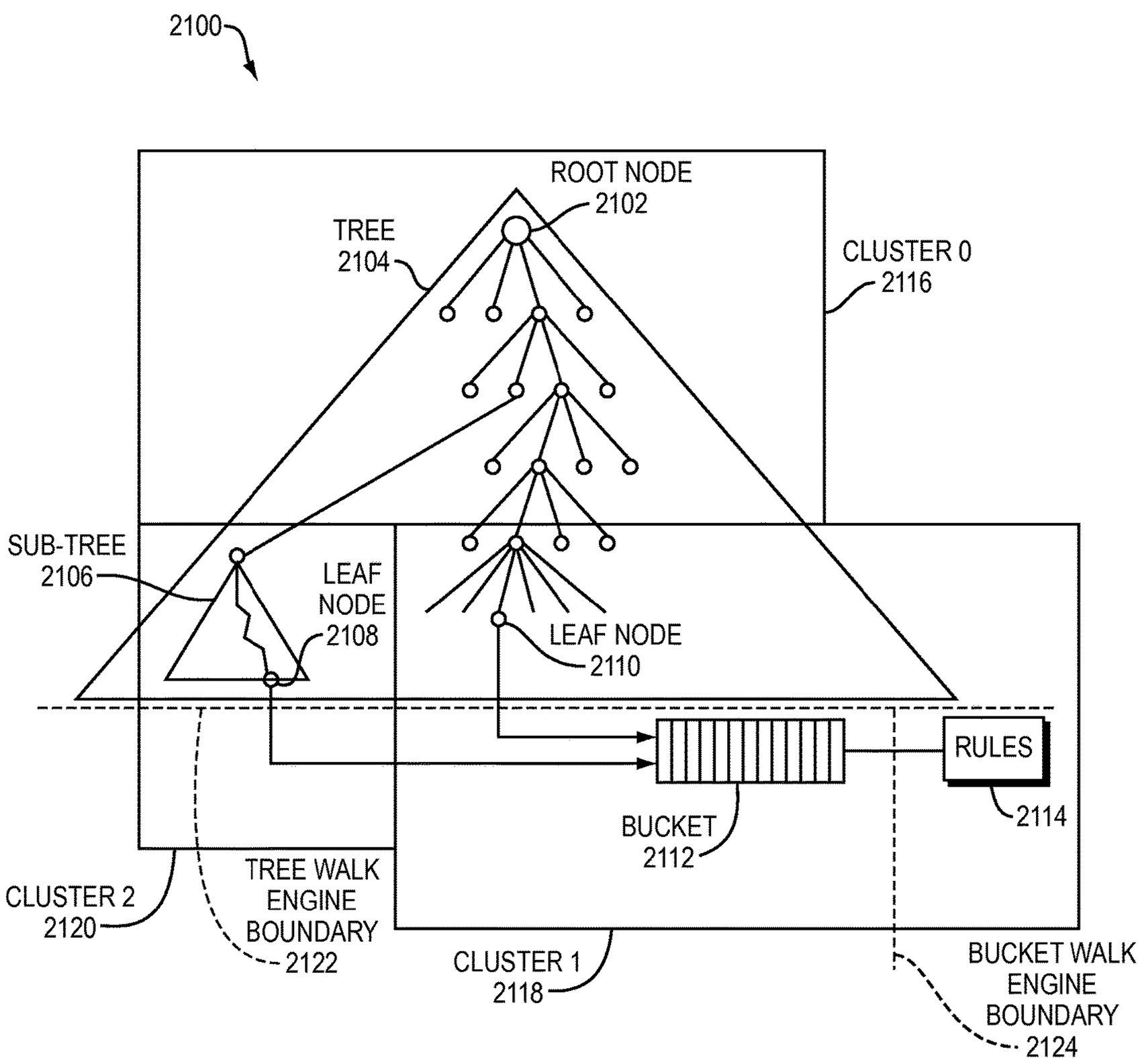
2100
ROOT NODE
2102
TREE
2104
CLUSTER 0
2116
SUB-TREE
2106
LEAF
NODE
2108
LEAF NODE
2110
RULES
2114
BUCKET
2112
CLUSTER 2
2120
TREE WALK
ENGINE
BOUNDARY
2122
CLUSTER 1
2118
BUCKET WALK
ENGINE
BOUNDARY
2124

FIG. 21

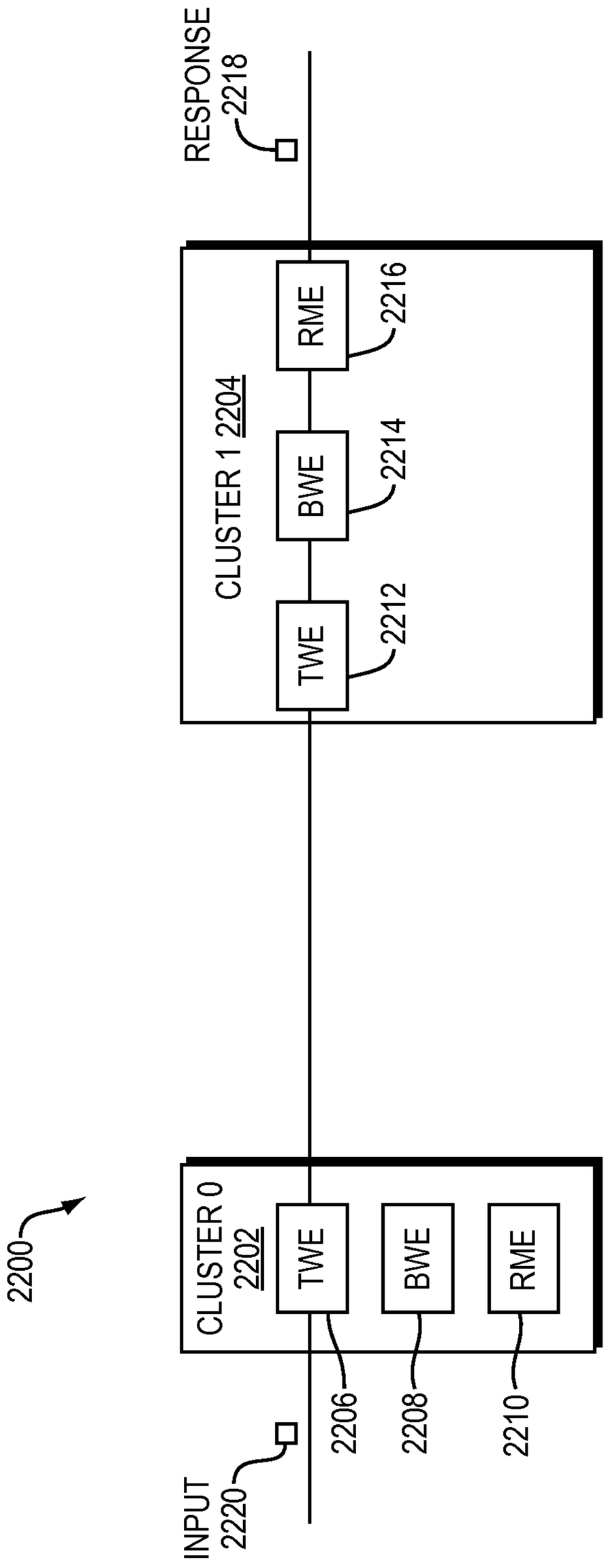
2200
CLUSTER 0
2202
TWE
BWE
RME
2206
2208
2210
INPUT
2220
CLUSTER 1 2204
TWE
BWE
RME
2212
2214
2216
RESPONSE
2218

FIG. 22A

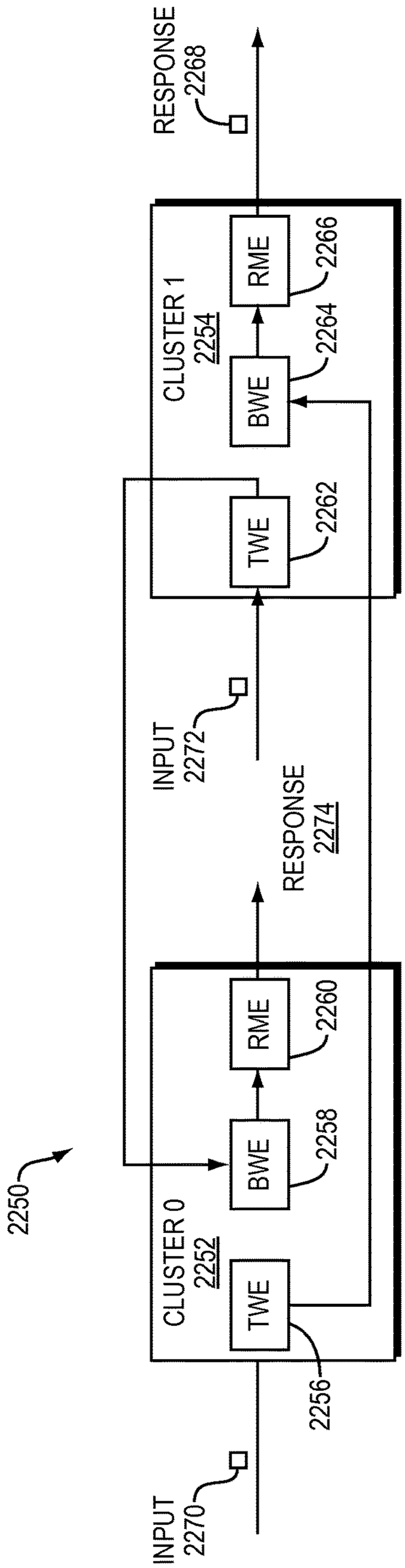
2250
CLUSTER 0
2252
TWE
BWE
RME
2256
2258
2260
INPUT
2270
RESPONSE
2274
INPUT
2272
CLUSTER 1
2254
TWE
BWE
RME
2262
2264
2266
RESPONSE
2268

FIG. 22B

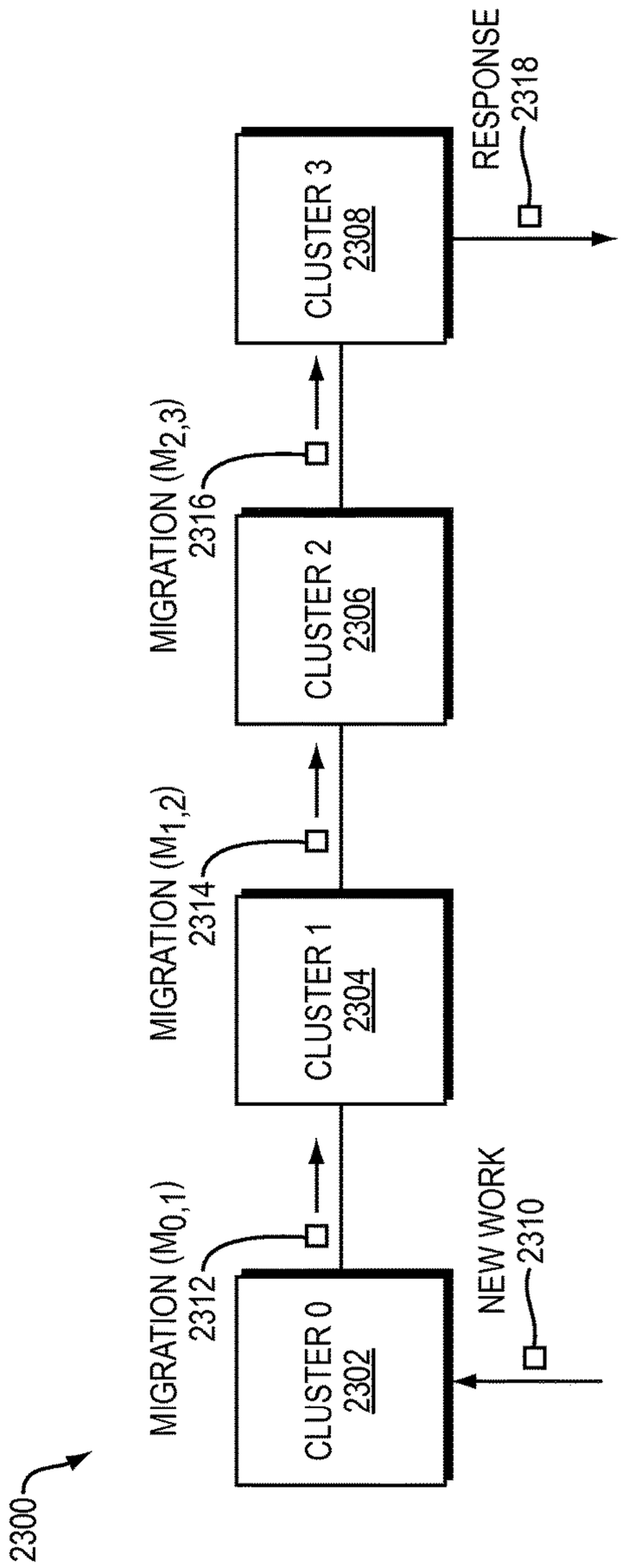

FIG. 23A

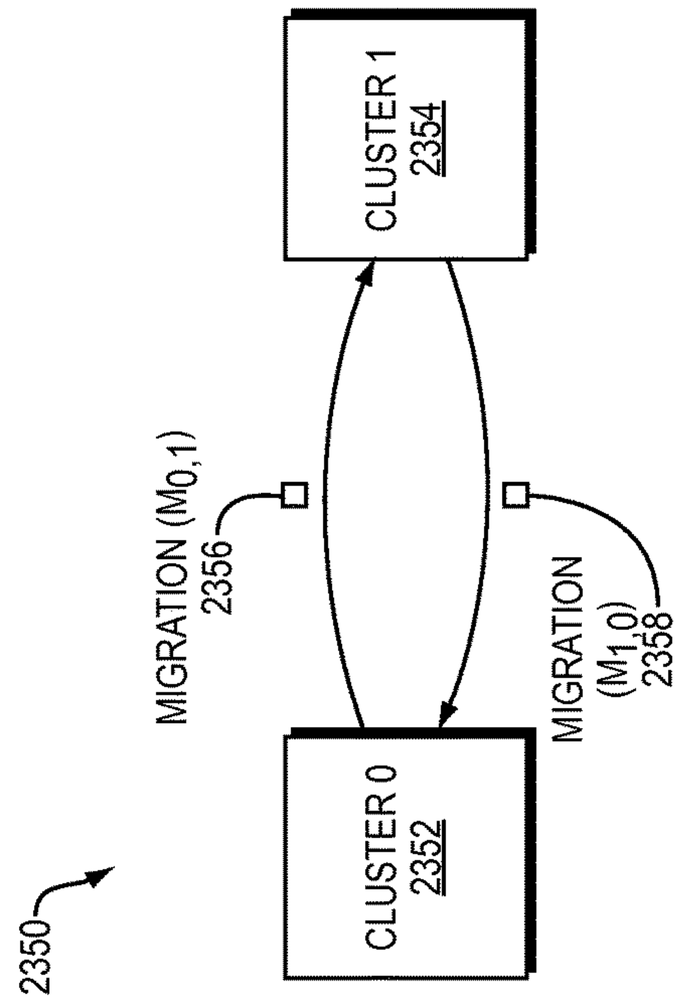

FIG. 23B

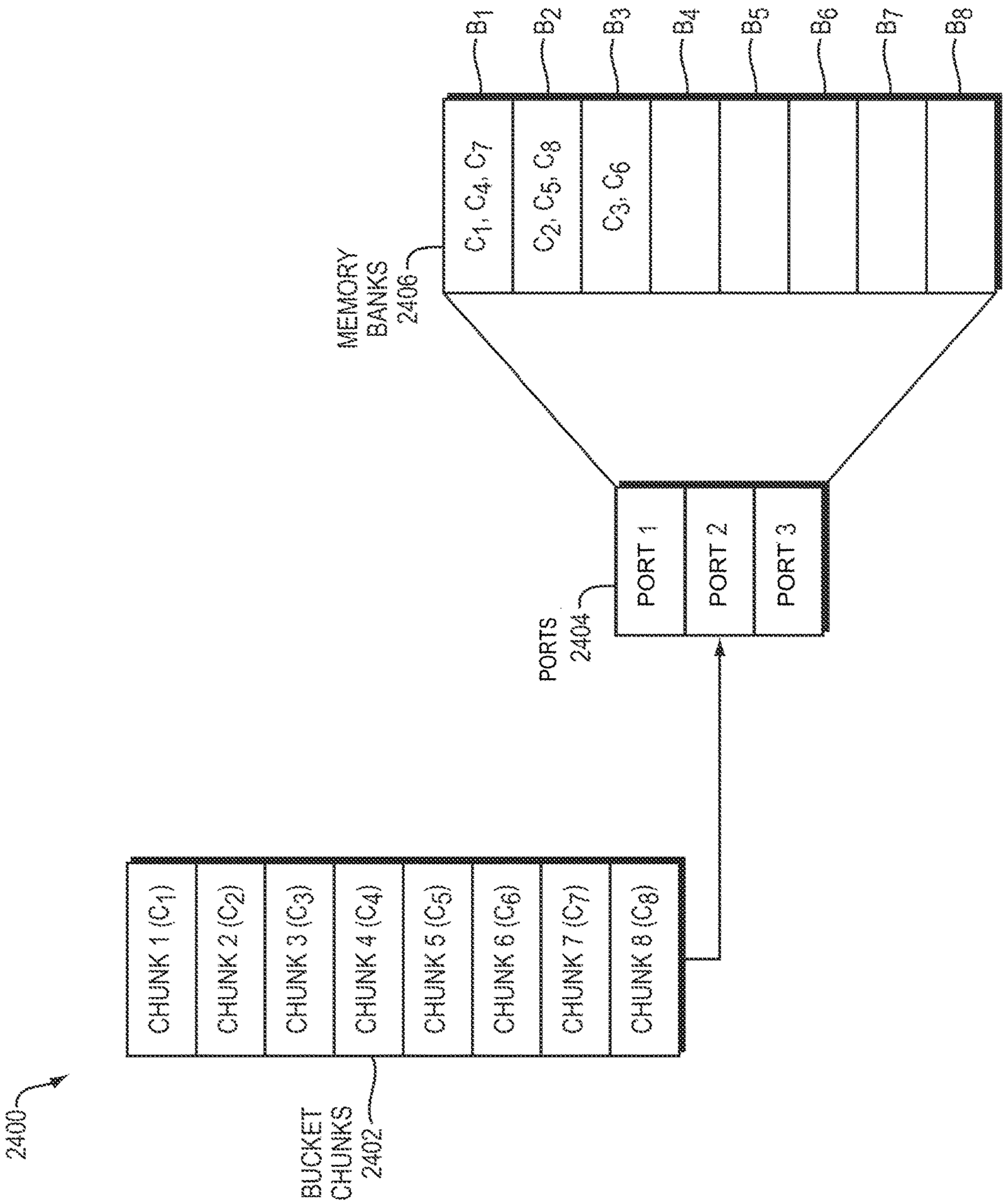
2400
BUCKET CHUNKS 2402
CHUNK 1 (C_1)
CHUNK 2 (C_2)
CHUNK 3 (C_3)
CHUNK 4 (C_4)
CHUNK 5 (C_5)
CHUNK 6 (C_6)
CHUNK 7 (C_7)
CHUNK 8 (C_8)
PORTS 2404
PORT 1
PORT 2
PORT 3
MEMORY BANKS 2406
C_1, C_4, C_7
C_2, C_5, C_8
C_3, C_6
B1
B2
B3
B4
B5
B6
B7
B8

FIG. 24

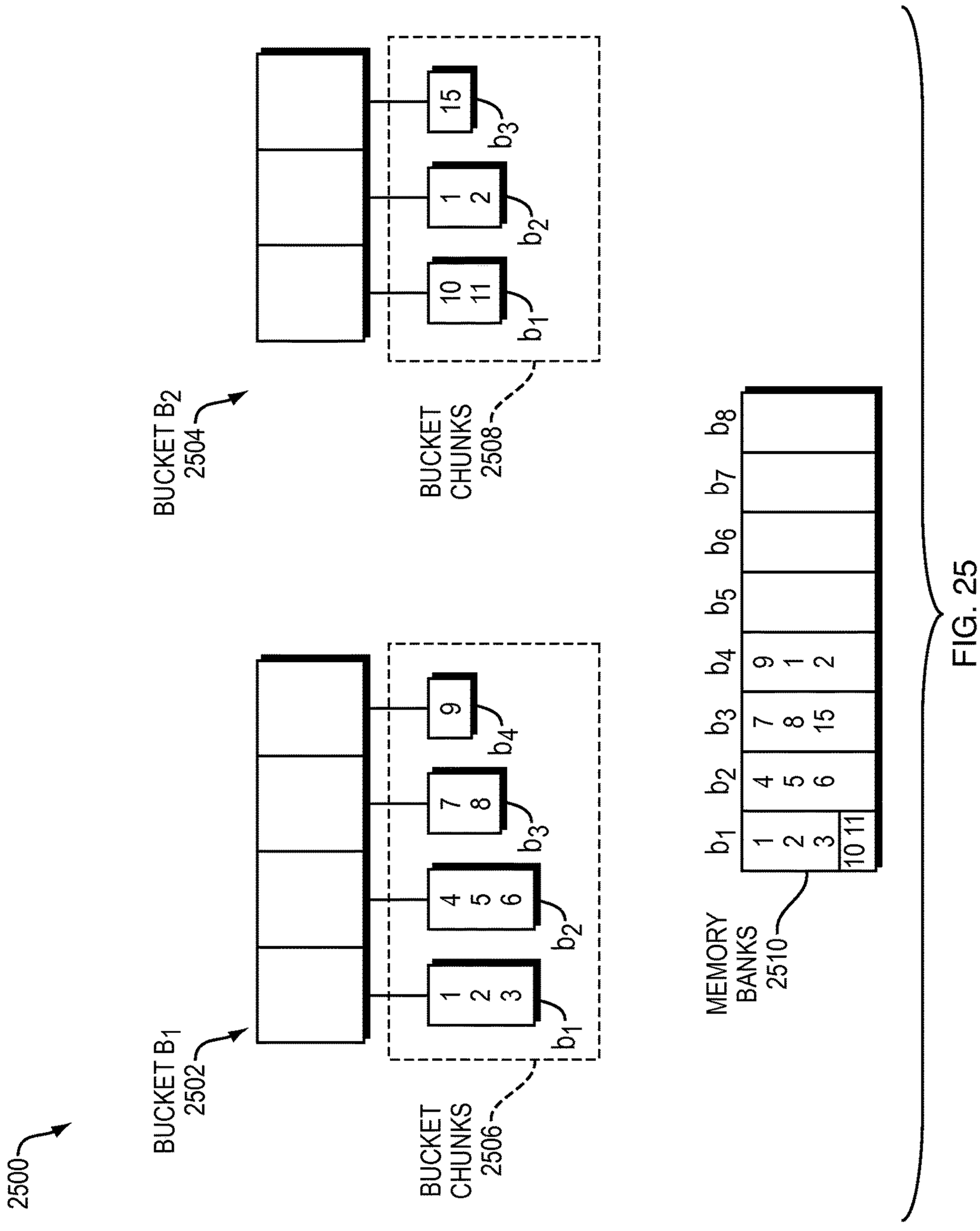
2500
BUCKET B1
2502
BUCKET CHUNKS
2506
1
2
3
b1
4
5
6
b2
7
8
b3
9
b4
BUCKET B2
2504
BUCKET CHUNKS
2508
10
11
15
MEMORY BANKS
2510
b1
b2
b3
b4
b5
b6
b7
b8

FIG. 25

NSP MANAGER

RELATED APPLICATIONS

This application claims the benefit of U.S. Provisional Application No. 61/799,013, filed on Mar. 15, 2013.

This application is being co-filed on Mar. 13, 2014 with U.S. patent application Ser. No. 14/207,928 entitled "Merging Independent Writes, Separating Dependent And Independent Writes, And Error Roll Back" by Satyanarayana Lakshmipathi Billa and Rajan Goyal.

This application is related to "A System And Method For Storing Lookup Request Rules In Multiple Memories," U.S. application Ser. No. 13/565,271, filed on Aug. 2, 2012, "A Method And Apparatus Encoding A Rule For A Lookup Request In A Processor," U.S. application Ser. No. 13/565,389, filed on Aug. 2, 2012, "A System And Method For Rule Matching In A Processor," U.S. application Ser. No. 13/565,406, filed on Aug. 2, 2012, "A Method And Apparatus For Managing Transfer Of Transport Operations From A Cluster In A Processor," U.S. application Ser. No. 13/565,743, filed on Aug. 2, 2012, and "Incremental Update," U.S. application Ser. No. 13/565,755, filed on Aug. 2, 2012.

The entire teachings of the above applications are incorporated herein by reference.

BACKGROUND

The Open Systems Interconnection (OSI) Reference Model defines seven network protocol layers (L1-L7) used to communicate over a transmission medium. The upper layers (L4-L7) represent end-to-end communications and the lower layers (L1-L3) represent local communications.

Networking application aware systems need to process, filter and switch a range of L3 to L7 network protocol layers, for example, L7 network protocol layers such as, HyperText Transfer Protocol (HTTP) and Simple Mail Transfer Protocol (SMTP), and L4 network protocol layers such as Transmission Control Protocol (TCP). In addition to processing the network protocol layers, the networking application aware systems need to simultaneously secure these protocols with access and content based security through L4-L7 network protocol layers including Firewall, Virtual Private Network (VPN), Secure Sockets Layer (SSL), Intrusion Detection System (IDS), Internet Protocol Security (IPSec), Anti-Virus (AV) and Anti-Spam functionality at wire-speed.

Improving the efficiency and security of network operation in today's Internet world remains an ultimate goal for Internet users. Access control, traffic engineering, intrusion detection, and many other network services require the discrimination of packets based on multiple fields of packet headers, which is called packet classification.

Internet routers classify packets to implement a number of advanced internet services such as routing, rate limiting, access control in firewalls, virtual bandwidth allocation, policy-based routing, service differentiation, load balancing, traffic shaping, and traffic billing. These services require the router to classify incoming packets into different flows and then to perform appropriate actions depending on this classification.

A classifier, using a set of filters or rules, specifies the flows, or classes. For example, each rule in a firewall might specify a set of source and destination addresses and associate a corresponding deny or permit action with it. Alternatively, the rules might be based on several fields of a packet header including layers 2, 3, 4, and 5 of the OSI model, which contain addressing and protocol information.

On some types of proprietary hardware, an Access Control List (ACL) refers to rules that are applied to port numbers or network daemon names that are available on a host or layer 3 device, each with a list of hosts and/or networks permitted to use a service. Both individual servers as well as routers can have network ACLs. ACLs can be configured to control both inbound and outbound traffic.

SUMMARY

In an embodiment, a method of managing a database including a tree, a plurality of buckets, and a plurality of rules, includes providing a memory with a plurality of cluster memories. Each cluster memory has a plurality of banks and a plurality of access ports. The memory stores the database across the plurality of cluster memories. The method also includes packing nodes of the tree in each of the plurality of cluster memories such that walking the tree accesses a minimal amount of cluster memories in the memory and walking the tree accesses each particular cluster memory no more than once.

In an embodiment, the method includes packing a first particular number of bucket chunks per bucket and a second particular number of rule pointers per bucket chunk based on addresses of the rules. Packing the first particular number of bucket chunks per bucket and the second particular number of rule pointers per bucket can include storing each of the first particular number of consecutively addressed bucket chunks in a respective unique memory bank, where the first particular number of bucket chunks is based on a quantity of the access ports.

In an embodiment, the method includes allocating the rules in the memory in a same order as an order of the rules in bucket chunks of the buckets.

In an embodiment, the method includes replicating a rule or a chunk of rules across a first and second bank in a particular cluster memory such that the rule or chunk of rules can be accessed on the second bank when the first bank has a memory access conflict during a particular clock cycle.

In an embodiment, the method includes distributing rules and buckets across a first cluster memory and a second cluster memory within the memory and enabling processing of the rules by a rule match engine within the first cluster memory and a rule match engine within the second cluster memory.

In an embodiment, the method includes storing a first rule of a particular bucket of the plurality of buckets in a particular cluster memory of the plurality of cluster memories and storing any other rules of the particular bucket in the particular cluster memory.

In an embodiment, the method includes allocating a node of the tree in a particular cluster memory of the memory and allocating a bucket in the memory that the node of the tree points to in the particular cluster memory.

In an embodiment, the method includes allocating a bucket in a particular cluster memory of the memory and allocating a rule associated with the bucket in the particular cluster memory storing the bucket.

In an embodiment, the method includes allocating the rules in chunks according to an order of the buckets, determining a need to replicate the rules across the plurality of cluster memories and replicating the rules across the plurality of cluster memories, if necessary.

In an embodiment, the method can include determining at least one division of the database, the database including the tree, the plurality of buckets, and the plurality of rules. The division can be based on either a horizontal division or a

vertical division. The horizontal division can separate the tree based on a depth of data of the tree. The vertical division can separate the tree based on sub-trees of the tree. The method can further include generating at least one memory request to store each division of the database in a respective cluster memory.

In an embodiment, the method can also include maintaining relationships among the tree, the plurality of buckets, and the plurality of rules. Allocating the memory can include allocating the tree, the plurality of buckets, and the plurality of rules to respective memory blocks based on the relationships to avoid migrations, remote reads, and bank conflicts by storing rules in a bucket of an optimal size and storing the tree, the plurality of buckets, and the plurality of rules within the same cluster across multiple banks.

In an embodiment, a system for managing a database including a tree, a plurality of buckets, and a plurality of rules includes a memory with a plurality of cluster memories. Each cluster memory has a plurality of banks and a plurality of access ports. The memory stores the database across the plurality of cluster memories. The system also includes a tree packing module configured to pack nodes of the tree in each of the plurality of cluster memories such that walking the tree accesses a minimal amount of cluster memories in the memory and walking the tree accesses each particular cluster memory no more than once.

In another embodiment, the system can include a bucket packet module configured to pack a first particular number of bucket chunks per bucket and a second particular number of rule pointers per bucket by storing each of the first particular number of consecutively addressed bucket chunks in a respective unique memory bank, where the first particular number of bucket chunks is based on a quantity of the access ports.

In an embodiment, a non-transitory computer-readable medium is configured to store instructions for managing a database including a tree, a plurality of buckets, and a plurality of rules. The instructions, when loaded and executed by a processor, can cause the processor to provide a memory with a plurality of cluster memories. Each cluster memory has a plurality of banks and a plurality of access ports. The memory stores the database across the plurality of cluster memories. The instructions can further cause the processor to pack nodes of the tree in each of the plurality of cluster memories such that walking the tree accesses a minimal amount of cluster memories in the memory and walking the tree accesses each particular cluster memory no more than once.

BRIEF DESCRIPTION OF THE DRAWINGS

The foregoing will be apparent from the following more particular description of example embodiments of the invention, as illustrated in the accompanying drawings in which like reference characters refer to the same parts throughout the different views. The drawings are not necessarily to scale, emphasis instead being placed upon illustrating embodiments of the present invention.

FIG. **1** is a block diagram of a typical network topology including network elements where a search processor may be employed.

FIGS. **2**A-**2**C show block diagrams illustrating example embodiments of routers employing a search processor.

FIG. **3** shows an example architecture of a search processor.

FIG. **4** is a block diagram illustrating an example embodiment of loading rules, by a software compiler, into an on-chip memory (OCM).

FIG. **5** shows a block diagram illustrating an example embodiment of a memory, or search, cluster.

FIGS. **6**A-**6**B show block diagrams illustrating example embodiments of transport operations between two search clusters.

FIG. **7** shows an example hardware implementation of the OCM in a search cluster.

FIGS. **8**A to **8**E show block and logic diagrams illustrating an example implementation of a crossbar controller (XBC).

FIGS. **9**A to **9**D show block and logic diagrams illustrating an example implementation of a crossbar (XBAR) and components therein.

FIGS. **10**A and **10**B show two example tables storing resource state information in terms of credits.

FIGS. **11**A to **11**C illustrate examples of interleaving transport operations and partial transport operations over consecutive clock cycles.

FIGS. **12**A and **12**B show flowcharts illustrating methods of managing transport operations between a first memory cluster and one or more other memory clusters performed by the XBC.

FIG. **13** shows a flowchart illustrating a method of assigning resources used in managing transport operations between a first memory cluster and one or more other memory clusters.

FIG. **14** shows a flow diagram illustrating a deadlock scenario in processing thread migrations between two memory clusters.

FIG. **15** shows a graphical illustration of an approach to avoid deadlock in processing thread migrations.

FIG. **16** is a flowchart illustrating a method of managing processing thread migrations within a plurality of memory clusters.

FIG. **17** is a block diagram illustrating an example embodiment of a search processor manager employed with a search processor.

FIG. **18** is a flow diagram illustrating an example embodiment of a control plane stack.

FIG. **19** is a block diagram illustrating an example embodiment of incremental update as managed by a search processor manager for a search processor.

FIG. **20**A is a block diagram illustrating an example embodiment of a tree, a bucket and rules.

FIG. **20**B is a block diagram illustrating an example embodiment of storing the tree, buckets and rules in multiple clusters in the memory.

FIG. **21** is a block diagram illustrating an example embodiment of storing the database in a plurality of clusters.

FIG. **22**A is a block diagram illustrating an example embodiment of two clusters receiving work in generating a response.

FIG. **22**B is a block diagram illustrating a more optimal setup for two clusters.

FIG. **23**A is a block diagram illustrating an example embodiment of migration from cluster to cluster.

FIG. **23**B is a block diagram illustrating an example embodiment of a loop formed during migration.

FIG. **24** is a diagram illustrating an example embodiment of writing bucket chunks to memory banks using access ports.

FIG. **25** is a diagram illustrating an example embodiment of memory replication.

DETAILED DESCRIPTION

A description of example embodiments of the invention follows.

Although packet classification has been widely studied for a long time, researchers are still motivated to seek novel and efficient packet classification solutions due to: i) the continued growth of network bandwidth, ii) increasing complexity of network applications, and iii) technology innovations of network systems.

Explosion in demand for network bandwidth is generally due to the growth in data traffic. Leading service providers report bandwidths doubling on their backbone networks about every six to nine months. As a consequence, novel packet classification solutions are required to handle the exponentially increasing traffics on both edge and core devices.

Complexity of network applications is increasing due to the increasing number of network applications being implemented in network devices. Packet classification is widely used for various kinds of applications, such as service-aware routing, intrusion prevention and traffic shaping. Therefore, novel solutions of packet classification must be intelligent to handle diverse types of rule sets without significant loss of performance.

In addition, new technologies, such as multi-core processors provide unprecedented computing power, as well as highly integrated resources. Thus, novel packet classification solutions must be well suited to advanced hardware and software technologies.

Existing packet classification algorithms trade memory for time. Although the tradeoffs have been constantly improving, the time taken for a reasonable amount of memory is still generally poor.

Because of problems with existing algorithmic schemes, designers use ternary content-addressable memory (TCAM), which uses brute-force parallel hardware to simultaneously check packets against all rules. The main advantages of TCAMs over algorithmic solutions are speed and determinism. TCAMs work for all databases.

A TCAM is a hardware device that functions as a fully associative memory. A TCAM cell stores three values: 0, 1, or 'X,' which represents a don't-care bit and operates as a per-cell mask enabling the TCAM to match rules containing wildcards, such as a kleene star '*'. In operation, a whole packet header can be presented to a TCAM to determine which entry, or rule, it matches. However, the complexity of TCAMs has allowed only small, inflexible, and relatively slow implementations that consume a lot of power. Therefore, a need continues for efficient algorithmic solutions operating on specialized data structures.

Current algorithmic methods remain in the stages of mathematical analysis and/or software simulation, that is observation based solutions.

Proposed mathematical solutions have been reported to have excellent time/spatial complexity. However, methods of this kind have not been found to have any implementation in real-life network devices because mathematical solutions often add special conditions to simplify a problem and/or omit large constant factors which might conceal an explicit worst-case bound.

Proposed observation based solutions employ statistical characteristics observed in rules to achieve efficient solution for real-life applications. However, these algorithmic methods generally only work well with a specific type of rule sets. Because packet classification rules for different applications have diverse features, few observation based methods are able to fully exploit redundancy in different types of rule sets to obtain stable performance under various conditions.

Packet classification is performed using a packet classifier, also called a policy database, flow classifier, or simply a classifier. A classifier is a collection of rules or policies. Packets received are matched with rules, which determine actions to take with a matched packet. Generic packet classification requires a router to classify a packet on the basis of multiple fields in a header of the packet. Each rule of the classifier specifies a class that a packet may belong to according to criteria on 'F' fields of the packet header and associates an identifier, e.g., class ID, with each class. For example, each rule in a flow classifier is a flow specification, in which each flow is in a separate class. The identifier uniquely specifies an action associated with each rule. Each rule has 'F' fields. An ith field of a rule R, referred to as R[i], is a regular expression on the ith field of the packet header. A packet P matches a particular rule R if for every i, the ith field of the header of P satisfies the regular expression R[i].

Classes specified by the rules may overlap. For instance, one packet may match several rules. In this case, when several rules overlap, an order in which the rules appear in the classifier determines the rules relative priority. In other words, a packet that matched multiple rules belongs to the class identified by the identifier, class ID, of the rule among them that appears first in the classifier.

Packet classifiers may analyze and categorize rules in a classifier table and create a decision tree that is used to match received packets with rules from the classifier table. A decision tree is a decision support tool that uses a graph or model of decisions and their possible consequences, including chance event outcomes, resource costs, and utility. Decision trees are commonly used in operations research, specifically in decision analysis, to help identify a strategy most likely to reach a goal. Another use of decision trees is as a descriptive means for calculating conditional probabilities. Decision trees may be used to match a received packet with a rule in a classifier table to determine how to process the received packet.

In simple terms, the problem may be defined as finding one or more rules, e.g., matching rules, that match a packet. Before describing a solution to this problem, it should be noted that a packet may be broken down into parts, such as a header, payload, and trailer. The header of the packet, or packet header, may be further broken down into fields, for example. So, the problem may be further defined as finding one or more rules that match one or more parts of the packet.

A possible solution to the foregoing problem(s) may be described, conceptually, by describing how a request to find one or more rules matching a packet or parts of the packet, a "lookup request," leads to finding one or more matching rules.

FIG. **1** is a block diagram **100** of a typical network topology including network elements where a search processor may be employed. The network topology includes an Internet core **102** including a plurality of core routers **104***a-h*. Each of the plurality of core routers **104***a-h* is connected to at least one other of the plurality of core routers **104***a-h*. Core routers **104***a-h* that are on the edge of the Internet core **102**, e.g., core routers **104***b-e* and **104***h*, are coupled with at least one edge router **106***a-f*. Each edge router **106***a-f* is coupled to at least one access router **108***a-e*.

The core routers **104***a*-**104***h* are configured to operate in the Internet core **102** or Internet backbone. The core routers

104*a*-**104***h* are configured to support multiple telecommunications interfaces of the Internet core **102** and are further configured to forward packets at a full speed of each of the multiple telecommunications protocols.

The edge routers **106***a*-**106***f* are placed at the edge of the Internet core **102**. Edge routers **106***a*-**106***f* bridge access routers **108***a*-**108***e* outside the Internet core **102** and core routers **104***a*-**104***h* in the Internet core **102**. Edge routers **106***a*-**106***f* may be configured to employ a bridging protocol to forward packets from access routers **108***a*-**108***e* to core routers **104***a*-**104***h* and vice versa.

The access routers **108***a*-**108***e* may be routers used by an end user, such as a home user or an office, to connect to one of the edge routers **106***a*-**106***f*, which in turn connects to the Internet core **102** by connecting to one of the core routers **104***a*-**104***h*. In this manner, the edge routers **106***a*-**106***f* may connect to any other edge router **106***a*-**106***f* via the edge routers **106***a*-**106***f* and the interconnected core routers **104***a*-**104***h*.

The search processor described herein may reside in any of the core routers **104***a*-**104***h*, edge routers **106***a*-**106***f*, or access routers **108***a*-**108***e*. The search processor described herein, within each of these routers, is configured to analyze Internet protocol (IP) packets based on a set of rules and forward the IP packets along an appropriate network path.

FIG. **2**A is a block diagram illustrating an example embodiment of an edge router **106** employing a search processor **202**. An edge router **106**, such as a service provider edge router, includes the search processor **202**, a first host processor **204** and a second host processor **214**. Examples of the first host processor include processors such as a network processor unit (NPU), a custom application-specific integrated circuit (ASIC), an OCTEON® processor available from Cavium Inc., or the like. The first host processor **204** is configured as an ingress host processor. The first host processor **204** receives ingress packets **206** from a network. Upon receiving a packet, the first host processor **204** forwards a lookup request including a packet header, or field, from the ingress packets **206** to the search processor **202** using an Interlaken interface **208**. The search processor **202** then processes the packet header using a plurality of rule processing engines employing a plurality of rules to determine a path to forward the ingress packets **206** on the network. The search processor **202**, after processing the lookup request with the packet header, forwards the path information to the first host processor **204**, which forwards the processed ingress packets **210** to another network element in the network.

Likewise, the second host processor **214** is an egress host processor. Examples of the second host processor include processors such as a NPU, a custom ASIC, an OCTEON processor, or the like. The second host processor **214** receives egress packets **216** to send to the network. The second host processor **214** forwards a lookup request with a packet header, or field, from the egress packets **216** to the search processor **202** over a second Interlaken interface **218**. The search processor **202** then processes the packet header using a plurality of rule processing engines employing a plurality of rules to determine a path to forward the packets on the network. The search processor **202** forwards the processed egress packets **221** from the host processor **214** to another network element in the network.

FIG. **2**B is a block diagram **220** illustrating another example embodiment of an edge router **106** configured to employ the search processor **202**. In this embodiment, the edge router **106** includes a plurality of search processors **202**, for example, a first search processor **202***a* and a second search processor **202***b*. The plurality of search processors **202***a*-**202***b* are coupled to a packet processor **228** using a plurality of Interlaken interfaces **226***a*-*b*, respectively. Examples of the packet processor **228** include processors such as NPU, ASIC, or the like. The plurality of search processors **202***a*-**202***b* may be coupled to the packet processor **228** over a single Interlaken interface. The edge router **106** receives a lookup request with a packet header, or fields, of pre-processed packets **222** at the packet processor **228**. The packet processor **228** sends the lookup request to one of the search processors **202***a*-**202***b*. The search processor, **202***a* or **202***b*, searches a packet header for an appropriate forwarding destination for the pre-processed packets **222** based on a set of rules and data within the packet header, and responds to the lookup request to the packet processor **228**. The packet processor **228** then sends the post processed packets **224** to the network based on the response to the lookup request from the search processors **202***a*-**202***b*.

FIG. **2**C is a block diagram **240** illustrating an example embodiment of an access router **246** employing the search processor **202**. The access router **246** receives an input packet **250** at an ingress packet processor **242**. Examples of the ingress packet processor **242** include OCTEON processor, or the like. The ingress packet processor **242** then forwards a lookup request with a packet header of the input packet **250** to the search processor **202**. The search processor **202** determines, based on packet header of the lookup request, a forwarding path for the input packet **250** and responds to the lookup requests over the Interlaken interface **252** to the egress packet processor **244**. The egress packet processor **244** then outputs the forwarded packet **248** to the network.

FIG. **3** shows an example architecture of a search processor **202**. The processor includes, among other things, an interface, e.g., Interlaken LA interface, **302** to receive requests from a host processor, e.g., **204**, **214**, **228**, **242**, or **244**, and to send responses to the host processor. The interface **302** is coupled to Lookup Front-end (LUF) processors **304** configured to process, schedule, and order the requests and responses communicated from or to the interface **302**. According to an example embodiment, each of the LUF processors is coupled to one of the super clusters **310**. Each super cluster **310** includes one or more memory clusters, or search clusters, **320**. Each of the memory, or search, clusters **320** includes a Lookup Engine (LUE) component **322** and a corresponding on-chip memory (OCM) component **324**. A memory, or search, cluster may be viewed as a search block including a LUE component **322** and a corresponding OCM component **324**. Each LUE component **322** is associated with a corresponding OCM component **324**. A LUE component **322** includes processing engines configured to search for rules in a corresponding OCM component **324**, given a request, that match keys for packet classification. The LUE component **322** may also include interface logic, or engine(s), configured to manage transport of data between different components within the memory cluster **320** and communications with other clusters. The memory clusters **320**, in a given super cluster **310**, are coupled through an interface device, e.g., crossbar (XBAR), **312**. The XBAR **312** may be viewed as an intelligent fabric enabling coupling LUF processors **304** to different memory clusters **320** as well as coupling between different memory clusters **320** in the same super cluster **310**. The search processor **202** may include one or more super clusters **310**. A lookup cluster complex (LCC) **330** defines the group of super clusters **310** in the search processor **202**.

The search processor **202** may also include a memory walker aggregator (MWA) **303** and at least one memory block controller (MBC) **305** to coordinate read and write operations from/to memory located external to the processor. The search processor **202** may further include one or more Bucket Post Processors (BPPs) **307** to search rules, which are stored in memory located external to the search processor **202**, that match keys for packet classification.

FIG. **4** is a block diagram **400** illustrating an example embodiment of loading rules, by a software compiler, into OCM components. According to an example embodiment, the software compiler **404** is software executed by a host processor or control plane processor to store rules into the search processor **202**. Specifically, rules are loaded to at least one OCM component **324** of at least one memory cluster, or search block, **320** in the search processor **202**. According to at least one example embodiment, the software compiler **404** uses multiple data structures, in storing the rules, in a way to facilitate the search of the stored rules at a later time. The software compiler **404** receives a rule set **402**, parameter(s) indicative of a maximum tree depth **406** and parameter(s) indicative of a number of sub-trees **408**. The software compiler **404** generates a set of compiled rules formatted, according at least one example embodiment, as linked data structures referred to hereinafter as rule compiled data structure (RCDS) **410**. The RCDS is stored in at least one OCM component **324** of at least one memory cluster, or search block, **320** in the search processor **202**. The RCDS **410** includes at least one tree **412**. Each tree **412** includes nodes **411***a*-**411***c*, leaf nodes **413***a*-**413***b*, and a root node **432**. A leaf node, **413***a*-**413***b*, of the tree **412** includes or points to one of a set of buckets **414**. A bucket **414** may be viewed as a sequence of bucket entries, each bucket entry storing a pointer or an address, referred to hereinafter as a chunk pointer **418**, of a chunk of rules **420**. Buckets may be implemented, for example, using tables, linked lists, or any other data structures known in the art adequate for storing a sequence of entries. A chunk of rules **420** is basically a chunk of data describing or representing one or more rules. In other words, a set of rules **416** stored in one or more OCM components **324** of the search processor **202** include chunks of rules **420**. A chunk of rules **420** may be a sequential group of rules, or a group of rules scattered throughout the memory, either organized by a plurality of pointers or by recollecting the scattered chunk of rules **420**, for example, using a hash function.

The RCDS **410** described in FIG. **4** illustrates an example approach of storing rules in the search engine **202**. A person skilled in the art should appreciate that other approaches of using nested data structures may be employed. For example, a table with entries including chunk pointers **418** may be used instead of the tree **412**. In designing a rule compiled data structure for storing and accessing rules used to classify data packets, one of the factors to be considered is enabling efficient and fast search or access of such rules.

Once the rules are stored in the search processor **202**, the rules may then be accessed to classify data packets. When a host processor receives a data packet, the host processor forwards a lookup request with a packet header, or field, from the data packet to the search processor **202**. On the search processor side, a process of handling the received lookup request includes:

1) The search processor receives the lookup request from the host processor. According to at least one example embodiment, the lookup request received from the host processor includes a packet header and a group identifier (GID).

2) The GID indexes an entry in a group definition table (GDT). Each GDT entry includes n number of table identifiers (TID), a packet header index (PHIDX), and key format table index (KFTIDX).

3) Each TID indexes an entry in a tree location table (TLT). Each TLT entry identifies which lookup engine or processor will look for the one or more matching rules. In this way, each TID specifies both who will look for the one or more matching rules and where to look for the one or more matching rules.

4) Each TID also indexes an entry in a tree access table (TAT). TAT is used in the context in which multiple lookup engines, grouped together in a super cluster, look for the one or more matching rules. Each TAT entry provides the starting address in memory of a collection of rules, or pointers to rules, called a table or tree of rules. The terms table of rules or tree of rules, or simply table or tree, are used interchangeably hereinafter. The TID identifies which collection or set of rules in which to look for one or more matching rules.

5) The PHIDX indexes an entry in a packet header table (PHT). Each entry in the PHT describes how to extract n number of keys from the packet header.

6) The KFTIDX indexes an entry in a key format table (KFT). Each entry in the KFT provides instructions for extracting one or more fields, e.g., parts of the packet header, from each of the n number of keys, which were extracted from the packet header.

7) Each of the extracted fields, together with each of the TIDs are used to look for subsets of the rules. Each subset contains rules that may possibly match each of the extracted fields.

8) Each rule of each subset is then compared against an extracted field. Rules that match are provided in responses, or lookup responses.

The handling of the lookup request and its enumerated stages, described above, are being provided for illustration purposes. A person skilled in the art should appreciate that different names as well as different formatting for the data included in a look up request may be employed. A person skilled in the art should also appreciate that at least part of the data included in the look up request is dependent on the design of the RCDS used in storing matching rules in a memory, or search, cluster **320**.

FIG. **5** shows a block diagram illustrating an example embodiment of a memory, or search, cluster **320**. The memory, or search, cluster **320** includes an on-chip memory (OCM) **324**, a plurality of processing, or search, engines **510**, an OCM bank slotter (OBS) module **520**, and a cross-bar controller (XBC) **530**. The OCM **324** includes one or more memory banks According to an example implementation, the OCM **324** includes two mega bytes (MBs) of memory divided into 16 memory banks According to the example implementation, the OCM **324** includes 64 k, or 65536, of rows each 256 bits wide. As such, each of the 16 memory banks has 4096 contiguous rows, each 256 bits wide. A person skilled in the art should appreciate that the described example implementation is provided for illustration and the OCM may, for example, have more or less than 2 MBs of memory and the number of memory banks may be different from 16. The number of memory rows, the number of bits in each memory row, as well as the distribution of memory rows between different memory banks may be different from the illustration in the described example implementation. The OCM **324** is configured to store, and provide access to, the RCDS **410**. In storing the RCDS **410**, the distribution of the data associated with the RCDS **410**

among different memory banks may be done in different ways. For example, different data structures, e.g., the tree data structure(s), the bucket storage data structure(s), and the chunk rule data structure(s), may be stored in different memory banks. Alternatively, a single memory bank may store data associated with more than one data structure. For example, a given memory bank may store a portion of the tree data structure, a portion of the bucket data structure, and a portion of the chunk rule data structure.

The plurality of processing engines **510** include, for example, a tree walk engine (TWE) **512**, a bucket walk engine (BWE) **514**, one or more rule walk engines (RWE) **516**, and one or more rule matching engines (RME) **518**. When the search processor **202** receives a request, called a lookup request, from the host processor, the LUF processor **304** processes the lookup request into one or more key requests, each of which has a key **502**. The LUF processor **304** then schedules the key requests to the search cluster. The search cluster **320** receives a key **502** from the LUF processor **304** at the TWE **512**. A key represents, for example, a field extracted from a packet header. The TWE **512** is configured to issue requests to access the tree **412** in the OCM **324** and receive corresponding responses. A tree access request includes a key used to enable the TWE **512** to walk, or traverse, the tree from a root node **432** to a possible leaf node **413**. If the TWE **512** does not find an appropriate leaf node, the TWE **512** issues a no match response to the LUF processor **304**. If the TWE **512** does find an appropriate leaf node, it issues a response that an appropriate leaf node is found.

The response that an appropriate leaf node is found includes, for example, a pointer to a bucket passed by the TWE **512** to the BWE **514**. The BWE **514** is configured to issue requests to access buckets **414** in the OCM **324** and receive corresponding responses. The BWE **514**, for example, uses the pointer to the bucket received from the TWE **512** to access one or more buckets **414** and retrieve at least one chunk pointer **418** pointing to a chunk of rules. The BWE **514** provides the retrieved at least one chunk pointer **418** to at least one RWE **516**. According to at least one example, BWE **514** may initiate a plurality of rule searched to be processed by one RWE **516**. However, the maximum number of outstanding, or on-going, rule searches at any point of time may be constrained, e.g., maximum of 16 rule searches. The RWE is configured to issue requests to access rule chunks **420** in the OCM **324** and receive corresponding responses. The RWE **416** uses a received chunk pointer **418** to access rule chunks stored in the OCM **324** and retrieve one or more rule chunks. The retrieved one or more rule chunks are then passed to one or more RMEs **518**. An RME **518**, upon receiving a chunk rule, is configured to check whether there is a match between one or more rules in the retrieved rule chunk and the field corresponding to the key.

The RME **518** is also configured to provide a response, to the BWE **514**. The response is indicative of a match, no match, or an error. In the case of a match, the response may also include an address of the matched rule in the OCM **324** and information indicative of a relative priority of the matched rule. Upon receiving a response, the BWE **514** decides how to proceed. If the response is indicative of a no match, the BWE **514** continues searching bucket entries and initiating more rule searches. If at some point the BWE **514** receives a response indicative of a match, it stops initiating new rule searches and waits for any outstanding rule searches to complete processing. Then, the BWE **514** provides a response to the host processor through the LUF processor **304**, indicating that there is a match between the field corresponding to the key and one or more rules in the retrieved rule chunk(s), e.g., a "match found" response. If the BWE **514** finishes searching buckets without receiving any "match found" response, the BWE **514** reports a response to the host processor through the LUF processor **304** indicating that there is no match, e.g., "no-match found" response. According to at least one example embodiment, the BWE **514** and RWE **516** may be combined into a single processing engine performing both bucket and rule chunk data searches. According to an example embodiment the RWEs **516** and the RMEs **518** may be separate processors. According to another example embodiment, the access and retrieval of rule chunks **420** may be performed by the RMEs **518** which also performs rule matching. In other words, the RMEs and the RWEs may be the same processors.

Access requests from the TWE **512**, the BWE **514**, or the RWE(s) are sent to the OBS module **520**. The OBS module **520** is coupled to the memory banks in the OCM **324** through a number of logical, or access, ports, e.g., M ports. The number of the access ports enforce constraints on the number of access requests that may be executed, or the number of memory banks that may be accessed, at a given clock cycle. For example, over a typical logical port no more than one access request may be executed, or sent, at a given clock cycle. As such, the maximum number of access requests that may be executed, or forwarded to the OCM **324**, per clock cycle is equal to M. The OBS module **520** includes a scheduler, or a scheduling module, configured to select a subset of access requests, from multiple access requests received in the OBS module **520**, to be executed in at least one clock cycle and to schedule the selected subset of access requests each over a separate access port. The OBS module **520** attempts to maximize OCM usage by scheduling up to M access requests to be forwarded to the OCM **324** per clock cycle. In scheduling access requests, the OBS module **520** also aims at avoiding memory bank conflict and providing low latency for access requests. Memory bank conflict occurs, for example, when attempting to access a memory bank by more than one access request at a given clock cycle. Low latency is usually achieved by preventing access requests from waiting for a long time in the OBS module **520** before being scheduled or executed.

Upon data being accessed in the OCM **324**, a response is then sent back to a corresponding engine/entity through a "Read Data Path" (RDP) component **540**. The RDP component **540** receives OCM read response data and context, or steering, information from the OBS. Read response data from each OCM port is then directed towards the appropriate engine/entity. The RDP component **540** is, for example, a piece of logic or circuit configured to direct data responses from the OCM **324** to appropriate entities or engines, such as TWE **512**, BWE **514**, RWE **516**, a host interface component (HST) **550**, and a cross-bar controller (XBC) **530**. The HST **550** is configured to store access requests initiated by the host processor or a respective software executing thereon. The context, or steering, information tells the RDP component **540** what to do with read data that arrives from the OCM **324**. According to at least one example embodiment, the OCM **324** itself does not contain any indication that valid read data is being presented to the RDP component **540**. Therefore, per-port context information is passed from the OBS module **520** to the RDP component **540** indicating to the RDP component **540** that data is arriving from the OCM **324** on the port, the type of data being received, e.g., tree data, bucket data, rule chunk data, or host data, and the destination of the read response data, e.g., TWE **512**, BWE **514**, RWE **516**, HST **550** or XBC **530**. For example, tree

data is directed to TWE **512** or XBC **530** if remote, bucket data is directed to BWE **514** or XBC if remote, rule chunk data is directed to RWE **516** or XBC **530** if remote, and host read data is directed to the HST **550**.

The search cluster **320** also includes the crossbar controller (XBC) **530** which is a communication interface managing communications, or transport operations, between the search cluster **320** and other search clusters through the crossbar (XBAR) **312**. In other words, the XBC **530** is configured to manage pushing and pulling of data to, and respectively from, the XBAR **312**.

According to an example embodiment, for rule processing, the processing engines **510** include a tree walk engine (TWE) **512**, bucket walk engine (BWE) **514**, rule walk engine (RWE) **516** and rule match engine (RME) **518**. According to another example embodiment, rule processing is extended to external memory and the BPP **307** also includes a RWE **516** and RME **518**, or a RME acting as both RWE **516** and RME **518**. In other words, the rules may reside in the on-chip memory and in this case, the RWE or RME engaged by the BWE, e.g., by passing a chunk pointer, is part of the same LUE as BWE. As such, the BWE engages a "local" RWE or RME. The rules may also reside on a memory located external to the search processor **202**, e.g., off-chip memory. In this case, which may be referred to as rule processing extended to external memory or, simply, "rule extension," the bucket walk engine does not engage a local RWE or RME. Instead, the BWE sends a request message, via the MWA **303** and MBC **305**, to a memory controller to read a portion, or chunk, of rules. The BWE **514** also sends a "sideband" message to the BPP **307** informing the BPP **307** that the chunk, associated with a given key, is stored in external memory.

The BPP **307** starts processing the chunk of rules received from the external memory. As part of the processing, if the BPP **307** finds a match, the BPP **307** sends a response, referred to as a lookup response or sub-tree response, to the LUF processor **304**. The BPP **307** also sends a message to the LUEs component **322** informing the LUEs component **322** that the BPP **307** is done processing the chunk and the LUEs component **322** is now free to move on to another request. If the BPP **307** does not find a match and the BPP **307** is done processing the chunk, the BPP **307** sends a message to the LUEs component **322** informing the LUEs component **322** that the BPP **307** is done processing and to send the BPP **307** more chunks to process. The LUEs component **322** then sends a "sideband" message, through the MWA **303** and MBC **305**, informing the BPP **307** about a next chunk of rules, and so on. For the last chunk of rules, the LUEs component **322** sends a "sideband" message to the BPP **307** informing the BPP **307** that the chunk, which is to be processed by the BPP **307**, is the last chunk. The LUEs component **322** knows that the chunk is the last chunk because the LUEs component **322** knows the total size of the set of rule chunks to be processed. Given the last chunk, if the BPP **307** does not find a match, the BPP **307** sends a "no-match" response to the LUF processor **304** informing the LUF processor **304** that the BPP **307** is done with the set of rule chunks. In turn, the LUEs component **322** frees up the context, e.g., information related to the processed key request or the respective work done, and moves on to another key request.

FIG. **6**A shows a block diagram illustrating an example embodiment of processing a remote access request between two search clusters. A remote access request is a request generated by an engine/entity in a first search cluster to access data stored in a second search cluster or memory outside the first search cluster. For example, a processing engine in cluster 1, **320***a*, sends a remote access request for accessing data in another cluster, e.g., cluster N **320***b*. The remote access request may be, for example, a tree data access request generated by a TWE **512***a* in cluster 1, a bucket access request generated by a BWE **514***a* in cluster 1, or a rule chunk data access request generated by a RWE **516***a* or RME in cluster 1. The remote access request is pushed by the XBC **530***a* of cluster 1 to the XBAR **312** and then sent to the XBC **530***b* of cluster N. The XBC **530***b* of cluster N then forwards the remote access request to the OBS module **520***b* of cluster N. The OBS module **520***b* directs the remote access request to OCM **324***b* of cluster N and a remote response is sent back from the OCM **324***b* to the XBC **530***b* through the RDP **540***b*. The XBC **530***b* forwards the remote response to the XBC **530***a* through the XBAR **312**. The XBC **530***a* then forwards the remote response to the respective processing engine in the LUEs component **322***a*.

FIG. **6**B shows a block diagram illustrating an example embodiment of a processing thread migration between two search clusters. Migration requests originate from a TWE **512** or BWE **514** as they relate mainly to a bucket search/access process or a tree search/access process, in a first cluster, that is configured to continue processing in a second cluster. Unlike remote access where data is requested and received from the second cluster, in processing thread migration the process itself migrates and continues processing in the second cluster. As such, information related to the processing thread, e.g., state information, is migrated to the second cluster from the first cluster. As illustrated in FIG. **6**B, processing thread migration requests are sent from TWE **512***a* or BWE **514***a* directly to the XBC **530***a* in the cluster 1, **320***a*. The XBC **530***a* sends the migration request through the crossbar (XBAR) **312** to the XBC **530***b* in cluster N, **320***b*. At the receiving cluster, e.g., cluster N **320***b*, the XBC **530***b* forwards the migration request to the proper engine, e.g., TWE **512***b* or BWE **514***b*. According to at least one example embodiment, the XBC, e.g., **530***a* and **530***b*, does not just forward requests. The XBC arbitrates which, among remote OCM requests, OCM response data, and migration requests, to be sent at a clock cycle.

FIG. **7** shows an example hardware implementation of the OCM **324** in a cluster **320**. According to the example implementation shown in FIG. **7**, the OCM includes a plurality, e.g., 16, single-ported memory banks **705***a*-**705***p*. Each memory bank, for example, includes 4096 memory rows, each of 256 bits width. A person skilled in the art should appreciate that the number, e.g., 16, of the memory banks and their storage capacity are chosen for illustration purposes and should not be interpreted as limiting. Each of the memory banks **705***a*-**705***p* is coupled to at least one input multiplexer **715***a*-**715***p* and at least one output multiplexer **725***a*-**725**-*p*. Each input multiplexer, among the multiplexers **715***a*-**715***p*, couples the input logical ports **710***a*-**710***d* to a corresponding memory bank among the memory banks **705***a*-**705***p*. Similarly, each output multiplexer, among the multiplexers **725***a*-**725***p*, couples the output logical ports **720***a*-**720***d* to a corresponding memory bank among the memory banks **705***a*-**705***p*.

The input logical ports **710***a*-**710***d* carry access requests' data from the OBS module **520** to respective memory banks among the memory banks **705***a*-**705***p*. The output logical ports **720***a*-**720***d* carry access responses' data from respective memory banks, among the memory banks **705***a*-**705***p*, to RDP component **540**. Given that the memory banks **705***a*-**705***p* are single-ported, at each clock cycle a single access is

permitted to each of the memory banks **705***a*-**705***p*. Also given the fact that there are four input logical/access ports, a maximum of four requests may be executed, or served, at a given clock cycle because no more than one logical port may be addressed to the same physical memory bank at the same clock cycle. For a similar reason, e.g., four output logical/access ports, a maximum of four responses may be sent out of the OCM **324** at a given clock cycle. An input multiplexer is configured to select a request, or decide which request, to access the corresponding physical memory bank. An output multiplexer is configured to select an access port on which a response from a corresponding physical memory bank is to be sent. For example, an output multiplexer may select an output logical port, to send a response, corresponding to an input logical port on which the corresponding request was received. A person skilled in the art should appreciate that other implementations with more, or less, than four ports may be employed.

According to an example embodiment, an access request is formatted as an 18 bit tuple. Among the 18 bits, two bits are used as wire interface indicating an access instruction/command, e.g., read, write, or idle, four bits are used to specify a memory bank among the memory banks **705***a*-**705***p*, and 12 bits are used to identify a row, among the 4096 rows, in the specified memory bank. In the case of a "write" command, 256 bits of data to be written are also sent to the appropriate memory bank. A person skilled in the art should appreciate that such format/structure is appropriate for the hardware implementation shown in FIG. **7**. For example, using 4 bits to specify a memory bank is appropriate if the total number of memory banks is 16 or less. Also the number of bits used to identify a row is correlated to the total number of rows in each memory bank. Therefore, the request format described above is provided for illustration purpose and a person skilled in the art should appreciate that many other formats may be employed.

The use of multi-banks as suggested by the implementation in FIG. **7**, enables accessing multiple physical memory banks per clock cycle, and therefore enables serving, or executing, more than one request/response per clock cycle. However, for each physical memory bank a single access, e.g., read or write, is allowed per clock cycle. According to an example embodiment, different types of data, e.g., tree data, bucket data, or rule chunk data, are stored in separate physical memory banks. Alternatively, a physical memory bank may store data from different types, e.g., tree data, bucket data, and rule chunk data. Using single-ported physical memory banks leads to more power efficiency compare to multi-port physical memory banks. However, multi-port physical memory banks may also be employed.

Processing operations, e.g., tree search, bucket search, or rule chunk search, may include processing across memory clusters. For example, a processing operation running in a first memory cluster may require accessing data stored in one or more other memory clusters. In such a case, a remote access request may be generated, for example by a respective processing engine, and sent to at least one of the one or more other memory clusters and a remote access response with the requested data may then be received. Alternatively, the processing operation may migrate to at least one of the one or more other memory clusters and continue processing therein. For example, a remote access request may be generated if the size of the data to be accessed from another memory cluster is relatively small and therefore the data may be requested and acquired in relatively short time period. However, if the data to be accessed is of relatively large size, then it may be more efficient to proceed with a processing thread migration where the processing operation migrates and continue processing in the other memory cluster. The transfer of data, related to a processing operation, between different memory clusters is referred to hereinafter as a transport operation. Transport operations, or transactions, include processing thread migration operation(s), remote access request operation(s), and remote access response operation(s). According to an example embodiment, transport operations are initiated based on one or more instructions embedded in the OCM **324**. When a processing engine, fetching data within the OCM **324** as part of a processing operation, reads an instruction among the one or more embedded instructions, the processing engine responds to the read instruction by starting a respective transport operation. The instructions are embedded, for example, by software executed by the host processor, **210**, **216**, **228**, **242**, **244**, such as the software compiler **404**.

The distinction between remote access request/response and processing thread migration is as follows: When a remote request is made, a processing engine is requesting and receiving the data (RCDS) that is on a remote memory cluster to the memory cluster where work is being executed by the processing engine. The same processing engine in a particular cluster executes both local data access and remote data access. For processing thread migrations, work is partially executed on a first memory cluster. The context, e.g., state and data, of the work is then saved, packaged and migrated to a second memory cluster where data (RCDS) to be accessed exists. A processing engine in the second memory cluster picks up the context and continues with the work execution.

FIG. **8**A shows a block diagram illustrating an overview of the XBC **530**, according to at least one example embodiment. The XBC **530** is an interface configured to manage transport operations between the corresponding memory, or search, cluster and one or more other memory, or search, clusters through the XBAR **312**. The XBC **530** includes a transmitting component **845** configured to manage transmitting transport operations from the processing engines **510** or the OCM **324** to other memory, or search, cluster(s) through the XBAR **312**. The XBC **530** also includes a receiving component **895** configured to manage receiving transport operations, from other memory, or search, cluster(s) through the XBAR **312**, and directing the transport operations to the processing engines **510** or the OCM **324**. The XBC **530** also includes a resource, or credit, state manager **850** configured to manage states of resources allocated to the corresponding memory cluster in other memory clusters. Such resources include, for example, memory buffers in the other memory clusters configured to store transport operations data sent from the memory cluster including the resource state manager **850**. The transmitting component **845** may be implemented as a logic circuit, processor, or the like. Similarly, the receiving component **895** may be implemented as a logic circuit, processor, or the like.

FIGS. **8**B and **8**C show logical diagrams illustrating an example implementation of the transmitting component **845**, of the XBC **530**, and the resource state manager **850**. The transmitting component **845** is coupled to the OCM **324** and the processing engines **510**, e.g., TWEs **512**, BWEs **514**, and RWEs **516** or RMEs **518**, as shown in the logical diagrams. Among the processing engines **510**, the TWEs **512** make remote tree access requests, the BWEs **514** make remote bucket access requests, and the RWEs **516** make remote rule access requests. The remote requests are stored in one or more first in first out (FIFO) buffers **834** and then pushed into per-destination FIFO buffers, **806***a* . . . **806***g*, to avoid

head-of-line blocking. The one or more FIFO buffers **834** may include, for example, a FIFO buffer **832** for storing tree access requests, FIFO buffer **834** for storing bucket access requests, FIFO buffer **836** for storing rule chunk access requests, and an arbitrator/selector **838** configured to select remote requests from the different FIFO buffers to be pushed into the per-destination FIFO buffers, **806***a*-**806***g*. Similarly, remote access responses received from the OCM **324** are stored in a respective FIFO buffer **840** and then pushed into a per-destination FIFO buffers, **809***a*-**809***g*, to avoid head-of-line blocking.

The remote requests for all three types of data, e.g., tree, bucket and rule chunk, are executable in a single clock cycle. The remote access responses may be variable length data and as such may be executed in one or more clock cycles. The size of the remote access response is determined by the corresponding remote request, e.g., the type of the corresponding remote request or the amount of data requested therein. Execution time of a transport operation, e.g., remote access request operation, remote access response, or processing thread migration operation, refers herein to the time duration, e.g., number of clock cycles, needed to transfer data associated with transport operation between a memory cluster and the XBAR **312**. With respect to a transport operation, a source memory cluster, herein, refers to the memory cluster sending the transport operation while the destination memory cluster refers to the memory cluster receiving the transport operation.

The TWEs **512** make tree processing thread migration requests, BWEs **514** make bucket processing thread migration requests. In the following, processing thread migration may be initiated either by TWEs **512** or BWEs **514**. However, according to other example embodiments the RWEs **516** may also initiate processing thread migrations. When TWEs **512** or BWEs **514** make processing thread migration requests, the contexts of the corresponding processing threads are stored in per-destination FIFO buffers, **803***a*-**803***g*. According to an example embodiment, destination decoders, **802**, **805**, and **808**, are configured to determine the destination memory cluster for processing thread migration requests, remote access requests, and remote access responses, respectively. Based on the determined destination memory cluster, data associated with the respective transport operation is then sent to a corresponding per-destination FIFO buffer, e.g., **803***a*-**803***g*, **806***a*-**806***g*, and **809***a*-**809***g*. The logic diagrams in FIGS. **8**B and **8**C assume a super cluster **310** including eight memory, or search, clusters **320**. As such, each transport operation in a particular memory cluster may be destined to at least one of seven memory clusters referred to in the FIGS. **8**B and **8**C with the letters a . . . g.

According to an example embodiment, a per-destination arbitrator, **810***a*-**810***g*, is used to select a transport operation associated with the same destination memory cluster. The selection may be made, for example, based on per-type priority information associated with the different types of transport operations. Alternatively, the selection may be made based on other criteria. For example, the selection may be performed based on a sequential alternation between the different types of transport operations so that transport operations of different types are treated equally. In another example embodiment, data associated with a transport operation initiated in a previous clock cycle may be given higher priority by the per-destination arbitrators, **810***a*-**810***g*. As shown in FIG. **8**C, each per-destination arbitrator, **810***a*-**810***g*, may include a type selector, **812***a*-**812***g*, a retriever, **814***a*-**814***g*, and a destination FIFO buffer, **816***a*-**816***g*. The type selector, **812***a*-**812***g*, selects a type of a transport operation and passes information indicative of selected type to the retriever, **814***a*-**814***g*, which retrieves the data at the head of a corresponding per-destination FIFO buffer, e.g., **803***a*-**803***g*, **806***a*-**806***g*, or **809***a*-**809***g*. The retrieved data is then stored in the destination FIFO buffer, **816***a*-**816***g*.

The transmitting component **845** also includes an arbitrator **820**. The arbitrator **820** is coupled to the resource state manager **850** and receives or checks information related to the states of resources, in destination memory clusters, allocated to the source memory cluster processing the transport operations to be transmitted. The arbitrator **820** is configured to select data associated with at least one transport operation, or transaction, among the data provided by the arbitrators, **810***a*-**810***g*, and schedule the at least one transport operation to be transported over the XBAR **312**. The selection is based at least in part on the information related to the states of resources and/or other information such as priority information. For example, resources in destination memory clusters allocated to the source memory cluster are associated with remote access requests and processing thread migrations but no resources are associated with remote access responses. In other words, for a remote access response a corresponding destination memory cluster is configured to receive the remote access response at any time regardless of other processes running in the destination memory cluster. For example, resources in the destination memory clusters allocated to the source memory cluster include buffering capacities for storing data associated with transport operations received at the destination memory clusters from the source memory cluster. As such no buffering capacities, at the destination memory clusters, are associated with remote access responses.

Priority information may also be employed by the arbitrator **820** in selecting transport operations or corresponding data to be delivered to respective destination memory clusters. Priority information, for example, may prioritize transport operations based on respective types. The arbitrator may also prioritize data associated with transport operations that were initiated at a previous clock cycle but are not completely executed. Specifically, data associated with a transport operation executable in multiple clock cycles and initiated in a previous clock cycle may be prioritized over data associated with transport operations to be initiated. According to at least one example embodiment, transport operations, or transactions, executable in multiple clock cycles are not required to be delivered in back to back clock cycles. Partial transport operations, or transactions, may be scheduled to be transmitted to effectively use the XBAR bandwidth. The arbitrator **820** may interleave partial transport operations, corresponding to different transport operations, over consecutive clock cycles. At the corresponding destination memory cluster, the transport operations, or transactions, are pulled from the XBAR **312** based on transaction type, transaction availability from various source ports to maximize the XBAR bandwidth.

The selection of transport operations, or partial transport operations, by the arbitrator **820** may also be based on XBAR resources associated with respective destination memory clusters. XBAR resources include, for example, buffering capacities to buffer data to be forwarded to respective destination memory clusters. As such, the resource state manager **850** in a first memory cluster keeps track of XBAR resources as well as the resources allocated to the first memory cluster in other memory clusters.

According to an example embodiment, the arbitrator **820** includes a destination selector **822** configured to select a

destination FIFO buffer, among the destination FIFO buffers **816***a*-**816***g*, from which data to be retrieved and forwarded, or scheduled to be forwarded, to the XBAR **312**. The destination selector passes information indicative of the selected destination to a retriever **824**. The retriever **824** is configured to retrieve transport operation data from the respective destination FIFO buffer, **814***a*-**814***g*, and forward the retrieved transport operation data to the XBAR **312**.

The resource state manager **850** includes, for example, a database **854** storing a data structure, e.g., a table, with information indicative of resources allocated to the source memory cluster in the other clusters. The data structure may also include information indicative of resources in the XBAR **312** associated with destination memory clusters. The resource state manager **850** also includes a resource state logic **858** configured to keep track and update state information indicative of available resources that may be used by the source memory cluster. In other words, the resource state logic **858** keeps track of free resources allocated to the sources memory cluster in other memory clusters as well as free resources in the XBAR **312** associated with the other memory clusters. Resource state information may be obtained by updating, e.g., incrementing or decrementing, the information indicative of resources allocated to the source memory cluster in the other clusters and the information indicative of resources in the XBAR **312** associated with destination memory clusters. Alternatively, state information may be stored in a separate data structure, e.g., another table. Updating the state information is, for example, based on information received from the other memory clusters, the XBAR **312**, or the arbitrator **820** indicating resources being consumed or freed in at least one destination resources or the XBAR **312**.

According to an example embodiment, a remote access request operation is executed in a single clock cycle as it involves transmitting a request message. A processing thread migration is typically executed in two or more clock cycles. A processing thread migration includes the transfer of data indicative of the context, e.g., state, of the search associated with the processing thread. A remote access response is executed in one or more clock cycle depending on the amount of data to be transferred to the destination memory cluster.

FIGS. **8**D and **8**E show logical diagrams illustrating an example implementation of the receiving component **895**, of the XBC **530**. According to at least one example embodiment, the receiving component **895**, e.g., in a first memory cluster, includes a type identification module **860**. The type identification module **860** receives information related to transport operations destined to the first memory cluster with data in the XBAR **312**. The received information, for example, includes indication of the respective types of the transport operations. According to the example implementation shown in FIG. **8**E, the type identification module **860** includes a source decoder **862** configured to forward the received information, e.g., transport operation type information, to per-source FIFO buffers **865***a*-**865***g* also included in the type identification module **860**. For example, received information associated with a given source memory cluster is forwarded to a corresponding per-source memory FIFO buffer. An arbitrator **870** then acquires the information stored in the per-source FIFO buffers, **865***a*-**865***g*, and selects at least one transport operation for which data is to be retrieved from the XBAR **312**. Data corresponding to the selected transport operation is then retrieved from the XBAR **312**.

If the selected transport operation is a remote access request, the retrieved data is stored in the corresponding FIFO buffer **886** and handed off to the OCM **324** to get the data. That data is sent back as remote response to the requesting source memory cluster. If the selected transport operation is a processing thread migration, the retrieved data is stored in one of the corresponding FIFO buffers **882** or **884**, to be forwarded later to a respective processing engine **510**. The FIFO buffers **882** or **884** may be a unified buffer managed as two separate buffers enabling efficient management of cases where processing thread migrations of one type are more than processing thread migrations of another type, e.g., more tree processing thread migrations than bucket processing thread migrations. When a processing engine handling processing thread migration of some type, e.g., TMIG or BMIG, becomes available respective processing thread migration context, or data, is pulled from the unified buffer and sent to the processing engine for the work to continue in this first memory cluster. According to at least one example embodiment, one or more processing engines in a memory cluster receiving migration work are reserved to process received migrated processing threads. When a remote access response operation is selected, the corresponding data retrieved from the XBAR **312** is forwarded directly to a respective processing engine **510**. Upon forwarding the retrieved data to the OCM or a processing engine **510**, an indication is sent to the resource state manager **850** to cause updating of corresponding resource state(s).

In the example implementation shown in FIG. **8**E, the arbitrator **870** includes first selectors **871**-**873** configured to select a transport operation among each type and a second selector **875** configured to select a transport operation among the transport operations of different types provided by the first selectors **871**-**873**. The second selector **875** sends indication of the selected transport operation to the logic operators **876***a*-**876***c*, which in turn pass only data associated with the selected transport operation. The example receiving component **895** shown in FIG. **8**D also includes a logic operator, or type decoder, **883** configured direct processing thread migration data to separate buffers, e.g., **882** and **884**, based on processing thread type, e.g., tree or bucket. Upon forwarding a transport operation to the OCM **324** or a respective processing engine **510**, a signal is sent to a resource return logic **852**. The resource return logic **852** is part of the resource state manager **850** and is configured to cause updating of resource state information.

FIG. **9**A is a block diagram illustrating an example implementation of the XBAR **312**. A person skilled in the art should appreciate that the XBAR **312** as described herein is an example of an interface device coupling a plurality of memory clusters. In general, different interface devices may be used. The example implementation shown in FIG. **9**A is an eight port fully-buffered XBAR that is constructed out of modular slices **950***a*-**950***d*. For example, the memory clusters are arranged in two rows, e.g., north memory clusters, **320***a*, **320***c*, **320***e*, and **320***g*, are indexed with even numbers and south memory clusters, **320***b*, **320***d*, **320***f*, and **320***h*, are indexed with odd numbers. The XBAR **312** is constructed to connect these clusters. To match the cluster topology, the example XBAR **312** in FIG. **9**A is built as a 2×4 (8-port) XBAR **312**. Each slice connects a pair of North-South memory clusters to each other and to its neighboring slice(s).

FIG. **9**B is a block diagram illustrating implementation of two slices, **950***a* and **950***b*, of the XBAR **312**. Each slice is built using half-slivers **910** and full-slivers **920**. The half-slivers **910** and the full-slivers **920** are, for example, logic circuits used in coupling memory clusters to each other. For an N-port XBAR **312**, each slice contains N−2 full-slivers

920 and 2 half-slivers 910. The full-slivers 920 correspond to memory cluster ports that are used to couple memory clusters 320 belonging to distinct slices 950. For the slice 950*a*, for example, full-slivers 920 correspond to ports 930*c* to 930*h* which couple memory clusters in the slice 950*a* to the memory clusters 950*b*-950*d*, respectively, in other slices 950. For the memory cluster ports coupling memory cluster within the same slice, the slivers are optimized to half-slivers 910. For the slice 950*a*, for example, half-slivers correspond to ports 930*a* and 930*b*.

FIG. 9C shows an example logic circuit implementation of a full-sliver 920. The full-sliver 920 contains two FIFO buffers, 925*a* and 925*b*, for storing data from other ports through a neighboring slice. One FIFO buffer, e.g., 925*a*, is for storing data destined to the north memory cluster and one FIFO buffer, e.g., 925*b*, is for storing the data destined to the south memory cluster. The control (GRQs) signals 922*a* and 922*b* identify which port the data is destined to. The data (GRFs) 921 is pushed into the appropriate full-sliver FIFO 925*a* or 925*b*. For example, when data from the memory cluster_320*c* is destined to the memory cluster_320*b*, GRQ2 and GRF2 will signal to the south FIFO buffer 925*b* of the full sliver SLV2 in slice 950*a* to capture and keep the data until it is demanded by the memory cluster 320*b*. Continuing with the same example, if data was destined to the memory cluster_320*a*, GRQ2 will signal the north FIFO buffer 925*a* the full sliver SLV2 in slice 950*a* to capture and keep the data until demanded by the memory cluster_320*a*.

FIG. 9D shows an example logic circuit implementation of a half-sliver 910. Each half-sliver 910 contains one FIFO buffer 925 for storing data from one of two memory clusters within a given slice. The data in each half-sliver 910 is meant for the opposite memory cluster in the same slice. For example, in the slice 950*a*, the half-sliver HSLV0 gets data (GRF0) from the memory cluster_320*a* and is destined to the memory cluster_950*b*.

When a memory cluster decides to fetch the data from a particular FIFO buffer, e.g., 925, 925*a*, or 925*b*, it sends a pop signal, 917, 927*a*, or 927*b*, to that FIFO buffer. When the FIFO buffer, e.g., 925, 925*a*, or 925*b*, is not selected by the memory cluster the logic AND operator 914, 924*a*, or 924*b*, outputs zeros. An OR operator, e.g., 916, 926*a*, or 926*b*, in each sliver is applied to the data resulting in a chain of OR operators either going north or going south. According to an example embodiment, one clock cycle delay between pop signal and data availability at the memory cluster that's pulling the data.

The XBAR 312 is the backbone for transporting various transport transactions, or operations, such as remote requests, remote responses, and processing thread migrations. The XBAR 312 provides the transport infrastructure, or interface. According to at least one example embodiment, transaction scheduling, arbitration and flow control is handled by the XBC 530. In any given clock cycle multiple pairs of memory clusters may communicate. For example, the memory cluster 320*a* communicates with the memory cluster 320*b*, the memory cluster 320*f* communicates to the memory cluster 320*c*, etc. The transfer time for transferring a transport operation, or a partial transport operation, from a first memory cluster to a second memory cluster is fixed with no queuing delays in the XBAR 312 or any of the XBCs of the first and second memory clusters. However, in the case of queuing delays, the transfer time, or latency, depends on other transport operations, or partial transport operations, in the queue and the arbitration process.

Resources are measured in units, e.g., "credits." For example, a resource in a first memory cluster, e.g., destination memory cluster, allocated to a second memory cluster, e.g. source memory cluster, represented by one credit corresponds to one slot in a respective buffer, e.g., 882, 884, or 886. According to another example, one credit may represent storage capacity equivalent to the amount of data transferrable in a single clock cycle. XBAR resources refer, for example, to storage capacity of FIFO buffers, e.g., 915, 925*a*, 925*b*, in the XBAR. In yet another example, one credit corresponds to storage capacity for storing a migration packet or message.

The XBAR 312 carries single- and multi-cycle packets, and/or messages, from one cluster to another over, for example, a 128 bit crossbar. These packets, and/or messages, are for either remote OCM access or processing thread migration. Remote OCM access occurs when a processing thread, e.g., the TWE and/or BWE, on one cluster encounters Rule Compiled Data Structure (RCDS) image data that redirects a next request to a different memory cluster within the same super-cluster. Processing thread migration occurs for two forms of migration, namely, a) Tree-Walk migration and b) Bucket-Walk migration. In either case, the processing thread context, e.g., details of the work done so far and where to start working, for the migrated thread is transferred to a different memory cluster within the same super-cluster, which continues processing for the thread.

FIGS. 10A and 10B show two example tables storing resource state information in terms of credits. The stored state information is employed in controlling the flow of transport transactions. Both tables, in FIGS. 10A and 10B, illustrate two examples of resource credits allocated to the memory cluster indexed with 0 in the memory clusters indexed with 1 through 7. In FIG. 10A, the first column shows unified migration, the second column shows remote request credits and the third column shows XBAR credits allocated to the memory cluster indexed with 0. In FIG. 10B, the migration credits are separated based on the type of processing thread migration, e.g., tree processing thread migration and bucket processing thread migration. Migration credits track the migration buffer(s) availability at a particular destination. Remote request credits track the remote request buffer(s) availability at the destination. XBAR credit tracks the resources inside the XBAR to a particular destination. There are no separate credits for responses. The response space is pre-allocated in the respective engine.

When a remote access request is sent from a first memory cluster, e.g., a source cluster, to a second memory cluster, e.g., destination cluster, the resource state manager 850 of the first memory cluster decrements, e.g., by a credit, the credits defining the remote request resources allocated to the first memory cluster in the second memory cluster. The resource state manager 850 may also decrement, e.g., by a credit, the credits defining the state of XBAR resources associated with the second memory cluster and allocated to the first memory cluster. When the remote access request is passed from the XBAR 312 to the destination memory cluster, the resource state manager 850, at the source memory cluster, receives a signal from the XBAR 312 indicating the resource represented by the decremented credit is now free. The resource state manager 850, in the first cluster, then increments the state of the XBAR resources associated with the second memory cluster and allocated to the first memory cluster by a credit. When the corresponding remote access response is received from the second memory cluster and is passed to a corresponding engine in the first cluster, a signal is sent to the resource return logic 852 which

in turn increments, e.g., by a credit, the state of resources allocated to the first memory cluster in the second memory cluster.

When a processing thread is migrated from the first memory cluster to the second memory cluster, the resource state manger **850** of the first memory cluster decrements, e.g., by a credit, the credits defining the migration resources allocated to the first memory cluster in the second memory cluster. The resource state manager **850** of the first memory cluster may also decrement, e.g., by a credit, the credits defining the state of XBAR resources associated with the second memory cluster and allocated to the first memory cluster. When the migrated processing thread is passed from the XBAR **312** to the destination memory cluster, the resource state manager **850** of the first memory cluster receives a signal from the XBAR **312** indicating the resource represented by the decremented credit is now free. The resource state manager **850**, in the first memory cluster, then increments the state of the XBAR resources associated with the second memory cluster and allocated to the first memory cluster by a credit. When the migrated processing thread is passed to a corresponding engine, a signal is sent to the resource return logic **852** of the second memory cluster, which in turn forwards the signal to the resource state manager **850** of the first memory cluster. The resource state manager **850** of the first memory cluster then increments, e.g., by a credit, the migration resources allocated to the first memory cluster in the second memory cluster. Decrementing or incrementing migrations credits may be performed based on the type of processing thread being migrated, e.g., tree processing thread or bucket processing thread, as shown in FIG. **10**B.

FIGS. **11**A to **11**C illustrate examples of interleaving transport operations and partial transport operations over consecutive clock cycles. In FIG. **11**A, a processing thread migration is executed in at least four non-consecutive clock cycles with remote access requests executed in between. Specifically, the processing thread migration is executed over the clock cycles indexed with 0, 2, 3 and 5 while two remote access requests are executed, respectively, over the clock cycles indexed with 1 and 4. The interleaved transport operations in FIG. **11**A are executed by a source memory cluster destined to the same memory cluster. FIG. **11**B shows an example of interleaving transport operations executed by a source memory cluster destined to two destination memory clusters, e.g., indexed with 0 and 1. FIG. **11**C shows an example of interleaving transport operations and partial transport operations executed by a destination memory cluster. Specifically, a remote access response, received from the memory cluster indexed with 0, is executed over the clock cycles indexed with 0, 1, and 3, while two remote access requests destined to two distinct memory clusters over the clock cycles indexed with 2 and 4.

FIG. **12**A shows a flowchart illustrating a method of managing transport operations between a source memory cluster and one or more other memory clusters performed by the XBC **530**. Specifically the method is performed by the XBC in a source memory cluster. At block **1210** at least one transport operation from one or more transport operations is selected, at a clock cycle in the source memory cluster, the at least one transport operation is destined to at least one destination memory cluster based at least in part on priority information associated with the one or more transport operations or current states of available processing resources allocated to the source memory cluster in each of a subset of the one or more other clusters. At block **1220**, the transport of the selected at least one transport operation is initiated. The one or more transport operations are received from processing engines **510** and/or OCM **324**. The method may be implemented through an implementation of the XBC as shown in FIGS. **8**B and **8**C. However, a person skilled in the art should appreciate that the method may be implemented a different implementation of the XBC. For example, the priority information may be based on the type, latency, or destination, of the one or more transport operations. The selection may further be based on XBAR resources associated with the destination memory cluster.

FIG. **12**B shows a flowchart illustrating another method of managing transport operations between a destination memory cluster and one or more other memory clusters performed by the XBC **530**. Specifically the method is performed by the XBC in a destination memory cluster. At block **1260**, information related to one or more transport operations with related data buffered in an interface device is received, in the source memory cluster, the interface device coupling the destination memory cluster to the one or more other memory clusters. At block **1270**, at least one transport operation, from the one or more transport operations, is selected to be transported to the destination memory cluster based at least in part on the received information. At block **1280** the transport of the selected at least one transport operation is initiated.

According to at least one example embodiment, resource credits are assigned to memory clusters by software of the host processor, e.g., **204**, **214**, **228**, **242**, or **244**. The software may be, for example, the software compiler **404**. The assignment resource credits may be performed, for example, when the search processor **202** is activated or reset. The assignment of the resource credits may be based on the type of data stored in each memory cluster, the expected frequency of accessing the stored data in each memory cluster, or the like.

FIG. **13** shows a flowchart illustrating a method of assigning resources used in managing transport operations between a first memory cluster and one or more other memory clusters. At block **1310**, information indicative of allocation of a subset of processing resources in each of the one or more other memory clusters to the first memory cluster is received, for example, by the resource state manager **850** of the first memory cluster. At block **1320**, information indicative of resources allocated to the first cluster is stored in the first memory cluster, specifically in the respective resource state manager **850**. The allocated processing resources may be stored as credits. The processing resources may be allocated per type of transport operations as previously shown in FIGS. **10**A and **10**B. The allocated processing resources may be stored in the form of a table or any other data structure. At block **1330**, the information indicative of resources allocated to the first memory cluster, stored in the resource state manager **850**, is then used to facilitate managing of transport operations between the first memory cluster and the one or more other memory clusters. For example, the stored information is used as resource state information indicative of availability of the allocated processing resources to the first memory cluster and is provided to the arbitrator **820** to manage transport operations between the first memory cluster and the one or more other memory clusters. The resource state information is updated in real time, as described above, to reflect which among the processing resources are free and which are in use. The processing resources represent, for example, buffering capacities in the each memory cluster, and as such the sum of processing resources in a given memory cluster allocated to

other memory clusters is equal to or less than the total number of respective processing resources of the given memory cluster.

The host processor, e.g., **204**, **214**, **228**, **242**, or **244**, may modify allocation of processing resources to the first memory cluster on the fly. For example, the host processor may increase or decrease the processing resources, or number of credits, allocated to the first memory cluster in a second memory cluster. In reducing processing resources, e.g., number of migration resources, allocated to the first memory cluster in the second memory cluster, the host processor indicates to the search processor a new value of processing resources, e.g., number of credits, to be allocated to the first memory cluster in the second memory cluster. The search processor determines, based on the state information, whether a number of free processing resources allocated to the first memory cluster in the second memory cluster is less than a number of processing resources to be reduced. Specifically, such determination may be performed by the resource state manager **850** in the first memory cluster. For example, let 5 credits be allocated to the first memory cluster in the second memory cluster, and the host processor, e.g., **204**, **214**, **228**, **242**, or **244**, decides to reduce the allocated credits by 3 so that the new allocated credits would be 2. The host processor sends the new credits value, e.g., 2, to the search processor **202**. The resource state manager **850** in the first memory cluster checks whether the current number of free credits, e.g., m, that are allocated to the first memory cluster from the second memory cluster is less than the number of credits to be reduce, e.g., 3. Upon determining that the number of free processing resources, e.g., m, is less than the number of processing resources to be reduced, e.g., 3, The XBC **530** in the first memory cluster blocks, initiation of new transport operations between the first memory cluster and the second memory cluster until the number of free processing resources, e.g., m, allocated to the first memory cluster in the second memory cluster is equal to or greater than the number of resource to be reduced. That is, the transfer of transport operations between the first and second memory clusters are blocked until, for example, m≥3. According to one example, only initiation of transport of new transport operations is blocked. According to another example, initiation of transport of new transport operations and partial transport operation is blocked. Once the number of free processing resources, e.g., m, allocated to the first memory cluster in the second memory cluster is equal to or greater than the number of resource to be reduced, the information indicative of allocated processing resources is updated, for example, by the resource state manager **850** in the first memory cluster to reflect the reduction, e.g., changed from 5 to 2. In another example, the checking may be omitted and the blocking of transport operations and partial transport operations may be applied until all allocated credits are free and then the modification is applied.

In increasing the number of processing resources allocated to the first memory cluster from the second memory cluster, the host processor determines whether a number of non-allocated processing resources, in the second memory cluster, is larger than or equal to a number of processing resources to be increased. For example if the number of allocated processing resources is to be increased from 5 to 8 in the first memory cluster, the number of non-allocated resources in the second memory cluster is compared to 3, i.e., 8-5. Upon determining that the number of non-allocated processing resources, in the second memory cluster, is larger than or equal to the number of processing resources to be increased, the host processor sends information, to the search processor **202**, indicative of changes to be made to processing resources allocated to the first memory cluster from the second memory cluster. Upon the information being received by the search processor **202**, the resource state manager **850** in the first memory cluster modifies the information indicative of allocated processing resources to reflect the increase in processing resources, in the second memory cluster, allocated to the first memory cluster. The resource state manager then uses the updated information to facilitate management of transport operations between the first memory cluster and the second memory cluster. According to another example, the XBC **530** of the first memory cluster may apply blocking of transport operations and partial transport operations both when increasing or decreasing allocated processing resources.

FIG. **14** shows a flow diagram illustrating a deadlock scenario in processing thread migrations between two memory clusters. Assume two migration credits are allocated to memory cluster **320***a* from memory cluster **320***b* and two migration credits are allocated to the memory cluster **320***b* from the memory cluster **320***a*. Also assume that a single processing engine is handling migration work in each of the memory clusters **320***a* and **320***b*. Two processing threads, **1410** and **1420**, are migrated from the memory cluster **320***a* to **320***b* and two other processing threads, **1415** and **1425**, are migrated from the memory cluster **320***b* to **320***a*. The processing threads **1410** and **1420** want to migrate back to the memory cluster **320***a*, while the processing threads **1415** and **1425** want to migrate back to the memory cluster **320***b*. Also the processing thread **1430** wants to migrate to the memory cluster **320***b* and the processing thread **1435** wants to migrate to the memory cluster **320***a*. However each memory cluster, **320***a* or **320***b*, can handle a maximum of three processing threads at any point in time, e.g., one by the processing engine and two in the buffers indicated by the credits. Given that there are three processing threads in each memory cluster, none of the processing threads, **1410**, **1415**, **1420**, **1425**, **1430**, or **1435**, can migrate. As such, a deadlock occurs with none of the migration works proceeding. The deadlock is mainly caused by allowing migration loops where a processing may migrate back to memory cluster that it migrated from previously.

According to an example embodiment, the deadlock may be avoided by limiting the number of processing threads, of a given type, being handled by a super cluster at any given point of time. Regardless of the number of memory clusters, e.g., N, in a super cluster, if a processing thread may migrate to any memory cluster in the super cluster, or in group of memory clusters, then there is a possibility that all processing threads in the super cluster may end up in two memory clusters of the super cluster, that is similar to the case of FIG. **14**. Consider that each memory cluster has k processing engines for processing migration work of the given type and that each destination memory cluster has M migration credits, e.g., for migration work of the given type, to be distributed among N−1 memory clusters. The maximum number of processing threads, of a given type, that may be handled by the super cluster without potential deadlock is defined as:

$$W_{max} = \left(\mathrm{Int}\left(\frac{M}{N-1}\right) + k\right) * 2 - 1,$$

where "Int" is a function providing the integer part of a number.

For example, let the number of processing engines for processing migration work of the given type per memory cluster be k=16. Let the number of total credits, for migration of the given type, in any destination memory cluster be M=15 and the total number of memory cluster in the super cluster, or the group of clusters, be N=4. As such the number of migration credits allocated to any source memory cluster in any destination memory cluster is 15 divided by (4-1), which is equal to 5. According to the equation above, the maximum number of processing threads that the super cluster may handle is 41. Applying the example of processing thread ending up distributed between only two memory clusters as in FIG. **14**, then a first memory cluster, having 16 processing engines and 5 migration credits, may end up with 21 processing threads. That is, the 16 processing engines and the buffering capacity represented by the 5 credits are being consumed. A second memory cluster, having 16 processing engines and 5 migration credits, then ends up with the 20 other processing threads. As such a processing thread may migrate from the first memory cluster to the second memory. Given that any processing thread may either finish processing completely, migrate, or transforms into a different type of processing thread, e.g., from tree processing thread to bucket processing thread, then at a given point of time a processing thread in the first memory cluster would either transform into a processing thread of different type, finish processing completely and vanish, or migrate to the second memory cluster. In each of these cases it would become possible for a processing thread in the second memory cluster to migrate to the first memory cluster. Therefore, with such deadlock is avoided. However, if the total number of migration thread is more than the maximum indicated by the equation above, a potential deadlock may occur if the total processing threads end up being distributed between two clusters with all the engines and the migration credits therein being consumed.

FIG. **15** shows graphical illustration of another approach to avoid deadlock. The idea behind approach to avoid deadlocks is to prevent any migrations loops where a migrating processing thread may migrate to a memory cluster from which it previously migrated. In the example shown in FIG. **15**, migration of four different processing threads, **1510**, **1520**, **1530**, and **1540**, across the memory clusters **324***a*-**324***d* are illustrated, with the memory cluster **324***a* assigned as a drain, or sink, memory cluster. A sink, or drain, memory cluster prevents a processing thread that migrated to it from another memory cluster to migrate out. In addition, a processing thread that migrated to a particular memory cluster may not migrate to another memory cluster from which other processing threads, e.g., of the same type, migrate to the particular memory cluster and therefore preventing migration loops. In other words, migrated processing threads may migrate to a sink memory cluster or memory cluster in a path to a sink memory cluster. A path to a sink memory cluster may not have structural migration loops. As illustrated in FIG. **15**, such design, of migrations, prevent structural migration loops from occurring.

Contrary to migrated work, new work that originated in a particular memory cluster but did not migrate yet, may migrate to any other memory cluster even if the particular memory cluster is a sink memory cluster. Further, at least one processing engine is reserved to handle migration work in memory clusters receiving migration work.

FIG. **16** is a flowchart illustrating a method of managing processing thread migrations within a plurality of memory clusters. According to at least one example embodiment, instructions indicative of processing thread migrations are embedded at block **1610**, in memory components of the plurality of memory clusters. Such instructions are, for example, received from the host processor in the search processor **202** and embedded by the latter in memory components of the plurality of memory clusters of the same search processor. When a processing thread, fetching data in the OCM of a first memory cluster encounters one of such instructions, the corresponding processing engine make a migration request to migrate to a second memory cluster indicated in the encountered instruction. At block **1620**, data, configured to designate a particular memory cluster as a sink memory cluster, is stored in one or more memory components of the particular memory cluster. The particular memory cluster is one among the plurality of memory clusters of the search processor **202**.

A sink memory cluster may be designed, for example, through the way the data to be fetched by processing engines is stored across different memory clusters and by not embedding any migration instructions in any of the memory components of the sink memory cluster. In other words, by distributing data to be fetched in a proper way between the different memory clusters, a sink memory cluster stores all the data that is to be accessed by a processing thread that migrated to the sink memory cluster. Alternatively, if some data, that is to be accessed by a processing thread that migrated to the sink memory cluster, is not stored in the sink memory cluster, then such data is accessed from another memory cluster through remote access, but no migration is instructed. In another example, the data stored in the sink memory cluster is arranged to be classified into two parts. A first part of the data stored is to be searched or fetched only by processing threads originating in the sink memory cluster. A second part of the data is to be searched or fetched by processing threads migrating to the sink memory cluster from other memory clusters. As such, the first part of the data may have migration instructions embedded therein, while the second part of the data does not include any migration instructions. At block **1630**, one or more processing threads executing in one or more of the plurality of memory clusters, are processed, for example, by corresponding processing engines, in accordance with at least one of the embedded migration instructions and the data stored in the sink memory cluster. For example, if the processing thread encounters migration instruction(s) then it is caused to migrate to another memory cluster according to the encountered instruction(s). Also if the processing thread migrates to the sink memory cluster, then the processing thread does migrate out of the sink memory cluster.

According to at least one aspect, migrating processing threads include at least one tree search thread or at least one bucket search thread. With regard to the instructions indicative of processing thread migrations, such instructions are embedded in a way that would cause migrated processing threads to migrate to a sink memory cluster or to a memory cluster in the path to a sink memory cluster. A path to a sink memory cluster is a sequence of memory clusters representing a migration flow path and ending with the sink memory cluster. The embedded instructions are also embedded in a way to prevent migration of a processing thread to a memory cluster from which the processing thread migrated previously. The instructions may further be designed to prevent a migrating processing thread arriving to a first memory cluster to migrate to a second memory cluster from which other migration threads migrate to the first memory cluster.

A person skilled in the art should appreciate that the RCDS **410**, shown in FIG. **4**, may be arranged according to another example of nested data structures. As such the processing engines **510** are defined in accordance with respective fetched data structures. For example, if the nested data structures include a table, a processing engine may defined as, for example, table fetching engine or table walk engine. Processing engines **510**, according to at least one example, refer to separate hardware processors such as single-core processors or specialized processors included in the XBC **530**. Alternatively, processing engines **510** may be functions performed by one or more hardware processors included in the XBC **530**.

FIG. **17** is a block diagram **1700** illustrating an example embodiment of a search processor manager **1704** employed with a search processor **1714**. The search processor **1714** is coupled with a control plane processor **1702**, in a control plane **1708**, and also with a data plane processor **1712**, in a data plane **1716**. The search processor **1714** is also in the data plane **1716**. The search processor **1714** is coupled to the control plane processor **1702** by an I^2C/PCI-E **1706** connection, and also to the data plane processor **1712** over an Interlaken **1718** connection. The search processor manager **1704** is a module within the control plane processor **1702** in the control plane **1708**. The data plane processor **1712** is connected to the control plane processor **1702** over a PCI-E connection **1710**.

The data plane processor **1712** receives packets **1722** over a receive line **1724**. The data plane processor communicates with the control plane processor **1702** and the search processor **1714** to determine how to process the received packets. The control plane processor **1702** and search processor manager **1704** assist the search processor **1714**. The search processor manager **1704** provides an instant view of the storage within the memory (e.g., clusters) of the search processor **1714**. The search processor manager **1704** also provides a knowledge base, or guidelines, on processing data. For example, the search processor manager **1704** organizes the memory of the search processor **1714** to avoid migration deadlock of the processors (e.g., TWE, BWE, and RME) within the cluster of the search processor **1714**. The search processor manager **1704** can also organize the memory to avoid remote accesses of other clusters within the search processor **1714**. Further, the search processor manager **1704** can help better and more efficiently utilize hardware resources of the search processor **1714**, such as the memory within, and the processing resources, such as tree walk engines, bucket walk engines, and rule match engines. The search processor manager **1704** can also help avoid bank conflicts in the memory of the search processor **1714**. The search processor manager **1704** can also help avoid rule replication in the memory of the search processor **1714**.

In one embodiment, the search processor manager **1704** organizes the memory by using a database packer. The database packer can include a tree packer (e.g., tree packing module), a bucket packer (e.g., a bucket packing module) and a rule packer (e.g., a rule packing module). While the database packer **1704** can reside in the search process manager **1704**, in other embodiments it could reside in the search processor **1714** or within another processor or module.

The tree packer is configured to pack the nodes of the tree in each of the clusters so that, when the search processor walks the tree in its memory, it does not perform a migration or a remote cluster access, or if it does, it performs as few migrations and as few remote accesses as possible. The tree packer is also configured to prevent the search processor from performing a migration loop during a tree walk, that is, migrating from a first cluster to any number of clusters before migrating back to the first cluster.

The bucket packer is configured to pack a particular number of bucket chunks per bucket and a particular number of rule pointers per bucket chunk based on addresses of the rules. The bucket packer can also allocate the rules in the memory in a same order as an order of allocation of the buckets. The bucket packer can also replicate a rule (or a chunk of rules) across multiple banks, such as a first and second bank, in the clusters to improve the performance of the search processor when the search processor runs. The pre-runtime replication of rules by the bucket packer prevents bank conflict(s) by allowing the search processor to access the rule or a chunk of rules on the second bank when the first bank has a memory access conflict during a particular clock cycle. The bucket packer can also distribute rules across a first cluster and a second cluster within the memory to enable processing of the rules by a rule match engine within the first cluster and a rule match engine within the second cluster. The bucket packer can also store a first rule of a particular bucket of the plurality of buckets in a first cluster of the plurality of clusters, and then store any other rules of the particular bucket in a second cluster. The bucket packer can also allocate a bucket in the same memory cluster as the node of the tree which points to it.

The rule packer can be configured to allocate rules in a particular cluster of the memory. Then, it can allocate a rule associated with the bucket in the particular cluster storing the bucket. The rule packer can further allocate the rules in chunks according to an order of the buckets, determine a need to replicate the rules across the plurality of banks/clusters, and replicating the rules across the plurality of banks/clusters, if necessary.

FIG. **18** is a flow diagram **1800** illustrating an example embodiment of a control plane stack. The control plane stack has a processing flow **1802** flowing from an application layer **1804** to an incremental update layer **1806** to a search processor manager **1808** to a driver layer **1810**. The application layer can create incremental updates **1806**. When an incremental update **1806** needs to be applied to the search processor, the search processor manager **1808** can process the update and apply it to the search processor in a more efficient manner. Upon determining the most efficient manner to apply the update, based on hardware resources of the search processor, the control plane sends the instructions to the drivers **1810**.

FIG. **19** is a block diagram **1900** illustrating an example embodiment of incremental update as managed by a search processor manager **1902** for a search processor **1912**. The search processor manager **1902** includes a search processor manager front end **1904** and a search processor manager back end **1906**. The search processor manager front end **1904** interfaces with the application layer to create a new set of rules. The search processor manager back end **1906** receives the set of new rules and interfaces with a shadow image of the search processor memory **1908**. The shadow image of the search processor memory **1908** is a copy of the memory that is on the search processor **1914**. The shadow image **1908** has the same format as the search processor memory **1914**, such as the number of lines the memory has and the width of each line. In addition, the banks and the format of clusters within the shadow image **1908** are identical so that the search processor manager can detect any errors that may occur from the memory being divided into clusters or different banks.

The search processor manager **1902** first creates a full set of memory instructions to write the incremental update to the memory. The series of writes to the memory are typically a series of writes to individual lines within the memory. A given line within the memory can be written to multiple times. Many of the sequential writes to the same line can be either independent of each other or dependent of each other. If a write is independent of another write, the search processor manager **1902** can merge the two writes into one write command. However, if the writes are dependent on each other, the search processor manager cannot merge the writes after the dependency. However, merging independent writes up to any dependency can reduce the total number of writes and improve the speed of writing the incremental update to the memory. Further, independent writes can be issued sequentially, one after another, without waiting for previous writes to complete. Once the independent writes are complete, the dependent writes are issued.

Further, the search processor manager **1902** can determine whether the incremental update has any errors being written to memory, by including a shadow image **1908** and initially writing to the shadow image **1908**. Such an error could be running out of memory in a particular cluster or in memory in general. If there is such an error, the search processor manager can revert the shadow image **1908** to the original search processor memory **1914**, (or a copy thereof). This way, the search processor **1912** and search processor memory **1914** remain undisturbed and receive no error. Further, the search processor **1912** does not employ a procedure to restore the search processor memory **1914**.

In addition, the search processor manager writes data to the clusters in a manner that is optimized for the fact that during each clock cycle, no two ports can access the same bank. Such optimization guarantees the atomicity of write operations. Therefore, the search processor manager writes the rules to the memory first, then the buckets, and then the tree during an incremental update. The search processor manager can write buckets and rules simultaneously if the bucket that points to the rule is stored on the same line as the rule.

FIG. **20**A is a block diagram **2000** illustrating an example embodiment of a tree **2002**, a bucket **2006** and rules **2008**, **2010** and **2012**. The database, including the tree **2002**, bucket **2006**, and rules **2008**, **2010**, and **2012**, often cannot be stored in a single cluster of memory of the search processor. The database therefore should be divided into multiple clusters and memory banks. The search processor manager determines how to divide the database into these multiple clusters. In the embodiment shown in the block diagram **2000**, the tree **2002** is divided by a horizontal memory boundary **2014**. The search processor manager therefore directs all of the tree nodes above the memory boundary **2014** to be stored in a first cluster, and the remainder of the tree, the bucket **2006** and the rules **2008**, **2010** and **2012** to be stored in a second cluster. More than two clusters can be employed in the search processor, and the database can be stored in any number of clusters. Two clusters are used in this example embodiment for simplicity.

FIG. **20**B is a block diagram **2050** illustrating an example embodiment of storing the tree **2002**, buckets **2006** and **2054** and rules **2008**, **2010**, **2012**, **2056** and **2058** in multiple clusters in the memory. In FIG. **20**A, one cluster had only tree nodes, while a second cluster had tree nodes, a bucket, and rules. Such a setup can reduce the efficiency of hardware resources of the first cluster because the bucket walk engines and rule match engines of the first cluster either are not used or have to remotely access another cluster, which is a costly operation. In FIG. **20**B, the memory boundary **2052** is "vertical." The tree **2002**, divided by the memory boundary **2052** is stored in both a first cluster and a second cluster, with the first cluster being on the left, and the second cluster being on the right. However, because the memory boundary **2052** is vertical, the first cluster includes a bucket **2006**, and rules **2008**, **2010**, and **2012**. Similarly, the second cluster includes the right side of the tree **2002**, and a bucket **2054**, and corresponding rules **2056** and **2058**. This setup allows the tree walk engines, bucket walk engines, and rule match engines of both the first and second cluster to be utilized without having to remotely read from another cluster. This maximizes the hardware resources of each cluster. The search processor manager is configured to store the database in the clusters in a manner such that the resources of each cluster can be maximized for each clock cycle in this manner.

As described herein, the tree can be stored across a plurality of clusters. In addition, each cluster can store bucket data and rule data. Further, the search processor manager considers distributing data across multiple banks within a cluster. Such a distribution can improve efficiency because each bank can be read by a separate port. Ideally, the search processor manager distributes data such that each port is performing a read from its respective bank on every clock cycle to maximize the throughput of the memory. For this reason, rules can be replicated across multiple banks, so that if there is a bank conflict for a particular bank, another bank may store the same rule and be able to service the request.

For example, consider the example in FIG. **25**, which is a diagram **2500** illustrating an example embodiment of memory replication using a Bucket B_1 **2502** and a Bucket B_2 **2504**. Each Bucket B_1 **2502** and B_2 2504 points to respective bucket chunks **2506** and **2508**. Each bucket chunk **2506** and **2508** includes bucket chunks storing rules, which are to be stored in memory banks **2510** (b_1-b_8). For example, a first bucket chunk of a first bucket (Bucket B_1 **2502**) may include Rules 1, 2, and 3 and is stored in memory bank **2510** b_1. Other bucket chunks **2506** store rules 4, 5, and 6 (to be stored in memory bank b_2), rules 7 and 8 (to be stored in memory bank b_3) and rule 9 (to be stored in memory bank b_4). Bucket chunks **2508** of the Bucket B_2 **2504** may also have some rules which are also part of the Bucket B_1 **2502**. For example, a second bucket chunk of a second bucket (Bucket B_2 **2504**) may include rules 1 and 2, which are already stored as part of Bucket B_1 in memory bank **2510** b_1. A scheme to save memory would only store Rules 1, 2, and 3 with the first bucket chunk of Bucket Chunks **2506**, and would not store a separate copy of the second bucket chunk of Bucket Chunks **2508**. However, this can lead to memory conflicts while processing bucket chunks of Bucket B_2, **2504**. As shown in FIG. **25**, if the first bucket chunk (containing rules 10 & 11) of Bucket B_2, **2504** is stored in memory bank **2510** b_1, rules 1 & 2 of second bucket chunk of bucket B_2, **2504** may also need to be accessed in the same clock cycle as first bucket chunk (containing rules 10 & 11) and this leads to a bank conflict because both are stored in the same memory bank **2510** b_1. Therefore, the second bucket chunk of bucket B_2, **2504** containing rules 1 and 2 is replicated and stored in memory bank **2510** b_4. This eliminates any possible future memory conflict. However, if there is no potential for a memory conflict, a pointer can point to an already stored rule. Memory replication, if no memory conflict is possible, is not necessary. Therefore, before memory replication, the system determines if storing a pointer would create a memory conflict. The system determines this by determining

whether a bucket chunk of the current bucket is stored in the same bank where the pointer would point to. If so, a memory conflict is possible and replication is needed. If not, replication is not needed and a pointer can point to the rules and save memory.

Further, the search processor manager maintains an instant view of each bank of memory in each cluster by maintaining an allocation list and a free list. The allocation list is a list of all of the addresses in memory that have been allocated, and can be represented by a balanced tree or hash. The free list is a list of all of the addresses in memory that are free and can be represented by a linked list sorted by address and free space length. In one embodiment, each bank of memory is 4096 (4K) rows of 256-bit width. The allocation list begins as an empty list because the memory is empty, while the free list begins by designating the first free memory address of the bank (i.e., zero) and the length of the free space (i.e., the size of the bank (4096×256)). Providing a view of the memory in such a manner also allows the search processor manager to defrag the memory, for example, upon an allocation request to a cluster not having enough room. Alternatively, if defragmentation of the cluster is unsuccessful, the search processor manager can move data to another cluster.

Further, each memory allocation request can include a designation of the type of data to be stored (e.g., tree data, bucket data, or rule data). The search processor manager is aware of the data type of each allocation. The manager therefore knows the type of data stored, the cluster number, and bank number of each allocation. For example, when a child node in the tree is allocated, the child being pointed to by a respective parent node, the search processor manager is aware of the child-parent relationship as well. The search processor manager's employs its awareness of tree relationships during allocation to avoid migrations, remote reads, and bank conflicts.

The search processor manager optimizes the memory based on the following guidelines and principles. The RMEs can process rules faster when there are more rules per chunk. However, in search processors with different numbers of cache lines, the optimal number of rules per chunk may be increased or decreased. Therefore, a bucket packer in the search processor manager organizes the buckets to be optimally sized.

The optimization of memory helps prevent the search processor from performing migrations and bank conflicts. Reading from a same cluster in a same bank can take up to four clock cycles. Migrating work from one cluster to another cluster takes 10 clock cycles. However, reading from the same cluster in a different bank only takes one clock cycle. This is the ideal efficiency, and the search processor places tree nodes, buckets and rules within cluster to make this the most likely scenario. Other implementations use more or fewer clock cycles for reading from a same cluster in a same bank, migrating work from one cluster to another cluster, and reading from the same cluster in a different bank. However, even in other implementations, reading from the same cluster in a different bank is the most efficient operation and should therefore be used most often.

Further, within the memory, the search processor manager determines where to place rules and buckets. For example, consider a bucket containing a "Rule 0," "Rule 30," and "Rule 45." If none of the rules are allocated in memory, the search processor manager finds three contiguous memory slots and stores them all in order. However, suppose "Rule 0" is already stored in memory while "Rule 30" and Rule 45" are not stored. If "Rule 30" and "Rule 45" can be stored after "Rule 0," the search processor manager stores them in those locations. However, if there is no room after "Rule 0" to store "Rule 30" and "Rule 45," the search processor manager copies (e.g., replicates) "Rule 0" to another location where "Rule 30" and "Rule 45" can be stored afterwards. The above is one example of one scheme for determining where to place rules and buckets. Other schemes may be employed to place rules and buckets in memory.

FIG. **21** is a block diagram **2100** illustrating an example embodiment of storing the database in a plurality of clusters **2116**, **2118** and **2120**. A tree **2104** is stored across cluster 0 **2116**, cluster 1 **2118**, and cluster 2 **2120**. Cluster 0 **2116** stores only tree nodes of the tree **2104**. Cluster 0 **2116**, therefore, has the same problem as that shown in FIG. **20**A. With reference to FIG. **21**, cluster 1 **2118** includes tree nodes, including leaf node **2110**, pointing to a bucket **2112**, which subsequently points to rules **2114**. Cluster 1 **2118** includes tree nodes, bucket nodes and rules. Therefore, the tree walk engines can pass work to the bucket walk engines once they pass the tree walk engine boundary **2122**, and the bucket walk engines can pass work to the rule match engines once a pass the bucket walk engine boundary **2124**, all within the same cluster (cluster 1 **2118**). This is the ideal setup for a cluster.

Cluster 2 **2120** includes subtree **2106** with a leaf node **2108** pointing to the bucket **2112**. The subtree **2106**, for example, can be added to the database as an incremental update. The cluster 2 **2120** includes only tree nodes, and therefore has the same problem as cluster 0 **2116**.

FIG. **22**A is a block diagram **2200** illustrating an example embodiment of two clusters **2202** and **2204** receiving work in generating a response **2218**. Cluster 0 **2202** receives an input **2220** at a tree walk engine **2206**. Cluster 0 **2202** also has a bucket walk engine **2208** and a rule match engine **2210**. The bucket walk engine **2208** and rule match engine **2210** go underutilized because cluster 0 **2202** only includes tree nodes and can only utilize the tree walk engine **2206**. The tree walk engine **2206** passes the output of its work to the cluster 1 **2204** via a migration. The cluster 1 **2204** tree walk engine **2212** receives the output of the tree walk engine **2206** and finishes walking the tree. The tree walk engine **2212** then passes its output to a bucket walk engine **2214** which walks the buckets to find a rule, which is passed to a rule match engine **2216**, which outputs a response **2218**. Even though a response is generated **2218**, each clock cycle for the bucket walk engine **2208** and rule match engine **2210** go unused. The search processor manager is designed to make sure that all of the engines, or at least an optimal amount of engines, are being used every clock cycle.

FIG. **22**B is a block diagram **2250** illustrating a more optimal setup for two clusters **2252** and **2254**. Cluster 0 **2252** receives an input **2270** at its tree walk engine **2256**. While the tree walk engine **2256** has to migrate to the cluster 1 **2254**, it passes work to the bucket walk engine **2264** after finishing walking the tree. Although migrations are not optimal, the cluster 0 **2252** also includes buckets and rules which it can do work on while the tree is being walked through. Cluster 1 **2254** then processes the bucket at its bucket walk engine **2264** and then processes the rules at the rule match engine **2266** to produce a response **2268**. In parallel, cluster 1 also receives an input **2272** at its tree walk engine **2262**. The tree walk engine **2262** migrates the work to the bucket walk engine **2258** of cluster 0 after walking through its tree. The bucket walk engine **2258** walks to the bucket, and passes a rule or chunk of rules to the rule match engine **2260**. The rule match engine then issues a response **2274**. Both cluster 0 **2252** and cluster 1 **2254** utilize their tree

walk engines, bucket walk engines and rule match engines. Further, new work can be pipelined through the clusters, such that when the tree walk engines finish their work, new work comes in and the tree work engines stay utilized.

FIG. **23**A is a block diagram **2300** illustrating an example embodiment of migration from cluster to cluster. Cluster 0 **2302** receives new work **2310** and performs a migration ($M_{0,1}$) **2312** to cluster 1 **2304**. Cluster 1 **2304** then migrates work to cluster 2 **2306** by performing migration ($M_{1,2}$) **2314**. Then clustered 2 performs a migration **2316** ($M_{2,3}$) to cluster 3 **2308**. Cluster 3 then issues a response **2318**. While migration is not ideal, the migration shown in FIG. **23**A does not include any loops and is therefore acceptable. The job of the search processor manager is to allocate the memory such that loops between clusters do not form during migration.

FIG. **23**B is a block diagram **2350** illustrating an example embodiment of a loop formed during migration. Cluster 0 **2352** initially migrates work via migration ($M_{0,1}$) **2356** to cluster 1 **2354**. After performing work, cluster 1 **2254** then migrates work via migration ($M_{1,0}$) **2358** back to cluster 0 **2352**. In this way, a loop is formed between cluster 0 **2352** and cluster 1 **2354**. The job of the search processor manager is to avoid such a loop by configuring the memory in such a way that the loop never forms. This can be performed by confirming that a pointer from one cluster to another cluster does not point to any other nodes in the tree that eventually point back to the original cluster.

FIG. **24** is a diagram **2400** illustrating an example embodiment of writing bucket chunks **2402** to memory banks **2406** using access ports **2404**. After a write request to write bucket chunks **2402** ($C_{1\text{-}8}$) to memory banks **2406** ($B_{1\text{-}8}$), the system determines a distribution of bucket chunks **2402** over the memory banks **2406** to eliminate memory conflicts over the access ports **2404** (Ports **1-3**). Most memory accesses request the buckets **2402** in the order in which they are addressed in the database, such that they are accessed according to the order in which they are stored/addressed. For example, a memory request to access the buckets accesses C_1, C_2, C_3, then C_4, and so on. Therefore, ideally any group of consecutive bucket chunks of at most the number of access ports **2404** can be accessed during one clock cycle. Each respective access port **2404** cannot access the same bank twice during one clock cycle, so the bucket chunks **2402** should be distributed in the memory banks **2406** carefully.

FIG. **24** illustrates an example embodiment of distributing the bucket chunks **2402** in the memory banks **2406**. In this example, the bucket chunks **2402** are distributed among memory banks **2406** $B_{1\text{-}3}$. Bucket chunks C_1, C_4, and C_7 are stored in memory bank B_1, bucket chunks C_2, C_5, and C_8 are stored in memory bank B_2, and bucket chunks C_3 and C_6 are stored in memory bank B_3. This distribution scheme is a round-robin distribution scheme, but other distribution schemes can be employed, so long as no bucket chunks **2402** that can be accessed in the same clock cycle based on the number of ports (e.g., bucket chunks that are separated by a number of chunks that is less than the number of ports) are stored in the same memory bank **2406**. The bucket chunks **2402** are distributed such that the three access ports **2404** can access any three consecutive bucket chunks **2402** in one clock cycle. For example, accessing C_1, C_2, and C_3 can be performed on one clock cycle by utilizing all three ports **2404**, each respective port accessing a respective memory bank B_1, B_2, and B_3, where one respective chunk (C_1, C_2, or C_3) are stored. Likewise, accessing C_5, C_6 and C_7 can be performed on one cycle by utilizing all three ports **2404**, each respective port accessing a respective memory bank B_2, B_3, and B_1. Other distributions of the bucket chunks **2402** are possible. For example, the bucket chunks **2402**, in a system with three ports **2404** can be distributed over three, four, or any number of banks, as long as consecutively addressed chunks are assigned respectively unique memory banks, such that every memory access is able to access consecutively addressed bucket chunks **2402** stored in the memory banks **2406**. Different numbers of bucket chunks **2402**, ports **2404**, and memory banks **2406** can be employed. For example, the number of available access ports **2404** dictates the number of consecutively addressed bucket chunks that can be accessed during one clock cycle. For example, with four access ports **2404**, bucket chunks **2402** are distributed over at least four memory banks **2406**.

Embodiments may be implemented in hardware, firmware, software, or any combination thereof. It should be understood that the block diagrams may include more or fewer elements, be arranged differently, or be represented differently. It should be understood that implementation may dictate the block and flow diagrams and the number of block and flow diagrams illustrating the execution of embodiments of the invention.

The teachings of all patents, published applications and references cited herein are incorporated by reference in their entirety.

While this invention has been particularly shown and described with references to example embodiments thereof, it will be understood by those skilled in the art that various changes in form and details may be made therein without departing from the scope of the invention encompassed by the appended claims.

What is claimed is:

1. A method of managing a database including a tree, a plurality of buckets, and a plurality of rules, the method comprising:

managing a memory with a plurality of cluster memories, the managing including storing the database across the plurality of cluster memories of the memory and designating a given cluster memory of the plurality of cluster memories as a sink memory cluster; and

improving performance, of a search processor configured to walk the tree, by packing nodes of the tree in each of the plurality of cluster memories, such that walking the tree by the search processor accesses a minimal amount of cluster memories in the memory and walking the tree by the search processor accesses each particular cluster memory no more than once by configuring the memory in such a way that prevents a migration loop from forming, the configuring including embedding one or more migration instructions in the memory that cause a migrated processing thread of the search processor to migrate to the sink memory cluster, the sink memory cluster configured to end migration of the migrated processing thread.

2. The method of claim **1**, further comprising packing a first particular number of bucket chunks per bucket and a second particular number of rule pointers per bucket chunk based on addresses of the rules.

3. The method of claim **1**, further comprising allocating the rules in the memory in a same order as an order of the rules in bucket chunks of the buckets.

4. The method of claim **1**, further comprising:

replicating a rule or a chunk of rules across a first and second bank in a particular cluster memory such that the rule or chunk of rules can be accessed on the second bank when the first bank has a memory access conflict during a particular clock cycle.

5. The method of claim 1, further comprising:
distributing rules and buckets across a first cluster memory and a second cluster memory within the memory; and
enabling processing of the rules by a rule match engine within the first cluster memory and a rule match engine within the second cluster memory.
6. The method of claim 1, further comprising:
storing a first rule of a particular bucket of the plurality of buckets in a particular cluster memory of the plurality of cluster memories; and
storing any other rules of the particular bucket in the particular cluster memory.
7. The method of claim 1, further comprising:
allocating a node of the tree in a particular cluster memory of the memory; and
allocating a bucket in the memory that the node of the tree points to in the particular cluster memory.
8. The method of claim 1, further comprising:
allocating a bucket in a particular cluster memory of the memory; and
allocating a rule associated with the bucket in the particular cluster memory storing the bucket.
9. The method of claim 1, further comprising:
allocating the rules in chunks according to an order of the buckets;
determining a need to replicate the rules across the plurality of cluster memories; and
replicating the rules across the plurality of cluster memories, if necessary.
10. The method of claim 1, further comprising:
determining at least one division of the database, the database including the tree, the plurality of buckets, and the plurality of rules, the division based on either a horizontal division, the horizontal division separating the tree based on a depth of data of the tree or a vertical division, the vertical division separating the tree based on sub-trees of the tree; and
generating at least one memory request to store each division of the database in a respective cluster memory.
11. A system for managing a database including a tree, a plurality of buckets, and a plurality of rules, the system comprising:
a search processor configured to walk the tree;
a memory with a plurality of cluster memories, the memory configured to store the database across the plurality of cluster memories, a given cluster memory of the plurality of memories designated as a sink memory; and
a tree packing module configured to improve performance of the search processor by packing nodes of the tree in each of the plurality of cluster memories such that walking the tree by the search processor accesses a minimal amount of cluster memories in the memory and walking the tree by the search processor accesses each particular cluster memory no more than once by configuring the memory in such a way that prevents a migration loop from forming, wherein configuring the memory includes embedding one or more migration instructions in the memory that cause a migrated processing thread of the search processor to migrate to the sink memory cluster, the sink memory cluster configured to end migration of the migrated processing thread.
12. The system of claim 11, further comprising a bucket packing module configured to pack a first particular number of bucket chunks per bucket and a second particular number of rule pointers per bucket chunk based on addresses of the rules.
13. The system of claim 11, wherein a bucket packing module is configured to allocate the rules in the memory in a same order as an order of the rules in bucket chunks of the buckets.
14. The system of claim 11, further comprising:
a bucket packing module configured to replicate a rule or a chunk of rules across a first and second bank in a particular cluster memory such that the rule or chunk of rules can be accessed on the second bank when the first bank has a memory access conflict during a particular clock cycle.
15. The system of claim 11, further comprising a bucket packing module configured to distribute rules and buckets across a first cluster memory and a second cluster memory within the memory and enable processing of the rules by a rule match engine within the first cluster memory and a rule match engine within the second cluster memory.
16. The system of claim 11, further comprising a bucket packing module configured to store a first rule of a particular bucket of the plurality of buckets in a particular cluster memory of the plurality of cluster memories and store any other rules of the particular bucket in the particular cluster memory.
17. The system of claim 11, further comprising a bucket packing module configured to allocate a node of the tree in a particular cluster memory of the memory and allocate a bucket in the memory that the node of the tree points to in the particular cluster memory.
18. The system of claim 11, further comprising a bucket packing module configured to allocate a bucket in a particular cluster memory of the memory, and allocate a rule associated with the bucket in the particular cluster memory storing the bucket.
19. The system of claim 11, further comprising a bucket packing module configured to allocate the rules in chunks according to an order of the buckets, determine a need to replicate the rules across the plurality of cluster memories, replicate the rules across the plurality of cluster memories, if necessary.
20. The system of claim 11, wherein the tree packing module is further configured to determine at least one division of the database, the database including the tree, the plurality of buckets, and the plurality of rules, the division based on either a horizontal division, the horizontal division separating the tree based on a depth of data of the tree or a vertical division, the vertical division separating the tree based on sub-trees of the tree, and generate at least one memory request to store each division of the database in a respective cluster memory.
21. A non-transitory computer-readable medium configured to store instructions for managing a database including a tree, a plurality of buckets, and a plurality of rules, the instructions, when loaded and executed by a control plane processor, causes the control plane processor to:
manage a memory with a plurality of cluster memories used by a search processor to walk the tree the manage operation including storing the database across the plurality of cluster memories of the memory and designating a given cluster memory of the plurality of cluster memories as a sink cluster memory; and
improve performance of the search processor by packing nodes of the tree in each of the plurality of cluster memories such that walking the tree by the search processor accesses a minimal amount of cluster memo-

ries in the memory and walking the tree by the search processor accesses each particular cluster memory no more than once by configuring the memory in such a way that prevents a migration loop from forming, wherein configuring the memory includes embedding one or more migration instructions in the memory that cause a migrated processing thread of the search processor to migrate to the sink memory cluster, the sink memory cluster configured to end migration of the migrated processing thread.

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